

10 HW

$$R = \begin{bmatrix} 0 & 1 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \end{bmatrix}$$

Calculate Transitive Closure

a can communicate to b and c
both b and c can communicate to d
assume communication is one way,

this can be described by the relation

$$R = \{(a, b), (a, c), (b, d), (c, d)\}$$

we want to change the system so a can communicate with d and still maintain the previous system.

- Described as
 - Find the smallest relation R^+ which contains R as a sub set $R \subset R^+$
 - and is transitive
- Thus, $R^+ = \{(a, b), (a, c), (b, d), (c, d), (a, d)\}$

Hence the transitive closure of R ,

- denoted by R^+
 - is the smallest transitive relation that contains R as a subset

It follows that is $(a, b) \in S$ and $(b, c) \in S$ then, $(a, c) \in S^+$

$$A = \{1, 2, 3, 4\}$$

$S = \{(1, 2), (2, 3), (3, 4)\}$ be the relation on set A

- This is called the successor relation on A
 - since each element is related to its successor.

$$S^+ = S$$

$$S = \{(1, 2), (2, 3), (3, 4)\}$$

$$SS = S^2 = \{(1, 3), (2, 4)\}$$

$$S^+ = S \cup S^2 = \{(1, 2), (1, 3), (2, 3), (2, 4), (3, 4)\}$$

$$S^2S = S^3 = \{(1, 4)\}$$

- because $(1, 3) \in S^2$ and $(3, 4) \in S$
this shows that $S^3 \subseteq S^+$

Thus,

$$S^+ = S \cup S^2 \cup S^3 = \{(1, 2), (1, 3), (1, 4), (2, 3), (2, 4), (3, 4)\}$$

Let r be a relation on the set $A = \{\text{Kira, Frank, Lola, Sarah, Judy, Reuben}\}$

with the relation matrix,

$$R = \begin{bmatrix} 0 & 1 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \end{bmatrix}$$

The matrix of the transitive closure R^+ ,

- can be computed by the equation $R^+ = R + R^2 + R^3 + \dots + R^n$.
 - By using the polynomial evaluation methods,
 - we can compute R^+ with $n - 1$ matrix multiplications:
 - $R^+ = R(I + R(I + (\dots R(I + R) \dots)))$
 - for example, if $n = 3$
 - $R^+ = R(I + R(I + R))$

We can make use of the fact that if T is a relation matrix $T + T = T$ due to the fact that $1 + 1 = 1$ in boolean arithmetic.

- let, $S_k = R + R^2 + \dots + R^k$

- Then,

-

$$R = S_1$$

$$S_1(I + S_1) = R(I + R) = R + R^2 = S_2$$

$$\begin{aligned} S_2(I + S_2) &= (R + R^2)(I + R + R^2) \\ &= (R + R^2) + (R^2 + R^3) + (R^3 + R^4) \\ &= R + R^2 + R^3 + R^4 \\ &= S_4 \end{aligned}$$

- Similarly,

-

$$S_4(I + S_4) = S_8$$

- and by induction we can prove

-

$$S_{2^k}(I + S_{2^k}) = S_{2^{k+1}}$$

- Each matrix multiplication doubles the number of terms that have been added to the sum that is currently computed.

FOR OUR PROBLEM WE HAVE

Let r be a relation on the set $A = \{\text{Kira, Frank, Lola, Sarah, Judy, Reuben}\}$

$n = 6$

with the relation matrix,

$$R = \begin{bmatrix} 0 & 1 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \end{bmatrix}$$

The matrix of the transitive closure R^+ ,

- can be computed by the equation $R^+ = R + R^2 + R^3 + \dots + R^n$.

$$R^2 = \begin{bmatrix} 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \end{bmatrix}$$

$$R^3 = \begin{bmatrix} 0 & 0 & 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \end{bmatrix}$$

[illegible]

[illegible]

[illegible]

Thus, $R^+ = R^1 + R^2 + R^3 + R^4 + R^5 + R^6$

$$R^+ = \begin{bmatrix} 0 & 1 & 1 & 1 & 1 & 1 \\ 0 & 0 & 0 & 1 & 1 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \end{bmatrix}$$

Q1.2

In set notion,

$$R^+ = \{(2, 1), (3, 1), (4, 1), (4, 2), (5, 1), (5, 2), (5, 4), (6, 1), (6, 2), (6, 4)\}$$

Q1.3

the **transitive closure** of

S

,

- denoted by

S^+

or

S^*

- is the smallest transitive relation that contains

S

as a subset.

For our specific, problem,
relation

R

is the successor relation on set

A

.

The **Transitive Closure** specifies an ancestral family tree,

- if there is a edge between two vertices, then they are
 - either a parent and child,
 - grandparent, and child,
 - or great grandparent and child.

Q2.1

$$A = \{c, d, e\}$$

so total # of equivalence relations on the set is 5

$$R_1 = \{(a, a), (b, b), (c, c)\}$$

$$R_2 = \{(a, a), (b, b), (c, c), (a, b), (b, a)\}$$

$$R_3 = \{(a, a), (b, b), (c, c), (b, c), (c, b)\}$$

$$R_4 = \{(a, a), (b, b), (c, c), (a, c), (c, a)\}$$

$$R_5 = \{(a, a), (b, b), (c, c), (a, b), (b, a), (c, a), (b, c), (b, c), (c, b)\}$$

Each equivalent relation on a set yields a partition of that set in disjoint equivalence classes for a finite set.

The number of equivalence relations is the number of partitions.

Q 2.2

How many different partial ordering can we define on the set

$$A = \{x, y\}$$

?

$$A \times A = \{(x, x), (y, y), (x, y), (y, x)\}$$

partial order is reflexive, lets list the reflexive relations on

$$A$$

$$R_1 = \{(x, x), (y, y)\}$$

$$R_2 = \{(x, x), (x, y), (y, y)\}$$

$$R_3 = \{(x, x), (y, z), (y, y)\}$$

$$R_4 = \{(x, x), (x, y), (y, x), (y, y)\}$$

- Due to antisymmetric property, we cant use

$$R_4$$

so, there exist 3 different partial orderings for set

$$A$$

Q 2.3

How many different total orderings on set

$$A = \{p, q\}$$

?

$$A \times A = \{(p, p), (p, q), (q, p), (q, q)\}$$

Total ordering is

- partial ordering and two elements of

$$A$$

are comparable.

so, 2 total orderings,

$$\{p < q, q < p\}$$

