Introduction

The Linux IPC (Inter-process communication) facilities provide a method for multiple processes to communicate with one another. There are several methods of IPC available to Linux C programmers:

Half-duplex UNIX Pipes

Basic Concepts

Simply put, a pipe is a method of connecting the standard output of one process to the standard input of another. Pipes are the eldest of the IPC tools, having been around since the earliest incarnations of the UNIX operating system. They provide a method of one-way communications (hence the term half-duplex) between processes.

This feature is widely used, even on the UNIX command line (in the shell).

ls | sort | lp

The above sets up a pipeline, taking the output of **ls** as the input of **sort**, and the output of sort as the input of **lp**. The data is running through a half duplex pipe, traveling (visually) left to right through the pipeline.

When a process creates a pipe, the kernel sets up **two file descriptors** for use by the pipe. One descriptor is used to allow a path of input into the pipe (write), while the other is used to obtain data from the pipe (read).

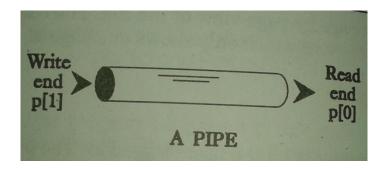
Creating Pipes in C

To create a simple pipe with C, we make use of the pipe() system call. It takes a single argument, which is an array of two integers, and if successful, the array will contain two new file descriptors to be used for the pipeline.

The first integer in the array (element 0) is set up and opened for reading, while the second integer (element 1) is set up and opened for writing. Visually speaking, the output of fd1 becomes the input for fd0. Once again, all data traveling through the pipe moves through the kernel.

```
See pipe0.c
main()
{
    int p[2];
    retval=pipe(p);
    printf("p[0] is %d and p[1] is %d\n ", p[0],p[1]);
}
```

Figure 1



We pass the starting address of array to pipe and values of two descriptors are returned in the array and those vales are printed. p[0] is read end and p[1] is write end.

Following program tells you how many pipes can be created at one time.

/* How many pipes can be created? */ See pipe3.c #include<stdio.h>

Accessing Pipe

To access a pipe directly, the same system calls that are used for low-level file I/O can be used .

To send data to the pipe, we use the write() system call, and to retrieve data from the pipe, we use the read() system call.

Reading Pipe

Reading pipe end can be done using read command which we have used already in files.

Syntax : size_t read(int fd, char *buffer , size_t bytes);

Example: read(fd1, buf, 100); include file: #include <sys/types.h>

Writing Pipe

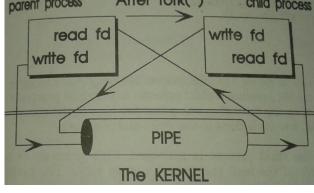
Writing pipe end can be done using write command which we have used already in files.

size_t write(int fd, char *buffer, size_t bytes);

Example: write(fd1,buf, 100); include file: #include <sys/types.h>

```
see pipe1.c
/* usage of pipe read, write with in same process */
#include<stdio.h>
#include<sys/types.h>
#define MSGSIZE 16
main()
        char* msg1="Hello World 1";
        char* msg2="Hello World 2";
        char* msg3="Hello World 3";
        char inbuf[MSGSIZE];
        int p[2],i;
        pipe(p);
        write(p[1],msg1,MSGSIZE);
        write(p[1] ,msg2 ,MSGSIZE);
        write(p[1] ,msg3 ,MSGSIZE);
              for (i=0;i<3;i++)
               read(p[0],inbuf,MSGSIZE);
               printf("%s\n",inbuf);
       exit(0);
}
Example:
    Fork and Pipe
/* usage of close(),read() and write(), fork() in pipe()*/
#include<stdlib.h>
#include<stdio.h>
#define MSGSIZE 16
main()
        char* msg1="Hello World 1";
        char inbuf[MSGSIZE];
        int p[2],pid;
       pipe(p);
       pid=fork();
       if (pid>0)
               printf("Parent : p[0]=%d and p[1]=%d\n", p[0],p[1]);
       if (pid==0)
               printf("Child:p[0]=%d and p[1]=%d\n", p[0],p[1]);
```

```
exit(0);
}
Example: See pipe???
/* parent process write data to pie and child reads from pipe, fork() used to create parent and child*/
#include<stdio.h>
#define MSGSIZE 16
main()
{
       int p[2], pid;
       char* msg1="Hello World 1";
       char inbuf[MSGSIZE];
       pipe(p);
       pid=fork();
       if(pid>0)
       {
               printf("In Praent \n");
               write(p[1],msg1,MSGSIZE);
       else
               printf("\nIn child \n");
               read(p[0],inbuf,MSGSIZE);
               printf("\n\% s\n",inbuf);
       }
}
       figure 2 and fig 3
                 After fork()
parent process
                                     child process
       read fd
                                write fd
```



By doing as above in the program, it is possible for the both parent and child to write to the write end of pipe (see fig 2). This will not be a communication, because we considered it as half duplex communication, i.e. parent writes first to write end of pipe and next child reads from read end of pipe or child writes and then parent reads. To ensure this we have take some precautions of closing and opening the write end, read end of pipe.

Parent has to close read end before writing to write end, and child has to close write end before reading from read end (see Figure 3).

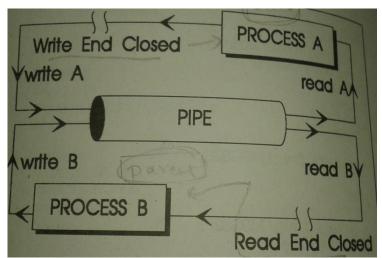


Figure 3 We need to use close()

(If the parent wants to receive data from the child, it should close fd1, and the child should close fd0.

If the parent wants to send data to the child, it should close fd0, and the child should close fd1. Since descriptors are shared between the parent and child, we should always be sure to close the end of pipe we aren't concerned with.)

```
int p[2];
pipe(p);
we can use close(p[0]) to close read end
close(p[1]); to close write end.
See pipe2.c
/* usage of close (), read (), fork () and write () in pipe ()*/
#include<stdio.h>
#define MSGSIZE 16
main()
{
        char* msg1="Hello World 1";
        char inbuf[MSGSIZE];
        int p[2],pid;
       pipe(p);
       pid=fork();
       if (pid>0)
              close(p[0]);
               write(p[1],msg1,MSGSIZE);
       }
       if (pid==0)
              close(p[1]);
```

```
read(p[0] ,inbuf ,MSGSIZE);
               printf(" In child %s",inbuf);
        exit(0);
}
Sort numbers in a array using parent child process created using fork() and use pipe() for
communication between parent and child process.
/* Sorting array of numbers using pipe */
#include <stdio.h>
#define MAXSIZE 10
void main()
{
  int pid,array[MAXSIZE],a[10];
  int i, j, num, temp;
  int p[2],o;
       pipe(p);
  printf("Enter the value of num \n");
  scanf("%d", &num);
  printf("Enter the elements one by one \n");
  for (i = 0; i < num; i++)
     scanf("%d", &array[i]);
  printf("Input array is \n");
  for (i = 0; i < num; i++)
     printf("%d\n", array[i]);
  pid=fork();
  /* Bubble sorting begins */
if (pid==0)
{ printf("Child -Sorting \n");
  for (i = 0; i < num; i++)
     for (j = 0; j < (num - i - 1); j++)
```

if (array[j] > array[j + 1])

array[j] = array[j + 1];array[j + 1] = temp;

temp = array[j];

```
}

for (i = 0; i < num; i++)
{
    printf("Child %d\n", array[i]);
}
close(p[0]);
write(p[1],array,sizeof(int)*num);
printf(" child ends \n");
}
else
{
    close(p[1]);
        read(p[0],a,sizeof(int)*num);
    printf("Sorted array is...\n");

    for (i = 0; i < num; i++)
{
        read(p[0],a[i],sizeof(int));
        printf("Parent %d\n", a[i]);
}
}
</pre>
```

Network Programming

Unix Pipes

A pipe is used for one-way communication of a stream of bytes. The command to create a pipe is *pipe()*, which takes an array of two integers. It fills in the array with two file descriptors that can be used for low-level I/O.

We can use two pipes to make Duplex communication.

Figure 4

```
#include<stdio.h>
#include<stdlib.h>
#include<sys/types.h>
main()
{
    int pp[2],pc[2], pid;
```

```
char msg1[20];
char msg2[20];
char msg3[20];
pipe(pp);
pipe(pc);
pid=fork();
if(pid==0)
{
       close(pc[0]);
       write(pc[1],"Hi Dad",6);
       msg2[6]='\0';
       close(pp[1]);
       read(pp[0],msg2,11);
       msg2[11]='\0';
       printf("chiild reading-%s\n",msg2);
       close(pc[0]);
       write(pc[1],"Thank you Dad",13);
}
else
{
       close(pc[1]);
       read(pc[0],msg1,6);
       msg1[6]='\0';
       printf("Parent \ Reading-\%s\n",msg1);
```

```
close(pp[0]);
    write(pp[1],"Hi My Child",11);
    close(pc[1]);
    read(pc[0],msg3,13);
    msg3[13]="\0';
    printf("Parent reading-%s\n",msg3);
}
```

EXAMPLE 2:

```
#include<stdio.h>
#include<stdlib.h>
#include<sys/types.h>
#include<unistd.h>
main(){
int pp[2],pc[2], pid;
int a[10],b[10],c[10],d[10];
 pipe(pp);
 pipe(pc);
 pid=fork();
 if(pid==0) {
 sleep(5);
 close(pc[1]); //pc - write close
 read(pc[0],b,5*sizeof(int));
 printf(" CHILD - received array b \n");
 for(int j=0; j<5; j++){
   printf(" %d \t",b[j]);
   c[j]=b[j]*b[j];
  }
 close(pp[0]);
 write(pp[1],c,5*sizeof(int));
```

```
else {
    printf(" Enter elements to array \n");
    for(int i=0;i<5;i++){
        scanf("%d",&a[i]);
    }
    close(pc[0]); // pc -read close
    write(pc[1],a,5*sizeof(int));
    close(pp[1]);
    read(pp[0],d,5*sizeof(int));
    printf("\n PARENT - Received squared arrayd \n");
    for(int k=0;k<5;k++){
        printf(" %d \t",d[k]);
    }
    printf("\n");
    wait(0);
}</pre>
```

NAMED PIPES (FIFOs)

In unnamed pipes (half duplex discussed previously) has some drawbacks, such as-

- 1. These pipes (half duplex) can be used only in process which have **common ancestry like parent** and child processes.
- 2. These **pipes are not permanent**, if a Process which creates them terminated then pipes also get terminated.

To overcome these Named Pipes are used.

These Named pipes are Permanent fixtures in Unix. Unix gives them size, owner, access permissions like files. It cane be opened ,closed or deleted.

Experiment with Named Pipe

We can create named pipe using mknod() or mkfifo() sytem call from command prompt. We use the command ls - l > cat > a1.txt to display list of files/folders details and it pipelined to cat which in turn store in a1.txt. Similar kind of activity we can do using pipes and displaying ls-l output to the screen instead of text file.

Type the following series of commands in terminal(first)-

\$ mknod testpipe p

mknod creates pipe named testpipe

You can check it

\$ ls -l testpipe*

Make the cat command which takes input from pipe as background process as below-

Open two terminals.

In first terminal type as below-

\$ cat < testpipe &

By making it background process, cat process will not hang if there is no input.

Now in the second terminal execute following command-

\$ ls -l >testpipe

This will send the display of files/folder listing to pipe and in-turn in second terminal(receives output) which was running **cat** in background receives output of **ls -l** through **testpipe** and displays output on screen.

Example: p217.c and p218.c

Initially create pipe testpipe.

\$ mknod testpipe p

Terminal 1

```
$ cc p217.c -o p217
$ cc p218.c -o p218
```

Terminal 1(client)

\$./p218 hai hello

Terminal 2(Server)

\$./p217

Message Received hai Message Received hello

```
#include<fcntl.h>
#include<stdio.h>
#include <strings.h>
#include<errno.h>
#define MSGSIZ 63

main()
{
    int fd;
    char msgbuf[MSGSIZ+1];
    int n;
```

```
bzero(msgbuf, MSGSIZ);
       if((fd=open("testpipe",O_RDWR))<0)
              perror("pipe open failed in p217\n");
       for(;;)
       {
             if((n=read(fd,msgbuf,MSGSIZ+1))>0)
                     printf(" Message Received %s \n",msgbuf);
}
  -----p218.c-----
#include<fcntl.h>
#include<stdio.h>
#include<stdlib.h>
#include<errno.h>
#include <strings.h>
#include<string.h>
#define MSGSIZ 63
main(argc,argv)
int argc;
char* argv[];
{
       int fd,j,nwrite;
       char msgbuf[MSGSIZ];
    bzero(msgbuf, MSGSIZ);
       if(argc<2)
             printf(" Type as - Filename msg1 msg2 ..\n");
              exit(1);
       if((fd=open("testpipe",O_WRONLY))<0)
             printf("pipe open failed in p218\n");
              exit(1);
       printf("FD:%d",fd);
       for(j=1;j<argc;j++)
              strcpy(msgbuf,argv[j]);
              if((nwrite=write(fd,msgbuf,MSGSIZ+1))<0)
                    printf("pipe wirte failed in p218 \n");
exit(0);
```

NAME MKFIFO() top

mkfifo - make a FIFO special file (a named pipe)

SYNOPSIS

top

#include <sys/types.h>
#include <sys/stat.h>

int mkfifo(const char *pathname, mode_t mode);

DESCRIPTION

top

mkfifo() makes a FIFO special file with name *pathname*. *mode* specifies the FIFO's permissions. It is modified by the process's **umask** in the usual way: the permissions of the created file are (*mode* & ~umask).

A FIFO special file is similar to a pipe, except that it is created in a different way. Instead of being an anonymous communications channel, a FIFO special file is entered into the file system by calling **mkfifo**().

Once you have created a FIFO special file in this way, any process can open it for reading or writing, in the same way as an ordinary file. However, it has to be open at both ends simultaneously before you can proceed to do any input or output operations on it. Opening a FIFO for reading normally blocks until some other process opens the same FIFO for writing, and vice versa.

```
RETURN VALUE top
```

On success **mkfifo**() returns 0. In the case of an error, -1 is returned (in which case, <u>errno</u> is set appropriately).

The *mode* argument is used to set the file's permissions.

The **file mode**, stored in the **st_mode** field of the file attributes, contains two kinds of information: the **file type** code, and the **access permission** bits. This section discusses only the access permission bits, which control who can read or write the file.

The file type must be one of **S_IFREG**, **S_IFCHR**, **S_IFBLK**, **S_IFIFO**, or **S_IFSOCK** to specify a regular file (which will be created empty), character special file, block special file, FIFO (named pipe), or UNIX domain socket, respectively. (Zero file type is equivalent to

```
type S_IFREG.)
```

Access permissions are the permissions bits like -0664 which indicates rwx to owner ;rwx to group and rw- to others. (See file attributes document)

RETURN VALUE top

// Creating pipe using mkfifo() and opening pipe for writing and writing some messgae , in another program n_pipe2.c we open pipe for reading and read the message and dipslay.

#include<fcntl.h>

#include<stdio.h>

#include<stdlib.h>

```
int main()
 int fd,fi,nwrite;
 char msgbuf[64]="Hello";
       fi=mkfifo("vmfifo",S_IRUSR|S_IWUSR|S_IRGRP|S_IROTH);
       if(fi<0)
              perror("vmfifo creation failed");
              exit(1);
       }
       fd=open("vmfifo",O_WRONLY);
       if (fd<0)
              perror("vmfifo open failed");
              exit(1);
       nwrite=write(fd,msgbuf,sizeof(msgbuf));
       if(nwrite<=0)
              perror("message write failed \n");
       close(fd);
return(0);
// In n_pipe1.c named pipe is created using mkfifo() and message is written into pipe, in this
program n_pipe2.c we open pipe for reading and read the message and dipslay.
#include<fcntl.h>
#include<stdio.h>
#include<stdlib.h>
main()
 int fd,fi,nread;
 char msgbuf[64];
//open named pipe created in n_pipe1.c
       fd=open("vmfifo",O_RDWR);
       if (fd<0)
              perror("vmfifo open failed");
              exit(1);
       fd=open("vmfifo",O_RDWR);
       nread=read(fd,msgbuf,sizeof(msgbuf));
```