

### Log Based Conquency Control

- In log based conquency control conflicting actions of differnt transctions in which the logs are opting and the log protocol extends this ordering on actions to transctions there by ensuring serializibility
- In optimistic control a time stamp ordering is imposed on transctions and validation checks that all conflicting actions occured in the same order timestamps can also be used in another way i.e each transction can be assigned a timestamp at startup and we can ensure at execution time if action  $a_i$  of transction  $T_i$  conflicts with action  $a_j$  of transction  $T_j$ ,  $a_i$  occurs before  $a_j$  if  $TS(T_i) < TS(T_j)$  if an action violates then the transction is aborted and restarted

### Multi Version Conquency Control

- This protocol represents yet another of using timestamps assigned at startup time to achieve serializibility the goal is to ensure that the transction never has to wait to read a database object and the idea is to maintain sereval versions of each database object each with a write timestamp and let transction  $T_i$  read the most recent version whose timestamp preceds  $TS(T_i)$
- If transction  $T_i$  wants to write an object we must ensure that the object has not already been read by some other transction  $T_j$  such that  $TS(T_i) < TS(T_j)$  if the allowed  $T_i$  to write such an object it's change should be seen by  $T_j$  for serializibility but obously  $T_j$  which read the object at sometime in the past will not see  $T_i$ 's change
- The drawbacks of this scheme are similar to those of timestamp conquency control and in addition their is the cost of maintaining versions on the other hand reads are never blocked which can be important for workloads dominatedd by transctions that only read values from the database