<u>Lab 2 data</u> <u>Red-Black Tree</u>

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Time analysis:

- Red-Black tree has elements n, so the space is O(n).
- Height is $O(\log n)$ -the worst case is $O(2 \log n_b)$ -where n_b is the number of black nodes.

• Rotation:

- is O(1) since a constant number of pointers are modified.

Insertion:

- To find the place to insert takes the height of the tree, or $O(\log n)$.
- To add the node is O(1).
- To fix double red, rotation is O(1).
- Worst case, the double red can cascade all the way to the root. The cascade is proportional to the height of the tree, so the fixing takes *O(logn)*, worst case.
- Therefore, insertion is $O(\log n)$.

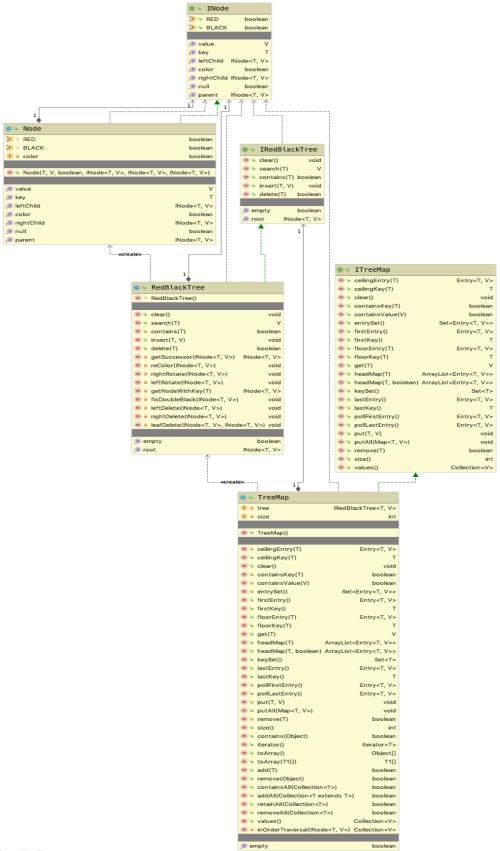
• Deletion:

- Finding the node to delete plus finding the leftmost right descendant is proportional to the height of the tree, so it is *O(log n)*.
- The swaping and deletion is O(1).
- Each individual fix (rotation, etc.) is O(1).

- In the worst case, a double-black may get passed up to the root. Since each rotation takes constant time, this would be proportional to the height of the tree, and thus is *O(log n)*.
- Therefore, the worst case cost of deletion is $O(\log n)$.
- Search & Contains:
- Take the height of the tree. Therefore, they are $O(\log n)$.
- CeilingEntry: height of the tree O(log n).
- FloorEntry: height of the tree O(log n).
- EntrySet: uses inorder traversal O(n).
- HeadMap: uses inorder traversal, so it's O(n).
- FirstEntry: height of the tree O(log n).
- LastEntry: height of the tree O(log n).

- pollFirstElement:
- Find least key in O(log n).
- Then remove it in O(log n).
- So it takes O(log n).
- PollLastEntry:
- Find greatest key in O(log n).
- Then remove it in O(log n).
- So it takes O(log n).
- Values: uses inorder traversal, so it's O(n).

UML diagram:



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