

Algebric Circuit to R1CS

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Definitions

Succinct

proof π is succinct if:

- 1. $|\pi| = \text{poly}(\lambda, \log |\mathcal{C}|)$ and
- 2. the verification time is $poly(\lambda, |x|, log|C|)$.

Proof of knowledge

A proof of knowledge for relation R with knowledge error κ is a two party protocol with a prover P and a verifier V with the following two properties:

- 1. Completeness: If $(x,w) \in R$, then the prover P who knows witness w for x succeeds in convincing the verifier V of his knowledge. More formally: $\Pr(P(x,w) \leftrightarrow V(x) \to 1) = 1$, i.e. given the interaction between the prover P and the verifier V, the probability that the verifier is convinced is 1.
- 2. Validity: Validity requires that the success probability of a knowledge extractor E in extracting the witness, given oracle access to a possibly malicious prover P

 in convincing the verifier. This property guarantees that no prover that doesn't know the witness can succeed in convincing the verifier.

Construction of SNARKS

 $\textbf{Computation} \rightarrow \textbf{Arithmetic Circuit} \rightarrow \textbf{R1CS} \rightarrow \textbf{QAP} \rightarrow \textbf{zk-SNARK}$

```
def qeval(x):
    y = x**3
    return x + y + 5
```

```
sym_1 = x * x

y = sym_1 * x

sym_2 = y + x

\sim out = sym_2 + 5
```

```
'~one', 'x', '~out', 'sym_1', 'y', 'sym_2'
```

```
a = [0, 1, 0, 0, 0, 0]
b = [0, 1, 0, 0, 0, 0]
c = [0, 0, 0, 1, 0, 0]
```

$$s.c = s.a * s.b \implies s = [0, 3, 0, 9, 0, 0]$$

```
a = [0, 1, 0, 0, 1, 0]
b = [1, 0, 0, 0, 0, 0]
c = [0, 0, 0, 0, 0, 1]
```



```
\begin{array}{l} a = [5, 0, 0, 0, 0, 0, 1] \\ b = [1, 0, 0, 0, 0, 0] \\ c = [0, 0, 1, 0, 0, 0] \end{array}
```

```
[0, 1, 0, 0, 0, 0]
[0, 0, 0, 1, 0, 0]
[0, 1, 0, 0, 1, 0]
[5, 0, 0, 0, 0, 1]
[0, 1, 0, 0, 0, 0]
[0, 1, 0, 0, 0, 0]
[1, 0, 0, 0, 0, 0]
[1, 0, 0, 0, 0, 0]
[0, 0, 0, 1, 0, 0]
[0, 0, 0, 0, 1, 0]
[0, 0, 0, 0, 0, 1]
[0, 0, 1, 0, 0, 0]
```

$$\implies$$
 $(A.s)*(B.s) = C.s$

