

Let  $x^*$  be the optimal sol<sup>n</sup> to the LP-relaxation of set-cover problem

We want to round it to some  $\hat{x}$  (Randomly)

High Level Idea:- Interpret the fractional value  $x_j^*$  as Probability that  $x_j$  should be set to 1.

Each subset  $S_j$  is included in our solution w.p.  $x_j^*$  and  $m$  events are independent.

$$x_j = \begin{cases} 1 & \text{if } S_j \text{ in sol}^n \\ 0 & \text{o/w} \end{cases}$$

$$\begin{aligned} E \left[ \sum_{j=1}^m w_j x_j \right] &= \sum_{j=1}^m w_j \Pr[x_j=1] \\ &= \sum_{j=1}^m w_j x_j^* \\ &= Z_{LP}^* \end{aligned}$$

But may NOT give a SET COVER!

$$\begin{aligned} \Pr[e_i \text{ is not covered}] &= \prod_{j: e_i \in S_j} (1 - x_j^*) \\ &\leq \prod_{j: e_i \in S_j} e^{-x_j^*} \quad \left( \because 1-x \leq e^{-x} \right) \\ &= e^{-\sum_{j: e_i \in S_j} x_j^*} \\ &\leq e^{-1} \quad \left( \sum_{j: e_i \in S_j} x_j^* \geq 1 \right) \end{aligned}$$

How the make failure probability smaller?

— Toss a coin for each  $j$   
 $c \ln n$  times instead of just once  
 If it comes heads even once, include  $S_j$  in sol<sup>n</sup>

$$\begin{aligned} \text{Now } \Pr[S_j \text{ not included}] &= (1 - x_j^*)^{c \ln n} \quad (c \geq 2) \end{aligned}$$

$$\begin{aligned} \Pr[e_i \text{ is not covered}] &= \prod_{j: e_i \in S_j} (1 - x_j^*)^{c \ln n} \\ &\leq \prod_{j: e_i \in S_j} e^{-x_j^* (c \ln n)} \\ &= e^{-(c \ln n) \sum_{j: e_i \in S_j} x_j^*} \\ &\leq \frac{1}{n^c} \end{aligned}$$

Thm :- Randomized Rounding Algo is a randomized  $O(\ln n)$  approx that produces set cover w.h.p

Proof :-  $p_j(x_j^*) = \text{prob } S_j \text{ is included in sol}^n$

$$p_j(x_j^*) = 1 - (1 - x_j^*)^{c \ln n}$$

If  $x_j^* \in [0, 1]$  and  $c \ln n \geq 1$

$$\begin{aligned} p_j'(x_j^*) &= c \ln n (1 - x_j^*)^{(c \ln n) - 1} \\ &\leq c \ln n \end{aligned}$$

$$p_j(0) = 0$$

Slope of  $p_j$  is bounded above by  $c \ln n$  on  $[0, 1]$

$$p_j(x_j^*) \leq (c \ln n) x_j^* \text{ on } [0, 1]$$

$$\begin{aligned} E \left[ \sum_{j=1}^m w_j x_j \right] &= \sum_{j=1}^m w_j \Pr[x_j=1] \\ &\leq \sum_{j=1}^m w_j (c \ln n) x_j^* \\ &= (c \ln n) Z_{LP}^* \end{aligned}$$

Not Enough !!

Have to check for set cover

Let  $F$  be the event that solution obtained is feasible set cover  
 $\overline{F}$  is the complement event

$$\Pr[F] \geq 1 - \frac{1}{n^{c-1}}$$

Now

$$\begin{aligned} E \left[ \sum_{j=1}^m w_j x_j \right] &= E \left[ \sum_{j=1}^m w_j x_j \mid F \right] \Pr[F] + E \left[ \sum_{j=1}^m w_j x_j \mid \overline{F} \right] \Pr[\overline{F}] \\ \text{Also } w_j &\geq 0 \quad \forall j \\ \therefore E \left[ \sum_{j=1}^m w_j x_j \mid \overline{F} \right] &\geq 0 \\ \therefore E \left[ \sum_{j=1}^m w_j x_j \mid F \right] &= \frac{1}{\Pr[F]} \left( E \left[ \sum_{j=1}^m w_j x_j \right] - E \left[ \sum_{j=1}^m w_j x_j \mid \overline{F} \right] \Pr[\overline{F}] \right) \\ &\leq \frac{1}{\Pr[F]} \cdot E \left[ \sum_{j=1}^m w_j x_j \right] \\ &\leq \frac{(c \ln n) Z_{LP}^*}{1 - \frac{1}{n^{c-1}}} \\ &\leq 2c (\ln n) Z_{LP}^* \quad \left( \text{for } n \geq 2, c \geq 2 \right) \end{aligned}$$