Control Flow

Processors do only one thing:

- From startup to shutdown, a CPU simply reads and executes (interprets) a sequence of instructions, one at a time
- This sequence is the CPU's control flow (or flow of control)

Physical control flow

Altering the Control Flow

- Up to now: two mechanisms for changing control flow:
 - Jumps and branches
 - Call and return

React to changes in *program state*

- Insufficient for a useful system: Difficult to react to changes in system state
 - Data arrives from a disk or a network adapter
 - Instruction divides by zero
 - User hits Ctrl-C at the keyboard
 - System timer expires
- System needs mechanisms for "exceptional control flow"



Exceptional Control Flow

- Exists at all levels of a computer system
- Low level mechanisms
 - 1. Exceptions
 - Change in control flow in response to a system event (i.e., change in system state)
 - Implemented using combination of hardware and OS software

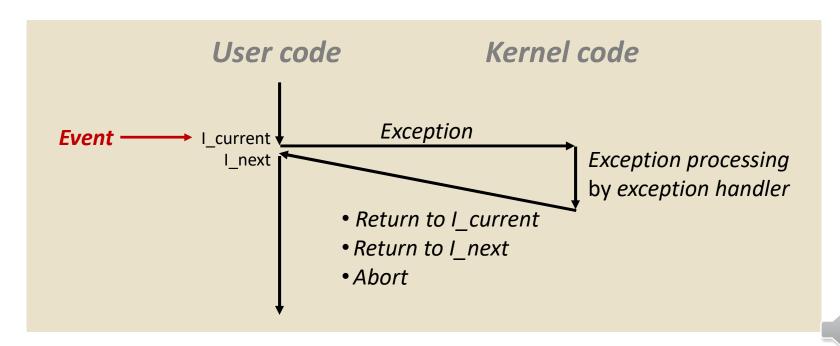
Higher level mechanisms

- 2. Process context switch
 - Implemented by OS software and hardware timer
- 3. Signals
 - Implemented by OS software
- 4. Nonlocal jumps: setjmp() and longjmp()
 - Implemented by C runtime library

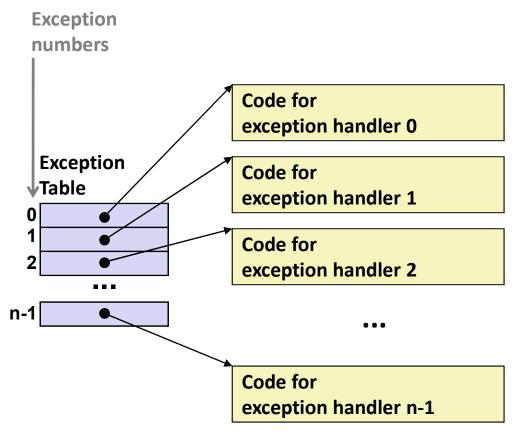


Exceptions

- An exception is a transfer of control to the OS kernel in response to some event (i.e., change in processor state)
 - Kernel is the memory-resident part of the OS
 - Examples of events: Divide by 0, arithmetic overflow, page fault, I/O request completes, typing Ctrl-C



Exception Tables



- Each type of event has a unique exception number k
- k = index into exception table (a.k.a. interrupt vector)
- Handler k is called each time exception k occurs

Asynchronous Exceptions (Interrupts)

Caused by events external to the processor

- Indicated by setting the processor's interrupt pin
- Handler returns to "next" instruction

Examples:

- Timer interrupt
 - Every few ms, an external timer chip triggers an interrupt
 - Used by the kernel to take back control from user programs
- I/O interrupt from external device
 - Hitting Ctrl-C at the keyboard
 - Arrival of a packet from a network
 - Arrival of data from a disk

Synchronous Exceptions

- Caused by events that occur as a result of executing an instruction:
 - Traps
 - Intentional
 - Examples: system calls, breakpoint traps, special instructions
 - Returns control to "next" instruction

Faults

- Unintentional but possibly recoverable
- Examples: page faults (recoverable), protection faults (unrecoverable), floating point exceptions
- Either re-executes faulting ("current") instruction or aborts

Aborts

- Unintentional and unrecoverable
- Examples: illegal instruction, parity error, machine check
- Aborts current program



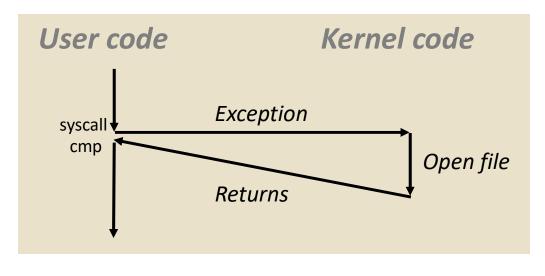
System Calls

- Each x86-64 system call has a unique ID number
- **Examples:**

Number	Name	Description
0	read	Read file
1	write	Write file
2	open	Open file
3	close	Close file
4	stat	Get info about file
57	fork	Create process
59	execve	Execute a program
60	_exit	Terminate process
62	kill	Send signal to process

System Call Example: Opening File

- User calls: open (filename, options)
- Calls __open function, which invokes system call instruction syscall



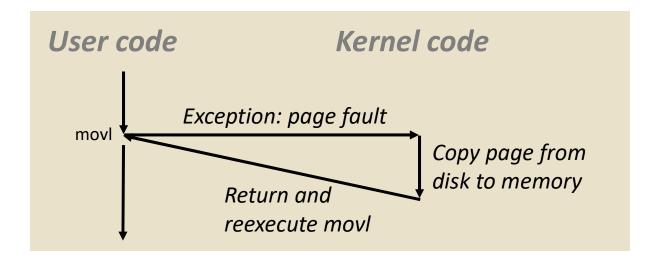
- %rax contains syscall number
- Other arguments in %rdi, %rsi, %rdx, %r10, %r8, %r9
- Return value in %rax
- Negative value is an error corresponding to negative errno

Fault Example: Page Fault

- User writes to memory location
- That portion (page) of user's memory is currently on disk

```
int a[1000];
main ()
{
    a[500] = 13;
}
```

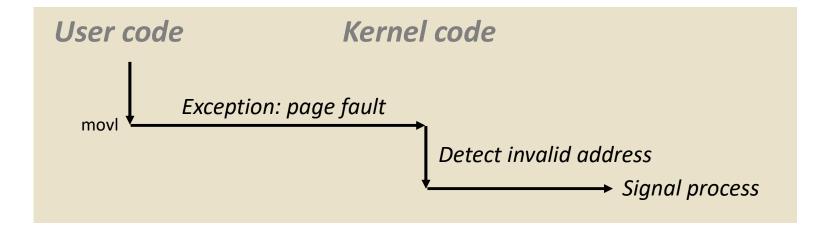
```
80483b7: c7 05 10 9d 04 08 0d movl $0xd,0x8049d10
```



Fault Example: Invalid Memory Reference

```
int a[1000];
main ()
{
    a[5000] = 13;
}
```

```
80483b7: c7 05 60 e3 04 08 0d movl $0xd,0x804e360
```



- Sends SIGSEGV signal to user process
- User process exits with "segmentation fault"

