Tuesday, 28 January 2025

Longe biportite subgraph

Evez graph with e edges contains biportite subgraph with at least $\frac{1}{2}$ e edges.

Proof: expected volue.

Diagonal Ramsey lower bound (Endos, 1947)

For any integer k and n if $\binom{n}{k} 2^{1-\binom{k}{k}} < 1$, then

any edge coloning at K_n contains monochromatic K_k . Hence $R(k,k) > n-\binom{n}{k}2^{1-\binom{k}{2}} > \left(\frac{1}{e\sqrt{k}} + o(1)\right)k2^{\frac{k}{2}}$

Proof: By Union-Bound such coloring exists. There exists a coloning with at most $\binom{n}{k} 2^{1-\binom{k}{k}}$ monochromatic K_k , hence there is a graph with at least $n-\binom{n}{k} 2^{1-\binom{k}{k}}$ vertices without monochromatic K_k .

Dependency graph

Set of events An, ..., An in probability spoce. Graph on [n] is dependency graph if A; is independent of events indexed by [n]\n(i).

Lovorz Local Lemma (symmetric) Let A; be a family of evals o. + IP[A:] < P and

each of them is dependent on at most of others. If $p(dt1)e \in 1$ then there exist our event that

hone of these events occur.

R(k,k) lower bound (Spencer, 1977) If $(\binom{k}{2}\binom{n}{k-2}+1) 2^{1-\binom{k}{2}}e \le 1$ then there exists

monochronotic K_k -free graph on n neutices. In posticular: $R(k,k) > \left(\frac{\sqrt{2}}{e} + o(1)\right) k 2^{\frac{k}{2}}$

Proof: LLL-events are dependent if they share an edge.

Lorosz Local Lema (Endős-Lonasz, 1975)

Let B1, ..., Bn be family at load events.

If there exist real numbers x_1, \dots, x_n s.t.

 $|P|B_{i} \leq x_{i} = \overline{\prod_{i \in N(i)}} (1-x_{i}) \quad \forall i \in [n]$ Then the event avoiding all bod events has probability at least II (1-x;)

Proof: We will show by induction that for each Sc[n]:

IP[B;] n B;] < X; if if S We proceed by induction on 151. ISI=O is trivial.

 $S_1 = S \cap N(i)$ $S_2 = S \setminus S_1$

 $\begin{aligned}
|P_{r}[B; \cap \overline{B}] &= \frac{|P_{r}[B; \cap \overline{B}, \cap \overline{B}, \cap \overline{B}, \cap \overline{B}]}{|P_{r}[n; \overline{B}, \cap \overline{B}, \cap \overline{B}]} &= \frac{|P_{r}[B; \cap \overline{B}, | \overline{B},$

 $\mathbb{P}\left[\bigcap_{j \in S_1} \overline{B}_i \mid \bigcap_{j \in S_2} \overline{B}_j\right] = \mathbb{P}\left[\overline{B}_{j_1} \mid \bigcap_{j \in S_2} \overline{B}_j\right] \cdot \mathbb{P}\left[\overline{B}_{j_2} \mid B_{j_1} \cap \bigcap_{j \in S_2} \overline{B}_j\right] \dots$ $\geq (1-x_1)(1-x_2)-(1-x_{1s_{11}}) \geq \overline{11}(1-x_i)$

Hence $\mathbb{R}[\overline{I}\overline{B}_i] = \mathbb{R}[\overline{B}_1] \cdot \mathbb{R}[\overline{B}_2|\overline{B}_2] \cdots \leq \overline{I}(1-x_i)$