

# Compiler

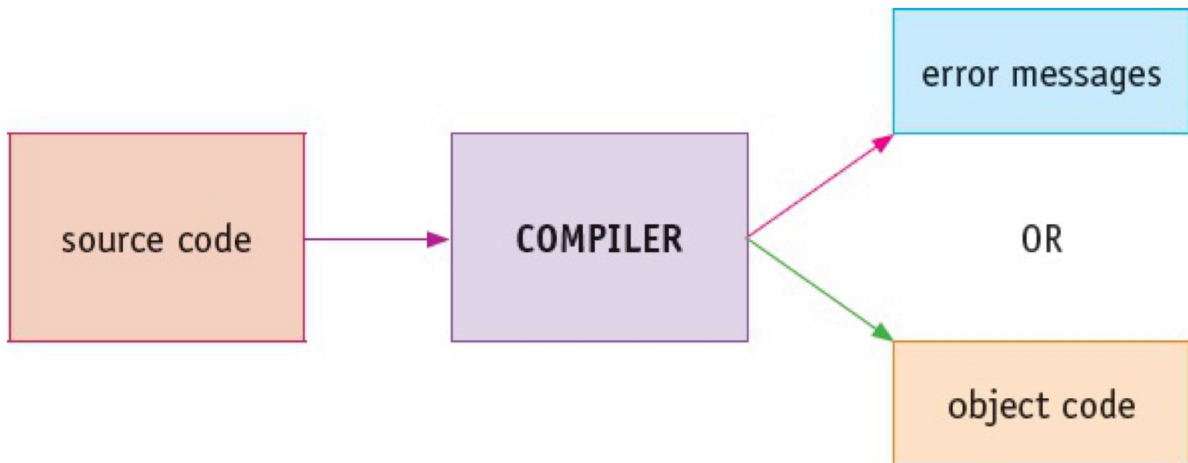



Figure 16.19

## Compilation stages

### Lexical analysis

- All necessary characters not required by the compiler, such as the white space and comments, are removed
- The white space removed // redundant characters are removed // removal of comments // identification of errors
- Tokenization
  - Using a keyword table that contains all the tokens for reserved keywords and symbols
  - Convert the source program into tokens
  - Keyword table
    - The reserved keyword used
    - The operator used

Keyword	Token
←	01
+	02
=	03
<>	04



IF	4A
THEN	4B
ENDIF	4C
ELSE	4D
REPEAT	4E
UNTIL	4F
TO	50
INPUT	51
OUTPUT	52
ENDFOR	53

- Their matching token
- Variables, constants, and identifiers added to a symbol table, and are then converted into locations/address
  - Symbol table
    - Identifier name used
    - The data type
    - Role: e.g. Constant, array, procedure
    - Location marker, value of constant

Symbol	Token	
	Value	Type
Counter	60	Variable
0	61	Constant
Password	62	Variable
"Cambridge"	63	Constant
1	64	Constant

Explain how the keyword table and symbol table are used to translate the source code program

- Keywords are looked up in the keyword table
- They are represented by tokens
- Identifiers are looked up in the symbol table
- They are converted into location/address
- Used to create a sequence of tokens (for the program)

## Syntax analysis

- Output from the lexical analysis is checked for grammatical/syntax errors - `parsing`
- The rules for `parsing` can be set out in Backus-Naur form(BNF) notation
- If errors are found: each statement and the associated error is outputted, but the next stage, code generation, will not be attempted
- If no error is found: passed to the next stage of compilation
- --
- Construction of a parse tree / parsing
- Checking that the rules of grammar/syntax have been obeyed
- Production of an error report

## Code generation

- Produces an object program to perform the task defined in the source code
- The object program is in machine-readable form(binary):
  - Either in machine code that can be directly executed by the CPU
  - Or in intermediate code that is converted into machine code when the program is loaded

# Optimization

- Performing the task using the minimum amount of resources
  - Execution time, storage space, memory, and CPU use.

A simple example of code optimisation is shown here:

Original code	$w = x + y$  $v = x + y + z$	Object code	LDD x ADD y STO w LDD x ADD y ADD z STO v
Optimised code	$w = x + y$  $v = w + z$	Object code	LDD x ADD y STO w ADD z STO v

Why code is optimized:

- Redundant code removed
- Program requires less memory
- Code reorganized to make it more efficient
- Program will complete task in a shorter time

Why optimization is necessary

- Optimisation means the code would have fewer instructions
- Optimised code occupies less memory in space
- Fewer instructions reduce the execution time of the program

Benefits

- Code has fewer instructions/occupies less memory in space
- Shortens the execution time for the program // time taken to execute whole program decreases

# Past-paper questions

(d) These lines of code are to be compiled:

```
X ← A + B  
Y ← A + B + C
```

Following the syntax analysis stage, object code is generated. The equivalent code, in assembly language, is shown below:

```
01 LDD 436 //loads value A  
02 ADD 437 //adds value B  
03 STO 612 //stores result in X  
04 LDD 436 //loads value A  
05 ADD 437 //adds value B  
06 ADD 438 //adds value C  
07 STO 613 //stores result in Y
```

Suggest what a compiler could do to optimise this code.

- Remove the second instances of LDD 436 // remove line 04
  - Remove the second instance s of ADD 437 // remove line 05
  - The value required is already stored in the accumulator
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Why compiled version helps protect the security of the source code

- Compiler produces executable version - not readable
  - Difficult for reverse engineering
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