Obfs2

Pluggable Transports documentation series.[*]

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ABSTRACT

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Tor, Bridge, Scary, Obscuration, Censorship, Circumvention, Pluggable Transport

PREAMBLE

This paper is part of a paper documentation series about Pluggable Transports (PTs) [3], how they work and their strengths and weaknesses.

1 INTRODUCTION

During the history of digital networks, we have been confronted more and more with the phenomenon of internet censorship and blocking by governments [4]. So, over the years, more and more circumvention tools were developed and also have been blocked due to deep packet inspections and the detailed analysis of its content. This lead us to, so-called, *Pluggable Transports (PTs)* [3], which help to bypass censorship attemps by transforming the traffic between client and server, in such a ways that it looks like innocent traffic.

In this paper, we will talk about *obfs2*, a protocol obfuscation layer for TCP protocols. We will show how it works, what we can do with it and which strengths and weaknesses it has.

1.1 Outline

TODO

2 OBFS2

Obfs2 is a protocol obfuscation layer for TCP protocols, to keep a third party from telling what protocol is in use based on message contents [2]. It's the continuation of brl's ssh obfuscation protocol [2] [1].

2.1 Overview

The protocol consists of two phases.

- First: The parties establish keys
- Second: The parties exchange superenciphered traffic.

2.2 Notation

Given are two parties: the 'initiator' (INIT), which opens the connection and can mostly be associated with a client, and the 'responder' (RESP), which accepts the connection and can mostly be associated with a server.

We use the following primitives,

- $H(x)$ is SHA256 of x
- $H^n(x)$ is $H(x)$ called iteratively n times
- $E(K,s)$ is the AES-CTR-128 encryption of s using K as key

notation

- x | y is the concatenation of x and y
- UINT32(n) is the 4 byte value of n in big-endian (network) order
- SR(n) is n bytes of strong random data
- WR(n) is n bytes of weaker random data
- "xyz" is the ASCII characters 'x', 'y', and 'z', not NUL-terminated
- s[:n] is the first n bytes of s
- s[n:] is the last n bytes of s

and constants

- MAGIC_VALUE = 0x2BF5CA7E
- SEED_LENGTH = 16
- MAX_PADDING = 8192
- HASH_ITERATIONS = 100000
- KEYLEN (16) is the length of the key used by $E(K,s)$
- IVLEN (16) is the length of the IV used by $E(K,s)$
- HASHLEN (32) is the length of the output of $H()$
- MAC(s, x) = $H(s | x | s)$

according to [2]. A "byte" is an 8-bit octet and we require that HASHLEN >= KEYLEN + IVLEN.

3 CONCLUSIONS

A APPENDIX

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