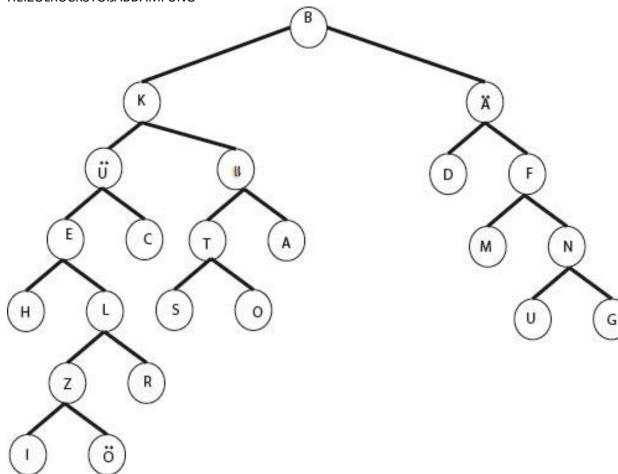
a. HEIZÖLRÜCKSTOßABDÄMFUNG



- c. Citation: hw-questions discord channel and project 5 from java 2 with Prof. Duke
- 2. A game with n terminals ends after (n-1) moves so to win the game you should go first if n is even and second if n is odd.
 - a. A game with two open terminals in the same section is the winning state. n is how many terminals there are that can be connected to another terminal. The base case is n=2. Using 2 for n in (n-1) returns 1 meaning that the game will end after 1 move which it will so the base case is true. For the induction hypothesis I assume that the formula (n-1) will hold for every integer k such that $2 \le k \le n$. For the induction step I'll show that it holds for k+1. After the first player draws their curve, the board is split into two sections with r and s many terminals in each. r+s=n. The formula should hold for both of those sections. The r section has a formula (r-1). The s section has a formula (s-1). Adding those two sections together plus 1 for the move that is used to split the first section is 1+(r-1)+(s-1) which equals (r+s)+1-2 which equals (r+s)-1. r+s is n so the final step says (n-1)=(n-1). So the induction step holds
 - b. Citation: classwork from 1/18/22

3.

b.

a. Citation: https://proofwiki.org/wiki/Fibonacci Number with Negative Index

- i. The base case is n = 1. Using 1 in $f(n) = (-1)^n(n+1)^*f(n)$ is equal to $f(1) = (-1)^n(1+1)^*f(1)$. We know that f(1) is 1 because that is one of the base cases for the Fibonacci sequence. So $1 = (-1)^n(2)^*1$ which is the same as $1 = 1^*1$ so 1 = 1 which is true and the base case is valid.
- ii. The induction hypothesis will be based on the backward Fibonacci sequence f(-n) = f(n+2) f(n+1). If (k+1) is used in place of n, then by the backward Fibonacci sequence it will be equal to f(-(k+1) + 2) f(-(k+1) + 1). Which simplifies to f(-(k-1)) f(-k). The formula $f(-n) = (-1)^n(n+1)^n$ (n) should hold for all of the other elements of the sequence and by extension to the values that are created by using (k+1) in place of f(-n). So the induction hypothesis is that f(-(k-1)) equals $f(-1)^n(k-1)^n$ ($f(k-1)^n$) which is the same as $f(-1)^n(k)^n$ and $f(-k)^n$ equals $f(-1)^n(k+1)^n$ ($f(k)^n$).
- iii. The induction step is that $f(-(k+1)) = (-1)^k(k+1+1)^*f(k+1) = f(-(k-1)) f(-k)$. As shown before in the induction hypothesis stage f(-(k-1)) f(-k) can be simplified to $(-1)^k(k)^*f(k-1) (-1)^k(k+1)^*f(k)$. The second part of that equation can be changed to be $(-1)^k$ instead to make it be added to the first part instead of subtracted. Then $(-1)^k$ can be factored out of both sides leaving $(-1)^k(k)^*(f(k) + f(k-1))$. In relation to f(k+1), f(k) + f(k-1) is the formula for the Fibonacci sequence. So f(k) + f(k-1) can be replaced with f(k+1). This leaves f(k) + f(k+1) which is the same as f(k) + f(k+1) + f(k+1) because adding two doesn't change the outcome at all. So the induction step is valid and $f(-n) = (-1)^k(n+1)^*f(n)$ is valid for every positive integer f(k) + f(k+1) + f(k+

b.

- i. The base case is n = 1. This is proven by $3^0 = 1$
- ii. Assume that the k and k+1 can be written as the sum of positive or negative three to some power i. Say $k = a1*3^{(i1)} + a2*3^{(i2)} + ... + am*3^{(im)}$ where a{x} is either 1 or -1 and i{x} are distinct nonnegative integers. In the case where i1 is 0, n is one less than n+1. So, n can be expressed as $n = (n + 1 3^0)$. Otherwise, n can be written as $n = (n + 1 + 3^{i1})$ which is recursive. So, every nonzero integer can be written in the form of a sum of +-3^i where the exponents i are distinct non negative integers