# Lists: ArrayList and LinkedList

#### List Interfaces Methods

```
add(Elm)
add(index, Elm)
addAll(Coll)
clear()
contains(Obj)
containsAll(Coll)
equals(Obj)
get(index)
hashCode()
index0f(0bj)
isEmpty()
```

```
iterator()
lastIndexOf(Obi)
remove(index)
removeAll(Coll)
• remove(0bj)
retainAll(Coll)
set(index, Elm)
size()
subList(from, to)
toArray()
toArray(T[] arr)
```

# List Implementations

- Two implementations:
  - ArrayList
  - LinkedList
- ArrayList:
  - Resizable array
    - Essentially a wrapper class around an array
  - Uses Contiguous Allocation
  - Usually has more memory allocated than is strictly necessary
- LinkedList
  - Uses Linked Allocation

# ArrayList: add(E)

Adds the element to the end of the list

myAL.add(1);

- What happens if we now myAL.add(5);?
  - List is full
  - Create a new list that is twice as big, copy the data, then add to the end

3		l .	1				
3	1	4	1	Ø	Ø	Ø	Ø
3	1	4	1	5	Ø	Ø	Ø

 Adding to the end is fast (except every once in a while, don't worry about that)

# ArrayList: add(index, E)

Adds the item at the specified index

myAL.add(1, 10); // Add the number 0 at
index 1

3	1	4	1	5	Ø	Ø	Ø
3	Ø	1	4	1	5	Ø	Ø
3	10	1	4	1	5	Ø	Ø

- Everything from the specified index on must be shifted to the right
  - Worst case, if inserting at index 0, every element must be shifted
- This operation is slow
  - The operation will take longer with 1000 elements than with 10

# Other ArrayList Operations

- get(index)
  - Does simple math, is fast
- index0f(0bj)
  - Must search through and visit every element (potentially), is slow
- lastIndexOf(Obj)
  - Like indexOf(Obj), but in reverse, so is slow
- remove(index) or remove(Obj)
  - Go to the index (fast ), shift everything left (slow), so slow
- set(index, E)
  - Go to the index (fast ), do assignment (fast ), so fast

# ArrayList: Summary

- Fast at/Good at
  - get(index), set(index, Elm) (Random Access)
  - Add and Remove at end (usually)

#### Slow at/Bad at

- Add at beginning/middle
- Remove from beginning/middle
- Finding elements (searching)
- Can use up to twice as much memory as needed to hold data

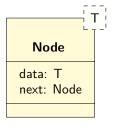
#### Linked Lists

- Linked Lists are built of nodes (like a chain)
  - A node consists of (at least) a data field and a reference to the next node in the list (Singly Linked List)
  - The node may also contain a reference to the previous node in the list (Doubly Linked List)

#### 

- The LinkedList maintains a reference to the beginning (head) and the end (tail) of the list
- All examples in the following slides will demonstrate with a Singly Linked List; it is relatively straightforward to expand your understanding to a Doubly Linked List

#### LinkedList Nodes



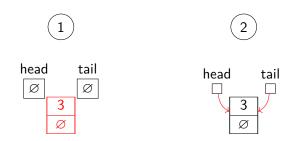
# size: int head, tail: Node

Note: A doubly linked node would have a prev member as well.

#### LinkedList

- myLL.add(3); // Add to empty list
- Create a new node with the data and the next reference set to null
- opoint both head and tail to point at the new node

# Adding to Empty LinkedList

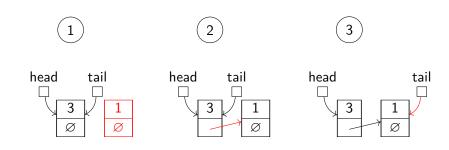


# Adding to End of LinkedList

```
myLL.add(1); // or myLL.addLast(1)
```

- Oreate a new node with the data (and a null next reference)
- Point tail.next at new node
- Point tail at new node

# Adding to End of LinkedList



# Adding to End of LinkedList

- The number of operations this takes is independent of the number of things in the list (and is small).
- This means the operation is fast
- Similarly, adding at the beginning is fast
  - Replace tail with head in pseudocode

#### LinkedList addLast Pseudocode

```
addLast(data):
    if size == 0:
        head = new Node(data, null)
        tail = head
else:
        temp = new Node(data, null)
        tail.next = temp
        // if doubly linked list, temp.prev = tail
        tail = temp
size ++
```

#### Random Access: LinkedList

- The only two nodes we know the positions of are head and tail
- In order to get to a node at a specific index (to either get the value or set the value), we must start at the head and iterate through
- The amount of time this can take grows as the list grows, thus it is a slow operation

#### Random Access Pseudocode: LinkedList

```
get(index):
    if index >= size:
        throw Exception
    Node it = head
    for (int i=0; i<index; i++)
        it = it.next
    return it.data
set(index, data):
    if index >= size:
        throw Exception
    Node it = head
    for (int i=0; i<index; i++)
        it = it.next
    it.data = data
```

#### Add to Middle Pseudocode: LinkedList

```
add(index, data):
    if index > size:
        throw Exception
    if index == 0:
        addFirst(data)
    else if index == size:
        addLast(data)
    else:
        Node it = head
        // iterate to node before index to add at
        for (int i=0; i<index - 1; i++)
            it = it.next
        temp = new Node(data, null)
        temp.next = it.next
        it.next = temp
        size ++
```

#### Addition and Removal from Middle of LinkedList

- Removal is similar to addition in the middle of a LinkedList
  - Iterate to node before index to add/remove
  - Link or Unlink node as appropriate
- The linking and unlinking is fast
- Iterating to the node before the node to modify, however, is slow
- Thus, these operations are slow
  - The exception: Removal from the beginning of a Doubly Linked List is fast

# Searching in a LinkedList

- Like the ArrayList, searching in a LinkedList requires visiting possibly all elements
- Thus, operations like index0f are slow

# Comparing ArrayList and LinkedList

	Rando	n Access	m Pdd/R	Nid Padde	End of	ind Search
ArrayList	fast	slow	slow	fast	fast	slow
Singly LinkedList	slow	fast	slow	fast	slow	slow
Doubly LinkedList	slow	fast	slow	fast	fast	slow

Note: Arrays are similar to ArrayLists, but are not resizable. They do have arguably more convenient indexing (citation needed).