Minimum Vertex Cover Problem:

1. fractional relaxation for an arbitrary graph G= (V, E)

Answer:

* Given a graph G = (V, E) and the vertex weights = Weights(V)

Minimum vertex cover problem for graph G as an ILP by using a variable xv for each vertex

vertex v, taking 0 or 1 values.

* xv = 0 means v ∉ Cover C, xv = 1 means v ∈ Cover C
* The weight of the solution, which we want to minimize, is Σv ∈V xv\*W(v) such that xu + xv >= 1 for each edge (u, v)
* This gives the Integer Linear Programming as

Minimize Σv ∈V xv\*W(v)

Subject to xu + xv >= 1 for all (u, v) ∈ Edges E

Xv <= 1 for all v ∈ Vertices V

Xv ∈ N for all v ∈ Vertices V

* Relaxing Integer Linear Programming to a linear program

Minimize Σv ∈V xv\*W(v)

Subject to xu + xv >= 1 for all (u, v) ∈ Edges E

Xv <= 1 for all v ∈ Vertices V

Xv >= 0 for all v ∈ Vertices V

* Linear program can be solved in polynomial time
* Suppose x\* is an optimal solution to the linear program, rounding each value to the closest integer

i.e.,

x’v = 1 if x\*v>= ½

x’v = 0 if x\*v < ½

* The set of valid vertex cover C, because of each edge (u, v) has the conditions

x\*u + x\*v >= 1

so at least one of x\*u or x\*v must be at least ½ and at least one of u or v belongs to Cover C

* The cost of cover is at most twice the optimum

1. Dual linear program for an arbitrary graph G= (V, E)

Answer:

• For the dual, the number of variables is equal to the number of vertices in the minimum vertex cover

• Using weak duality, we convert minimum vertex cover to maximum matching in a bipartite graph

• As we know, the LP for minimum vertex cover is

Minimize Σv ∈V xv\*W(v)

Subject to xu + xv >= 1 for all (u, v) ∈ Edges E

Xv >= 0 for all v ∈ Vertices V

• Its dual has one variable y(u,v) for edge (u, v) and it is

Maximize Σ (u, v) ∈V y(u,v)

Subject to Σ u: (u, v) ∈V y(u,v) <= Weights(v) for all v ∈ Vertices V

y(u,v) >= 0 for all (u, v) ∈ Edges E

• Assigning a non-negative charge to each edge such that the total charge

Overall the edges are as large as possible.

• The sum of charges is a lower bound to the weight of the minimum vertex cover

In the weighted graph, G = (V, E) with weights W

1. ILP corresponding to the dual LP (Maximum matching)

Answer:

* Using weak duality, we convert minimum vertex cover to maximum matching in a bipartite graph
* The input is a graph with each edge having a positive weight Wuv
* A maximum weighted bipartite matching is defined as a perfect matching where the sum of values of the edges in the matching have a maximum value.
* The size of such a matching is n2.
* graph is not complete bipartite then missing edges are inserted with the value zero
* The Variable xuv is defined as,

xuv = 0 if the edge (u, v) belongs to the matching

xuv = 1 if the edge (u, v) does not belong to the matching

* ILP for this graph is

Maximum ΣWuv xuv

For all u, Σxuw = 1

For all v, Σxpv = 1

For all u, v xuv >= 0 is integral

1. write in words the graph problem corresponding the dual ILP

* A matching in a [Bipartite Graph](https://www.geeksforgeeks.org/bipartite-graph) is a set of the edges chosen in such a way that no two edges share an endpoint.
* A maximum matching is a matching of maximum size (maximum number of edges).
* In a maximum matching, if any edge is added to it, it is no longer matching.
* There can be more than one maximum matchings for a given Bipartite Graph.
* Given a bipartite graph G = (A ∪ B, E), find an S ⊆ A × B that is a matching and is as large as possible.
* A matching *M* is maximal if and only if there does not exist an augmenting path with respect to *M*.
* Algorithm

**bipartiteMatch (u, visited, assign)**

Input: Starting node, visited list to keep track, assign the list to assign node with another node.

Output: Returns true when a matching for vertex u is possible.

Begin

for all vertex v, which are adjacent with u, do

if v is not visited, then

          mark v as visited

          if v is not assigned, or bipartiteMatch(assign[v], visited, assign) is true, then

            assign[v] := u

            return true

done

   return false

End

**maxMatch (graph)**

Input: The given graph.

Output: The maximum number of the match.

Begin

   initially no vertex is assigned

   count := 0

   for all applicant u in M, do

      make all node as unvisited

      if bipartiteMatch(u, visited, assign), then

         increase count by 1

   done

End

1. Given graph with 5 vertices and edges 1-2, 2-3, 3-4, 4-5, 5-1
2. **ILP for Minimum Weighted Vertex Cover**

2

1 3

5 4

1. Objective: Minimize Σ a ∈ V weight (V) \* xa

Assuming weights of all vertices are equal to 1

x1 + x2 + x3 + x4 + x5

We must add 0/1 constraints for the variables to state that the values come from 0 or 1. So, xi ∈ {0, 1}, ∀ i ∈ V is the additional constraint

1. Subject to:

x1 + x2 >= 1

x2 + x3 >= 1

x3 + x4 >= 1

x4 + x5 >= 1

x1 + x5 >= 1

Such that x1 = x3 = x4 = 1

X2 = x5 = 0

the above can be written in matrix form as where variables = vertices = columns in A

and constraints = edges = rows in A

AX >= b

we get x1 = x2 = x4 = 1 and x3 = x5 = 0 solution is 3

**ILP for Maximum Matching**

1. Objective: Maximize Σ (u, v) ∈V y(u,v)

y1 + y2 + y3 + y4 + y5 = yTB

Subject to Σ u: (u, v) ∈V y(u,v) <= Weights(v) for all v ∈ Vertices V

y(u,v) >= 0 for all (u, v) ∈ Edges E

y1 x1 + x2 <= 1

y2 x2 + x3 <= 1

y3 x3 + x4 <= 1

y4 x4 + x5 <= 1

y5 x1 + x5 <= 1

x1 \* (y1 + y5) + x2 \* (y1 + y2) + x3 \* (y2 + y3) + x4 \* (y3 + y4) + x5 \* (y4 + y5) is Maximized

y1 + y5 >= 1

y1 + y2 >= 1

y2 + y3 >= 1

y3 + y4 >= 1

y4 + y5 >= 1

y2 = y3 = 1 and y1 = y4 = y5 = 0 The matching set is ([1,2], [3,4]) Assuming all edge weights are equal to 1 we get the solution as 2

1. **Optimal Solution for Vertex Cover**

Considering the greedy approach, we get x1 = x2 = x4 = 1 and x3 = x5 = 0 so the sub-optimal cost is 3 whereas optimal x1 = x2 = x4 = 1 and x3 = x5 = 0 cost is 3

Greedy approach <= 2 \* optimal value

3 <= 2\*3

3 <= 6

**Optimal Solution for Maximum matching**

Considering the greedy approach here too, y2 = y3 = 1 and y1 = y4 = y5 = 0 The matching set is ([1,2], [3,4]) Assuming all edge weights are 1 we get 2

Greedy approach <= 2 \* optimal value

2 <= 2\*2

2 <= 4

1. **Fractional Optimal solution for Minimum Vertex Cover**

x1 + x2 >= 1

x2 + x3 >= 1

x3 + x4 >= 1

x4 + x5 >= 1

x1 + x5 >= 1

are the integral constraints

2x1 + 2x2 + 2x4 + 2x3 + 2x5 >= 5

x1 + x2 + x4 + x3 + x5 >= 5/2

1/2

1/2

2

1 3

x1 = 0.5, x2 = 0.5, x3 = 0.5, x4 = 0.5, x5 = 0.5

1/2

1/2

5 4

1/2

**Fractional Optimal solution for Maximum Matching**

y1 + y5 >= 1

y1 + y2 >= 1

y2 + y3 >= 1

y3 + y4 >= 1

y4 + y5 >= 1

are the integral constraints

2y1 + 2y2 + 2y4 + 2y3 + 2y5 >= 5

y1 + y2 + y4 + y3 + y5 >= 5/2

1/2

1/2

2

1 3

1/2

1/2

5 4

1/2

y1 = 0.5, y2 = 0.5, y3 = 0.5, y4 = 0.5, y5 = 0.5

Here, The fractional solutions for minimum vertex cover and its dual (maximum matching) are equal.