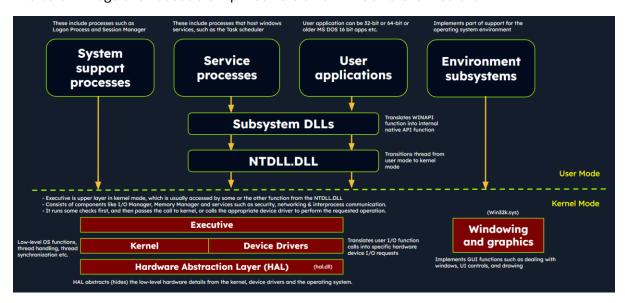
Windows Internals

To conduct effective malware analysis, a profound understanding of Windows internals is essential. Windows operating systems function in two main modes:

- User Mode: This mode is where most applications and user processes operate.
 Applications in user mode have limited access to system resources and must interact with the operating system through Application Programming Interfaces (APIs). These processes are isolated from each other and cannot directly access hardware or critical system functions. However, in this mode, malware can still manipulate files, registry settings, network connections, and other user-accessible resources, and it may attempt to escalate privileges to gain more control over the system.
- Kernel Mode: In contrast, kernel mode is a highly privileged mode where the Windows kernel runs. The kernel has unrestricted access to system resources, hardware, and critical functions. It provides core operating system services, manages system resources, and enforces security and stability. Device drivers, which facilitate communication with hardware devices, also run in kernel mode. If malware operates in kernel mode, it gains elevated control and can manipulate system behavior, conceal its presence, intercept system calls, and tamper with security mechanisms.

Windows Architecture At A High Level

The below image showcases a simplified version of Windows' architecture.



The simplified Windows architecture comprises both user-mode and kernel-mode components, each with distinct responsibilities in the system's functioning.

User-mode Components

User-mode components are those parts of the operating system that don't have direct access to hardware or kernel data structures. They interact with system resources through APIs and system calls. Let's discuss some of them:

• System Support Processes: These are essential components that provide crucial functionalities and services such as logon processes (winlogon.exe), Session Manager

(smss.exe), and Service Control Manager (services.exe). These aren't Windows services but they are necessary for the proper functioning of the system.

- Service Processes: These processes host Windows services like the Windows Update Service, Task Scheduler, and Print Spooler services. They usually run in the background, executing tasks according to their configuration and parameters.
- User Applications: These are the processes created by user programs, including both 32-bit and 64-bit applications. They interact with the operating system through APIs provided by Windows. These API calls get redirected to NTDLL.DLL, triggering a transition from user mode to kernel mode, where the system call gets executed. The result is then returned to the user-mode application, and a transition back to user mode occurs.
- Environment Subsystems: These components are responsible for providing execution environments for specific types of applications or processes. They include the <u>Win32</u> <u>Subsystem</u>, <u>POSIX</u>, and <u>OS/2</u>.
- Subsystem DLLs: These dynamic-link libraries translate documented functions into appropriate internal native system calls, primarily implemented in NTDLL.DLL. Examples include kernelbase.dll, user32.dll, wininet.dll, and advapi32.dll.

Kernel-mode Components

Kernel-mode components are those parts of the operating system that have direct access to hardware and kernel data structures. These include:

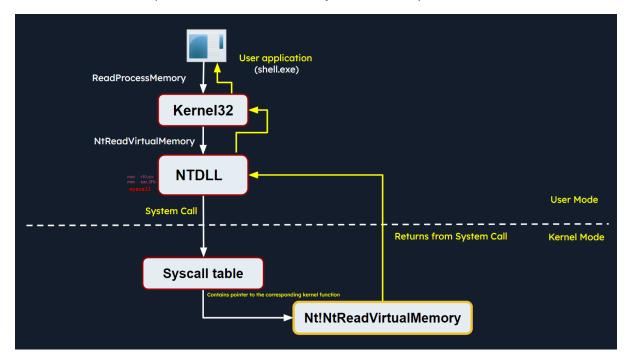
- Executive: This upper layer in kernel mode gets accessed through functions
 from NTDLL.DLL. It consists of components like the I/O Manager, Object
 Manager, Security Reference Monitor, Process Manager, and others, managing the core
 aspects of the operating system such as I/O operations, object management, security,
 and processes. It runs some checks first, and then passes the call to kernel, or calls the
 appropriate device driver to perform the requested operation.
- Kernel: This component manages system resources, providing low-level services like thread scheduling, interrupt and exception dispatching, and multiprocessor synchronization.
- Device Drivers: These software components enable the OS to interact with hardware devices. They serve as intermediaries, allowing the system to manage and control hardware and software resources.
- Hardware Abstraction Layer (HAL): This component provides an abstraction layer between the hardware devices and the OS. It allows software developers to interact with hardware in a consistent and platform-independent manner.
- Windowing and Graphics System (Win32k.sys): This subsystem is responsible for managing the graphical user interface (GUI) and rendering visual elements on the screen.

Now, let's discuss in what happens behind the scenes when an user application calls a Windows API function.

Windows API Call Flow

Malware often utilize Windows API calls to interact with the system and carry out malicious operations. By understanding the internal details of API functions, their parameters, and expected behavior, analysts can identify suspicious or unauthorized API usage.

Let's consider an example of a Windows API call flow, where a user-mode application tries to access privileged operations and system resources using the **ReadProcessMemory function**. This function allows a process to read the memory of a different process.



When this function is called, some required parameters are also passed to it, such as the handle to the target process, the source address to read from, a buffer in its own memory space to store the read data, and the number of bytes to read. Below is the syntax of ReadProcessMemory WINAPI function as per Microsoft documentation.

```
BOOL ReadProcessMemory(

[in] HANDLE hProcess,

[in] LPCVOID lpBaseAddress,

[out] LPVOID lpBuffer,

[in] SIZE_T nSize,

[out] SIZE_T *lpNumberOfBytesRead
);
```

```
C++

BOOL ReadProcessMemory(
   [in] HANDLE hProcess,
   [in] LPCVOID lpBaseAddress,
   [out] LPVOID lpBuffer,
   [in] SIZE_T nSize,
   [out] SIZE_T *lpNumberOfBytesRead
);
```

ReadProcessMemory is a Windows API function that belongs to the kernel32.dll library. So, this call is invoked via the kernel32.dll module which serves as the user mode interface to the Windows API. Internally, the kernel32.dll module interacts with the NTDLL.DLL module, which provides a lower-level interface to the Windows kernel. Then, this function request is translated to the corresponding Native API call, which is NtReadVirtualMemory. The below screenshot from x64dbg demonstrates how this looks like in a debugger.

```
48:83EC 48
48:804424 30
48:894424 20
48:FF15 F30416
0F1F4400 00
48:885424 70
48:85D2

sub rsp,48
lea rax,qword ptr ss:[rsp+30]
mov qword ptr ss:[rsp+20],rax
call qword ptr ds:[<a href="https://www.wightheadvirtualMemory">www.wightheadvirtualMemory</a>
nop dword ptr ds:[rax+rax],eax
mov rdx,qword ptr ss:[rsp+70]
test rdx,rdx
```

The NTDLL.DLL module utilizes system calls (syscalls).

```
mov r10, rcx
mov eax, 3F Syscall Number
4C:8BD1
                                                      NtReadVirtualMemory
B8 3F000000
                                                      3F: '?'
F60425 0803FE7 test byte ptr ds:[7FFE0308],1
75 03
                 jne ntdll.7FFD0C1CD7E5
0F05
                    Syscall instruction
C3
                 ret
                     2E
CD 2E
                 int
C3
                 ret
```

The syscall instruction triggers the system call using the parameters set in the previous instructions. It transfers control from user mode to kernel mode, where the kernel performs the requested operation after validating the parameters and checking the access rights of the calling process.

If the request is authorized, the thread is transitioned from user mode into the kernel mode. The kernel maintains a table known as the System Service Descriptor Table (SSDT) or the syscall table (System Call Table), which is a data structure that contains pointers to the various system service routines. These routines are responsible for handling system calls made by user-mode applications. Each entry in the syscall table corresponds to a specific system call number, and the associated pointer points to the corresponding kernel function that implements the requested operation.

The syscall responsible for ReadProcessMemory is executed in the kernel, where the Windows memory management and process isolation mechanisms are leveraged. The kernel performs necessary validations, access checks, and memory operations to read the memory from the target process. The kernel retrieves the physical memory pages corresponding to the requested virtual addresses and copies the data into the provided buffer.

Once the kernel has finished reading the memory, it transitions the thread back to user mode and control is handed back to the original user mode application. The application can then access the data that was read from the target process's memory and continue its execution.

Portable Executable

Windows operating systems employ the Portable Executable (PE) format to encapsulate executable programs, DLLs (Dynamic Link Libraries), and other integral system components. In the realm of malware analysis, an intricate understanding of the PE file format is indispensable. It allows us to gain significant insights into the executable's structure, operations, and potential malign activities embedded within the file.

PE files accommodate a wide variety of data types including executables (.exe), dynamic link libraries (.dll), kernel modules (.srv), control panel applications (.cpl), and many more. The PE file format is fundamentally a data structure containing the vital information required for the Windows OS loader to manage the executable code, effectively loading it into memory.

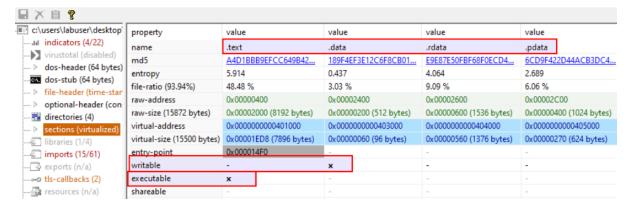
PE Sections

The PE Structure also houses a Section Table, an element comprising several sections dedicated to distinct purposes. The sections are essentially the repositories where the actual content of the file, including the data, resources utilized by the program, and the executable code, is stored. The .text section is often under scrutiny for potential artifacts related to injection attacks.

Common PE sections include:

- Text Section (.text): The hub where the executable code of the program resides.
- Data Section (.data): A storage for initialized global and static data variables.
- Read-only initialized data (.rdata): Houses read-only data such as constant values, string literals, and initialized global and static variables.
- Exception information (.pdata): A collection of function table entries utilized for exception handling.
- BSS Section (.bss): Holds uninitialized global and static data variables.
- Resource Section (.rsrc): Safeguards resources such as images, icons, strings, and version information.
- Import Section (.idata): Details about functions imported from other DLLs.
- Export Section (.edata): Information about functions exported by the executable.
- Relocation Section (.reloc): Details for relocating the executable's code and data when loaded at a different memory address.

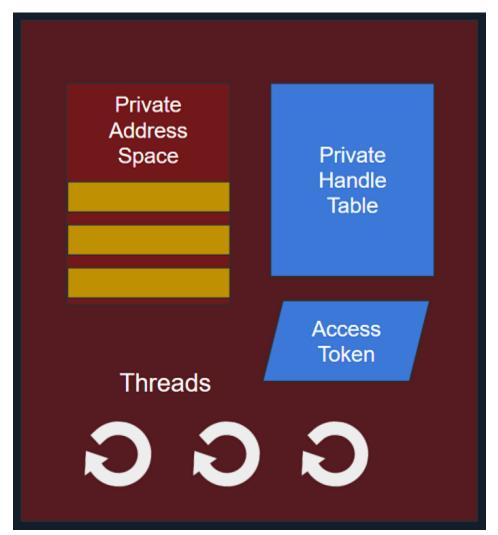
We can visualize the sections of a portable executable using a tool like pestudio as demonstrated below.



Delving into the Portable Executable (PE) file format is pivotal for malware analysis, offering insights into the file's structure, code analysis, import and export functions, resource analysis, anti-analysis techniques, and extraction of indicators of compromise. Our comprehension of this foundation paves the way for efficacious malware analysis.

Processes

In the simplest terms, a process is an instance of an executing program. It represents a slice of a program's execution in memory and consists of various resources, including memory, file handles, threads, and security contexts.



Each process is characterized by:

- A unique PID (Process Identifier): A unique Process Identifier (PID) is assigned to each
 process within the operating system. This numeric identifier facilitates the tracking and
 management of the process by the operating system.
- Virtual Address Space (VAS): In the Windows OS, every process is allocated its own
 virtual address space, offering a virtualized view of the memory for the process. The VAS
 is sectioned into segments, including code, data, and stack segments, allowing the
 process isolated memory access.
- Executable Code (Image File on Disk): The executable code, or the image file, signifies the binary executable file stored on the disk. It houses the instructions and resources necessary for the process to operate.
- Table of Handles to System Objects: Processes maintain a table of handles, a reference catalogue for various system objects. System objects can span files, devices, registry keys, synchronization objects, and other resources.
- Security Context (Access Token): Each process has a security context associated with it, embodied by an Access Token. This Access Token encapsulates information about the process's security privileges, including the user account under which the process operates and the access rights granted to the process.
- One or More Threads Running in its Context: Processes consist of one or more threads, where a thread embodies a unit of execution within the process. Threads enable concurrent execution within the process and facilitate multitasking.

Dynamic-link library (DLL)

A Dynamic-link library (DLL) is a type of PE which represents "Microsoft's implementation of the shared library concept in the Microsoft Windows OS". DLLs expose an array of functions which can be exploited by malware, which we'll scrutinize later. First, let's unravel the import and export functions in a DLL.

Import Functions

- Import functions are functionalities that a binary dynamically links to from external libraries or modules during runtime. These functions enable the binary to leverage the functionalities offered by these libraries.
- During malware analysis, examining import functions may shed light on the external libraries or modules that the malware is dependent on. This information aids in identifying the APIs that the malware might interact with, and also the resources such as the file system, processes, registry etc.
- By identifying specific functions imported, it becomes possible to ascertain the actions the malware can perform, such as file operations, network communication, registry manipulation, and more.
- Import function names or hashes can serve as IOCs (Indicators of Compromise) that assist in identifying malware variants or related samples.

Below is an example of identifying process injection using DLL imports and function names:



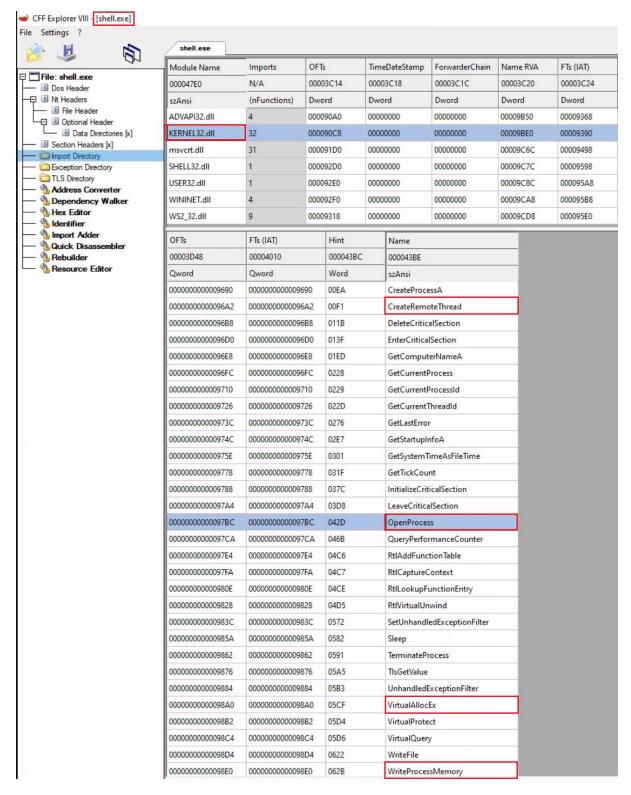
In this diagram, the malware process (shell.exe) performs process injection to inject code into a target process (notepad.exe) using the following functions imported from the DLL kernel32.exe:

- OpenProcess: Opens a handle to the target process (notepad.exe), providing the necessary access rights to manipulate its memory.
- VirtualAllocEx: Allocates a block of memory within the address space of the target process to store the injected code.
- WriteProcessMemory: Writes the desired code into the allocated memory block of the target process.
- CreateRemoteThread: Creates a new thread within the target process, specifying the entry point of the injected code as the starting point.

As a result, the injected code is executed within the context of the target process by the newly created remote thread. This technique allows the malware to run arbitrary code within the target process.

The functions above are WINAPI (Windows API) functions. Don't worry about WINAPI functions as of now. We'll discuss these in detail later.

We can examine the DLL imports of shell.exe (residing in the C:\Samples\MalwareAnalysis directory) using CFF Explorer (available at C:\Tools\Explorer Suite) as follows.

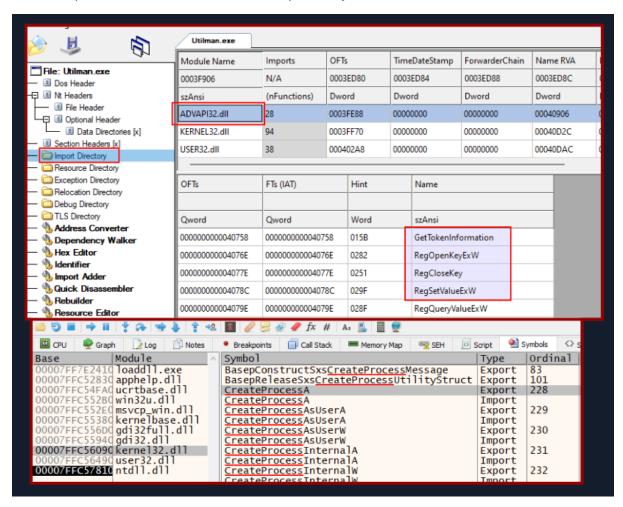


Export Functions

- Export functions are the functions that a binary exposes for use by other modules or applications.
- These functions provide an interface for other software to interact with the binary.

In the below screenshot, we can see an example of DLL imports (using CFF Explorer) and exports (using x64dbg - Symbols tab):

- Imports: This shows the DLLs and their functions imported by an executable Utilman.exe.
- Exports: This shows the functions exported by a DLL Kernel32.dll.



In the context of malware analysis, understanding import and export functions assists in discerning the behavior, capabilities, and interactions of the binary with external entities. It yields valuable information for threat detection, classification, and gauging the impact of the malware on the system.