

CMSC 474, Game Theory

5. Imperfect-Information Games

Dana Nau

University of Maryland

Motivation

- So far, we've assumed that players in an extensive-form game always know what node they're at
 - Know all prior choices
 - Both theirs and the others'
 - Thus “perfect information” games
- But sometimes players
 - Don't know all the actions the others took or
 - Don't recall all their past actions
- Extend extensive-form game representation to include this

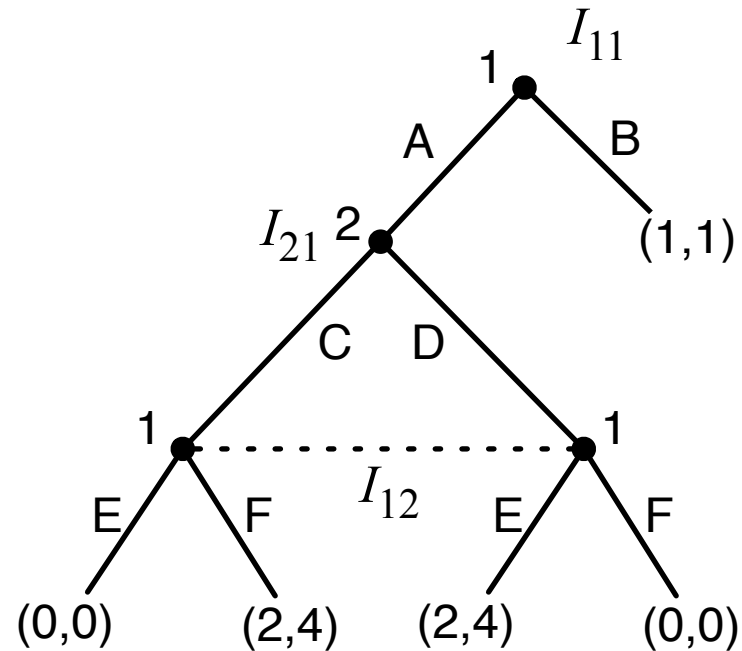
Definition

- **Imperfect-information** game: extensive-form game in which each agent's choice nodes are partitioned into **information sets**
 - An information set = {all choice nodes an agent *might* be at}
- Let $H = \{\text{all nodes where it's agent } i\text{'s move}\}$
 - Agent i 's information sets are I_{i1}, \dots, I_{im} for some m , where
 - $I_{i1} \cup \dots \cup I_{im} = H$
 - $I_{ij} \cap I_{ik} = \emptyset$ for all $j \neq k$

} What is this called?
- If h and h' are in the same information set, they are **indistinguishable** to i
 - So they must have the same set of actions
 - $\chi(h) = \chi(h')$
 - But the actions may have different outcomes
- A perfect-information game is a special case
 - Each I_{ij} contains just one node

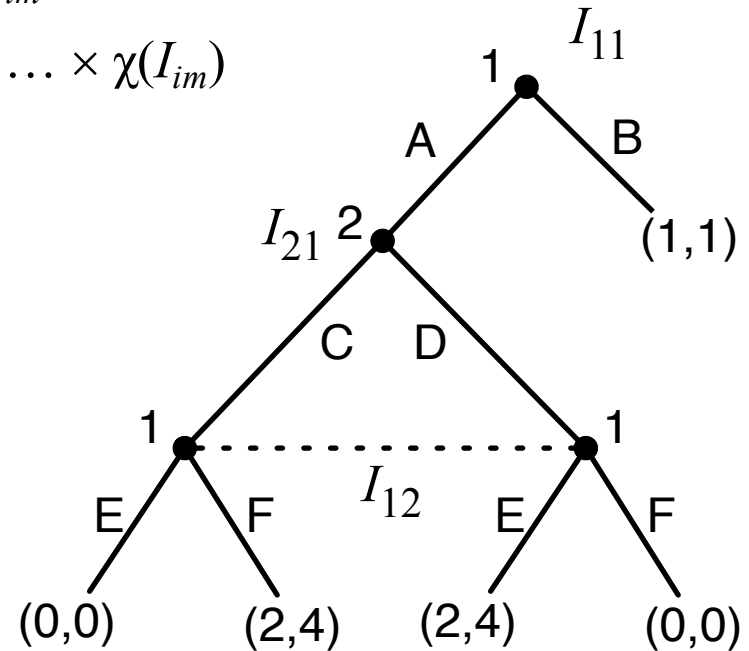
Example

- Agent 1 has two information sets
 - I_{11} and I_{12}
- In I_{12} , agent 1 doesn't know whether agent 2 chose C or D
- Agent 2 has just one information set
 - I_{21}



Strategies

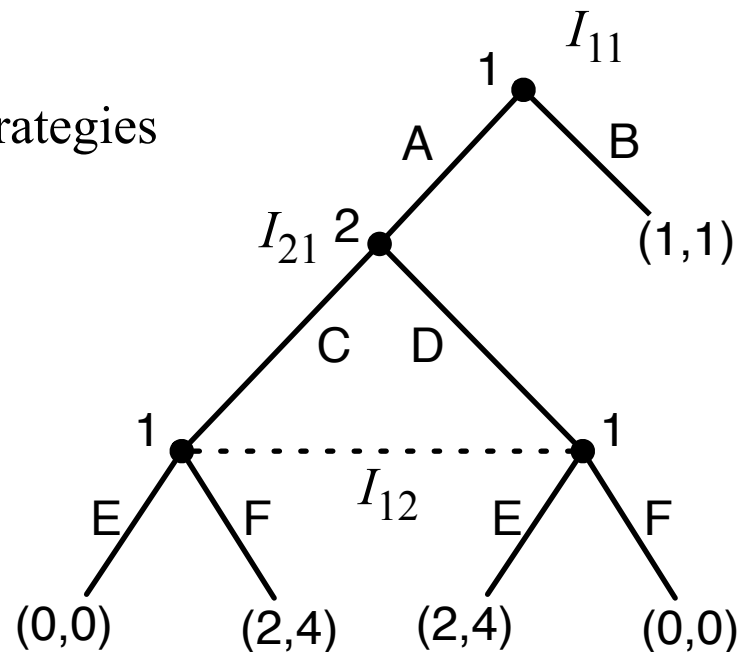
- **Pure strategy** for agent i
 - a function s_i telling what action to take in each of i 's information sets
 - $s_i(I) =$ agent i 's action in information set I
- Suppose i has information sets I_{i1}, \dots, I_{im}
 - $\{\text{all pure strategies for } i\} = \chi(I_{i1}) \times \dots \times \chi(I_{im})$
- Agent 1's pure strategies:
 - $\{A, B\} \times \{E, F\} = \{(A, E), (A, F), (B, E), (B, F)\}$
- Agent 2's pure strategies: $\{C, D\}$



Extensive Form \rightarrow Normal Form

- Can transform any extensive-form imperfect-information game into an equivalent normal-form game
 - Same strategies and same payoffs
 - Thus same Nash equilibria, same Pareto optimal strategy profiles, etc.
- Just like we did it for perfect-information games
 - n -dimensional payoff matrix
 - i 'th dimension \Leftrightarrow agent i 's pure strategies

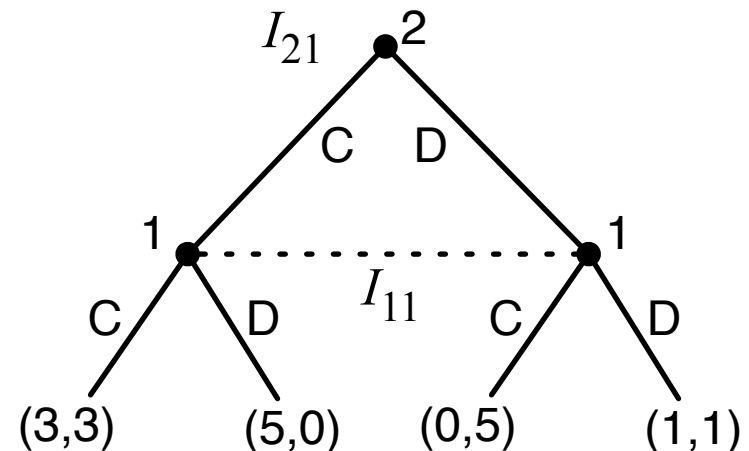
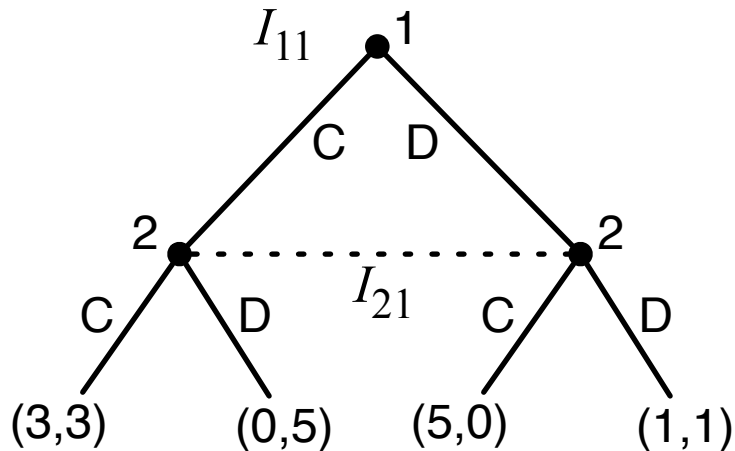
	C	D
(A,E)	0, 0	2, 4
(A,F)	2, 4	0, 0
(B,E)	1, 1	1, 1
(B,F)	1, 1	1, 1



Normal Form \rightarrow Extensive Form

- Can translate any normal-form game into an equivalent extensive-form imperfect-information game
 - n -level game tree, one level for each agent
 - each agent has exactly one information set
- Same strategies, payoffs, Nash equilibria, Pareto optimal strategy profiles, etc.
- Example: Prisoner's Dilemma
 - Two equivalent game trees

	C	D
C	3, 3	0, 5
D	5, 0	1, 1



Different Kinds of Strategies

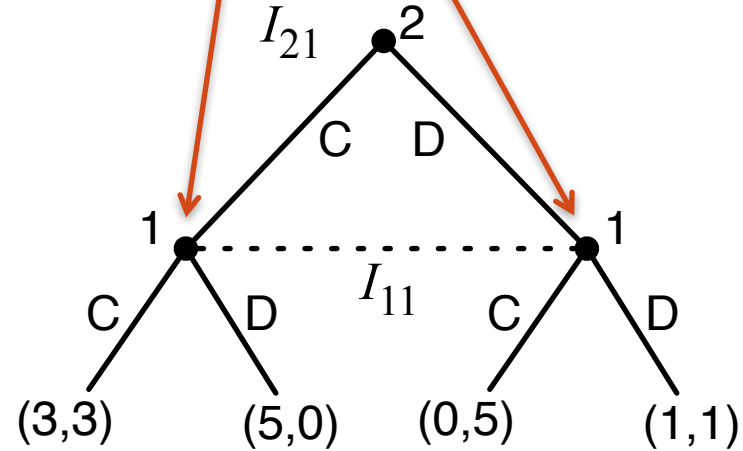
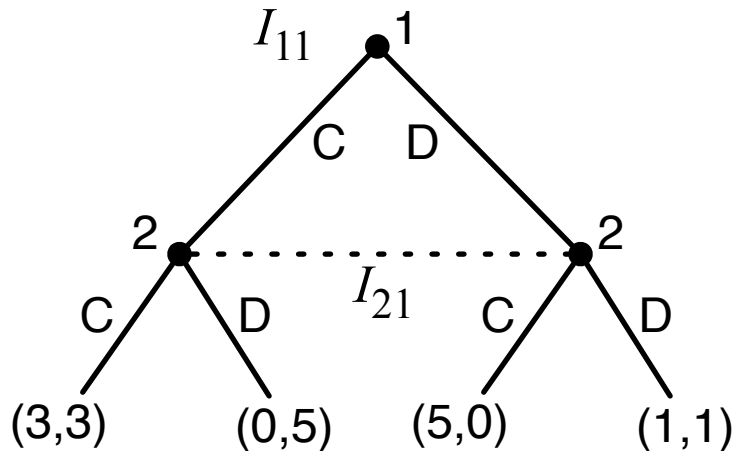
- **Pure strategy:**

- for each information set an agent makes the same move at all nodes in the information set
- E.g., agent 1 choosing D in the normal form game
 - same as choosing D at both of these nodes

	C	D
C	3, 3	0, 5
D	5, 0	1, 1

- **Mixed strategy:**

- Does the agent make the same move at all nodes in the information set?



Different Kinds of Strategies

- New class of strategies: **behavioral strategies**

- Suppose agent i has behavioral strategy s_i
- Each time i is in information set I_{ij} , he/she chooses from the same probability distribution $s_i(I_{ij})$
 - *independently* of i 's choices at other nodes in I_{ij}
- E.g., flip coin every time you're in I_{ij}

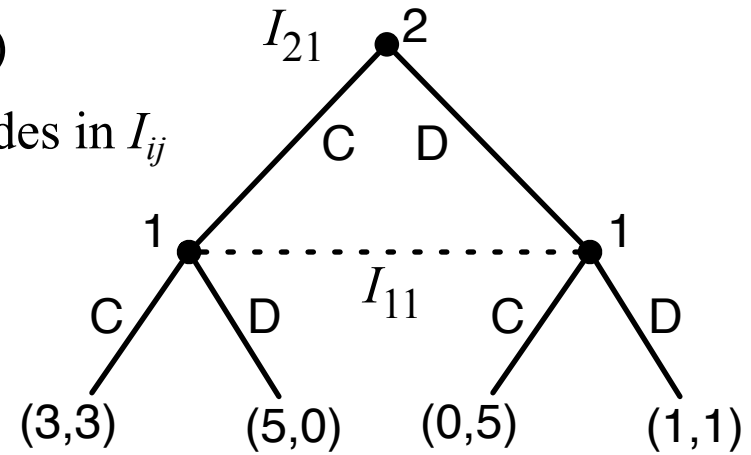
- In some games, every behavioral strategy has an **equivalent** mixed strategy, and vice versa

- Same probability distributions over outcomes

- Example

- Behavioral strategy for Agent 1: $(I_{11}, \{(0.3, C), (0.7, D)\})$
- Mixed strategy is basically the same: $\{(0.3, C), (0.7, D)\}$

- More examples later



Example 2

- A behavioral strategy for Agent 1:

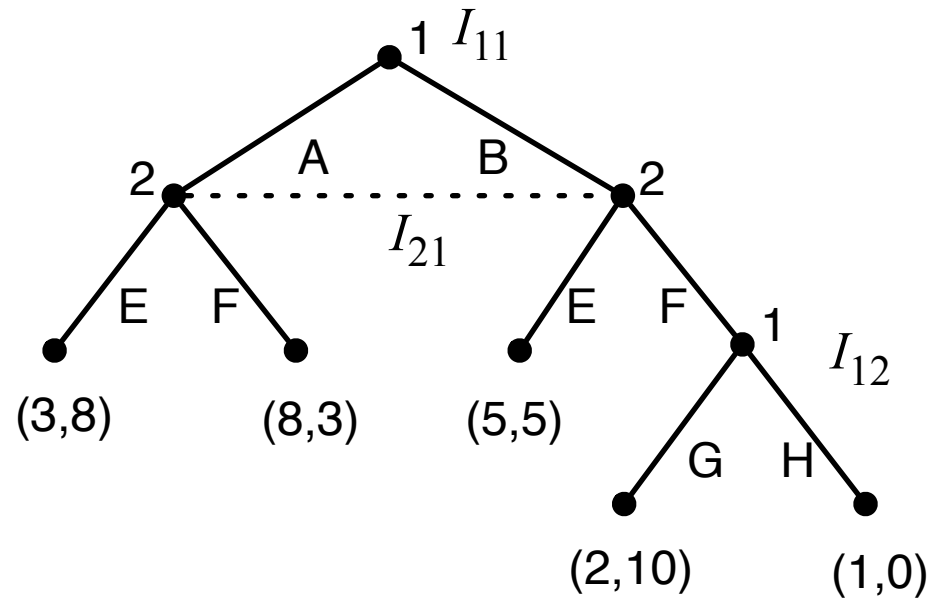
- In I_{11} , $\{(0.4, A), (0.6, B)\}$

- In I_{12} , $\{(0.3, G), (0.7, H)\}$

- An equivalent mixed strategy:

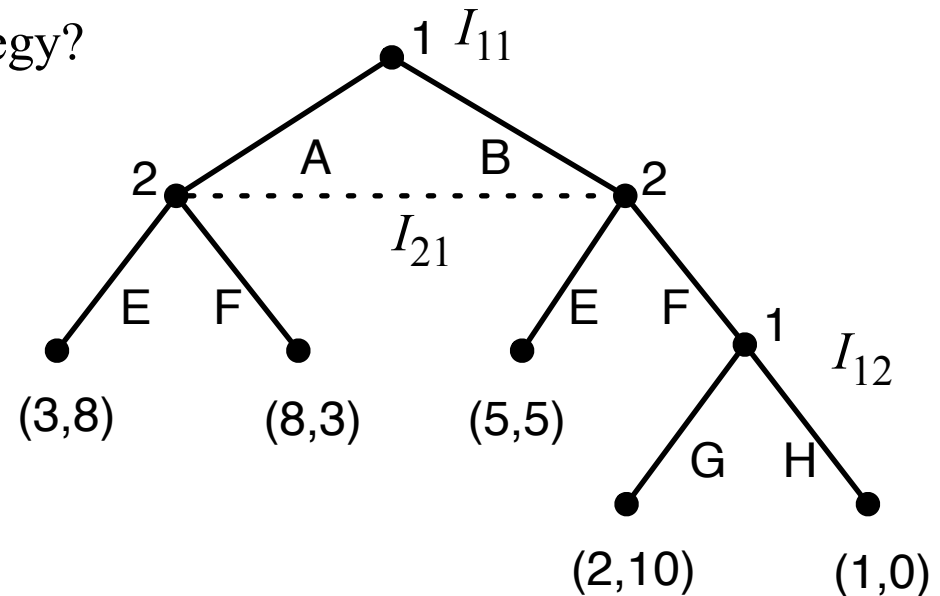
- $\{(0.12, (A,G)),$
 $(0.28, (A,H)),$
 $(0.18, (B,G)),$
 $(0.42, (B,H))\}$

- How did I get those numbers?



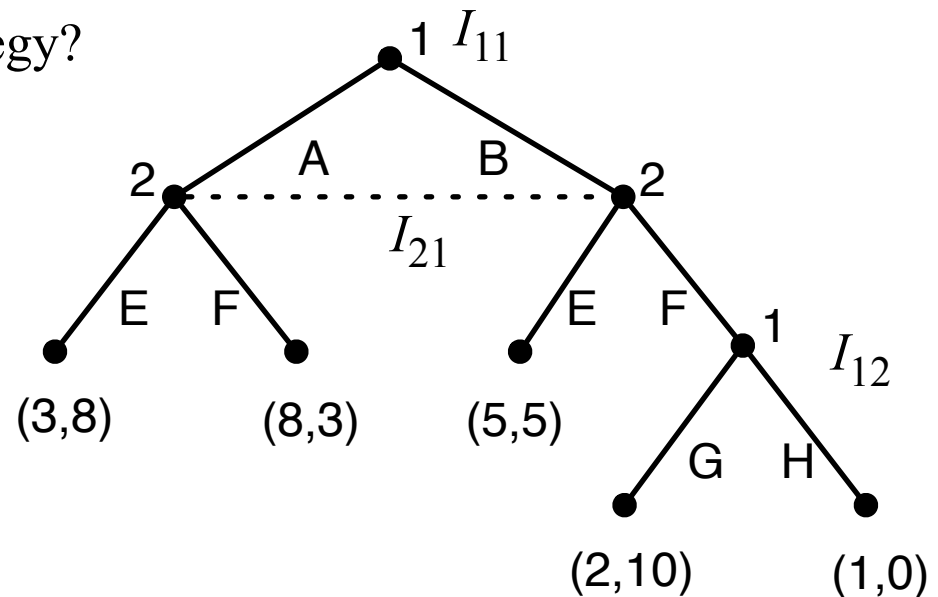
Example 3

- A mixed strategy for agent 1:
 - $\{(0.6, (A,G)), (0.4, (B,H))\}$
- The choices in the two information sets aren't independent
 - Choose A in $I_{11} \Leftrightarrow$ choose G in I_{12}
 - Choose B in $I_{11} \Leftrightarrow$ choose H in I_{12}
- Is there an equivalent behavioral strategy?



Example 3

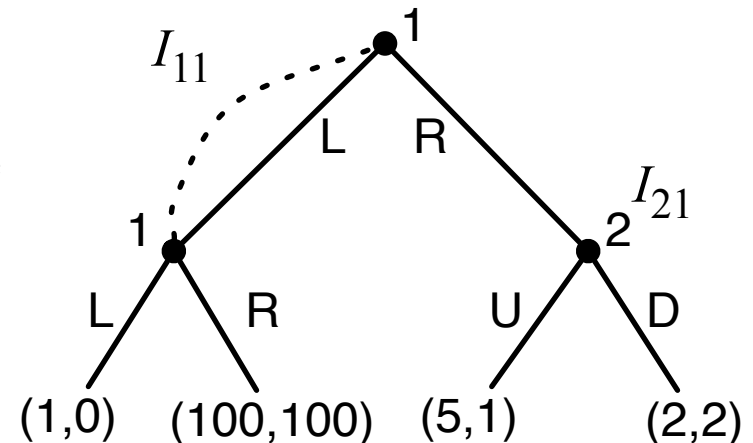
- A mixed strategy for agent 1:
 - $\{(0.6, (A,G)), (0.4, (B,H))\}$
- The choices in the two information sets aren't independent
 - Choose A in $I_{11} \Leftrightarrow$ choose G in I_{12}
 - Choose B in $I_{11} \Leftrightarrow$ choose H in I_{12}
- Is there an equivalent behavioral strategy?
 - In I_{11} , $\{(p, A), (1-p, B)\}$
 - In I_{12} , $\{(q, G), (1-q, H)\}$
- Look for p and q that give the same probabilities of outcomes
 - $\Pr[A] = p = 0.6$
 - $\Pr[(B,G)] = (1-p) q = 0$
 - $\Pr[(B,H)] = (1-p) (1-q) = 0.4$



Behavioral vs. Mixed Strategies

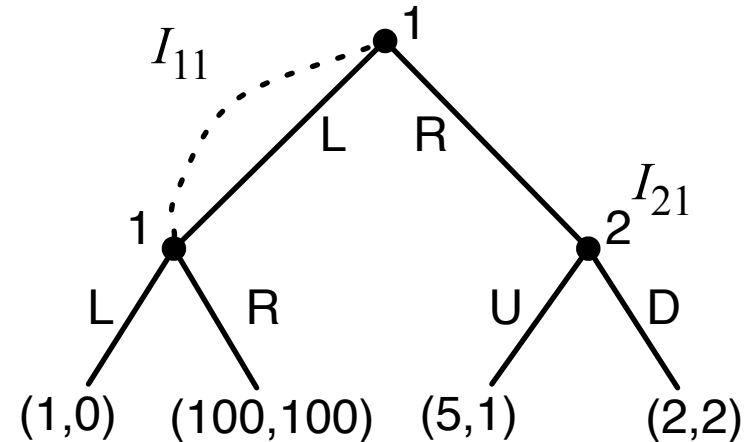
- In some games, there are
 - mixed strategies that have no equivalent behavioral strategy
 - behavioral strategies that have no equivalent mixed strategy
- Thus mixed and behavioral strategies can produce different Nash equilibria

- Example:
 - Both of Agent 1's choice nodes are in the same information set, I_{11}
 - How could this ever happen?



Behavioral vs. Mixed Strategies

- Mixed strategy $\{(p, L), (1-p, R)\}$
 - agent 1 chooses L or R randomly, but commits to it
 - Choose L \Rightarrow history $\langle L, L \rangle$
 - Choose R \Rightarrow history $\langle R, U \rangle$ or $\langle R, D \rangle$
 - **never** $\langle L, R \rangle$
- Nash equilibrium in mixed strategies:
 - For agent 1, R is strictly dominant
 - For agent 2, D is strictly dominant
 - So (R,D) is the unique Nash equilibrium



Behavioral vs. Mixed Strategies

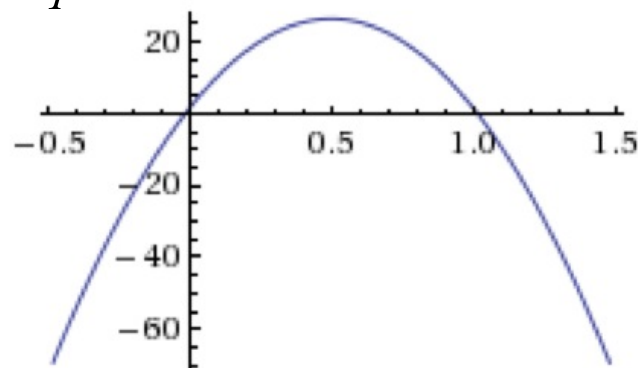
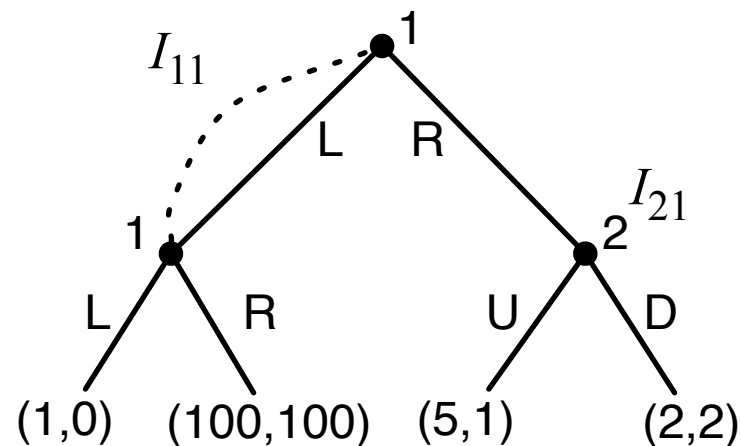
- Behavioral strategy (I_{11} , $\{(q, L), (1 - q, R)\}$)

- Remake the choice each time agent 1 is in I_{11}
 - If $p = q = 0$, have the pure strategy L
 - If $p = q = 1$, have the pure strategy R
 - In all other cases, $\Pr[\langle L, R \rangle] = q(1 - q) > 0$

- Nash equilibrium in behavioral strategies:

- For 2, D is strictly dominant
 - Find 1's best response among behavioral strategies
 - Suppose 1's behavioral strategy is (I_{11} , $\{(q, L), (1 - q, R)\}$)
 - $u_1 = 1 q^2 + 100 q(1 - q) + 2 (1 - q) = -99q^2 + 98q + 2$
 - Maximum is where $du_1/dq = 0$
 - $-198q + 98 + 0 = 0$
 - $q = 49/99$

- Equilibrium is ($\{(49/99, L), (50/99, R)\}$, D)



Games of Perfect Recall

- The reason the strategies weren't equivalent was because agent 1 could be in the same information set more than once
 - Mixed strategy \Rightarrow agent 1 will make the same move every time
 - Behavioral strategy \Rightarrow agent 1 may make a different move each time
 - Like mixed strategy + faulty memory
- Look at games where agents have perfect memories
 - Agent i has **perfect recall** if i never forgets anything i knew earlier
 - G is a **game of perfect recall** if every agent in G has perfect recall

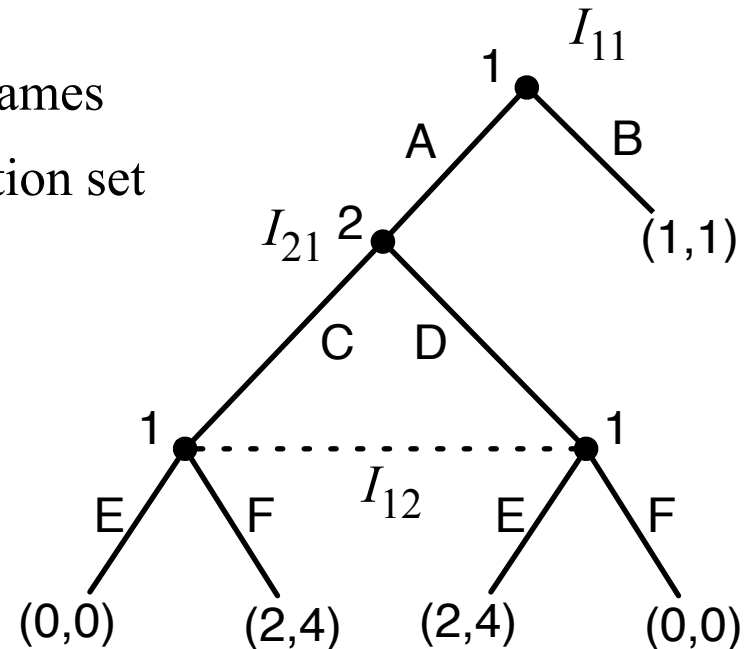
Theorem: For every history in a game of perfect recall, no agent will have the same information set more than once

Games of Perfect Recall

- **Theorem:** For every history in a game of perfect recall, no agent will have the same information set more than once
- **Proof:** Let h be any history for G . Suppose that
 - At one point in h , i 's information set is I
 - At another point later in h , i 's information set is J
 - Then i must have made at least one move in between
 - If i remembers all his/her moves, then
 - At J , i remembers a longer sequence of moves than at I
 - Thus I and J are different information sets
- **Theorem** (Kuhn, 1953). In a game of perfect recall, for every mixed strategy there is an equivalent behavioral strategy, and vice versa
- **Corollary:** In a game of perfect recall, the set of Nash equilibria doesn't change if we consider behavioral strategies instead of mixed strategies

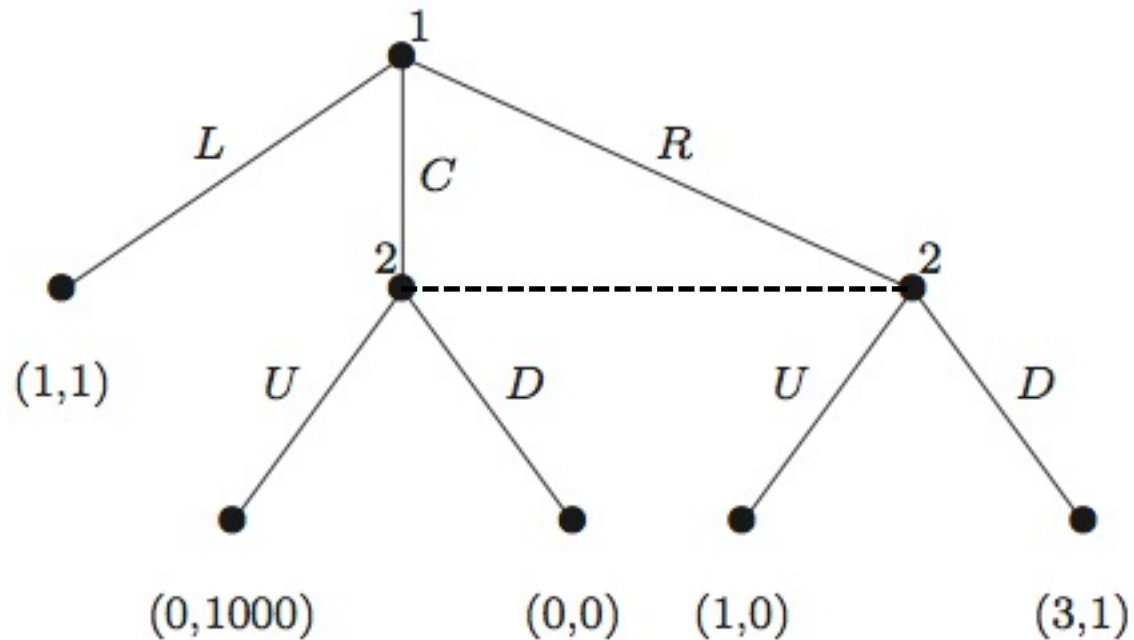
Sequential Equilibrium

- For perfect-information games, subgame-perfect equilibria were useful
 - Avoided non-credible threats; could be computed more easily
 - Each agent's strategy must be a best response in every subgame
- Generalize to imperfect-information games?
- Information set \Leftrightarrow a set of possible subgames
 - One for each element of the information set
 - Could we require an agent's strategy to be a best response in all of the subgames?



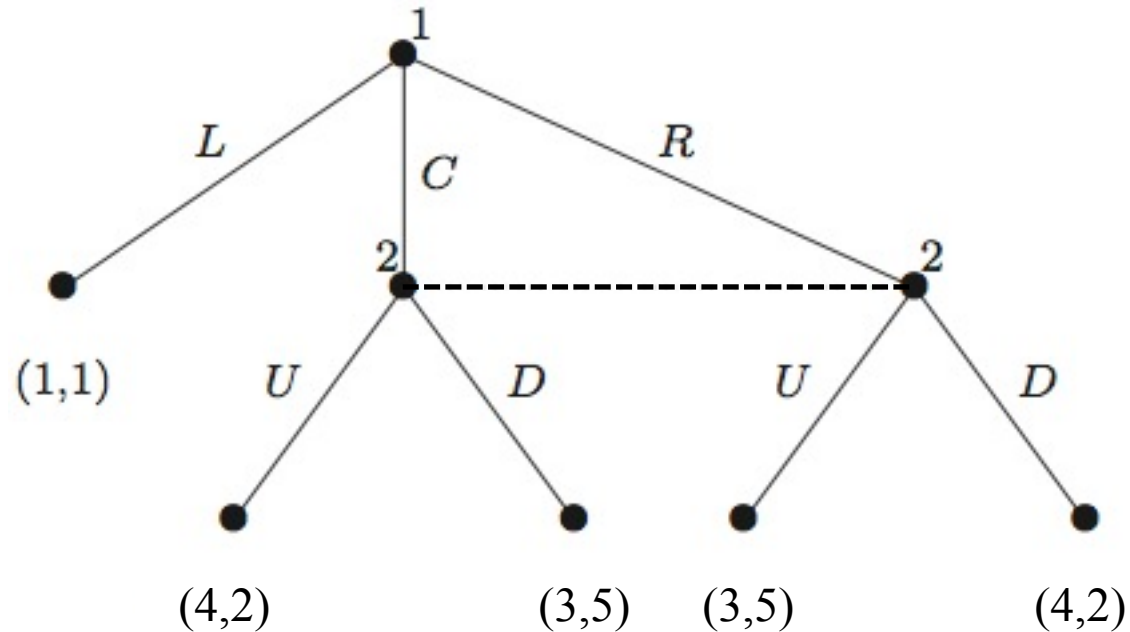
Example

- No strategy is a best response to both C and R
- Assume common knowledge of rationality
 - 1 will never choose C
 - 2 only needs a best response to R



Example

- No strategy is a best response to both C and R
- Assume common knowledge of rationality
 - 1 will never choose L
- Suppose the agents' mixed strategies are
 - $s_1 = \{(p, C), (1-p, R)\}$
 - and
 - $s_2 = \{(q, U), (1-q, D)\}$
- Can show there is one Nash equilibrium, at $p = q = \frac{1}{2}$
 - But $q = \frac{1}{2}$ is not a best response to either C or R



Sequential Equilibrium

- In general, need Bayesian reasoning about the players' strategy profiles
- This leads to a complicated solution concept called **sequential equilibrium**
 - A little like a trembling-hand perfect equilibrium, but with additional complications to deal with the tree structure

Definition 5.3.1 (Sequential equilibrium). *A strategy profile S is a sequential equilibrium of an extensive-form game G if there exist probability distributions $\mu(h)$ for each information set h in G , such that the following two conditions hold:*

1. *$(S, \mu) = \lim_{n \rightarrow \infty} (S^n, \mu^n)$ for some sequence $(S^1, \mu^1), (S^2, \mu^2), \dots$, where S^n is fully mixed, and μ^n is consistent with S^n (in fact, since S^n is fully mixed, μ^n is uniquely determined by S^n); and*
 2. *For any information set h belonging to agent i , and any alternative strategy S'_i of i , we have $u_i(S | h, \mu(h)) \geq u_i(S'_i, S_{-i} | h, \mu(h))$.*
- Every finite game of perfect recall has a sequential equilibrium
 - Every subgame-perfect equilibrium is a sequential equilibrium, but not vice versa
 - I won't discuss it further

Summary

- Topics covered:
 - information sets
 - behavioral vs. mixed strategies
 - games of perfect recall
 - equivalence between behavioral and mixed strategies in such games
 - very brief discussion of sequential equilibrium