

Assignment No. 01

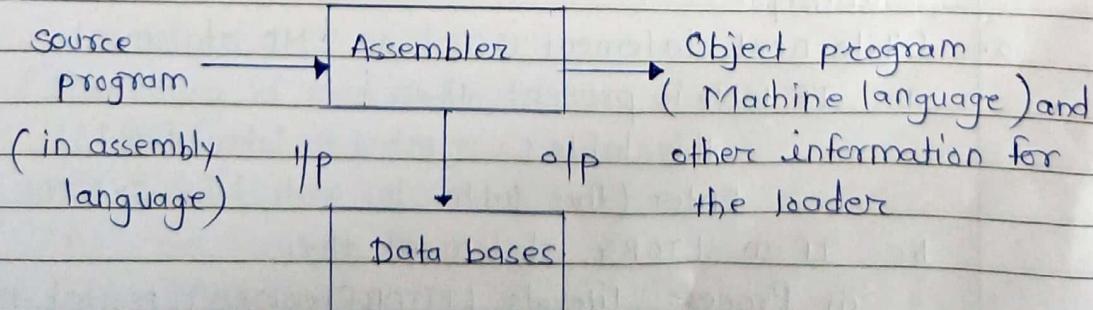
Title: - To design a Pass-1 of two Pass assembler

Problem Statement: - Design suitable data structures and implement pass-1 of two-pass assembler for pseudo-machine in Java using object oriented feature.

Implementation should consist of a few instructions from each category and few assembler directives.

Objective: - Design suitable data structures and develop a subset of an assembler. Subset should consist of a few instructions from each category and few assembler directives.

Theory: - An assembler is a computer program for translating assembly language essentially, a mnemonic representation of machine language - into object code. To make writing and reading programs written in machine language more convenient, mnemonic (symbol) were used for each machine instruction. These mnemonics were then later translated into machine language. Such a mnemonic language is called assembly language. Assemblers are programs to automatically translate assembly language into machine language.



Other than producing the machine language, the assembler must also produce the information for the loader to use eg. externally defined symbols must be noted and passed to the loader.

Data structures:-

Pass T data bases;

1. Input source program
2. A Location Counter (LC), used to keep track of each instruction's location
3. A machine operation table (MOT), that indicates the symbolic mnemonic for each and its length.
4. A pseudo-operation table (POT), that indicates the symbolic mnemonic and action to be taken for each pseudo-op in pass I
5. A symbol table (ST), that is used to store each literal encountered and its corresponding value.
6. A literal table (LT) that is used to store each literal encountered that its corresponding assigned location.
7. A copy of the input, to be passed to pass 2. This is to be stored in storage device.

Algorithm of Pass I:

1. loc_ctr := 0; (default value)
 pooltab_ptr := 1; POOLTAB[1] := 1;
 littab_ptr := 1;
2. While next statement is not an END statement
 - a. If label is present then
 this_label := symbol in label field;
 Enter (this_label, loc_ctr) in SYMTAB
 - b. If an LTORG statement then
 - (i) Process literals LTTAB[POOLTAB[pooltab_ptr]:LTTAB[littab_ptr-1]] to allocate memory and put the address in the address field. Update loc_ctr accordingly.

- (ii) $\text{Pooltab_ptr} := \text{pooltab_ptr} + 1;$
 - (iii) $\text{POOLTAB}[\text{pooltab_ptr}] := \text{littab_ptr};$
 - c. If a START or ORIGIN statement then
 $\text{loc_cntr} := \text{value specified in operand field}$
 - d. If an EQU statement then
 - (i) $\text{this_addr} := \text{value of <address spec>};$
 - (ii) correct the symtab entry for this-table to $(\text{this_table}, \text{this_addr})$
 - e. If a declaration statement then
 - (i) $\text{code} := \text{code of the declaration statement};$
 - (ii) $\text{size} := \text{size of memory area required by DC/DS}.$
 - (iii) $\text{loc_cntr} = \text{loc_cntr} + \text{size};$
 - (iv) Generate $\text{IC}^e(\text{BL}, \text{code}) \dots$.
 - f. If an imperative statement then
 - (i) $\text{code} := \text{machine opcode from OPTAB};$
 - (ii) $\text{loc_cntr} := \text{loc_cntr} + \text{instruction length from OPTAB};$
 - (iii) If operand is a literal then
 $\text{this_literal} := \text{literal in operand field};$
 $\text{LITTAB}[\text{littab_ptr}] := \text{this_literal};$
 $\text{littab_ptr} := \text{littab_ptr} + 1;$
 Else (i.e. operand is a symbol)
 $\text{this_entry} := \text{SYMTAB entry number of operand};$
 Generate $\text{IC}^e(\text{IS}, \text{code}) (\text{S}, \text{this_entry})$;
3. (Processing of END statement)
- (a) Perform step 2(b)
 - (b) Generate $\text{IC}'(\text{AP}, \text{o2})$.
 - (c) Go to Pass II.

Conclusion:- Hence we studied To design a Pass-I of two pass assembler.

PROGRAM:

```

}
}
else
{
lc++;
}
}

System.out.print("**SYMBOL TABLE*\n"); System.out.println("_____");
for(i=0;i<5;i++)
{
for(cnt=0;cnt<2;cnt++)
{
System.out.print(st[i][cnt]+"\t");
}
System.out.println();
}

String
inst[]={ "STOP","ADD","SUB","MULT","MOVER","MOVEM","COMP","BC","DIV","READ",
",PRINT"};
String reg[]={ "NULL","AREG","BREG","CREG","DREG"};
int op[][]=new int[12][3];
int j,k,p=1,cnt1=0;
for(i=1;i<11;i++)
{
for(j=0;j<11;j++)
{
if(a[i][1].equalsIgnoreCase(inst[j]))
{
op[cnt1][0]=j;
}
else
if(a[i][1].equalsIgnoreCase("DS"))
{
p=Integer.parseInt(a[i][2]);
}
else if(a[i][1].equalsIgnoreCase("DC"))
{
op[cnt1][2]=Integer.parseInt(a[i][2]);
}
}
}

```

```
        }
    }
    for(k=0;k<5;k++)
    {
        if(a[i][2].equalsIgnoreCase(reg[k]))
        {
            op[cnt1][1]=k;
        }
    }
    for(l=0;l<5;l++)
    {
        if(a[i][3].equalsIgnoreCase(st[l][0]))
        {
            int mn=Integer.parseInt(st[l][1]);
            op[cnt1][2]=mn;
        }
    }
    cnt1=cnt1+p;
} System.out.println("\n **OUTPUT*\n"); System.out.println("****MOT TABLE***");
int dlc=Integer.parseInt(a[0][2]);
for(i=0;i<12;i++)
{
    System.out.print(dlc+++"\t");
    for(j=0;j<3;j++)
    {
        System.out.print(" "+op[i][j]+" ");
    }
    System.out.println();
} System.out.println("");
}
```

OUTPUT

```
Output  
java -cp /tmp/ENh950LMbI P1  
**SYMBOL TABLE*  
_____  
AGAIN 102 N 106  
RESULT 108  
ONE 110  
TERM 111  
**OUTPUT*  
  
****MOT TABLE***  
101 4 2 110  
102 3 2 111 103 4 3 111  
104 1 3 106  
105 5 3 111  
106 0 0 0  
107 0 0 0  
108 0 0 0  
109 0 0 0  
110 0 0 1  
111 0 0 0  
112 0 0 0
```