

Exam 2. Problem 2

Base case

when $n=1$, $A = [\alpha_{11}]$ and bordered algorithm will run for one iteration. $\check{L} = [1]$ $\check{U} = [\alpha_{11}]$ as the result.

$$\Delta A = [0] \leq \gamma_1 |\check{L}| |\check{U}| \text{ since } \gamma_1 \geq 0$$

Induction step

Assume $A_{00} + \Delta A_{00} = \check{L}_{00} \check{U}_{00}$ with $\Delta A_{00} \leq \gamma_n |\check{L}_{00}| |\check{U}_{00}|$

then for $A = \mathbb{R}^{m+1} \times \mathbb{R}^{n+1}$

$$\begin{bmatrix} \check{L}_{00} & 0 \\ \check{L}_{10}^T & 1 \end{bmatrix} \begin{bmatrix} \check{U}_{00} & \check{U}_{01} \\ 0 & \check{U}_{11} \end{bmatrix} = \begin{bmatrix} \check{L}_{00} \check{U}_{00} & \check{L}_{00} \check{U}_{01} \\ \check{L}_{10}^T \check{U}_{00} & \check{L}_{10}^T \check{U}_{01} + \check{U}_{11} \end{bmatrix}$$

To prove

$$\gamma_{n+1} \left| \begin{bmatrix} \check{L}_{00} & 0 \\ \check{L}_{10}^T & 1 \end{bmatrix} \right| \left| \begin{bmatrix} \check{U}_{00} & \check{U}_{01} \\ 0 & \check{U}_{11} \end{bmatrix} \right| = \begin{bmatrix} \gamma_{n+1} |\check{L}_{00}| |\check{U}_{00}| & \gamma_{n+1} |\check{L}_{00}| |\check{U}_{01}| \\ \gamma_{n+1} |\check{L}_{10}^T| |\check{U}_{00}| & \gamma_{n+1} |\check{L}_{10}^T| |\check{U}_{01}| + \gamma_{n+1} |\check{U}_{11}| \end{bmatrix}$$

We need to prove the following conclusion.

1. $A_{00} + \Delta A_{00} = \check{L}_{00} \check{U}_{00}$ with $|\Delta A_{00}| \leq \gamma_{n+1} |\check{L}_{00}| |\check{U}_{00}|$
this is proved by the fact that $\gamma_{n+1} \geq \gamma_n$

2. $A_{01} + \Delta A_{01} = \check{L}_{00} \check{U}_{01}$ with $|\Delta A_{01}| \leq \gamma_{n+1} |\check{L}_{00}| |\check{U}_{01}|$

We first need to prove a lemma:

~~Lemma~~

Lemma 1

Let $L \in \mathbb{R}^{n \times n}$ be a non-singular lower triangular matrix and \check{x} be the computed result of $Lx = y$ then there is δy such that:

$$y + \delta y = L\check{x} \text{ with } |\delta y| \leq \gamma_n |L| |\check{x}|$$

According to Corollary 6.4.1.4 of the book.

$$(L + \delta L) \tilde{x} = y \text{ where } |\delta L| \leq \gamma_n |L|$$

$$\text{let } \delta y = -\delta L \tilde{x} \text{ then we have } y + \delta y = L \tilde{x} \text{ with } |\delta y| = |\delta L| |\tilde{x}| \leq \gamma_n |L| |\tilde{x}|$$

Then for conclusion 2. $L_{00} \in \mathbb{R}^{n \times n}$ is a non-singular lower triangular matrix and \tilde{u}_0 is the computed result when this bordered algorithm

Solve $L_{00} \tilde{u}_0 = a_0$ according to Lemma 1

$$a_0 + \delta a_0 = L_{00} \tilde{u}_0 \text{ with } |\delta a_0| \leq \gamma_n |L_{00}| |\tilde{u}_0| \leq \gamma_{n+1} |\tilde{L}_0| |\tilde{u}_0|$$

Conclusion 3

$$a_0^T + \delta a_0^T = \tilde{L}_0^T \tilde{u}_0 \text{ with } |\delta a_0^T| \leq \gamma_{n+1} |\tilde{L}_0^T| |\tilde{u}_0|$$

As \tilde{L}_0^T is the computed result by solving $L_{00}^T \tilde{u}_0 = a_0^T$

Thus \tilde{L}_0 is the result of $U_{00}^T \tilde{L}_0 = a_0$

since U_{00}^T is non-singular lower triangular matrix

then according to Lemma 1:

$$a_0 + \delta a_0 = U_{00}^T \tilde{L}_0 \text{ with } |\delta a_0| \leq \gamma_n |U_{00}^T| |\tilde{L}_0|$$

Finally we have:

$$a_0^T + \delta a_0^T = \tilde{L}_0^T \tilde{u}_0 \text{ with } |a_0^T| \leq \gamma_n |\tilde{L}_0^T| |\tilde{u}_0| \leq \gamma_{n+1} |\tilde{L}_0^T| |\tilde{u}_0|$$

Conclusion 4.

$$\alpha_{11} + \delta \alpha_{11} = \tilde{L}_0^T \tilde{u}_0 + \tilde{u}_{11} \text{ with } |\delta \alpha_{11}| \leq \gamma_{n+1} (|\tilde{L}_0^T| |\tilde{u}_0| + |\tilde{u}_{11}|)$$

\tilde{u}_{11} is the computed result of bordered algorithm step $\tilde{u}_{11} = \alpha_{11} - \tilde{L}_0 \tilde{u}_0$

According to Reference [1], there exist $\delta \alpha_{11}$ such

$$\tilde{u}_{11} = \alpha_{11} - \tilde{L}_0^T \tilde{u}_0 + \delta \alpha_{11} \text{ with } |\delta \alpha_{11}| \leq \gamma_n (|\tilde{L}_0^T| |\tilde{u}_0| + |\tilde{u}_{11}|)$$

From

$$\text{Lastly, } \alpha_{11} + \delta \alpha_{11} = \tilde{L}_0^T \tilde{u}_0 + \tilde{u}_{11} \text{ with } |\delta \alpha_{11}| \leq \gamma_{n+1} (|\tilde{L}_0^T| |\tilde{u}_0| + |\tilde{u}_{11}|)$$

By proving the 4 conclusion we prove the induction step for our question.

Reference:

- [1] Paolo Bientinesi, Robert A. van de Geijn, Goal oriented and Modular Stability Analysis. SIAM Journal on Matrix Analysis and Applications, Volume 32 Issue 1, Feb 2011