

# SVD Computation

Samprit Chakraborty , Samahriti Mukherjee , Arunsoumya Basu , Aytijhya Saha  
Indian Statistical Institute, Kolkata

Fall 2021

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## 1 Introduction

We are interested to write a program to find SVD of a matrix from scratch.

## 2 Algorithm 1

We used QR Decomposition to find the SVD of a matrix.

Let  $A$  is a matrix and we want to find its eigenvalues. Let  $A = A_0$ , then let

$$A_0 = Q_0 R_0,$$

where  $Q_0 R_0$  is the QR Decomposition of  $A_0$ .

Now we define  $A_1 = R_0 Q_0$  and find the QR Decomposition of  $A_1$ , call it  $Q_1 R_1$ , then we define  $A_2 = R_1 Q_1$ . We continue like this and finally we will see that for large  $n$ ,  $A_n$  converges to an upper

triangular matrix. Now let  $P_n = Q_1 Q_2 \dots Q_n$  and  $P_n' A P_n = Q_n' Q_{n-1}' \dots Q_1' A Q_1 Q_2 \dots Q_n$

$$\implies P_n' A P_n = Q_n' Q_{n-1}' \dots Q_2' A_1 Q_2 Q_3 \dots Q_n$$

$\implies P_n' A P_n = A_n$  Now as  $A$  is similar to  $A_n$ , so eigenvalues of  $A$  will be same as eigenvalues of  $A_n$ .

Now as  $A_n$  is upper triangular, so diagonal elements of  $A_n$  will be the eigenvalues of  $A$ . Columns of  $P_k$  gives the eigenvectors to the corresponding eigenvalues which are orthogonal to each other. Proof of  $A_n$  converging to an upper triangular matrix:

We see  $P_{k+1} = P_k Q_{k+1}$

$\implies Q_{k+1} = P_k' P_{k+1}$

Now if  $P_k \rightarrow P$ , for large  $k$ , then  $Q_{k+1} \rightarrow I$ .

$\implies A_{k+1} \rightarrow R_k$ .

Now if we are given a matrix  $B$  for which we want the SVD Decomposition, we find the eigenvalues of  $BB^T$  and  $B^T B$ , which are the squares of the singular values of  $B$  and the corresponding orthogonal eigenvectors. The matrices formed by eigenvectors of  $BB^T$  and  $B^T B$  are  $U$  and  $V$ , say. We take a matrix  $C$  which has singular values of  $B$  in diagonals and all other elements are 0. Then  $UCV^T$  will be the SVD Decomposition of  $B$ .

## 2.1 R code based on this algorithm:

```
QR=function(A)
{
  m <- nrow(A)
  n <- ncol(A)
  count <- 1
  flag <- 0
  Q <- diag(1,n)
  unit = function(v) {v/sqrt(sum(v*v))} #
    Function to find the unit vector
  hmult = function(u, x) {x - 2* sum(u*x)*u} #
    Function to find the product of a Householder matrix with a vector
  shaver = function(x) #
    Function to shave a vector
  {
    x[1] = x[1] - sqrt(sum(x*x))
    unit(x)
  }
  u_vec=numeric((m*(m+1)/2-(m-n)*(m-n+1)/2)) #Vector
    to hold the values of the u_i s obtained at different steps
  for (i in 1:n)
  {
    if (i<m)
    {
```

```

        if (sum((A[(i+1):m,i])^2)>0.0000001) #
            Checking if the entries of A are already shaved, in case A is not of
            full column rank
        {
            u_vec[count:(count+(m-i))]=shaver(A[i:m,i]) #u_vec
                holds the values of the u_i s linearly
            for (j in 1:(n+1-i))
            {
                A[i:m,(i+j-1)]=hmult(u_vec[count:(count+(m-i))],A[i:m,(i+j-1)]) #
                    Multiplication with Householder matrix
            }
        }
        else
        {
            u_vec[count:(count+(m-i))]=0
        }
    }
    newmat1=diag(1,m)
    temp=diag(1,(m+1-i))-2*matrix(u_vec[count:(count+(m-i))])%*%t(matrix(u_vec[
        count:(count+(m-i))]))
    newmat1[i:m,i:m]=temp
    Q <- Q%*%newmat1
    count <- count+m+1-i
}

R=round(A,digits=6)
R <- R
return(list(Q=Q,R=R))
}

eig=function(M){
    X=M
    P=diag(1,dim(M)[1])
    for(i in 1:1000)
    {
        d=QR(X)
        Q=QR(X)$Q
        P=P%*%Q
        X=QR(X)$R%*%Q
    }
    return(list(diag(round(X,5)),round(P,5)))
}

SVDcomp=function(M){
    X=M
    D=t(M)%*%M

```

```

E=M%*%t(M)
V=eig(D)[2]
U=eig(E)[2]
a=unlist(eig(D)[1])
f=sqrt(as.numeric(a))
print(U)
print(f)
print(V)
}
SVDcomp(M)

```

## 2.2 Drawbacks of the method:

- We are blindly running the QR-flip loop (ie, computing QR and then finding RQ) a fixed number of times instead of running it until convergence.
- we are rounding off the result after each iteration.

## 3 Algorithm 2

First, we convert the given matrix  $M_{m \times n}$  into a bidiagonal matrix by a series of multiplications by Householder matrices. Here, we assume that  $m \geq n$ , otherwise all the operations are performed on  $\mathbf{M}^T$ . This is done in the following manner: We know that we can shave all but the first coordinate of a vector by multiplying it with a Householder matrix.

So, given distinct vectors  $\mathbf{u} \in \mathbb{R}^n$  and  $\mathbf{v} \in \mathbb{R}^n$  (with  $\|\mathbf{u}\| = \|\mathbf{v}\|$  and with the last  $(n-1)$  coordinates of  $\mathbf{v}$  0) we can multiply  $\mathbf{u}$  by a matrix  $\mathbf{A}$  where  $\mathbf{A} = \mathbf{I}_n - 2\mathbf{t}\mathbf{t}'$ , with  $\mathbf{t}$  being the unit vector in the direction of  $\mathbf{u} - \mathbf{v}$ .

This enables us to make the last  $(m-1)$  entries of the first column of  $M$  0 through a premultiplication with a Householder matrix  $\mathbf{U}_1$ . Next, we make the first  $(n-2)$  entries of the first row of this new  $\mathbf{M}$  zeroes, by postmultiplying the matrix with  $\mathbf{V}_1$ . Then the last  $(m-2)$  entries of the second column of  $\mathbf{M}$  are converted to 0 by another premultiplication with  $\mathbf{U}_2$ . After this, all but the first three entries of the second row of this new matrix are shaved by a multiplication with  $\mathbf{V}_2$ . This process of alternately shaving columns and rows continues till we have generated  $\mathbf{U}_1, \mathbf{U}_2, \dots, \mathbf{U}_n$  and  $\mathbf{V}_1, \mathbf{V}_2, \dots, \mathbf{V}_{n-2}$ .

We note that at each step, the matrix  $\mathbf{M}$  is overwritten after multiplication with the Householder matrix. Also, the Householder matrices are gradually decreasing in size, so they are packed into a  $m \times m$  identity matrix (for  $\mathbf{U}$ ) and into a  $n \times n$  identity matrix (for  $\mathbf{V}$ ) as a principal sub-matrix formed by the last few rows and columns. Thus, it makes sense to compute the product

$$\mathbf{U}^T = \mathbf{U}_n \mathbf{U}_{n-1} \cdots \mathbf{U}_1 \text{ and } \mathbf{V} = \mathbf{V}_1 \mathbf{V}_2 \cdots \mathbf{V}_{n-2}.$$

At the end of this process,  $\mathbf{M}$  has been reduced to a bidiagonal matrix and we have found the orthogonal matrices  $\mathbf{U}$  and  $\mathbf{V}$  which transform  $\mathbf{M}$  into a bidiagonal matrix.

Now we need to compute the Singular Value Decomposition of this bidiagonal  $\mathbf{M}$ , which turns out to be computationally simpler than directly finding the Singular Value Decomposition of the original matrix.

This is done through a series of alternate **QR** and **LQ** decompositions, carried out until we get an approximately diagonal matrix of singular values.

This diagonal matrix is the singular matrix of the original matrix  $\mathbf{M}$ , since orthogonal transformations do not alter the singular values. The  $\mathbf{U}$  and the  $\mathbf{V}$  matrices obtained from the Singular Value Decomposition of the bidiagonal matrix are then multiplied with the existing  $\mathbf{U}$  and  $\mathbf{V}$  matrices to get the final  $\mathbf{U}$  and  $\mathbf{V}$  matrices.

### 3.1 R code based on this algorithm:

```
SVD=function(B)      #A function to compute the SVD of a matrix with nrow>=2, ncol>=2
{
  flag <- 0

  if(nrow(B)<ncol(B))
  {
    A=matrix(0,(ncol(B)+1),(nrow(B)+1))
    A[1:ncol(B),1:nrow(B)]=t(B)
  }

  else
  {
    A=matrix(0,(nrow(B)+1),(ncol(B)+1))
    A[1:nrow(B),1:ncol(B)]=B
  }

  m <- nrow(A)-1
  n <- ncol(A)-1

  unit = function(v) {v/sqrt(sum(v*v))}      #Function to find the unit
                                              vector

  hmult = function(u, x) {x - 2* sum(u*x)*u}   #Function to find the product
                                              of a Householder matrix with a vector

  shaver = function(x)                        #Function to shave a vector
  {
```

```

x[1] = x[1] - sqrt(sum(x*x))
unit(x)
}

for (i in 1:n)
{
  if (sum((A[(i+1):m,i])^2)>10^(-12))      #Checking if the entries of A
    are already shaved
  {
    if (i !=m)
    {
      temp=shaver(A[i:m,i])
    }
    else
    {
      temp=0
    }
    for (j in i:n)
    {
      A[i:m,j]=hmult(temp,A[i:m,j])      #Multiplication with
      Householder matrix
    }
    A[(i+1):(m+1),i]=temp
  }

  else
  {
    A[(i+1):(m+1),i]=0
  }

  if (i<=(n-2))
  {
    if (sum((A[i,(i+2):n])^2)>10^(-12))
    {
      temp=shaver(A[i,(i+1):n])
      for (k in i:m)
      {
        A[k,(i+1):n]=hmult(temp,A[k,(i+1):n])
      }
      A[i,(i+2):(n+1)]=temp
    }
  }
  else

```

```

        {
            A[i,(i+2):(n+1)]=0
        }
    }
}

A <- A

#Calculating U_A and V_A, the orthogonal matrices used for bidiagonalization

U_A <- diag(1,m)

for (i in 1:n)
{
    newmat1=diag(1,m)
    temp=diag(1,(m+1-i))-2*matrix(A[(i+1):(m+1),i])%*%t(matrix(A[(i+1):(m+1),i]))
    newmat1[i:m,i:m]=temp
    U_A <- U_A%*%newmat1
}

V_A <- diag(1,n)

if (n>=3)
{
    for (i in 1:(n-2))
    {
        newmat1=diag(1,n)
        temp=diag(1,(n-i))-2*matrix(A[i,(i+2):(n+1)])%*%t(matrix(A[i,(i+2):(n+1)]))
        newmat1[(i+1):n,(i+1):n]=temp
        V_A <- V_A%*%newmat1
    }
}

#print(U_A)
#print(V_A)
#print(round(t(U_A)%*%U_A,digits=6))
#print(round(t(V_A)%*%V_A,digits=6))

if (nrow(B)>=ncol(B))
{
    bid <- t(U_A)%*%B%*%V_A
}

```

```

else
{
  bid <- t(U_A)%*%t(B)%*%V_A
}

bid <- round(bid,digits=6)      #The bidiagonal matrix

U3=diag(1,min(nrow(B),ncol(B)))
V3=diag(1,min(nrow(B),ncol(B)))

U_A <- U_A[1:m,1:n]

epsilon=max(abs(bid[1:n,1:n]-diag(diag(bid))))

while(epsilon>0.0000001)      #Alternating between QR and LQ decompositions to
  iteratively arrive at the SVD of the bidiagonal matrix
{
  Q=qr.Q(qr(bid[1:n,1:n]))
  R=qr.R(qr(bid[1:n,1:n]))
  L=t(qr.R(qr(t(R))))
  P=qr.Q(qr(t(R)))
  U3=U3%*%Q
  V3=t(P)%*%V3
  bid[1:n,1:n]=L
  epsilon=max(abs(bid[1:n,1:n]-diag(diag(bid))))
}
V3=t(V3)

U <- U_A%*%U3
V <- V_A%*%V3
sig <- round(bid[1:n,1:n],digits=6)

if (nrow(B)<ncol(B))
{
  mat1=t(U%*%sig%*%t(V))
  temp1 <- U
  U <- V
  V <- temp1
}
else
{
  mat1=U%*%sig%*%t(V)
}

```



```

for (i in 1:n)
{
  if (sig[i,i]<0)
  {
    U[,i]=U[,i]*-1
    sig[i,i]=-sig[i,i]
  }
}

cat('\nU is:\n')
print(round(U,digits=6))    #The final matrix U which may not be a square
                             matrix; orthonormal basis extension may be required to do so
cat('\nSigma is:\n')
print(sig)                  #The diagonal matrix of singular values
cat('\nV is:\n')
print(round(V,digits=6))    #The final matrix V which can be a non-square
                             matrix

cat('\nThe product U*%Sigma*%t(V) is:\n')
print(round(mat1,digits=5))
}

#Examples

mat=matrix(c(2,6,1,8,5,0,7,4,3),3,3)
SVD(mat)

mat=matrix(1:15,5,3)
SVD(mat)

mat=matrix(1:40,5,8)
SVD(mat)

```