# Notes on Set Theory

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### 1 Foreword

Notes for the entrance examination

### 2 Models of Set - Sertraline

## 2.1 Some mathematical logic

**Theorem 2.1** (Gödels second incompleteness theorem). If a consistent recursive axiom set T contains **ZFC**, then

$$T \not\vdash \operatorname{Con}(t)$$

especially, **ZFC**  $\not\vdash$  Con(**ZFC**)

**Definition 2.2.** Suppose  $(M, E_M)$  and  $(N, E_N)$  are two models of set theory, then

- 1. if for any formula  $\sigma$ ,  $M \models \sigma$  if and only if  $N \models \sigma$ , then M and N are **elementary equivalent**, denoted by  $M \equiv N$
- 2. If bijection  $f: M \to N$  satisfies: for any  $a, b \in M$ ,  $aE_Mb$  iff  $f(a)E_Nf(b)$ , then  $f: M \cong N$  is an **isomorphism**
- 3. If  $M \subseteq N$  and  $E_M = E_N \upharpoonright M$ , then M is N's submodel
- 4. If M is isomorphic to a submodel of N by injection f, and for any formula  $\varphi(x_1,\ldots,x_n)$ , for any  $a_1,\ldots,a_n\in M$ ,  $M\models\varphi[a_1,\ldots,a_n]$  iff  $N\models\varphi[f(a_1),\ldots,f(a_n)]$ , then f is called an **elementary embedding** from M to N, written as  $f:M\prec N$
- 5. If  $M \subseteq N$  and  $M \prec N$ , then M is a **elementary submodel** of N

**Lemma 2.3.** Suppose  $N \models \mathbf{ZFC}, M \subseteq N$ , then  $M \prec N$  iff  $\forall \varphi(x, x_1, \dots, x_n)$ ,  $\forall (a_1, \dots, a_n) \in M$ , if  $\exists a \in N \text{ s.t. } N \models \varphi[a, a_1, \dots, a_n]$ , then  $\exists a \in M \text{ s.t. } M \models \varphi[a, a_1, \dots, a_n]$ 

**Definition 2.4.** Suppose  $(M, E) \models \mathbf{ZFC}$ 

- 1.  $h_{\varphi}: M^n \to M$  is  $\varphi$ 's **Skolem function** if  $\forall a_1, \ldots, a_n \in M$ , if  $\exists a \in M$  s.t.  $M \models \varphi[a, a_1, \ldots, a_n]$ , then  $M \models \varphi[h_{\varphi}(a_1, \ldots, a_n), a_1, \ldots, a_n]$ . requires **AC**
- 2. Let  $\mathcal{H} = \{h_{\varphi} \mid \varphi \text{ is a formula on set theory}\}$ . For any  $S \subseteq M$ , **Skolem** hull  $\mathcal{H}(S)$  is the smallest set consisting of S and closed under  $\mathcal{H}$

**Lemma 2.5.**  $N \models \mathbf{ZFC}, S \subseteq N$ , if  $M = \mathcal{H}(S)$ , then  $M \prec N$ 

**Theorem 2.6** (Löwenheim-Skolem theorem). Suppose  $N \models \mathbf{ZFC}$  and is infinite, then there is a model M s.t.  $|M| = \omega$  and  $M \prec N$ 

## 2.2 Cumulative Hierarchy

This section works in  $\mathbf{ZF}$  (a.k.a.  $\mathbf{ZF}$  – axiom of foundation)

**Definition 2.7.** For any  $\alpha$ , define sequence  $V_{\alpha}$ 

- 1.  $V_0 = \emptyset$
- 2.  $V_{\alpha+1} = \mathcal{P}(V_{\alpha})$
- 3. For any limit ordinal  $\lambda$ ,  $V_{\lambda} = \bigcup_{\beta < \lambda} V_{\beta}$

And WF = 
$$\bigcup_{\alpha \in \mathbf{On}} V_{\alpha}$$

**Lemma 2.8.** For any ordinal  $\alpha$ 

- 1.  $V_{\alpha}$  is transitive
- 2. if  $\xi \leq \alpha$ , then  $V_{\xi} \subseteq V_{\alpha}$
- 3. if  $\kappa$  is inaccessible cardinal, then  $|V_{\kappa}|=\kappa$

*Proof.* 1. Obviously  $\kappa \leq V_{\kappa}$ . Since  $\kappa$  is inaccessible, then for any  $\alpha < \kappa$ ,  $|V_{\alpha}| < \kappa$ .

**Definition 2.9.** For any set  $x \in WF$ ,

$$\operatorname{rank}(x) = \min\{\beta \mid x \in V_{\beta+1}\}$$

**Lemma 2.10.** 1.  $V_{\alpha} = \{x \in \mathbf{WF} \mid \text{rank}(x) < \alpha\}$ 

- 2. **WF** is transitive
- 3. For any  $x, y \in \mathbf{WF}$  , if  $x \in y$ , then  $\mathrm{rank}(x) < \mathrm{rank}(y)$
- 4. for any  $y \in \mathbf{WF}$  ,  $\mathrm{rank}(y) = \sup\{\mathrm{rank}(x) + 1 \mid x \in y\}$

**Lemma 2.11.** Supoose  $\alpha$  is an ordinal

1.  $\alpha \in \mathbf{WF}$  and  $\mathrm{rank}(\alpha) = \alpha$ 

2. 
$$V_{\alpha} \cap \mathbf{On} = \alpha$$

**Lemma 2.12.** 1. If  $x \in \mathbf{WF}$ , then  $\bigcup x, \mathcal{P}(x), \{x\} \in \mathbf{WF}$ , and their ranks are all less than  $\mathrm{rank}(x) + \omega$ 

- 2. If  $x,y\in \mathbf{WF}$  , then  $x\times y, x\cup y, x\cap y, \{x,y\}, (x,y), x^y\in \mathbf{WF}$  , and their ranks are all less than  $\mathrm{rank}(x)+\mathrm{rank}(y)+\omega$
- 3.  $\mathbb{Z}, \mathbb{Q}, \mathbb{R} \in V_{\omega+\omega}$
- 4. for any set x,  $x \in \mathbf{WF}$  iff  $x \subset \mathbf{WF}$

## Lemma 2.13. Suppose AC

- 1. for any group G, there exists group  $G' \cong G$  in **WF**
- 2. for any topological space T, there exists  $T' \cong T$  in **WF**

**Definition 2.14.** Binary relation < on set A is **well-founded** if for any nonempty  $X \subseteq A$ , X has minimal element under <

**Theorem 2.15.** If  $A \in \mathbf{WF}$  , then  $\in$  is a well-founded relation on A

**Lemma 2.16.** If set *A* is transitive and  $\in$  is well-founded on *A*, then  $A \in \mathbf{WF}$ 

**Lemma 2.17.** For any set x, there is a smallest transitive set  $\operatorname{trcl}(x)$  s.t.  $x \subseteq \operatorname{trcl}(x)$ 

Proof.

$$x_0 = x$$

$$x_{n+1} = \bigcup_{n < \omega} x_n$$

$$\operatorname{trcl}(x) = \bigcup_{n < \omega} x_n$$

trcl(x) is called **transitive closure** of x

Lemma 2.18. Without axiom of power set

- 1. if x is transitive, then trcl(x) = x
- 2. if  $y \in x$ , then  $trcl(y) \subseteq trcl(x)$
- 3.  $\operatorname{trcl}(x) = x \cup \bigcup \{\operatorname{trcl}(y) \mid y \in x\}$

**Theorem 2.19.** For any set X, the following are equivalent

- 1.  $X \in \mathbf{WF}$
- 2.  $\operatorname{trcl}(X) \in \mathbf{WF}$
- 3.  $\in$  is a well-founded relation on trcl(X)

**Theorem 2.20.** The following propositions are equivalent

- 1. Axiom of foundation
- 2. For any set X,  $\in$  is a well-founded relation on X
- 3. V = WF

#### 2.3 Relativization

**Definition 2.21.** Let **M** be a class  $\varphi$  a formula, the **relativization** of  $\varphi$  to **M** is  $\varphi^{\mathbf{M}}$  defined inductively

$$(x \in y)^{\mathbf{M}} \leftrightarrow x = y$$
$$(x \in y)^{\mathbf{M}} \leftrightarrow x \in y$$
$$(\varphi \to \psi)^{\mathbf{M}} \leftrightarrow \varphi^{\mathbf{M}} \to \psi^{\mathbf{M}}$$
$$(\neg \varphi)^{\mathbf{M}} \leftrightarrow \neg \varphi^{\mathbf{M}}$$
$$(\forall x \varphi)^{\mathbf{M}} \leftrightarrow (\forall x \in \mathbf{M}) \varphi^{\mathbf{M}}$$

Note  $\varphi^{\mathbf{V}} = \varphi$  and

$$f^{\mathbf{M}} = \{(x_1, \dots, x_n, x_{n+1}) \in \mathbf{M} \mid \varphi^{\mathbf{M}}(x_1, \dots, x_n, x_{n+1})\}$$

**Definition 2.22.** For any theory T, any class  ${\bf M}$  ,  ${\bf M} \models T$  iff for any axiom  $\varphi$  of T,  $\varphi^{\bf M}$  holds

Theorem 2.23 (ZF). WF  $\models$  ZF

*Proof.* • Axiom of existence

 $(\exists x(x=x))^{\mathbf{M}} \leftrightarrow \exists x \in \mathbf{M} \ (x=x)$ , which is equivalent to  $\mathbf{M}$  being nonempty

• Axiom of extensionality

$$\forall X \forall Y \forall u ((u \in X \leftrightarrow u \in Y) \to X = Y)^{\mathbf{M}} \Leftrightarrow$$
$$\forall X \in \mathbf{M} \ \forall Y \in \mathbf{M} \ \forall u \in \mathbf{M} \ ((u \in X \leftrightarrow u \in Y) \to X = Y)$$

**Lemma 2.24.** If  $\mathbf{M}$  is transitive, then axiom of extensionality holds in  $\mathbf{M}$ 

• Axiom schema of specification

$$\forall X \in \mathbf{M} \exists Y \in \mathbf{M} \ \forall u \in \mathbf{M} \ (u \in Y \leftrightarrow u \in X \land \varphi^{\mathbf{M}} (u))$$

Since for any  $X \in \mathbf{WF}$  ,  $\mathcal{P}(X) \subseteq \mathbf{WF}$ 

- Axiom of paring
- Axiom of union
- Axiom of power set

$$\forall X \in \mathbf{M} \ \exists Y \in \mathbf{M} \ \forall u \in \mathbf{M} \ (u \in Y \leftrightarrow (u \subseteq X)^{\mathbf{M}})$$

and

$$(u \subseteq X)^{\mathbf{M}} \leftrightarrow \forall x \in \mathbf{M} \ (x \in u \to x \in X) \leftrightarrow u \cap \mathbf{M} \subseteq X$$

- Axiom of foundation
- Axiom schema of replacement

#### 2.4 Absoluteness

**Definition 2.25.** For any formula  $\psi(x_1,\ldots,x_n)$  and any class  ${\bf M}$  ,  ${\bf N}$  ,  ${\bf M}\subseteq {\bf N}$  , if

$$\forall x_1 \dots \forall x_n \in \mathbf{M} \left( \psi^{\mathbf{M}} \left( x_1, \dots, x_n \right) \leftrightarrow \psi^{\mathbf{N}} \left( x_1, \dots, x_n \right) \right)$$

then  $\psi(x_1,\dots,x_n)$  is absolute for  ${\bf M}$  ,cn. If  ${\bf N}={\bf V}$  , then  $\psi$  is absolute for  ${\bf M}$ 

**Lemma 2.26.** Suppose  $\mathbf{M} \subseteq \mathbf{N}$  and  $\varphi, \psi$  are formulas, then

- 1. if  $\varphi$ , $\psi$  are absolute for **M** ,cn, then so are  $\neg \varphi$ ,  $\varphi \rightarrow \psi$
- 2. if  $\varphi$  doesn't contain any quantifiers, then  $\varphi$  is absolute for any **M**
- 3. if **M** ,**N** are transitive and  $\varphi$  is absolute for them, then so are  $\forall x \in y\varphi$

**Definition 2.27.**  $\Delta_0$  formula

- 1.  $x = y, x \in y$  are  $\Delta_0$  formulas
- 2. if  $\varphi$ , $\psi$  are  $\Delta_0$ , then so are  $\neg \varphi$ ,  $\varphi \rightarrow \psi$
- 3. if  $\varphi$  is  $\Delta_0$ , y is any set, then  $(\forall x \in y)\varphi$  is  $\Delta_0$

If  $\varphi$  is  $\Delta_0$ , then  $\exists x_1 \dots \exists x_n \varphi$  is  $\Sigma_1$  formula,  $\forall x_1 \dots \forall x_n \varphi$  is  $\Pi_1$ 

**Lemma 2.28.** M  $\subseteq$  N are both transitive,  $\psi(x_0,\ldots,x_n)$  is a formula, then

- 1. if  $\psi$  is  $\Delta_0$ , then it's absolute for **M**, cn
- 2. if  $\psi$  is  $\Sigma_1$ , then

$$\forall x_1 \dots x_n (\psi^{\mathbf{M}} (x_1, \dots, x_n) \to \psi^{\mathbf{N}} (x_1, \dots, x_n))$$

3. if  $\psi$  is  $\Pi_1$ , then

$$\forall x_1 \dots x_n (\psi^{\mathbf{N}}(x_1, \dots, x_n) \to \psi^{\mathbf{M}}(x_1, \dots, x_n))$$

**Lemma 2.29.** If  $\mathbf{M} \subseteq \mathbf{N}$ ,  $\mathbf{M} \models \Sigma$ ,  $\mathbf{N} \models \Sigma$  and

$$\Sigma \vdash \forall x_1 \dots \forall x_n (\varphi(x_1, \dots, x_n) \leftrightarrow \psi(x_1, \dots, x_n))$$

then  $\varphi$  is absolute for **M**, **N** if and only if  $\psi$  is absolute for **M**, **N** 

**Definition 2.30.** Suppose  $\mathbf{M} \subseteq \mathbf{N}$ ,  $f(x_1, \dots, x_n)$  is a function. f is **absolute** for  $\mathbf{M}$  and  $\mathbf{N}$  if and only if  $\varphi(x_1, \dots, x_n, x_{n+1})$  defining f is absolute.

**Theorem 2.31.** Following relations and functions can be defined in  $\mathbf{ZF}^-$  Pow – Inf and are equivalent to some  $\Delta_0$  formulas. So they are absolute for any transitive model  $\mathbf{M}$  on  $\mathbf{ZF}^-$  – Pow – Inf

- 1.  $x \in y$
- 2. x = y
- 3.  $x \subset y$
- 4.  $\{x, y\}$
- 5. {*x*}
- 6. (x, y)
- 7. Ø

- 8.  $x \cup y$
- 9. x y
- 10.  $x \cap y$
- 11.  $x^+$
- 12. x is a transitive set
- 13. LJ*x*
- 14.  $\bigcap x (\bigcap \emptyset = \emptyset)$

Lemma 2.32. Absoluteness is closed under operation composition

**Theorem 2.33.** Following relations and functions are absolute for any transitive model M on  $\mathbf{Z}\mathbf{F}^-$  – Pow – Inf

- 1. z is an ordered pair
- 2.  $A \times B$
- 3. R is a relation
- 4. dom(R)
- 5. ran(R)
- 6. *f* is a function
- 7. f(x)
- 8. *f* is injective

### 2.5 Relative consistence of the axiom of foundation

**Lemma 2.34.** Suppose transitive class  $\mathbf{M} \models \mathbf{Z}\mathbf{F}^- - \mathrm{Pow} - \mathrm{inf}$  and  $\omega \in \mathbf{M}$ , then the axiom of infinity is true in  $\mathbf{M}$ . Hence the axiom of infinity is true in  $\mathbf{W}\mathbf{F}$ 

**Theorem 2.35.** Let T be a theory of set theory language and  $\Sigma$  a set of sentences. Suppose  $\mathbf{M}$  is a class and  $T \vdash \mathbf{M} \neq \emptyset$ , then if  $\mathbf{M} \models_T \Sigma$ , then

- 1. for any sentences  $\varphi$ , if  $\Sigma \vdash \varphi$ , then  $T \vdash \varphi^{\mathbf{M}}$
- 2. if *T* is consistent, then so is  $Cn(\Sigma)$

**Theorem 2.36.** The axiom of foundation is consistent with **ZF**.

*Proof.* By 2.35, let T be **ZF**, 
$$\Sigma$$
 be **ZF** and **M** be **WF**

**Lemma 2.37 (ZF**<sup>-</sup>). Suppose transitive model  $\mathbf{M} \models \mathbf{ZF}^- - \text{Pow} - \text{Inf.}$  If  $X, R \in \mathbf{M}$  and R is a well-order on X, then

$$(R \text{ is a well-order on } X)^{\mathbf{M}}$$

**Theorem 2.38 (ZF**<sup>-</sup>). 
$$V_{\omega} \models \mathbf{ZFC} - \mathbf{Inf} + \neg \mathbf{Inf}$$

*Proof.* For any  $X \in V_{\omega}$ , X is finite hence there is a well-ordering on X

Corollary 2.39. 
$$Con(\mathbf{Z}\mathbf{F}^-) \to Con(\mathbf{Z}\mathbf{F}\mathbf{C} - Inf + \neg Inf)$$

### 2.6 Induction and recursion based on well-order relation

**Definition 2.40.** R is a well-founded relation on X if and only if

$$\forall U \subset \mathbf{X}(U \neq \emptyset \to \exists y \in U(\neg \exists z \in U(z\mathbf{R}y)))$$

**Definition 2.41.** Relation  $\mathbf{R}$  is **set-like** on  $\mathbf{X}$  iff for any  $x \in \mathbf{X}$ ,  $\{y \in \mathbf{X} \mid y\mathbf{R}x\}$  is a set

**Definition 2.42.** If R is a set-like relation on X and  $x \in X$ , define

$$\operatorname{pred}^{0}(\boldsymbol{X}, x, \boldsymbol{R}) = \{ y \in \boldsymbol{X} \mid y\boldsymbol{R}x \}$$
$$\operatorname{pred}^{n+1}(\boldsymbol{X}, x, bR) = \bigcup \{ \operatorname{pred}(\boldsymbol{X}, y, \boldsymbol{R}) \mid y \in \operatorname{pred}^{n}(\boldsymbol{X}, x, \boldsymbol{R}) \}$$
$$\operatorname{cl}(\boldsymbol{X}, x, \boldsymbol{R}) = \bigcup_{n \in \omega} \operatorname{pred}^{n}(\boldsymbol{X}, x, \boldsymbol{R})$$

**Lemma 2.43.** If R is a set-like relation on X, then for any  $y \in cl(X, x, R)$ ,  $pred(X, y, R) \subseteq cl(X, x, R)$ 

**Theorem 2.44** (Induction on well-founded set-like relation). If R is a well-founded set-like relation on X, then every nonempty  $Y \subseteq X$  has minimal element under R

**Theorem 2.45.** Suppose R is a well-founded set-like relation on X. If  $F: X \times V \to V$ , then there is a unique  $G: X \to V$  s.t.

$$\forall x \in \boldsymbol{X}(\boldsymbol{G}(x) = \boldsymbol{F}(x, \boldsymbol{G} | \text{pred}(\boldsymbol{X}, x, \boldsymbol{R})))$$

**Definition 2.46.** If R is a set-like well-founded relation on X, define

$$rank(x, \boldsymbol{X}, \boldsymbol{R}) = \sup\{rank(y, \boldsymbol{X}, \boldsymbol{R}) + 1 \mid y\boldsymbol{R}x \wedge y \in \boldsymbol{X}\}\$$

Note that

$$F(x,h) = \sup\{\alpha + 1 \mid \alpha \in \operatorname{ran}(h)\}\$$

**Lemma 2.47** (**ZF**<sup>-</sup>). If X is transitive and  $\in$  is well-founded on X, then  $X \subseteq WF$  and for any  $x \in X$ , rank $(x, X, \in) = \operatorname{rank}(x)$ 

**Definition 2.48.** R is a set-like well-founded relation on X, **Mostowski** function G on (X,R) is

$$\mathbf{G}(x) = \{ \mathbf{G}(y) \mid y \in \mathbf{X} \land y\mathbf{R}x \}$$

 $\mathbf{M} = \operatorname{ran}(\mathbf{G})$  is called the **Mostowski collapse** of  $(\mathbf{X}, \mathbf{R})$ 

**Lemma 2.49.** 1. 
$$\forall x, y \in X(xRy \rightarrow G(x) \in G(y))$$

- 2. **M** is transitive
- 3. If the axiom of power set holds,  $\mathbf{M} \subseteq \mathbf{WF}$
- 4. if the axiom of power set holds and  $x \in X$ , then  ${\rm rank}(x,X,R) = {\rm rank}(G(x))$

**Definition 2.50.** R is extensional on X iff

$$\forall x, y \in X (\forall z \in X (zRx \leftrightarrow zRy) \rightarrow x = y)$$

**Lemma 2.51.** If X is transitive then  $\in$  is extensional on X

**Lemma 2.52.** Let R be a set-like well-founded relation on X, G is a Mostowski function on it. If R is extensional, then G is an isomorphism

**Theorem 2.53** (Mostowski collapse theorem). Suppose R is set-like well-founded extensional on X, then there are unique transitive class M and bijection  $G: X \to M$  s.t.  $G: (X, R) \cong (M, \in)$