

## Lecture 5.3: Deterministic Finite Automata

Presenter: Warida Rashid

Scribe: Warida Rashid

### Extended Transition Function

The extended transition function takes a state  $q$  and an input string  $w$  and yields the resulting state in which the processing of the string ends. The definition proceeds by induction over the length of  $w$ . The definition precedes by induction over length of  $w$ . It is sometimes represented as  $\hat{\delta}$  (delta hat) to distinguish it from the transition function.

**Induction basis:** When  $w$  is of length 0 or an empty string ( $\epsilon$ ), it is defined as  $\hat{\delta}(q, \epsilon) = q$

**Inductive Step:** In this step, we determine the transition of a string ( $w$ ) of length  $l+1$  from a string of length  $l$ . Let's say,  $w$  is of the form " $va$ " where  $v$  is a string of length  $l$  and " $a$ " is a symbol. Therefore,  $\hat{\delta}(q, va) = \hat{\delta}(\hat{\delta}(q, v), a)$

### Example

Let's consider the DFA that accepts strings with two consecutive 1s and we have to determine  $\hat{\delta}(q_0, 110)$ .

Transition table:

	0	1
$\rightarrow q_0$	$q_0$	$q_1$
$q_1$	$q_0$	$q_2$
$*q_2$	$q_2$	$q_2$

$$\begin{aligned}
& \delta (q_0, 110) \\
&= \delta ( \delta (q_0, 11), 0) \\
&= \delta ( \delta ( \delta (q_0, 1), 1), 0) \\
&= \delta ( \delta ( \delta ( \delta (q_0, \varepsilon ), 1), 1), 0) \\
&= \delta ( \delta ( \delta (q_0, 1), 1), 0) \text{ [The basis of the induction]} \\
&= \delta ( \delta (q_1, 1), 0) \\
&= \delta (q_2, 0) \\
&= q_2
\end{aligned}$$

### Language of a DFA in terms of the Extended Transition Function

The language of a DFA  $(Q, \Sigma, \delta, q_0, F)$  is a set that contains strings  $w$  such that  $\delta (q_0, w) \in F$