# COM S 352: Introduction to Operating Systems Midterm Practice Exam Spring 2023

### **Cover Sheet**

| Student Name: |  |  |
|---------------|--|--|
|               |  |  |

#### Format:

Time: 50 minsPoints: 100

• Question Types: matching, true/false and short answer

#### **Instructions:**

- You may use 1 (one) letter sized sheet of paper (front and back) or 2 one-side sheets, that you have prepared yourself with notes before the exam, as "cheat sheet" during the exam.
- You may not consult classmates, electronic devices or resources other than the cheat sheet during the exam.
- Questions of clarification should be asked directly to an instructor or TA.

| Question | Points |
|----------|--------|
| 1        |        |
|          | /24    |
| 2        |        |
|          | /30    |
| 3        |        |
|          | /10    |
| 4        |        |
|          | /10    |
| 5        |        |
|          | /8     |
| 6        |        |
|          | /8     |
| 7        |        |
|          | /10    |
| Total    |        |
|          | /100   |

**1. (24 pts, 3 pts each)** For each description on the left, select the best matching term on the right, each term is used only once but some will not be used.

fork makes a copy of the current process

CPU bound a process that frequently exceeds its time slice

semaphore mechanism that can provide both locking and signaling

TCB created as a result of calling pthread\_create()

Kernel mode system calls are implemented using traps because they require this

pipe a queue for inter-process communication

locality reason caching can improve performance

TestAndSet machine instruction that can be used to implement mutex locks

- a. CPU bound
- b. TestAndSet
- c. Heap
- d. TCB
- e. kernel mode
- f. page out
- g. pipe
- h. semaphore
- i. fork()
- j. swap
- k. MMU
- I. locality

| F In RR scheduling, best performance is when the time quantum is small with respect to context-switch time.                                 |
|---|
| F In the Linux Completely Fair Scheduler processes with low nice values always run before those with high nice values.                      |
| T LRU never experiences Belady's Anomaly.   |
| T The purpose of priority boost is to prevent starvation.   |
| F A context switch is preformed when changing a process from the blocked state to ready state.  |
| F Thrashing is a result of insufficient multiprogramming.   |
| T A segment can be shared by multiple processes.  |
| F If a page's contents in memory are different than they are on disk, the valid bit must be set to invalid.                                 |
| F An illegal memory access results in a page fault.   |
| F A translation lookaside buffer (TLB) is used to search for free space in physical memory.   |
| T Paging has the problem of internal fragmentation.   |
| T Tick based Operating Systems depend on a hardware interrupt to keep track of time.  |
| T The bounded-buffer problem can be solved using only semaphores to control concurrency.  |
| T A binary semaphore is equivalent to a mutex lock.   |
| T A thread requesting a resource never causes a deadlock as long as that thread does not currently have any other resources assigned to it. |
|   |

2. (30 pts, 2 pts each) Which of the following statements are true? Write T or F for true or false.

**3. (10 pts)** Consider the following set of jobs, with arrival times and the length of CPU bursts given in milliseconds.

|   | Arrival Time | CPU Burst |
|---|--------------|-----------|
| А | 0            | 10        |
| В | 1            | 12        |
| С | 4            | 10        |
| D | 5            | 2         |

Create Gantt charts and compute the average response time for each of the following scheduling algorithms. Show your calculations.

a) SJF (without preemption)

Gantt chart:

Average response time is \_8.5\_\_\_ (show calculation below)

$$(0 + (10 - 5) + (12 - 4) + (22 - 1))/4 = 8.5$$

b) STCF

Gantt chart:

Average response time is \_\_\_7.25\_\_\_\_ (show calculation below)

$$(0 + (22-1) + (12-4) + (5-5))/4 = 7.25$$

**4. (10 pts)** Given the reference string of page accesses: 4 3 2 1 4 3 5 1 and a system with 3 page frames, how many page faults result when using the following page replacement policies? Show your work.

## a) FIFO

Number of page faults is \_\_8\_\_ (show work below)

| Page # | 4 | 3 | 2 | 1 | 4 | 3 | 5 | 1 |
|--------|---|---|---|---|---|---|---|---|
| H/M?   | M | M | M | M | M | M | M | М |
| Frames |   | 4 | 4 | 4 | 3 | 2 | 1 | 5 |
|        |   |   | 3 | 3 | 2 | 1 | 4 | 4 |
|        |   |   |   | 2 | 1 | 4 | 3 | 3 |

## b) OPT

Number of page faults is \_\_5\_\_ (show work below)

| Page # | 4 | 3 | 2 | 1 | 4 | 3 | 5 | 1 |
|--------|---|---|---|---|---|---|---|---|
| H/M?   | M | M | M | M | Н | Н | M | Н |
| Frames |   | 4 | 4 | 4 | 4 | 4 | 4 | 5 |
|        |   |   | 3 | 3 | 3 | 3 | 3 | 3 |
|        |   |   |   | 2 | 1 | 1 | 1 | 1 |

5. (8 pts) Below is an excerpt from xv6 scheduler code encountered in Project 1.

```
void
sched(void)
{
    // ...
    struct proc *p = myproc();
    // ...
    swtch(&p->context, &mycpu()->context);
    // ...
}
```

Why does this function call swtch()? What does swtch() do? Answer in 2 to 3 sentences.

sched() calls swtch() to perform a context switch from the context of the kernel code that interrupted the currently running user process (I.e., the process returned by myproc()) into the context of the scheduler. The purpose is for the scheduler to possibly context switch to a different process. swtch() performs the actions of a context switch: saving the current CPU registers to the current user process' context and loading CPU registers from the saved context of the next user process.

**6. (8 pts)** Given a virtual address space of 1,024 pages and a 4,096 byte page size, how many bits are required to store a virtual address. Show calculations.

1024 pages => page number takes 10 bits
Page size = 4096 bytes => page offset takes 12 bits
Hence, 10+12=22 bits are required to store a virtual address.

Or, Log (1024 \* 4096) = Log(2^{10} \* 2^{12}) = 22.

**7. (10 pts)** Consider two threads, a ping\_thread and pong\_thread, that output the words "ping" and "pong" respectively. The required output of the program is as follows (it is required that ping always goes first):

ping pong ping pong

Provide a solution for ping\_thread and pong\_thread using only semaphores for concurrency control. Declare and initialize all variables. Exact pthread syntax is not required.

```
#include <pthread.h>
     // declare any required variables here
sem t ping sem, pong sem;
void init() { // perform initialization here
sem init(&ping sem, 1);
sem init(&pong sem, 0);
}
void* ping thread(void *arg) {
    while (1) { // complete the ping thread
           sem wait(&ping sem);
printf("ping\n");
sem post(&pong sem);
    }
}
void* pong thread(void *arg) {
   while (1) { // complete the pong thread
sem wait(&pong sem);
printf("pong\n");
sem post(&ping sem);
   }
}
```