

Lab 2 data
Red-Black Tree

Names: Ahmed Talaat Noser ID: 05
Salma Ragab Gad 32

Time analysis:

- Red-Black tree has elements n , so the space is $O(n)$.
- Height is $O(\log n)$ -the worst case is $O(2 \log n_b)$ - where n_b is the number of black nodes.
- **Rotation:**
 - is $O(1)$ since a constant number of pointers are modified.
- **Insertion:**
 - To find the place to insert takes the height of the tree, or $O(\log n)$.
 - To add the node is $O(1)$.
 - To fix double red , rotation is $O(1)$.
 - Worst case, the double red can cascade all the way to the root. The cascade is proportional to the height of the tree, so the fixing takes $O(\log n)$, worst case.
 - Therefore, insertion is $O(\log n)$.
- **Deletion:**
 - Finding the node to delete plus finding the left-most right descendant is proportional to the height of the tree, so it is $O(\log n)$.
 - The swaping and deletion is $O(1)$.
 - Each individual fix (rotation, etc.) is $O(1)$.

- In the worst case, a double-black may get passed up to the root. Since each rotation takes constant time, this would be proportional to the height of the tree, and thus is $O(\log n)$.
- Therefore, the worst case cost of deletion is $O(\log n)$.

- **Search & Contains:**

- Take the height of the tree. Therefore, they are $O(\log n)$.

- **CeilingEntry:** height of the tree $O(\log n)$.

- **FloorEntry:** height of the tree $O(\log n)$.

- **EntrySet:** uses inorder traversal $O(n)$.

- **HeadMap:** uses inorder traversal, so it's $O(n)$.

- **FirstEntry:** height of the tree $O(\log n)$.

- **LastEntry:** height of the tree $O(\log n)$.

- **pollFirstElement:**
 - Find least key in $O(\log n)$.
 - Then remove it in $O(\log n)$.
 - So it takes $O(\log n)$.
- **PollLastEntry:**
 - Find greatest key in $O(\log n)$.
 - Then remove it in $O(\log n)$.
 - So it takes $O(\log n)$.
- **Values:** uses inorder traversal, so it's $O(n)$.

UML diagram:

