

Operating System

dec D

Syllabus of Operating System

1. Basic Introduction → types of OS

Process diagram ✓ (Main)

System call

**

2. Process scheduling (CPU scheduling)

FIFO

SJF

Round Robin

3. Process synchronization → Semaphore ✓

4. Deadlocks and threads → Banker's algorithm

5. Memory management → Paging

Virtual memory

Segmentation

fragmentation

Page replacement algo ✓

6. Disk scheduling ← SCAN

CSCAN

FCFS

7. Unix commands ls

→ mkdir

→ cd

→ chmod

→ Open system call

sequential

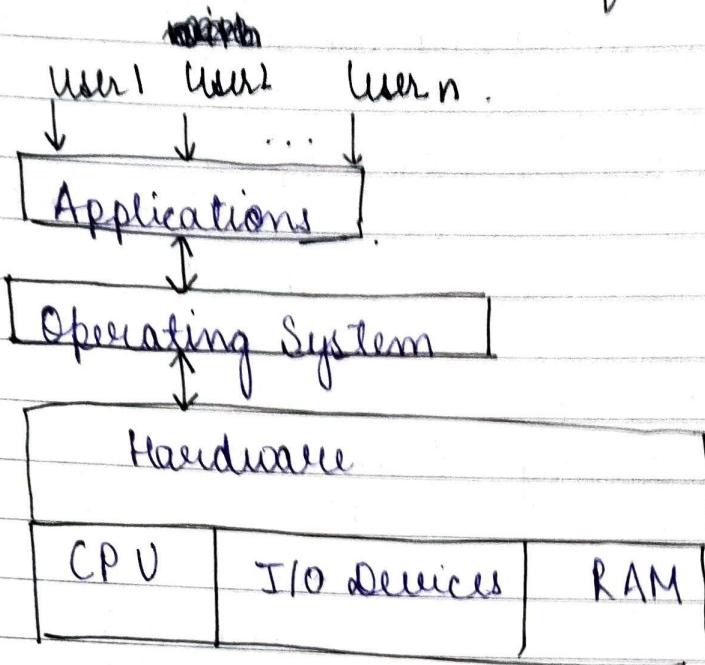
Random

linked

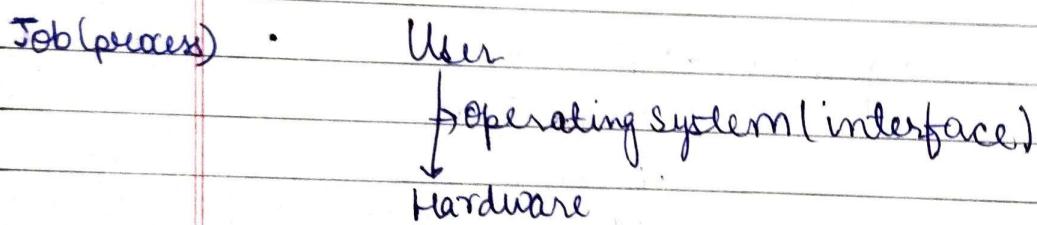
8. File management & Security

L-1.1

Introduction to OS and its function



Operating System :- It is a system software which works as an interface between users and hardware.



Eg: Windows, Macintosh (Mac), Linux

Primary goal → To provide convenience to user.

Throughput → no of tasks executed per unit time. Eg - Linux

Functionality of Operating System

- 1) Resource management (kis user ko kitne particular time ke liye hardware ko provide krsana hai).

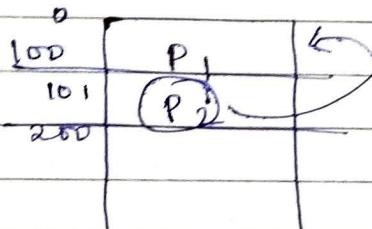
2. Process management - Managing & executing multiple processes at a time. we use CED scheduling

program execute process

3. Storage management \rightarrow How to store the data permanently using file system.
Eg - Hardisk.

4. Memory management (RAM) \rightarrow RAM is limited in system. Every process before execution comes to RAM. Allocation & deallocation takes place.

5. Security and privacy - (Password protection)
Windows uses two security protocol.



It provides security b/w the process also.

Windows (cmd)

Linux (Terminal)

OS works only through system calls.
(Read, open, write)



L-1.2

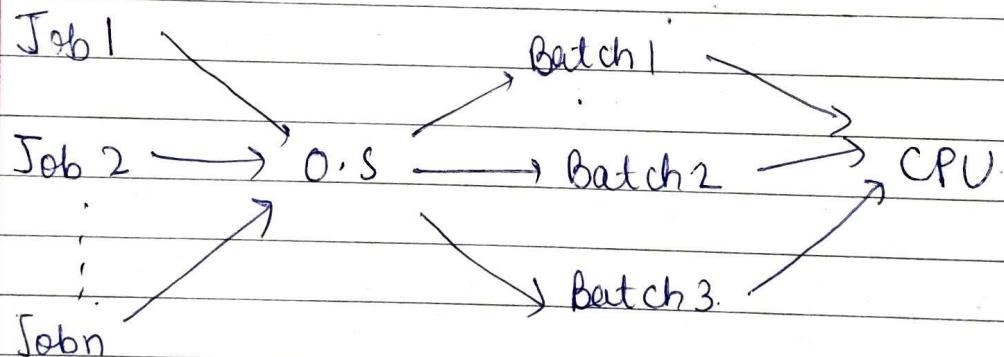
Batch Operating System

Types of Operating System

- 1) Batch OS.
- 2) Multiprogramming OS.
- 3) Multitasking OS / Time Sharing
- 4) Real time OS
- 5) Distributed OS
- 6) Clustered OS
- 7) Embedded OS.

1) Batch OS

Similar (batch) kind of jobs.

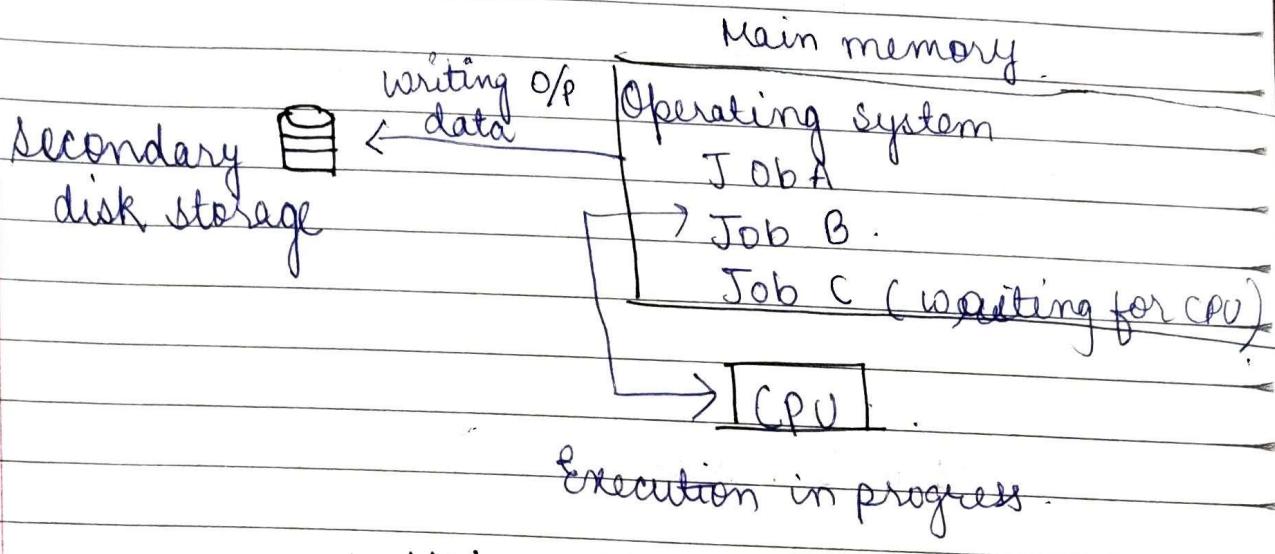


Same type of Jobs batch together and execute at a time.

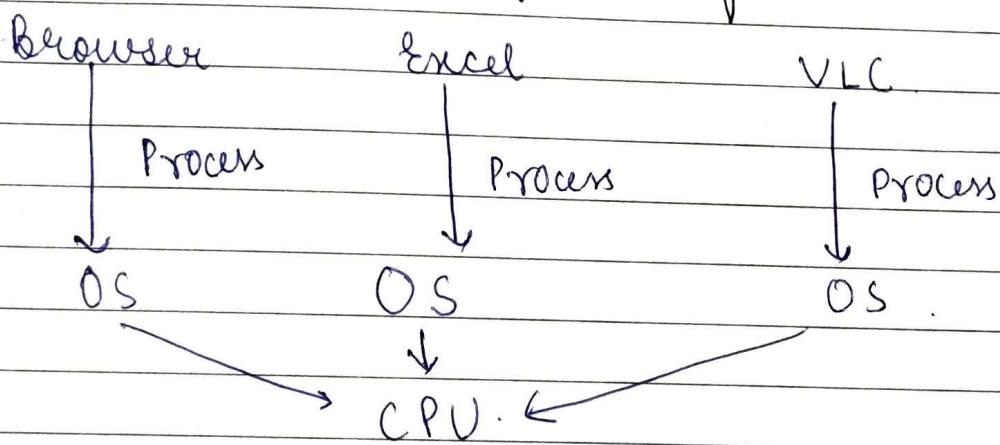
Demerit: It processes one job at a time till the time CPU remains idle.



L-1.3 (Multiprogramming & Multitasking system)



Multiprogramming OS / Time sharing



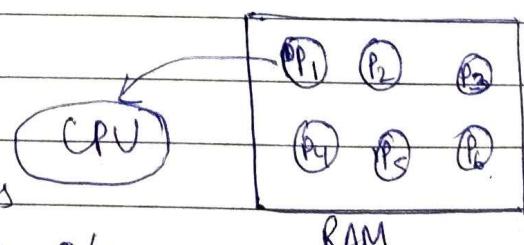
Multitasking OS

Multiprogramming OS.

Non preemptive

CPU executes a process completely until or unless process demands any I/O or P operation.

Idleness should be removed



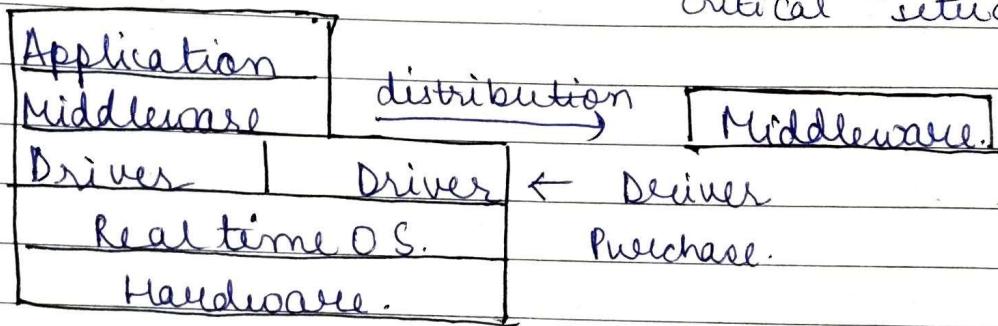


Multitasking / Time sharing - Provides particular time for every process. If its complete its great but if not it moves to other process P₂ & schedule P₁ for further.

Response time is reduced for every process. Idleness is not there for CPU. Job.

Dec 1.4 Types of OS (Real time)

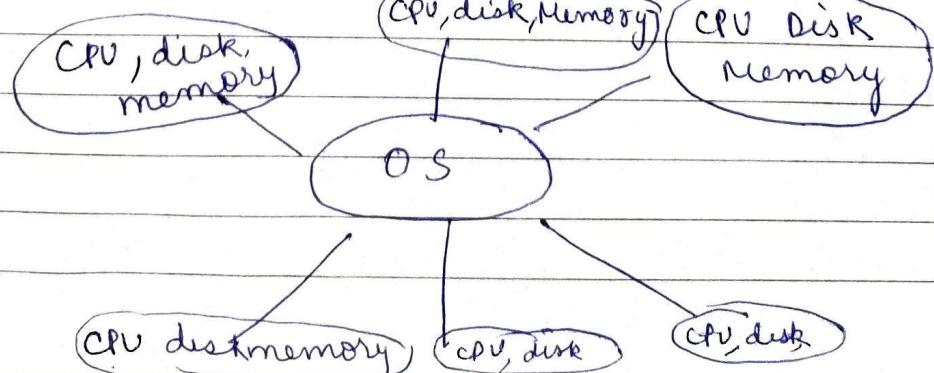
Real time OS → Hard (Restriction on time is more)
soft Eg - In gaming no critical situation



Immediate output

- i) No delays

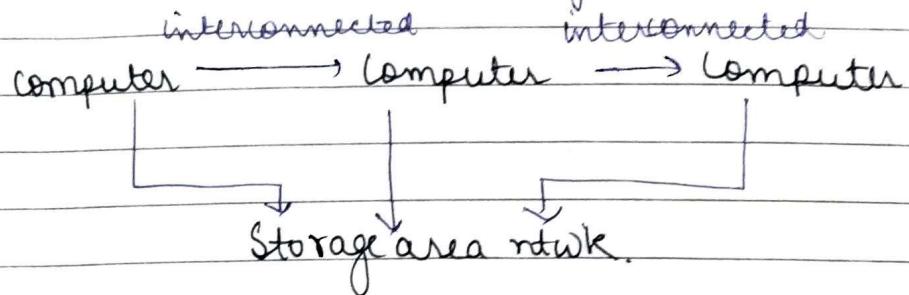
Distributed → Processing environment is distributed all over the world. Every process has its own environment.



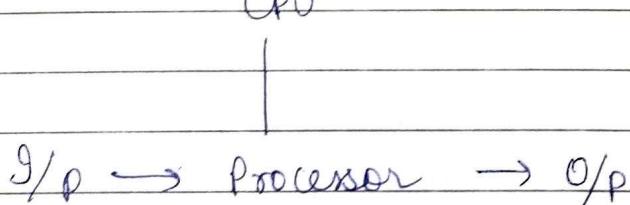
It uses multiple central processors to serve multiple real-time applications & multiple users using telephone wires.

Clustered OS. → ~~the~~ Clustered OS share storage and are closely linked via lan or a faster inter connection.

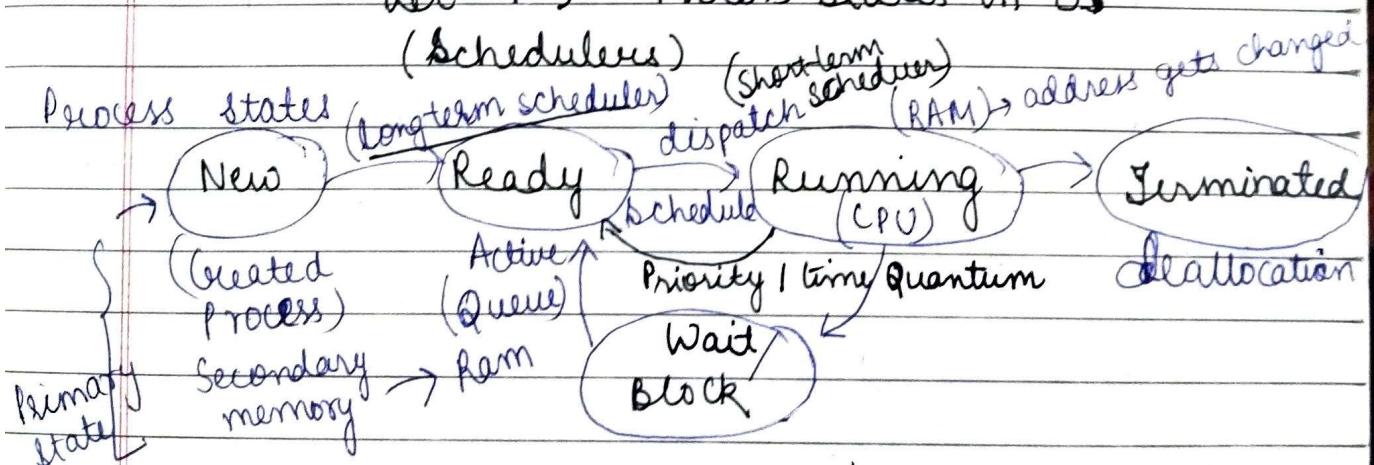
Group of connected computers
* fault tolerance, scalability



Embedded OS - It is a special purpose computer designed to perform one or few dedicated fun'n's. with real time computing. Works on a fixed functionality.



sec - 1.5 Process states in OS



long term scheduler - Bring more & more process in ready state.

LTS - long term scheduler

STS - short term scheduler.

Non Preemptive.

Preemptive.

Non preemptive - Running process isn't stopped at middle.

Preemptive - running process stopped at middle.

X

X

Syllabus of OS.

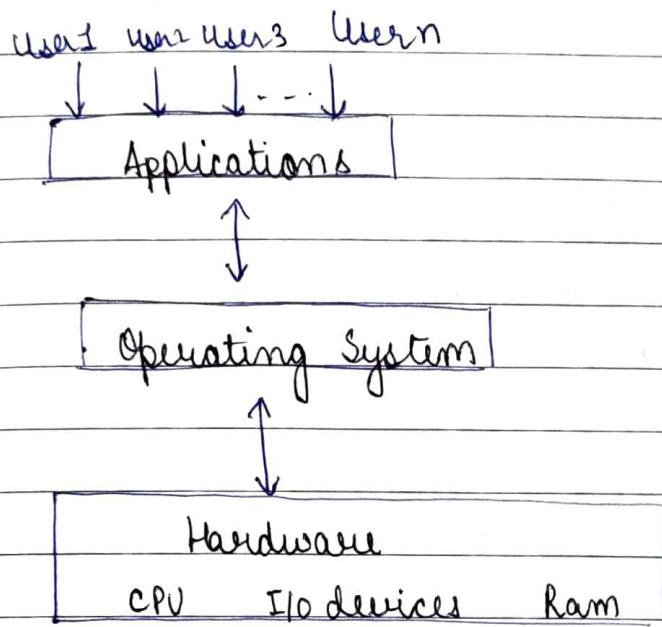
1. Basic Introduction : what is O/S, goals, need, system calls, diff' n OS.
2. Process Mdg : what is process, Algo - FIFO, SJF, Preemptive, Non preemptive, Round Robin
3. Deadlock - Banker Algo
4. Concurrency control - Semaphore, msg passing, race cond'n.
5. Memory - Paging, Segmentation, Fragmentation, virtual memory, Page Replacement.
6. Disk Mdg - C-look, Algo, PCFS, SSCF, look.
7. File mg : NIFS, FAT 32, Seg., Random
8. Unix System : commands.
9. Case study : MS DOS
10. Security - Attacks, Firewalls, exceptions.



L-1.1

Operating system acts as an interface between user and hardware. If there is no OS user has to write a program to interact with hardware.

Primary goal - Its primary goal is to provide convenience to the user.

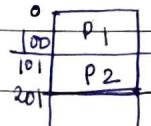


Throughput - No of tasks executed per unit time.
Eg - Linux

Functionalities of OS.

- i) Resource management - In case of multi-user managing, OS acts as key to provide hardware among the users for a given time. When many users are sending request at a time so we use OS so that there is no load on system.

2. Storage / Process management → Performing multiple process at a time, OS uses CPU scheduling algorithms for processing & executing in a proper & efficient way
3. Storage management (Hard disk) - How to store data in hard disk using file system management is done by OS.
4. Memory management (RAM) - In RAM we have limited no of size. Every process which gets executed goes in RAM there is allocation & deallocation done by OS in RAM after the process is done.



5. Security & privacy
It provides security & privacy between the process such that it doesn't get blocked. Password protection

Windows (cmd) } system calls
Linux (terminal) } (read, open, write)

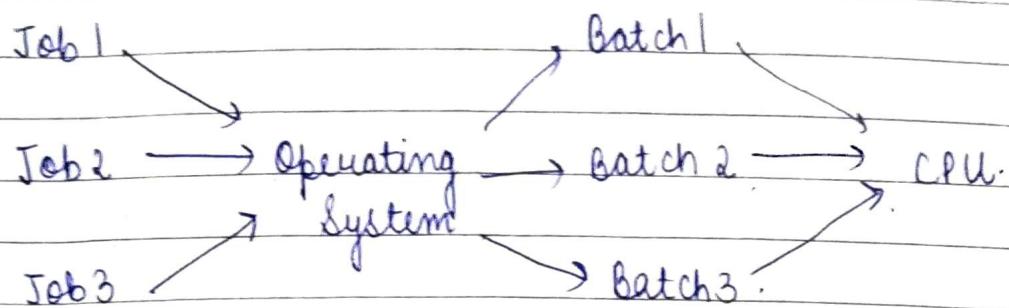
1.2.

Types of OS

- 1) Batch OS
- 2) Multiprogramming
- 3) Multitasking
- 4) Real time OS
- 5) distributed
- 6) Clustered
- 7) Embedded



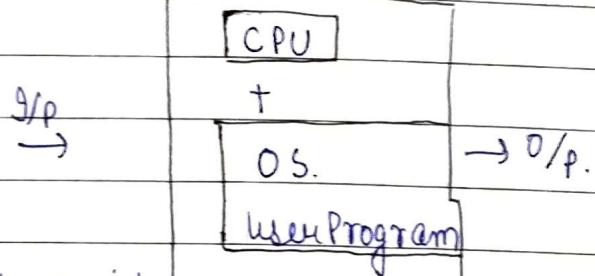
i) Batch OS



~~before~~

before batch OS., we used mainframe computers

- 1) Common I/O & O/P devices were card readers & tape drives.



- 2) User prepare a job which consist of the program I/O data & control instructions.

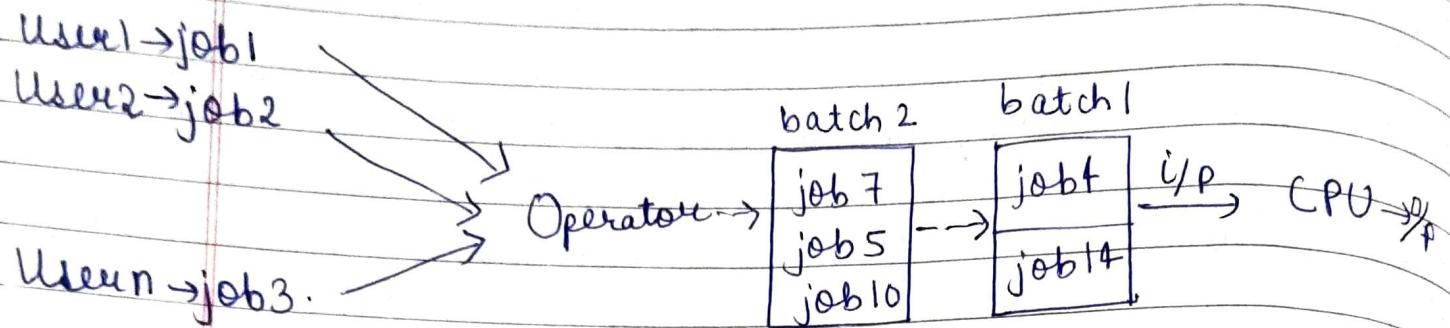
Job	
program	
I/O data	
Control instruction	

- 3) I/O ~~Punch~~ O/P (Punch card)

Major problem was.

- 1) Speed mismatch (user & CPU)
- 2) Less memory
- 3) Every job has diffⁿ req^m

1. Batch processing (e.g.: payroll systems)



- 1) Jobs with similar needs are batched together and executed through the processor as a group.
- 2) Operator sorts job as a deck of punch cards in a batch with similar needs.
e.g. - FORTAN batch, COBOL batch etc.
- 3) first jobs → batch (with same req) → CPU

Advantages

- 1) In batch job execute one after another saving time for activities like loading compiler.
- 2) During a batch execution no manual intervention is needed.

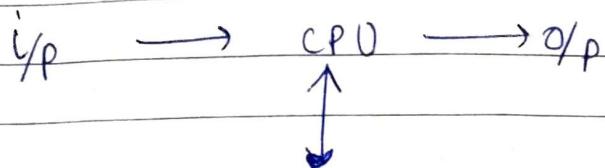
Disadvantage

- 1) Memory limitation
- 2) Interaction of i/p & o/p devices directly with CPU.
- 3) CPU remains idle during loading & unloading



Spooling (Simultaneous peripheral device)

i/p or o/p



secondary storage.

temporarily process is stored in secondary storage.

Advantages

- 1) Increase the system performance.
- 2) Spooling resolve the problem of speed mismatch of diffⁿ. device
- 3) % of one job is overlapped with the computation of other jobs.
- 4) Spooling use the disk as a huge buffer.

Spooling

Uses the hard disk as a large spool
(spool is a temporary storage area in hard disk)

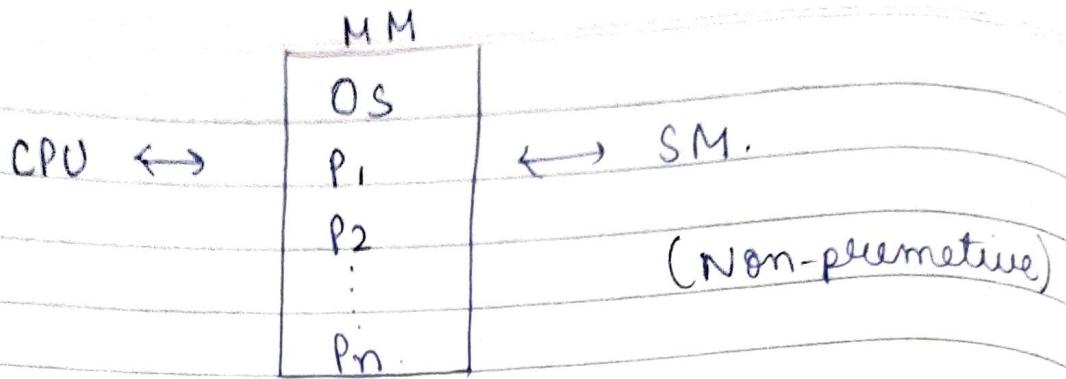
Buffering

Uses limited memory space in RAM usually called buffer
(Buffer is a temporary storage area in RAM)

1.3

2.

Multiprogramming OS

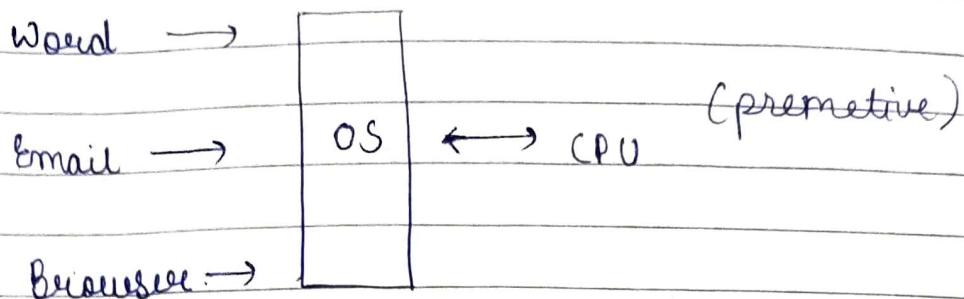


- 1) Maximize CPU utilization
- 2) Multiprogramming means more than one process in main memory which are ready to execute.
- 3) Process generally requires CPU time & I/O time. So if running process performs I/O or some other event which don't require CPU then instead of sitting idle, CPU makes context switch & pick some other process & this idea will continue.
- 4) CPU never ^{remains} idle unless there is no process ready to execute or at the time of context switch

Advantage	Disadvantage
1) High CPU utilization.	1) Difficult scheduling
2) less waiting time, response time etc	2) Main memory management is required
3) May be executed to multiple users	3) Memory fragmentation
4) Now a days useful when load is more	4) Paging (non contiguous memory allocation)

3o Multi-tasking OS. (Eg Unix)

Time sharing / Fair share / R.R



1. Multitasking is multiprogramming with time sharing
 2. Only one CPU but switches betn processes so quickly that it gives illusion that all executing at same time
 3. the task in multitasking may refer to multiple threads of the same program
 4. Main idea is better response time & executing multiple process together.
4. Real-time OS - It is used in environment where a large no of events mostly external to system must be accepted & processed in a short time within certain deadlines. The time interval required to process & response to input is very small.
Hard RTOS - It guarantees critical task to be completed within a range of time.
Eg - A robot hand to weld a car body
- Soft RTOS - It provides some relaxation in time.

Advantages

- 1) Max^{m.} utilization of devices & system
- 2) Better task shifting
- 3) Systems are error free



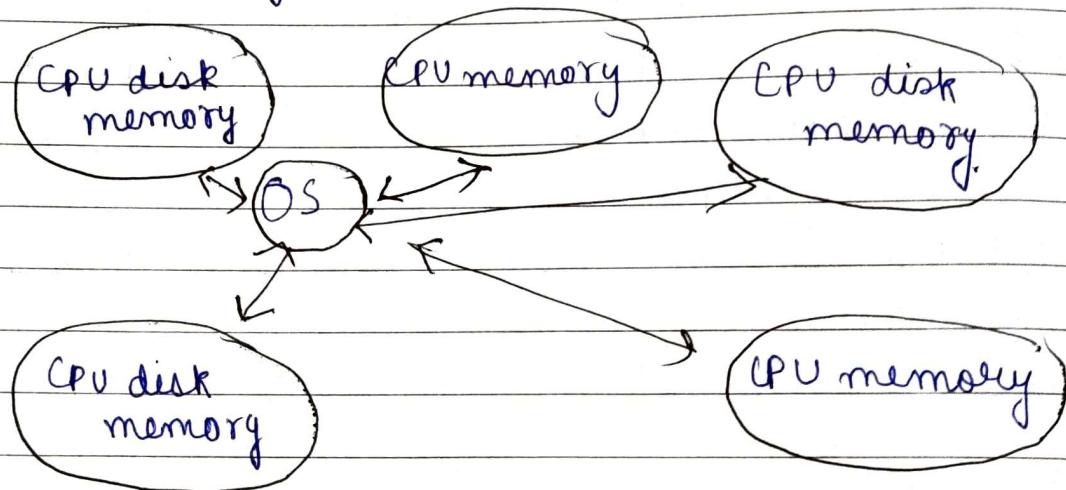
Disadvantages

- 1) limited tasks
- 2) Complex algorithms
- 3) Use heavy system resources

5) distributed OS. (Eg - LOCUS)

Distributed system uses multiple central processors to serve multiple real time application & multiple users. Data processing jobs are distributed among the processor accordingly.

Processors communicate with each other through various communication lines (such as high speed buses or telephone lines). These are referred as loosely coupled systems or distributed system.



Advantages

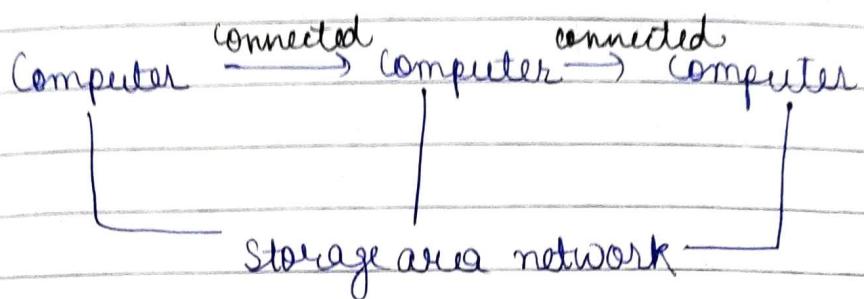
- 1) Failure of one network connection doesn't affect other
- 2) Delay in data processing decreases

Disadvantages

- Failure of main network will stop entire communication.

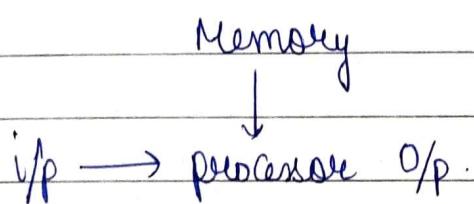
6) Clustered OS

- Clustered computers share storage & are closely linked via LAN or faster connection.
- Combination of multiprocessor + distributed OS



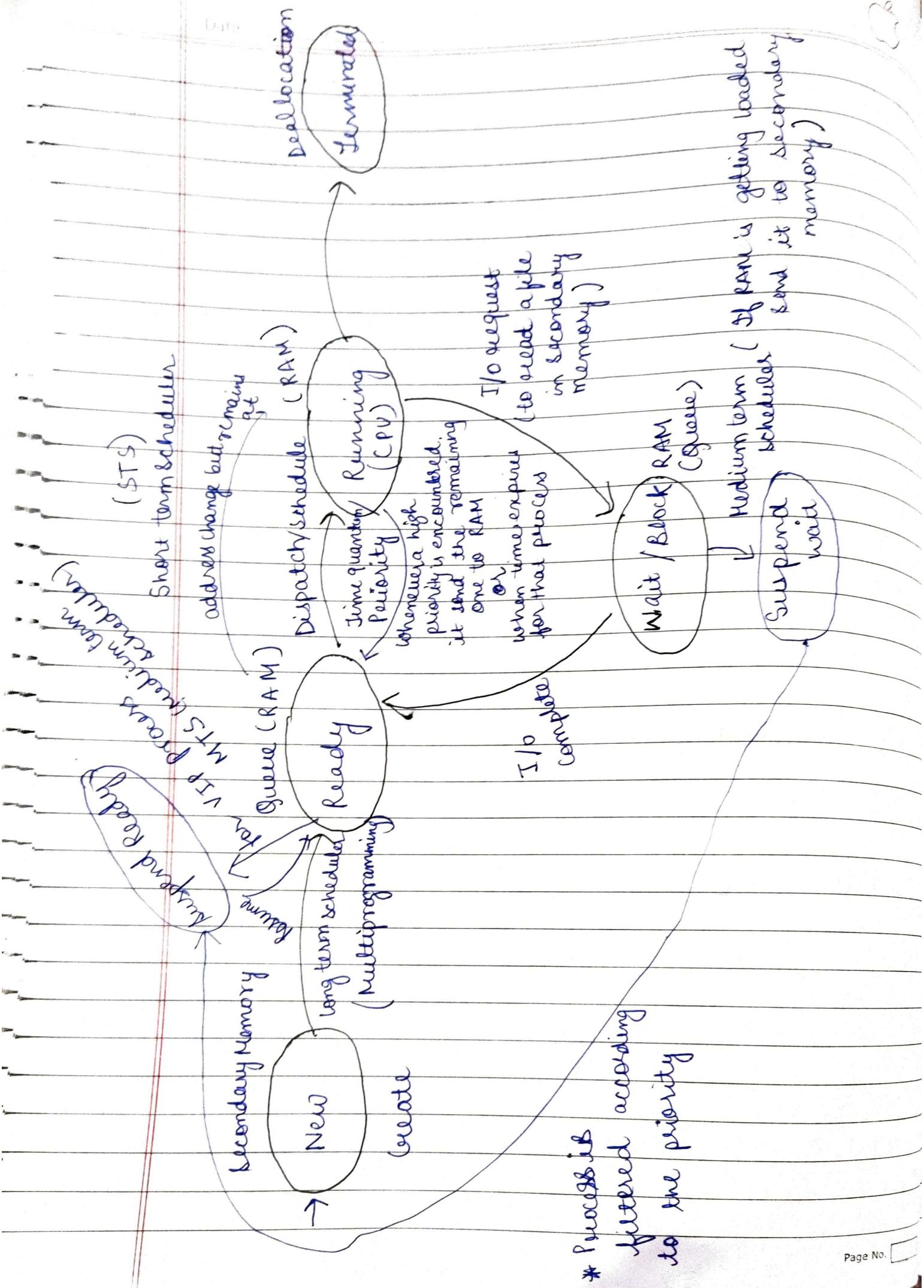
7) Embedded OS. - An OS in which we work on embedded systems. It is designed to perform a specific task for a device that is not in computer.

Eg - ATM



1.5

Process states in OS





- * Non preemptive - NO time quantum, no process can jump in between.
- * preemptive - We stop the running process in midway due to high priority process.

L-1.6

Linux Commands

Q. which command is used to assign only Read permission to all three categories of file 'note'.

- chmod a-~~rwx~~ rwx
- chmod go+r note
- chmod ugo = r note
- chmod. u+r, g+r, o-x note

chmod - change mode

<u>u</u>	<u>w</u>	<u>x</u>	<u>u</u>	<u>w</u>	<u>x</u>	<u>u</u>	<u>w</u>	<u>x</u>
user	group	other	user	group	other	user	group	other

u - user
g - group
o - others

rx own	u - read 4
6 7	w - write 2
	x - execute 1

Q. 'chmod ugo+rwx note' command can be represented in octal notation as

- 1) chmod 555 node
 2) chmod 666 node
 3) chmod 333 node
 4) chmod 444 node

$$7+2 = 9 + 2 = 6.$$

Q. Suppose you have a file "f" whose contents are

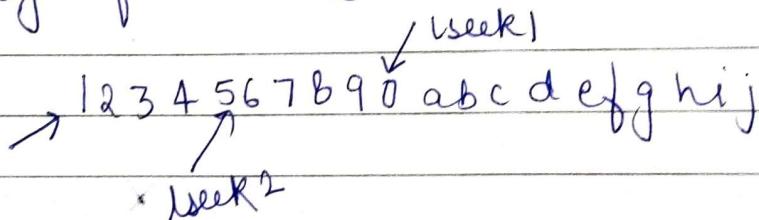
1 2 3 4 5 6 7 8 9 0 abcdefghij
 Here seek is used 2 times sequentially.

lseek (n, 10, SEEK_CUR);
 lseek (n, 5, SEEK_SET);

n is file descriptor. After applying seek Q times, what will be the current position of R/W head? (Index starts from 0)

- a) 0.
 b) 5
 c) 10
 d) 15.

Seek - system call (To move read write head we use seek command)
 by default read write head is at 0.



10, Seek_CUR → move to the given no from the current

5, Seek_Set → set the current pos at given

2, Seek_End → Count from the end towards left.

L - 10.7

System Calls in OS.

User mode System call → Kernel mode

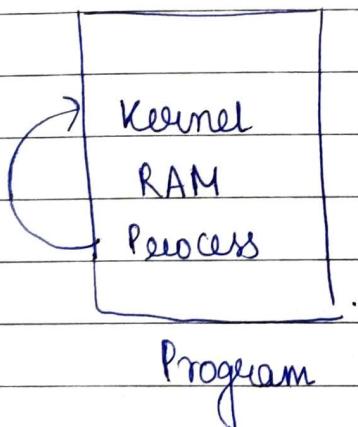
2+4 monitor, print(6)

In Linux we can use system call directly.

printf() → accessing system call

System calls.

- 1) File related SC → open(), Read(), write(), close()
create file etc.



- 2) Device related → read, write, reposition
ioctl (I/O control), fcntl (file related control)

for hardware devices

- 3) Information → Process or regarding device attributes

For e.g.: getpid, attributes, get system time data.

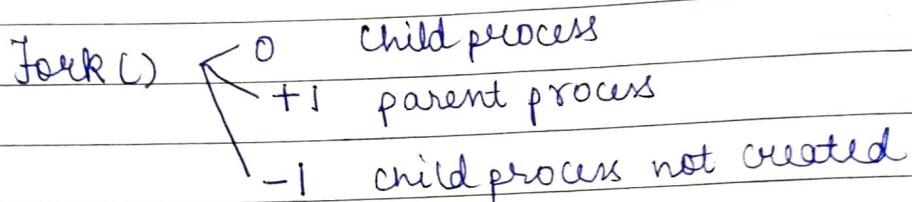


- 4) Process control \rightarrow Program being loaded
 \rightarrow load, execute, abort, fork, wait,
 signal, allocate etc.
- 5) Communication control - Interprocess communication
 \rightarrow Pipe(), create/delete connections, shmget()

L-1.8

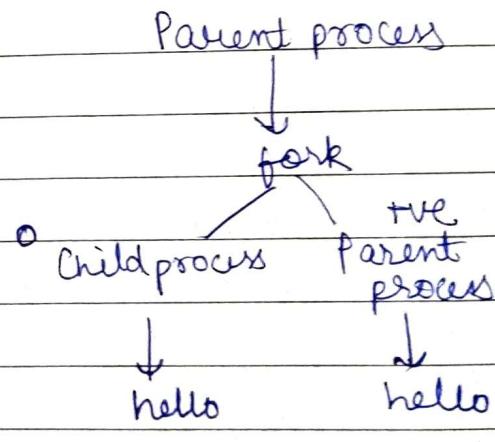
Fork system call

To create a child process from parent process
 we can use fork() or thread.



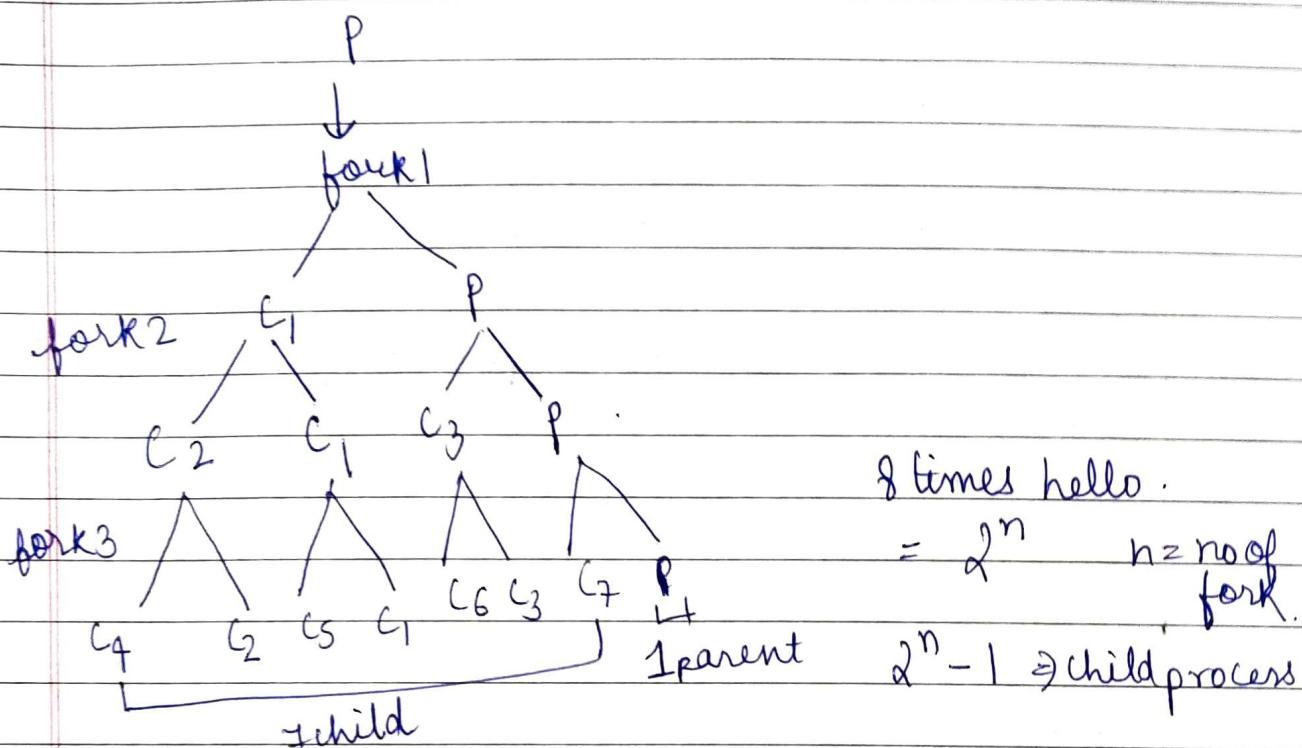
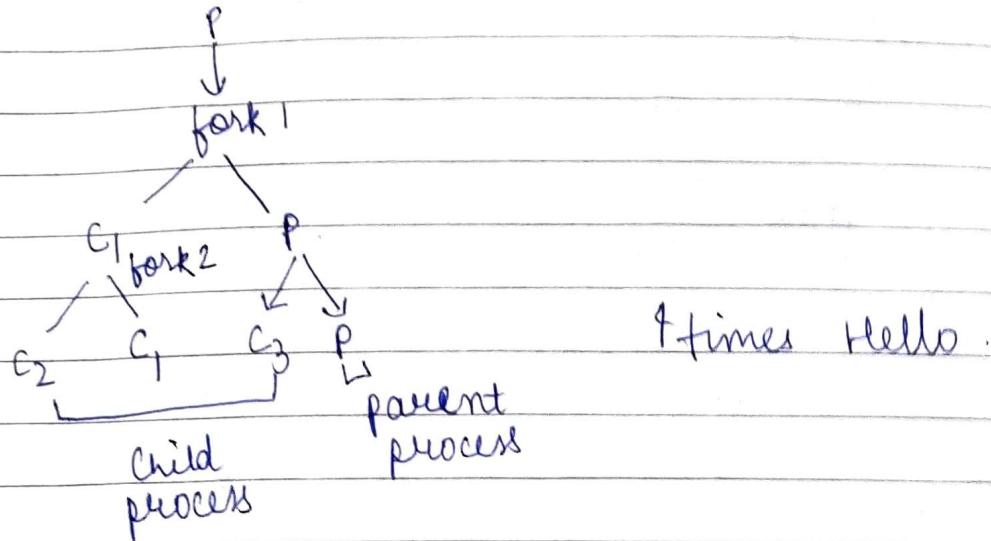
```

main() {
    fork();
    printf("hello");
}
  
```



```

main() {
    fork();
    fork();
    printf("hello");
}
  
```



L-1.9

Questions on Fork System Call

include <stdio.h>

include <unistd.h>

int main ()

{

if (fork() & fork())

fork();

printf ("Hello");

return 0;

~~fork() & fork()~~

Date 07/01/2023

~~fork() & fork()~~



~~fork() & fork()~~

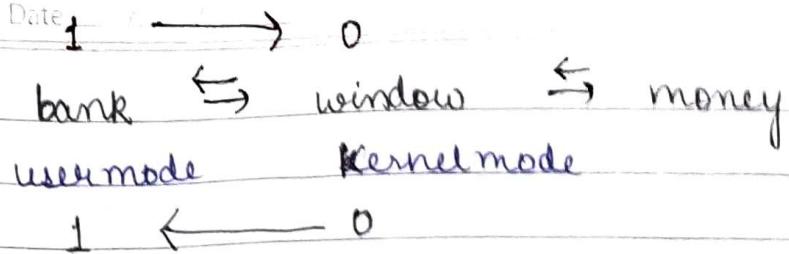
+ve

+ve

(false)

0

+ve



L-1.11

Process Vs Threads in OS

Process - It is a heavy weight task] multitasking

Thread - It is a light weight task] environment

Process

- 1) System calls involved
- 2) OS treats diffⁿ process differently.
- 3) Diffⁿ process have diffⁿ copies of data, files, cache
- 4) Context switching is slow
- 5) Blocking a process-not block other
- 6) Independent

Thread (User level)

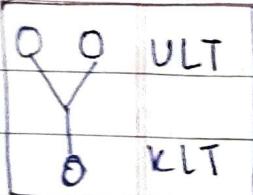
- 1) No system call involved
- 2) All user level threads treated as single task for OS.
- 3) Threads share same copy of code
- 4) Context switch is fast
- 5) Blocking a thread-block entire process
- 6) Interdependent

	Stack
Code	code
Data	data

fork()

T ₁	T ₂
Stack	Stack
1	2
Code	
data	

Hybrid environment



L-1.12

User level & kernel level thread
code & data is shared
have own stack but

User level thread

- 1) User level threads are managed by user level library.

- 2) User level threads are typically fast

- 3) Context switching is faster.

- 4) If one user level thread performs blocking operation then entire process get blocked

(context switch) Process $\gtreqless;$ KLT $\gtreqless;$ ULT

Kernel level thread

- 1) Kernel level threads are managed by OS (system calls)

- 2) Kernel level threads are slower than user level

- 3) Context switching is slower.

- 4) If one kernel level thread blocked no affect on others.

L - 2.1

Process scheduling Algorithms

A way of selecting an algorithm from ready queue & putting it on the CPU.

more processes

P₁

P₂

⋮

P_n

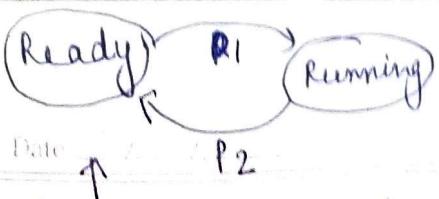
Ready Queue

RAM

scheduling
algorithm

CPU

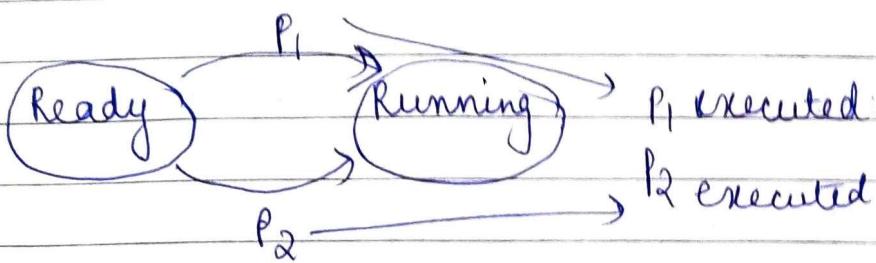
uniprocessor



- Reasons
- 1) Time quantum (provided time for P_1)
 - 2) Priority process

* Preemptive \rightarrow If a process is taken out from ready queue of RAM and is put in running process (queue) (CPU). Then we can stop the process P_1 & give another process P_2 & send P_1 to ready queue. It provides responsiveness.

* Non preemptive \rightarrow If a process is taken out from ready queue it would be processed till its burst time, after then only new process can be introduced.



Scheduling Algorithms

Preemptive

Non-preemptive

BT

1) SRTF (shortest remaining time first)

1) FCFS (first come first serve)

AT

BT

BT 2) LRTF (longest remaining time first) 2) SJF (Shortest job first)

BT

TQ 3) Round Robin

placed in
ready queue
in
descending
order of
burst
time

3) LJF (longest job first)

4) Priority based

des
of
burst
time

4) HRRN (High response ratio next).

Priority given

5) Multilevel Queue

6) Priority based

L-2.2

Different times in CPU scheduling

Arrival time - The time at which process enters the ready queue or stack.

Burst time- Time required by a process to get executed on CPU.

Completion time - The time at which process completes its execution

Two-second time - Completion - Arrival time

Waiting time - Turnaround - Burst time

Response time - (The time at which process get CPU) -
(Arrival-time))

```

graph LR
    A[Arrival time  
(point of time)] --> B[process  
execution  
duration]
    B --> C[completion time  
(point of time)]
  
```

~~Arrived Bank~~ → Create new ~~Bank~~ amount.

L-2-3

First Come First Serve (FCFS)



FCFS is an OS scheduling algorithm that automatically executes queued requests and process in order of their arrival.

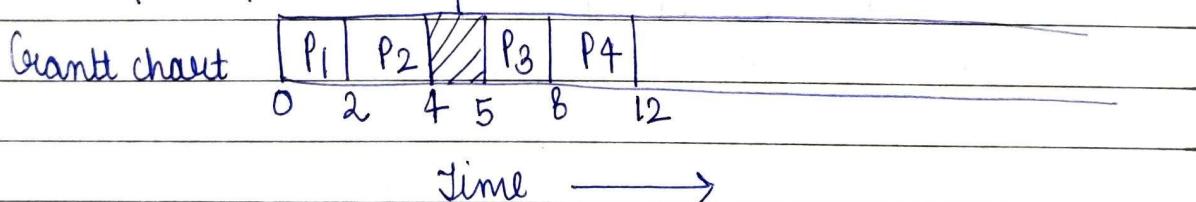
- i) The process which requests the CPU first get the CPU allocation first.
- ii) Managed using FIFO queue.

CPU idle Kb/milli	Process No	AT	BT	CT	TAT	WT	RT
		Arrival time	Burst time	Completion time	(CT-AT)	(TAT-BT)	
0	P ₁	0	2	2	2	0	0
2	P ₂	1	2	4	3	1	1
5	P ₃	5	3	8	3	0	0
8	P ₄	6	4	12	6	2	2

Criteria of arrival time

Mode : Non pre-emptive

cpu idle



At t=12, all the processes got executed.

$$\text{Avg TAT} = \frac{14}{4}$$

$$\text{Avg WAT} = \frac{3}{4}$$



SJF - (Shortest Job first) \rightarrow SJF

is an algorithm in which the process having the smallest execution time is chosen for next execution.

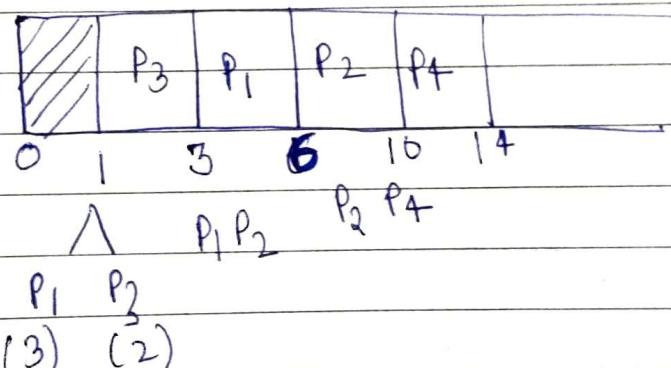
- ii) It can be preemptive or non-preemptive
- iii) It reduces the waiting time for other process.
- w) It is a greedy Algo.

Job Queue No	Process No	AT	BT	UT	TAT	WT	RT
3	P ₁	1	3	8	5.	2.	2
6	P ₂	2	4	10	8.	4.	4
1	P ₃	1	2	3	2	0	0
10	P ₄	4	4	14	10	6	6

Criteria - Burst time

Mode - Non-preemptive

Gantt chart



when BT is same see the one having less arrival time else can use process id.

L-2.5

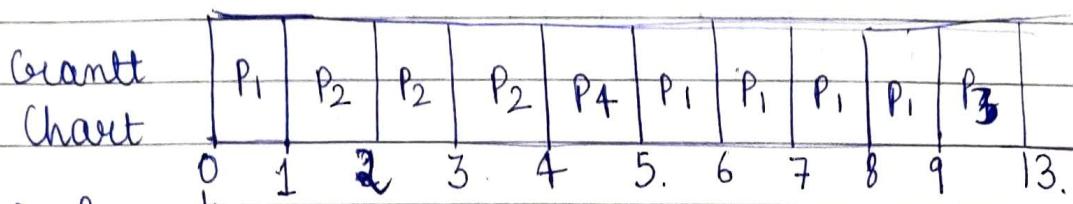
Shortest Remaining Time First

(Shortest job first + pre emptive)

Process no	AT	BT	Compt	TAT	WT	RT	CPU first time
P ₁	0	3 5 4	9	9	4	0	0
x P ₂	1	1 3 2	4	3.	0.	0	1
P ₃	2	4	13	11	7	.7	9
x P ₄	4	8	5	1	0	0	7

Criteria : "Burst time"

Mode : "Preemptive"

0 → P₁1 → P₁, P₂2 → P₁, P₂, P₃.3 → P₁, P₂, P₃.4 → P₁, P₃, P₄5 → P₁, P₃.6 → P₁, P₃7 → P₁, P₃8 → P₁, P₃9 → P₁, P₃

$$\text{Avg TAT} = \frac{24}{4} = 6$$

$$\text{Avg WT} = \frac{11}{4} = 2.75$$

$$\text{Avg RT} = \frac{7}{4} = 1.75$$

P₁ → 8 4P₂ → 1 3 2 0P₃ → 9P₄ → 1

P ₁	P ₂	P ₂	P ₄	P ₁	P ₃	
0	1	2	4	5	9	13



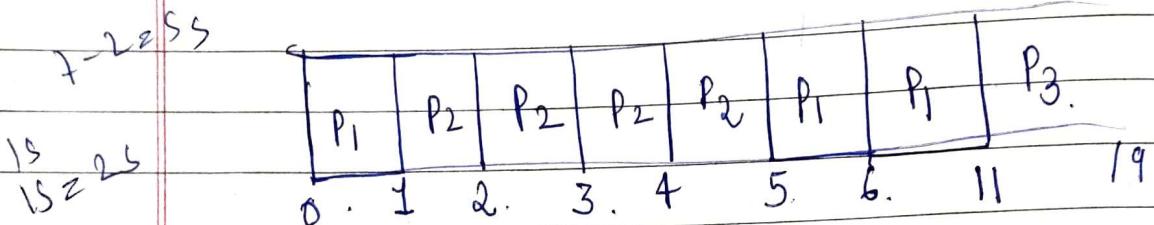
L-206

SJF with preemption ex:

process	AT	BT	CT	TAT	WT	RT
P ₁	0	4	7.68			
P ₂	1	4	2.1			
P ₃	2	8	.			

Criteria : Burst time
Mode : preemptive

Gantt chart



0 → P₁

1 → P₁, P₂

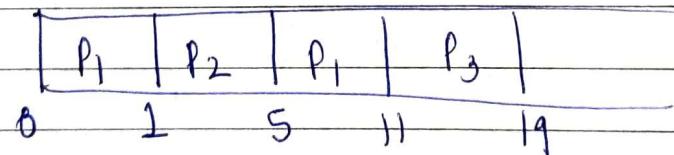
2 → P₁, P₂, P₃

3 → P₁, P₂, P₃

4 → P₁, P₂, P₃

5 → P₁, P₂, P₃

5 → P₁, P₂, P₃





L-2.7

Round Robin Scheduling Algo.

Round robin is preemptive process scheduling algorithm. Each process is provided a fix time to execute, it is called quantum. Once a process is executed for a given time period, it is pre-empted & other process executes for a given time period.

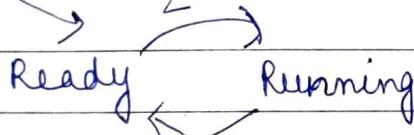
Process no	AT	BT	CT	TAT	WT	RT	cpu first time
P ₁	0	8.31	12	12	7	0	0
→ P ₂	1	0.42	11	10.	6	1	2
→ P ₃	2	2.0	6	4	2	2	4
P ₄	4	1	9	5.	4	4	8

TQ=2

Criteria : "Time Quantum"

(We have to resume the process not restart it)

Mode : "Preemptive"

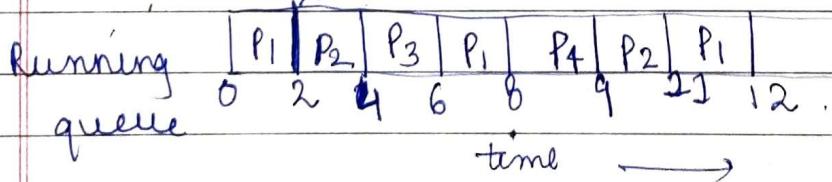


$$\begin{matrix} 3 \\ 4 \\ 2 \\ 1 \end{matrix}$$

Sequence of
process in
Ready queue

P ₁	P ₂	P ₃	P ₁	P ₄	P ₂	P ₁
----------------	----------------	----------------	----------------	----------------	----------------	----------------

context switching (save running process & welcome new process)



1 - 2.8 Scheduling algo.

Preemptive Priority ~~Process~~

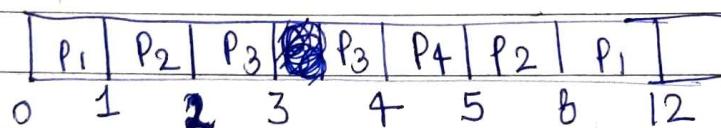
In pre-emptive scheduling, the tasks are mostly assigned with their priorities.

Priority scheduling is a method of scheduling process that is based on priority. In this algorithm, the scheduler selects the task to work as per the priority.

Equal priority - Round Robin or FCFS.

Priority	Process No	AT	BT	CT	TAT	WT	RT	CPU first
10	P ₁	0	48	12	12	7	0	0
20	P ₂	1	34	8	7	3	0	1
30	P ₃	2	120	4	2	0.	0	2
40	P ₄	4	01	5	1	0	0	4

Higher the no. higher the priority.



0 → P₁ time →

1 → P₁, P₂

2 → P₁, P₂, P₃

3 → P₁, P₂, P₃

4 → P₁, P₂, P₃, P₄

5 → P₁, P₂



L - 2.9

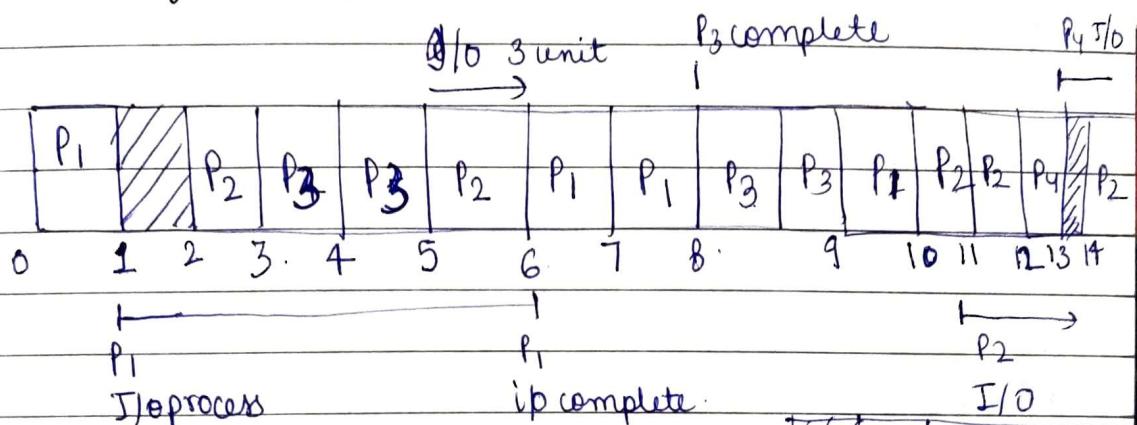
Example of Mix Burst Time (CPU & I/O both) in CPU

Process	AT	Priority	CPU	I/O	CPU
P ₁	0	2	X 0	5 0	8 2 1 0
P ₂	2	3	3 2 1	3	1
P ₃	3	1	2 X 0	3 0	8 0
P ₄	3	4	2	4	1

Lowest the no. highest the priority.

Mode: Preemptive

Criteria: Priority based

find CT of P₁, P₂, P₃, P₄0 → P₁

1 → no process

2 → P₂3 → P₃, P₄

↓
high
priority

4 → P₁, P₂, P₃, P₄

5 →

6 → P₁, P₂, P₄

	CT
P ₁	10
P ₂	15
P ₃	9
P ₄	18

$$\text{CPU idleness} = \frac{4}{18}$$

$$\text{Usage} = \frac{14}{18}$$

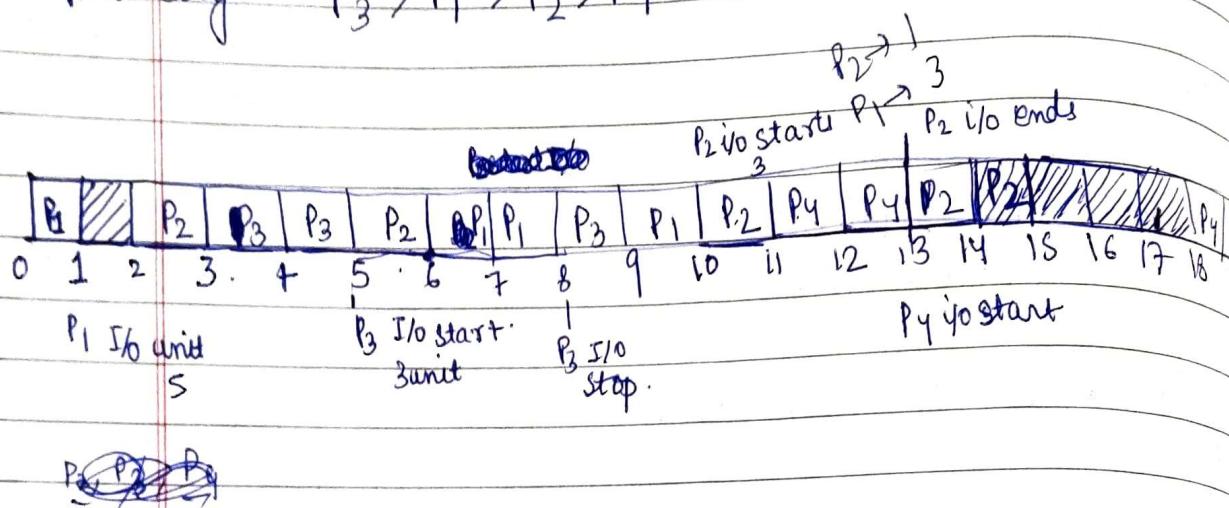


Date

Priority	AT	Priority	CPU	I/O	CPU
P ₁	0	2	20	5	
P ₂	2	3	320	3	1
P ₃	3	1	210	3	10
P ₄	3	4	210	4	1

32x0 Mode: Preemptive
Criteria Priority based.

Priority $P_3 > P_1 > P_2 > P_4$



In Ready queue.

I/O Process

$$P_1 \rightarrow 6$$

$$P_3 \rightarrow 5$$

$$6 \rightarrow P_1, P_2, P_4$$

$$7 \rightarrow P_1, P_2, P_4$$

$$8 \rightarrow P_1, P_2, P_3, P_4$$

$$9 \rightarrow P_1, P_2, P_4$$

$$10 \rightarrow P_2, P_4$$

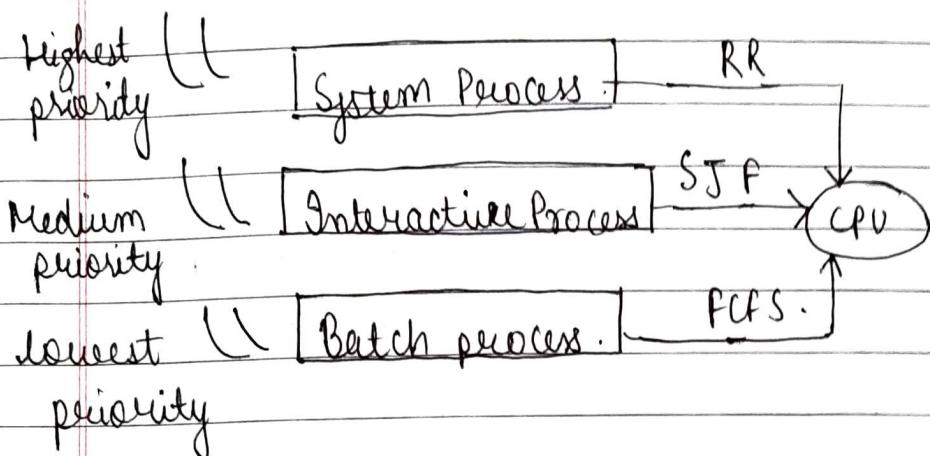
$$13 \rightarrow P_2$$

$$L-2 \cdot 10$$

Multilevel queue Scheduling

In every scheduling algorithm till now we have taken one queue.

- i) There must be different queue for diff' n process
- ii) Every process can have their own algorithms.
- i) Systems calls - Round robin
- ii) Interactive calls - SJFS.



Problem :

Lower level or lower priority may experience starvation due to more no of calls on higher priority

It can be resolved using multilevel feedback queue.

L - 2.11

Multilevel Feedback Queue

More system process - lowest priority waits

lowest priority - feedback.

lowest priority

(Batch Process)

RQ₁

P₁

TQ = 2

If completed then out

RQ₂

TQ = 4

If completed then out

RQ₂

TQ = 8

If completed then out

RQ₁

FCFS.

Remaining all
executed in FCFS
manner.

highest priority

(System process).

P₁

19

TQ = 2

17

TQ = 4

13

TQ = 8

5

Process Synchronization

Processes

Cooperative process

Independent process

One's execution

affects the other

as they share same
or common variable,
memory/buffer, code,
resources

One's execution

doesn't affect other

They have nothing
in common.

Ex - IRCTC



Process Synchronization is used to solve the problems of co-operative process.

int shared = 5; +

```

graph LR
    subgraph P1_Context [P1]
        X1[int X = shared; 5]
        X1p[X++ ; X=6]
        S1[Sleep(1); pause]
        S1p[shared = X; ←]
        E1["shared = 6  
terminated"]
        E1_in_PCB["JCS shared  
in PCB"]
    end

    subgraph P2_Context [P2]
        Y1[int y = shared; 5]
        Y1p[y-- ; Y=4]
        S2[Sleep(1)]
        S2p[shared = y;]
    end

    X1 --> X1p
    X1p --> S1
    S1 --> S1p
    S1p -.-> E1
    E1 -.-> E1_in_PCB

    Y1 --> Y1p
    Y1p --> S2
    S2 --> S2p
    S2p -.-> E1

```

When they are parallel.

Uni processor

Processes aren't synchronized as the process at last gives either 4 & 6. but it should give 5.

This is called trace condⁿ.

L-3.2

Producers = consumer Problem

Co-operative process

There is one producer that is producing something and there is one consumer that is consuming the products produced by producer. The producer & consumer share the same memory buffer.

Consumer

void consumer(void)
{

```

    ind item c;
    while (true)
    {
        localv
        ← out
        Buffer
empty while (count == 0);
        item c = Buffer(out);
        Out = (Out + 1) mod n;
        count = count - 1;
    }
}

```

Process-item (itemc)

1. load $R_{C,m}[\text{count}]$ $n = 8$

2. ~~DECR R_c~~

3. Store $m[\text{count}]$, R_C

Buffer [0.....n-1]
index

D - 2

2

COUNT

O

COUNT

OUT

1

2

4

1

20

丁

Rp → Registers

INCR → increment 6

$n = \text{size of buffer}$

Case I. x_1 (producer is processing x_1)

Out

1

$$(0+1) \bmod 8$$

$$1 \bmod 8 = 1$$

m

Count

∅ 10

$$(0+) \bmod b = 1$$

$$\text{count} = 1 - 1 = 0$$

It forms a circular queue so we used mod n.

Consumer.

	0	x_1
Out	1	x_2
I_1	2	x_3
	3	x_4
	4	
	5	
	6	
	7	

Producer

Count	In
3	4

Let x_4 be inserted in the in value.

Producer:

$$R_p = 3. \quad \text{load } R_p, m[\text{count}] \quad I_1$$

$$R_p = 4. \quad \text{INCR } R_p. \quad I_2$$

If after INCR Rp the process gets preempt

$$\text{Resuming } R_p = 4 \quad I_3$$

Let x_1 be consumed.

$$\text{load } R_c, m[\text{count}] \quad R_c = 3 \quad I_1$$

$$\text{DEC } R_c. \quad R_c = 2 \quad I_2$$

$$\text{Resuming } R_p = 2 \quad I_3$$

$I_1 I_2$ consumer $I_1 I_2$ Producer. I_3 consumer I_3 .

There are 3 items in buffer but count says only 2 values is there, there is race condn. so we can't achieve process synchronization.

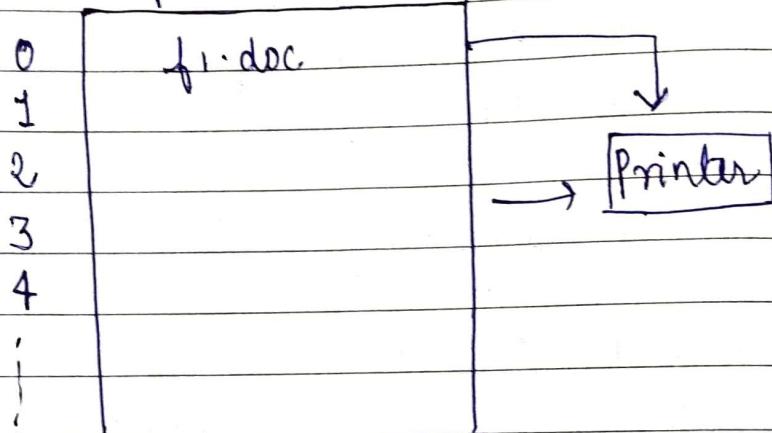
L - 3. 3.

Printer-Spooler Problem

multiple programs → spoolers → take items sequentially → gives the print

1. Load $R_i, m[i:n]$
2. Store $SD[R_i], "F-N"$
 $f_1.doc$
3. JNCR R_i
4. Store $m[i:n], R_i$

Spooler Directory



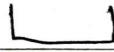
IN [01]

Empty slot

Case 1: $f_1.doc$ P_1
 R_1 [01]

$f_4.doc$
 \uparrow $f_5.doc$
 \uparrow

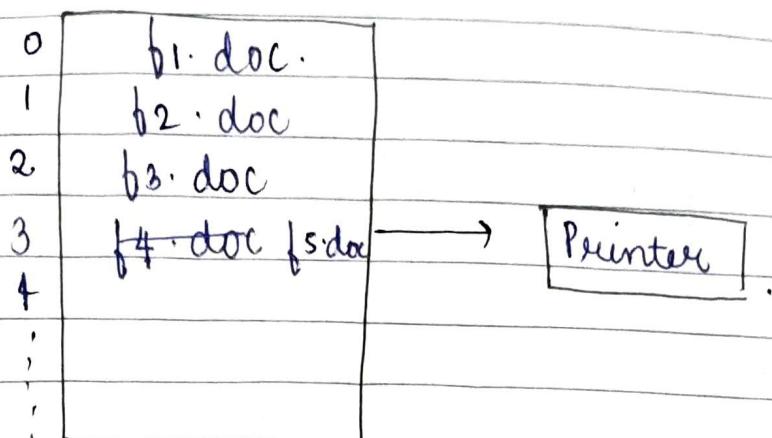
Case 2: Two processes P_1 & P_2



cooperative process



In



In [84] 4

R₁ [3] 4R₂ [84]

P ₁ . I ₁ I ₂ I ₃	I ₄
P ₂ I ₁ I ₂ I ₃	I ₄

P₁ & P₂ due to non synchronisation put the file in the same index. So there is loss of data. When 2 processes share the same code, same data there is a problem created.

3.4

Critical Section Problem

"It is the part of program where shared resources are accessed by various processes"

↓
co operative

P₁

#include

main()

A,B → non critical section

P₂

#include

main() {

X,Y

↳ common code

→ non critical section

Common code

→ critical section

count → critical section → count

main()

A, B

Non critical
section

Should
clean
entry
section

Entry section

count++

Critical section

Entry section

shouldn't clean
critical section

P₁

P₂

(Exit section)

(Exit section)

Different solⁿ for process synchronization

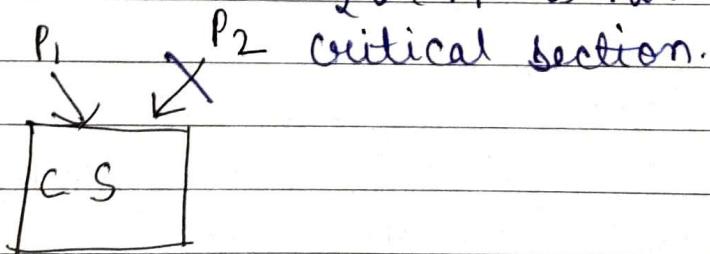
4 conditions

1) Mutual Exclusion }
2) Progress } Primary

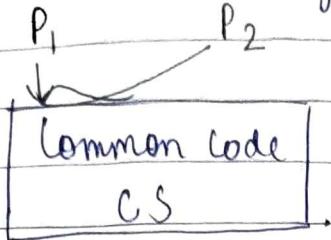
{ 3) Bounded wait

Secondary
(4) No assumption related to H/W speed

Mutual Exclusion → If P₁ or P₂ is inside critical section the P₂ or P₁ is not allowed to enter in

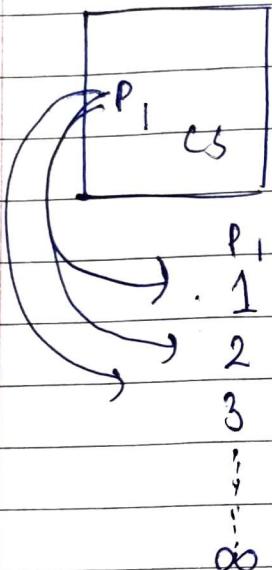


2. Progress → when both can use common code there's progress.



P₁ can't enter due to some code written in P₂ so there is no progress.

3. Bound Wait - There shouldn't be a condition of starvation as one process uses the CPU every time. Let every process use the CPU.



P ₁	CS
1	0
2	0
3	0
⋮	⋮
∞	⋮

If P₁ uses only the CS it will go in infinite loop

4) H/w speed - Every solution should be portable and universal. There shouldn't be anything dependent on H/w speed.

Now Solution to Critical Problem

L-3.5

Critical section solⁿ using lock.

CS - critical section

Date _____ / _____ / _____

do {
 acquire lock.
 CS
 release lock
}.

↑ entry code

1. while (LOCK == 1);
2. LOCK = 1
3. Critical section
4. LOCK = 0.

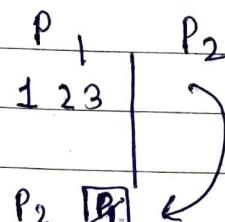
↓ exit code.

- * Execute in user mode.
- * Multiple process solution.
- * No mutual exclusion
- * Guarantee

Case 1 : P₁ P₂

lock = 0 → CS is vacant

lock = 1 → CS is full



lock ≠ 0 ≠ 1

Se

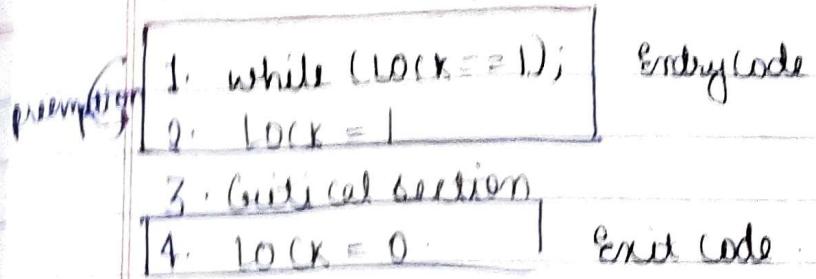
Case 2:

P ₁	P ₂
1 2 3	1 2

L = 0 X 1

P₂ P₁ → No mutual exclusion
CS

3.6 Test and Set Instruction



It combines 1 & 2

`while (test_and_set (&lock));`

CS

`lock = false;`

`boolean test_and_set (boolean *target)`

`boolean r = *target ;`

`*target = TRUE ;`

`return r ;`

}

lock	target	re	P ₁	P ₂
false	1000	false		
1000		1	2	1

True

Mutual exclusion as well as Progress achieved.

3.7 Turn variable

1) 2 process solution

2) Run in user mode.

Process "P ₀ "	Process "P ₁ "
<code>Entry code → while (turn != 0)</code>	<code>while (turn != 1)</code>
<code>CS</code>	<code>CS</code>

`Exit code → turn = 1 ;`

`turn = 0 ;`

int turn = 0; 1

P0

CS

P1

CS

P1 X ,]

P0 X ,]

Mutual exclusion is

satisfied

2) Progress (there is no progress)

CS

P1 X

when turn = 0

3) Bounded wait \rightarrow so there is
no bound wait as turn
value changes.

4) Not H/W dependent

3.8 Semaphore

Semaphore (It is a tool to prevent race cond)

Counting
from to ∞)

Binary.
(0,1)

Running cooperative
process side
by side

Semaphore is an integer
variable which is used in
exclusive manner by
various concurrent
cooperative process in order
to achieve synchronization

Entry section

Date _____ / _____

Down (~~lock~~ Semaphores)

{

s.value = s.value - 1;

if (s.value < 0)

{

put process (P(B)) in
suspended list sleep();

}

else.

freeturn;

}

Exit section



UP (Semaphores)

{

s.value = s.value + 1

if (s.value < 0)

{

Select a process

from suspended list

wake up();

{

.

P(), Down Wait
V(), UP, (signal, post, release)

P₁ P₂ P₃

entry code

CS :- :-

exit section

P₁ P₂ P₃ P₄ (process control block)
P₅

CS	P ₁
	P ₂
	P ₃

Block list (sleep())

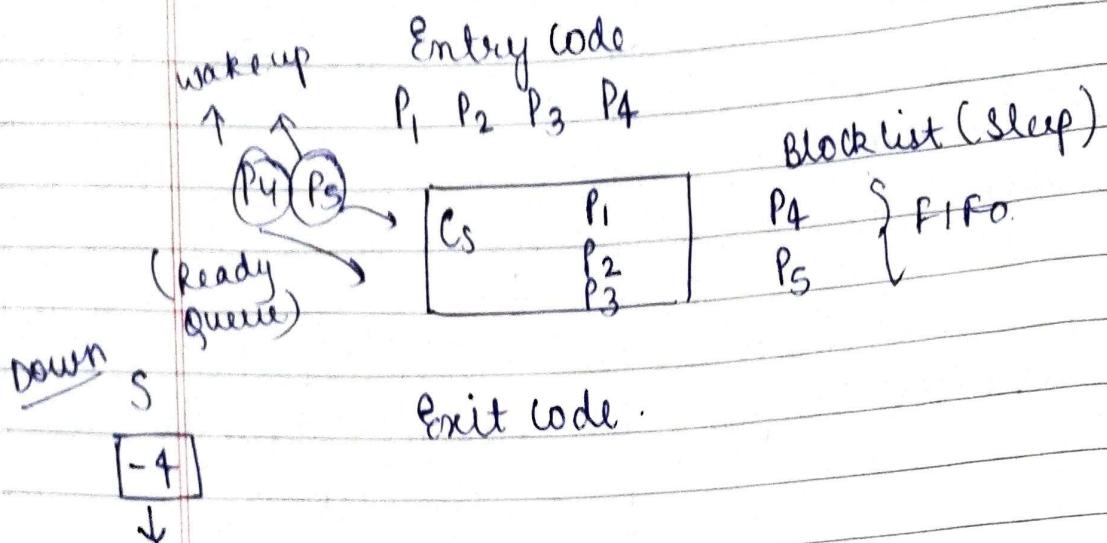
P₄

P₅,

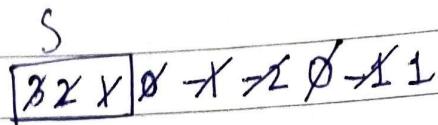
S

3 X X X X (s.value < 0)

-2



4 processes are there which is in CS.



S
0

~~Further no process can come in CS.~~

up S [0] No process in suspend list

S
10 How many process can successfully come in CS.

$$S = 10$$

Successful operation = 10

10 6 P final value of
 4 V \$1

$$10 - 6 = 4 + 4 = 8$$

S
17

SP 3V 1P

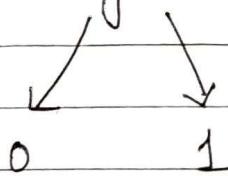
$$= 17 - 5 + 3 - 1$$

$$= 12 + 3 - 1$$

$$= 15 - 1$$

$$= 14$$

3.9 Binary Semaphore (integer variable



shared by a lot
of process in
mutual exclusive
manner).

Semaphore $\rightarrow \infty$ to ∞ .

Binary Semaphore $\rightarrow 0, 1$

Down (Semaphores)

{

if (value == 1)
{

s.value = 0;

{

else
{

Block this Process.

And place in suspend

list sleep();

{

{

Up (Semaphores)

{

if (suspend list is
empty)

{

s.value = 1;

{

else {

{

Select a process from
suspend list &

wake up();

{

{

Down, p, wait

Up, v, signal

$S = 1$
Down } successful
 $S = 0$ openⁿ.

$S = 0$
sleep(); } unsuccessful
block openⁿ.

$S = 0$ $S = 1$
Up Up
 $S = 1$ $S = 1$

In up we check if the block queue is empty or not ($S = \emptyset 1$)

If empty ($S = \emptyset 1$)

If not empty \rightarrow It will bring process from block state to unblock state.

$S = 01$

P_1
Down(s)

[CS]

Up(s)

P_2

Down(s)

[CS]

Up(s)

- Q. Each process P_i { $i = 1$ to 9} execute the following code.

Date

 $P_i \{ i=1 \text{ to } 9 \}$ P_{10} 

entry-section → repeat
 $P(\text{mutex})$
 [CS]
 $V(\text{mutex})$
 exit-section → forever.

mutex
 [1]

repeat
 $V(\text{mutex})$
 [CS]
 $V(\text{mutex})$
 forever.

What is the max^m no of process that may present in CS at any point of time?

(Case 1) mutex 1
 $P(\text{mutex})$
 mutex 0.

[CS.]

$V(\text{mutex})$ mutex = 1.

(Case 2) mutex

$P(\text{mutex})$ 1

0

P₂ P₃
 P₄
 (block) → [CS.
 P₁]

for P₁₀

mutex : & V & X & Z P₁ P₁₀ P₂ Block

CS
 P₁ P₁₀ P₂ ↓ P₃ P₃
 P₄
 P₅

P₁₀
 P₁ P₂ P₃ P₄ P₅ ← P₅
 P₁ P₂ P₃ P₄ P₅ ← P₁₀ P₁ P₂ P₃ P₄ P₁₀ P₉ No. □

maxⁿ no of process can be 10 but putting P₁₀ in and out of critical section.

Soultion of Producer Consumer using Binary Semaphore

Counting Semaphore

- full 0 = No of filled slots
- Empty N → No of empty slots

Produce item(item p);
 1) down(empty);
 2) down(s);
 3) Buffer [IN] = item p;
 4) IN = (IN + 1) mod n;
 5) up(s);
 6) up(full);

Consumer.
 1) down(full);
 2) down(s);
 3) item c = Buffer[OUT];
 4) OUT = (OUT + 1) mod n;
 5) up(s);
 6) up(empty);

N=8

IN	0	a	OUT
	1	b	
	2	c	
	3	d	
3	4		0
	5		
	6		
	7		

$$\text{Empty} = 5 \quad S = X \neq 1 \quad IN = (3+1) \bmod 8 \\ \text{full} = 7 \quad 4 \bmod 8 \\ \text{full} = 7 \quad 4$$



for consumer.

Empty = $\emptyset \neq 5$ In = 4

full = $\neq 3$ Out = $\emptyset \neq 1$ $((0+1) \bmod 8)$

0	
1	b
2	c
3	d
4	
5	
6	
7	

With preemption (context switching)

Empty = $5 \neq 5$

full = $3 \neq 2$ (consumer in CS)

S = $\emptyset \neq \emptyset$

Out = $\emptyset \neq 1$ \hookrightarrow consumer can't do down of S.

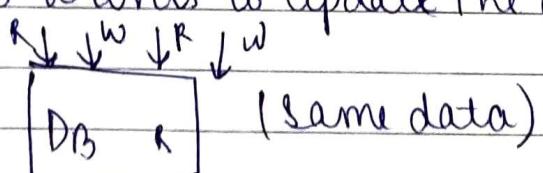
In = $3 \neq 4$

3.12 \rightarrow Solution of Readers writers problem

We have a database used by Reader & writer

Reader - who only reads the data

Writer - who wants to update the data also.



R-W (problem)

W-R (problem)

G-W (problem)

R-R (no problem)

read count
no of readers in buffer.

```

int rc=0;
Semaphore mutex=1;
Semaphore db=1;
Void Reader(void)
{
    while (true)
    {
        down(mutex);
        rc=rc+1;
        if (rc==f) then
            down(db);
        up(mutex);
    }
}
Void writer(void)
{
    while (true)
    {
        down(db);
        DB
        up(db);
    }
}

```

C S. DB

To synchronise problem we used Semaphores.

Case 1: R₁ comes first

$RC = 0 \oplus 1$ mutex = 1 $\oplus 1$ db = 1 $\oplus 0$

down(mutex)

DB R₁

For writer

$RC = 1$ mutex = 1 db = 0

down(db)

\rightarrow db ≠ -1 so the process gets blocked.

Case 2: W₁ comes first



Date: _____

$$R_c = \emptyset \times 1 \quad \text{muter} = X \emptyset \cdot \quad db = X \emptyset$$

down(db)

$db \neq -1$ can't be negative
so r, gets blocked.

Case 3: write-write problem

when writer w, comes down(db) $db = X \emptyset$

when w_2 comes $db \neq -1$ so it gets blocked

Case 4: Read -Read

$$R_c = \emptyset \times 2 \quad \text{muter} = X \emptyset \times \emptyset \cdot \quad db = X \emptyset$$

$[DB \ R_1 \ R_2]$

$2 = 1$ no

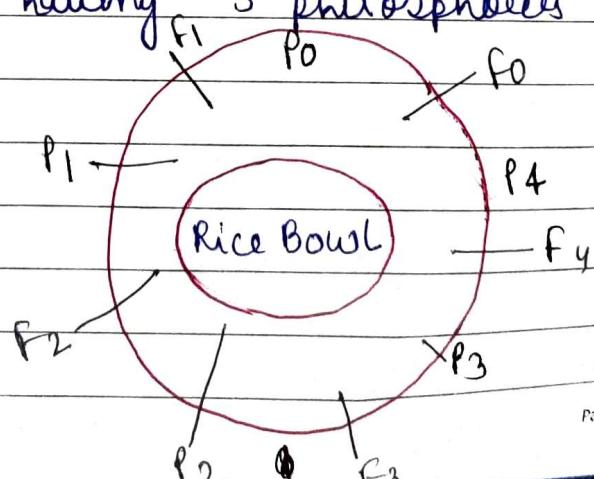
down(~~muter~~)

$db); X$

up(muter)

3.13 Dining Philosophers Problem.

A dining table having 5 philosophers & 5 forks



void philosopher (void)

{

while (true)

{

Thinking();

table_fork(i); \leftarrow Left fork

table_fork((i+1) % N); \leftarrow Right fork

EAT();

put_fork(i);

put_fork((i+1) % N);

}

}.

Philosopher \rightarrow think
 \rightarrow eat

Case 1: P_0 comes

i = 0

left fork = 0 (f_0)

right fork = $(0+1) \% N = 1 (f_1)$

Eating

f_0

f_1

end.

Case 2: P_1

i = 1

left fork = 1 (f_1)

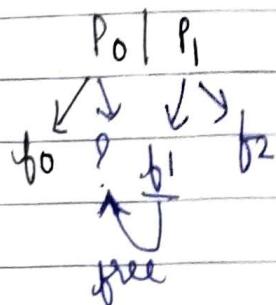
right fork = 2 (f_2)

Eating

f_1

f_2

Case 3:



"problem of race cond" occurs.

$S[i]$ → use array of semaphores

Initially → $\begin{matrix} S_0 & S_1 & S_2 & S_3 & S_4 \\ 1 & x & 1 & 1 & 1 \end{matrix}$

(Initialize every semaphore with 1 when initialized with 0, it waits & get blocked)

Now

void philosopher(void)
{

while (True)
{

Thinking();

Wait (table-fork(i));

Wait (table-fork((i+1) mod N))

EAT();

Signal (Put fork(i));

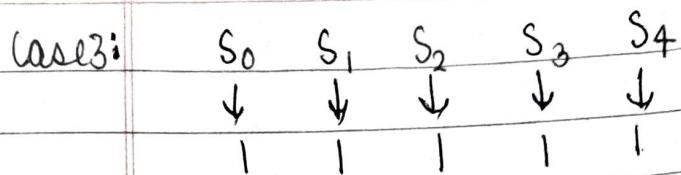
Signal (Put fork((i+1) mod N))

}

.

$P_0 \rightarrow S_0$	S_1
$P_1 \rightarrow S_1$	S_2
$P_2 \rightarrow S_2$	S_3
$P_3 \rightarrow S_3$	S_4
$P_4 \rightarrow S_4$	S_0
Process 6	

But in db P_0 & P_2 can come. It is a special case of mutual exclusion as P_1 & P_2 are independent of each other.



P_0 comes first

S_0
↓
0

It gets preempt

P_1 comes

S_1
↓
0

preempt

P_2 comes

S_2
↓
0

preempt

P_3 comes

S_3
↓
0

preempt

P_4 comes
↓
 S_4
↓
0

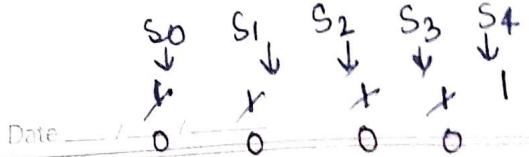
preempt

for eight fork (blocked)

The take left hand side fork belt gets blocked for the right one. This situation is called deadlock (All process get blocked)

Q: How to remove deadlock?
On changing the sequence of one process.

P_4 S_0 S_4



blocked $\rightarrow P_4$

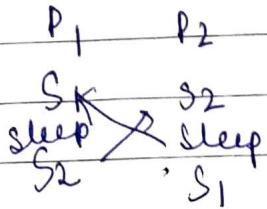
$s_0 \quad s_4$

We can change sequence of any philosopher.

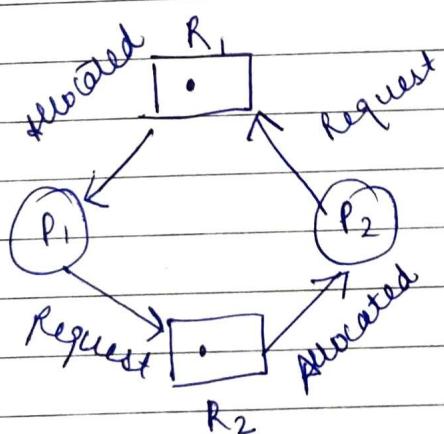
N^{th} philosopher.

wait (take_fork ($s_{(i+1) \bmod N}$))
 wait (take_fork (s_i));

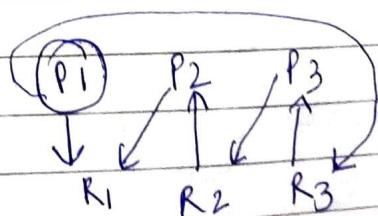
4.1 Deadlock concept



If two or more process are waiting on happening of some event which never happened we say these processes are involved in a deadlock then that state is called deadlock.



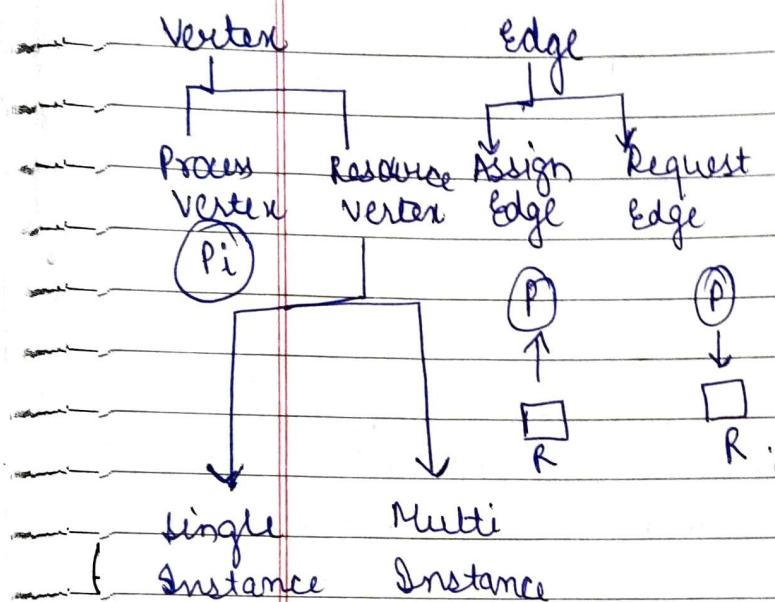
$R_1 \rightarrow P_1$
 R_1 is allocated to P_1
 P_2 request for R_1



Necessary conditions for deadlock.

- 1) Mutual Exclusion
- 2) No Preemption
- 3) Hold & Wait
- 4) Circular Wait

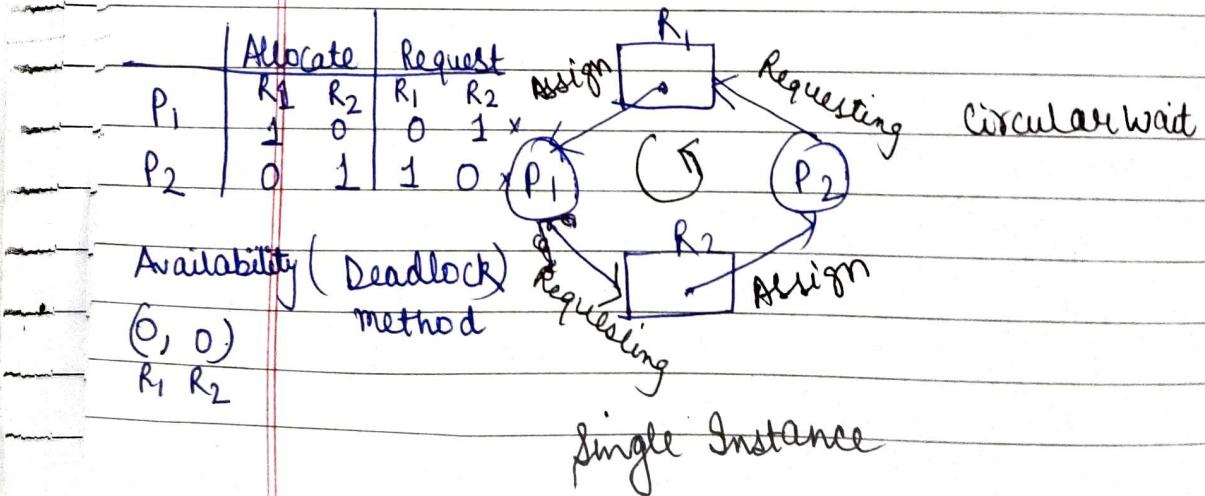
4.2 Resource Allocation Graph in Deadlock (RAG)

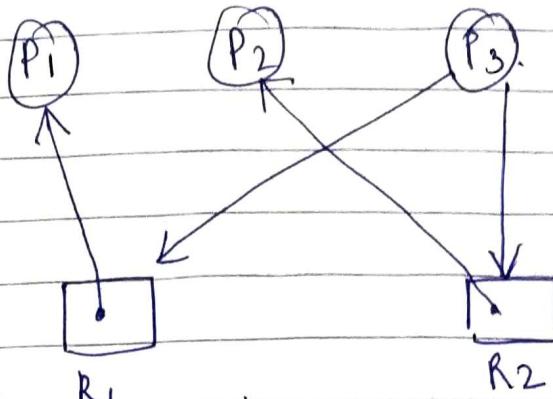


eg: CPU

Monitor

eg : Register





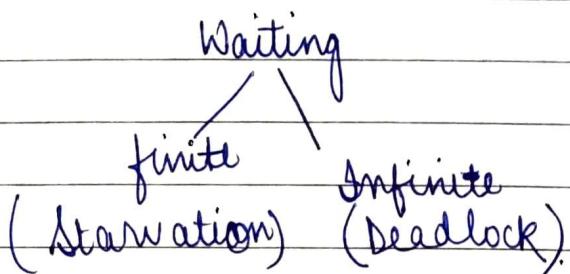
	Allocation		Request		
	R1	R2	R1	R2	
P1	1	0	0	0	✓ Terminate
P2	0	1	0	0	✓ Terminate
P3	0	0	1	1	✓

Availability (0, 0)

(1, 0) → after P1 terminate

(1, 1) → after P2 terminate

No deadlock occurs.



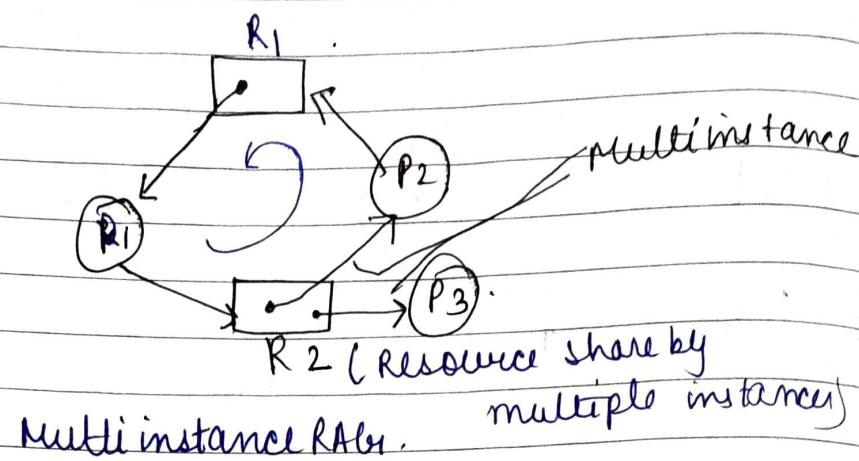
If RAG (Circular wait) cycle → Always in Deadlock

↓
A single instance

Single instance + Circular wait = Deadlock.
Multi instance + CW = Possible (may or may not)

Date _____ / _____ / _____

4.3 Multiinstance RAGs.

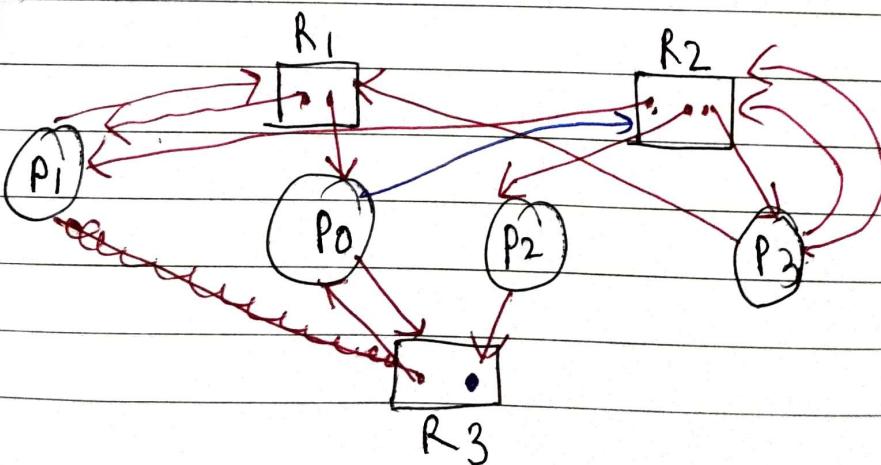


Process	Allocate		Request	
	R ₁	R ₂	R ₁	R ₂
P ₁	1	0	0	1 ✓
P ₂	0	1	1	0
P ₃	0	1	0	0 ✓

Current avai $\rightarrow (0, 0)$ -

0 L -
1 0 ✓

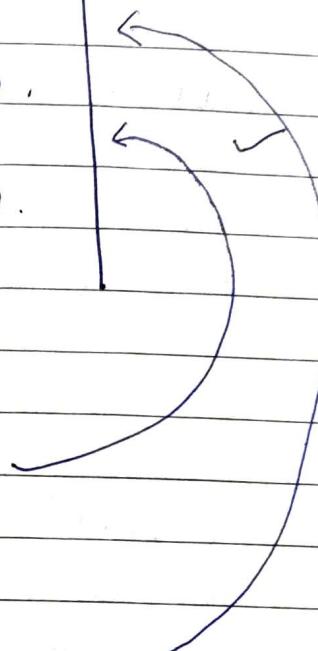
No deadlock



	Allocate		Request	
	R ₁	R ₂	R ₁	R ₂
P ₀	1	0	1	
P ₁				
P ₂				
P ₃				

	Allocate			Request		
	R ₁	R ₂	R ₃	R ₁	R ₂	R ₃
P ₀	1	0	1	0	0	1
P ₁	1	1	0	1	1	0
P ₂	0	1	0	0	0	1
P ₃	0	1	0	1	2	0

curr aw = $\begin{pmatrix} R_1 & R_2 & R_3 \\ 0 & 0 & 1 \end{pmatrix}$



$$P_2 \rightarrow \underline{\quad 0 \quad 0 \quad 1 \quad 0 \quad}$$

$$\underline{\quad 0 \quad 1 \quad 1 \quad}$$

$$P_0 \rightarrow \underline{\quad 1 \quad 0 \quad 1 \quad}$$

\rightarrow No deadlock in the system.

	R ₁	R ₂	R ₃
	1	1	2
P ₁	1	1	0

	R ₁	R ₂	R ₃
	2	2	2
P ₃	0	1	0

$\underline{\quad 2 \quad 3 \quad 2 \quad}$ - curr availability

$P_2 \rightarrow P_0 \rightarrow P_1 \rightarrow P_3$

4.4 Deadlock Handling methods

1. Deadlock ignorance (Ostrich method)
2. Deadlock prevention
3. Deadlock avoidance (Banker's Algo)
4. Deadlock detection & Recovery

Deadlock ignorance
↓

(Just ignore the deadlock.)

Deadlock occurs very rare. Windows has lots of code in this. We want more & more speed so we don't write any code for deadlock, so it's easy to avoid as it's rare.

Why ostrich method called?

Whenever there is a sand storm Ostrich puts its head in the sand assuming no sand. In the same way we ignore deadlock so that the performance & speed doesn't get degraded.

Deadlock prevention

Before the deadlock occurs find the prevention 4 necessary conditions for deadlock

- 1) Mutual Exclusion
- 2) No preemption
- 3) Hold & wait
- 4) Circular wait

either remove all the conditions or deny to false any one of the situation.

- 1) Mutual exclusion (there should be no sharing betⁿ diffⁿ process at same time).

So, if we make all resources sharable we can remove mutual exclusion but some resources like printer can't be made shareable.

- 2) No preemption (there should be any preemption in betⁿ the resources)

So, if we make the process preempted using timestamp or TQ we can prevent deadlock.

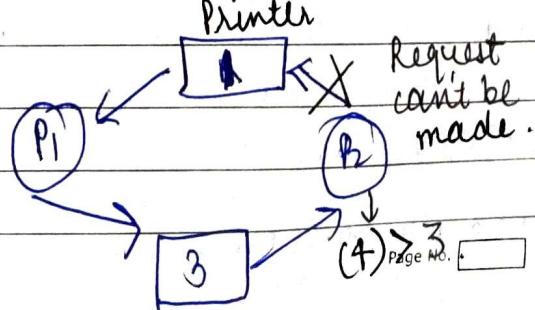
- 3) Hold wait (Holding some resource & waiting for some resources).

Before the process starts, give all the resources to the process.

- 4) Circular wait

Order all the resources (like numbering). Whenever a process request for a resource it would be ordered in increasing order. A process can request in increasing order:

1. Printer
2. Scanner
3. CPU
4. Register



Deadlock Avoidance

while giving the resource to the process check if its safe or not.

Deadlock detection & Recovery

Step 1: Firstly check whether there is a deadlock or not using RAG (Resource Allocation graph).

Step 2: After detection do some recovery

- 1) Kill the processes or a process (lower priority resources)
- 2) Resource preemption (preempt the resource it is holding)

4.5 Deadlock Avoidance (Banker's Algorithm)

Deadlock Avoidance Algorithm

while providing the resources to the process we check if deadlock occurs or not. It is also called Deadlock prevention.

$$\text{CPU (A)} = 10 \text{ (7)}$$

$$\text{Memory (B)} = 5 \text{ (2)}$$

$$\text{Printer (C)} = 7 \text{ (5)}$$

(Total AT) Allocation \rightarrow how much is allocated

Max Need \rightarrow how much is needed

(Total AT - A) Availability \rightarrow how much resource is available

(Max Need - Allocation) Remaining Need

Process	Allocation			Max Need			Available			Remaining need Max Allocation		
	(10-7)	(5-2)	(5-2)	A	B	C	A	B	C	A	B	C
→ P ₁	0	1	0	7	5	3	3 ₍₂₎	3 ₍₁₎	2 ₍₀₎	7	4	3
→ P ₂	2	0	0	3	2	2	5 ₍₂₎	3 ₍₁₎	2 ₍₁₎	1	2	2
→ P ₃	3	0	2	9	0	2	7 ₍₀₎	4 ₍₀₎	3 ₍₂₎	6	0	0
→ P ₄	2	1	1	4	2	2	7 ₍₀₎	4 ₍₁₎	5 ₍₀₎	2	1	1
→ P ₅	0	0	2	5	3	3	7 ₍₀₎	5 ₍₀₎	5 ₍₂₎	5	3	1
+							10	5	7			
	7	2	5									

If we can't fulfil need of every process it is called deadlock.

P₂ → P₄ → P₅ → P₁ → P₃ (safe sequence)
 ↓
 (sequence having no deadlock)

In real life it is not possible.

4.6. Gate Question.

Process	Allocation			Max			Available			Remaining need		
	A	B	C	A	B	C	A	B	C	A	B	C
P ₀	1	0	1	4	3	1	3 ₍₁₎	3 ₍₀₎	0 ₍₁₎	3	3	0
P ₁	1	1	2	2	1	4	4 ₍₁₎	3 ₍₀₎	1 ₍₂₎	1	0	2
P ₂	1	0	3	1	3	3	5 ₍₁₎	4 ₍₀₎	3 ₍₀₎	0	3	0
P ₃	2	0	0	5	4	1	6 ₍₂₎	4 ₍₀₎	6 ₍₀₎	3	4	1
+							8	4	6			
	5	1	6									

4.7 Question Explanation on Deadlock

Q A system requires 2 units of resources having 3 process. The min^m. no of units of 'R' such that no deadlock will occur.

- a) 3
- b) 5
- ~~c)~~ 6
- d) 4

$P_1 \quad P_2 \quad P_3$

$$3 \times 2 = 6 \text{ resources}$$

but we need minimum no of resources.

$$R = 2$$

$P_1 \quad P_2 \quad P_3$.

1 1

↑
need
1 more ↑
need
1 more \rightarrow deadlock.

$P_1 \quad P_2 \quad P_3$
1 1
1

$$P_1 \rightarrow P_2 \rightarrow P_3$$

but not always true

$P_1 \quad P_2 \quad P_3$
~~1~~ 1 1
~~1~~

~~X~~

$R = 3$

P_1	P_2	P_3
1 1 1	1 1 1	1 1 1
$\{x\}$	$\{x\}$	$\{x\}$

P_1	P_2	P_3
1	1	1

\rightarrow deadlock can occur

$R = 4$

P_1	P_2	P_3
1 1 x	1 x	1 x
$\{x\}$	$\{x\}$	$\{x\}$

\rightarrow Deadlock never occurs.

$$\begin{aligned} & (\text{min req}) + 1 \\ & 3 + 1 = 4 \end{aligned}$$

Max^m resources allocated still deadlock - 1

$$\begin{array}{ccc} P_1 & P_2 & P_3 \\ 3 & 4 & 5 \\ 2 & 3 & 4 \end{array} \quad \begin{array}{l} (\text{min req}) \\ = (4+3+2) + 1 \\ = 10 \end{array}$$

Max^m resources allocated still deadlock = 9

4.8 Create Question on Deadlock.

- Q. Consider a system with 3 processes with 3 instances of same resource type. Each process can request a max. of 'K' instances. The largest value of 'K' that will always avoid deadlock is $\rightarrow 2$.

$$R = 4 \text{ units of } R$$

	P_1	P_2	P_3
$K=1$	1	1	1

	P_1	P_2	P_3
$K=2$	1 1	1 1	1 1

	P_1	P_2	P_3
$K=3$	1 1 1	1 1 1	1 1 1
	\cancel{x}	\cancel{x}	\cancel{x}

	P_1	P_2	P_3
$K=3$	1 1 -	1 - -	1 - -
			Deadlock

Ans. 2.

Another method to find

'R' resources

'n' processes ($P_1, P_2, P_3 \dots P_n$)

'd' demand ($d_1, d_2, d_3 \dots d_n$)

max^m resources but still deadlock

$$\begin{array}{ccc} P_1 & P_2 & P_3 \\ 2 & 2 & 2 \\ 1 & 1 & 1 \end{array} = 6.$$

$= 3 + 1 = 4$ (free from deadlock)

$$(d_1 - 1) (d_2 - 1) (d_3 - 1)$$

$$R \leq \sum_{i=1}^n d_i - n. \quad (\text{there will be a deadlock})$$

$$R > \sum_{i=1}^n d_i - n \quad (\text{there will be no deadlock})$$

$$R + n > \sum_{i=1}^n d_i$$

↓ ↗ total demand

total
resource total
processes

$$4 + 3 > 3 \times 2$$

$$7 > 6 \rightarrow \text{True}$$

4 processes
4 resources

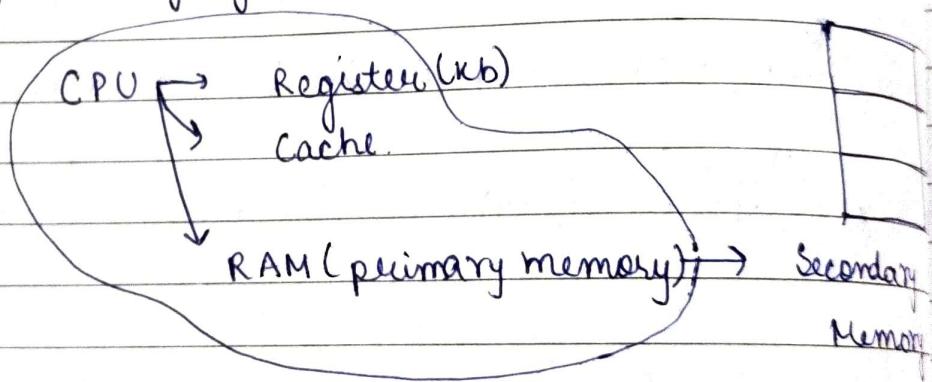
$$4+4 > 4 \times 1 \\ 8 > 4 \quad (\text{Valid})$$

$$4+4 > 4 \times 2 \\ 8 > 8 \quad (\text{false})$$

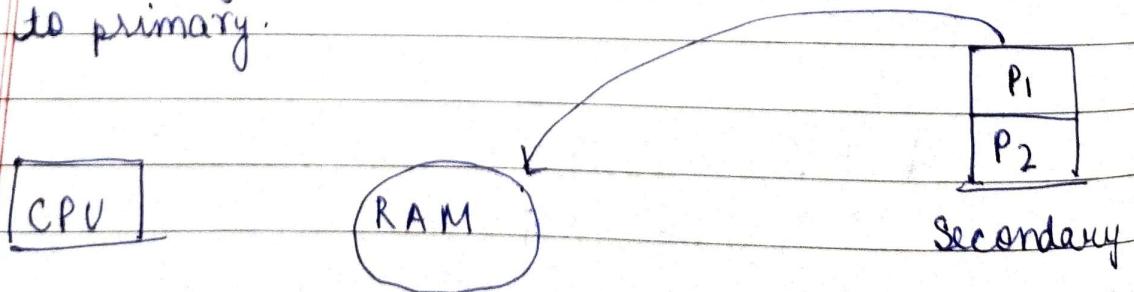
5.1 Memory Management and Degree of Multi programming

It is a kind of functionality to manage all the memory resources in a more & more efficient manner.

Method of managing primary memory.

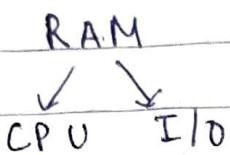


Secondary memory is not directly connected with CPU because secondary memory is a bit slower in comparison to primary.

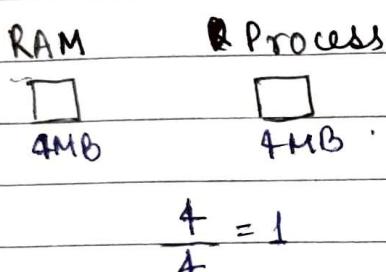


Multiprogramming - Whenever we are giving the programs in RAM being more process in main memory.

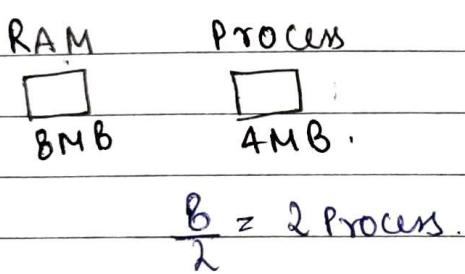
More process in RAM, degree of multiprogramming increases. (more & more processes in RAM to increase efficiency)



When we bring a processes P_1 & if it wants to have some I/O then CPU shouldn't be idle.

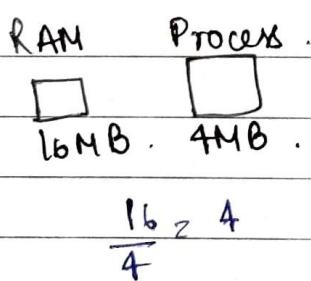


$$k \rightarrow \text{I/O operation (70\%)} \\ CPU = (1-k) \quad (30\%) \\ \text{utilization}$$



$$k^2 \rightarrow \text{I/O operatn.} \\ CPU utilization = (1 - k^2) \\ = 1 - (0.7)^2 \\ = 76\%$$

$$k = 70\%$$



$$k^4 \rightarrow \text{I/O operatn.} \\ CPU utilization = 1 - (0.7)^4 \\ = 93\% \text{ approx.} \\ \downarrow \\ \text{Efficiency increases}$$

- * Increase the no of processes as well as RAM of the system
- * OS needs to pay attention to allocation & deallocation

5.2 Contiguous and non contiguous.

Memory management Techniques

(continuous memory address allocation)

Contiguous

fixed
Partition
(static)
(make fixed sizes)

variable
partition.
(Dynamic)
(provide at run time)

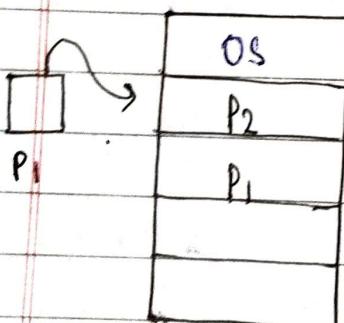
(Non contiguous memory allocation)

Non-contiguous

- Paging
- Multilevel Paging
- Inverted paging
- Segmentation
- Segmented Paging

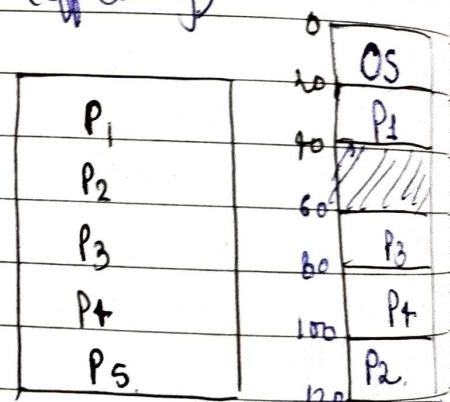
Readystate process of CPU utilization

(efficiency)



Memory Blocks
(RAM)

contiguous allocation



Memory Block

Non-contiguous allocation



5.3 Internal fragmentation | Fixed Size Partitioning

Whenever the process are coming into RAM, how we are allocating it the space.

Points to remember.

- 1) No of partitions are fixed
- 2) size of each partition may or may not be same.
- 3) Contiguous allocation so spanning is not allowed.

- * No of partitions may differ or be same. No of partition is always same.



We can put process in any partition taking into consideration its size should be less than partition size.

	OS .		OS .	
0		4MB.		8MB.
1		6MB.		8MB
2		8MB.		8MB
3		16MB .		8MB

When we allocate any process to the partition. But after allocating process we get some extra memory which is wasted. It is called ~~process~~ Internal fragmentation.

Although we having extra extra memory for diff. partition
we couldn't accomodate a process in pieces for that
fragments.

Limitations

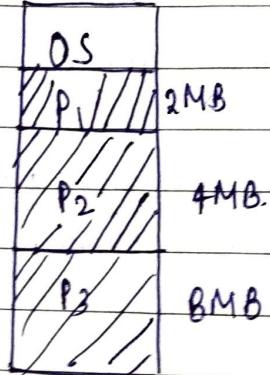
- 1) Internal fragmentation
- 2) limit in process size.
- 3) limitation on degree of multiprogramming.
We can't bring more and more processes.
- 4) External fragmentation

5.4 Variable Size Partitioning

Ram is empty. At run time processes are allocated to
the Ram at runtime.

$$P_1 = 2 \text{ MB}$$

$$P_2 = 4 \text{ MB}$$

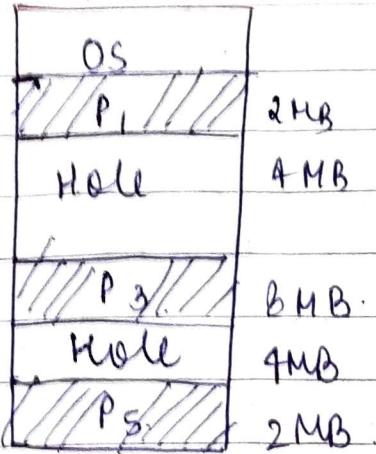


Notes

- 1) No internal fragmentation
- 2) No limitation on number of processes.
- 3) There is no limitation on the process size.

Limitations

- 1) External fragmentation occurs. We can remove it using compaction (one process is copied & pasted to other location) but it takes very much time.
- 2) Deallocation creates Hole. Allocation & deallocation is difficult.
- 3) lot of holes is created



In which we should put which process.

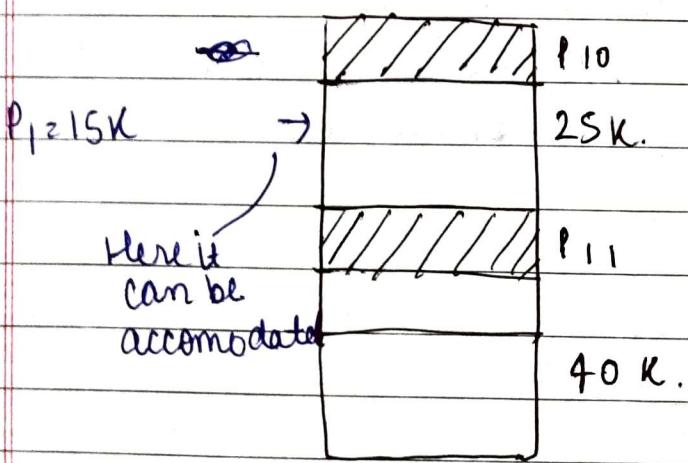
5.5 First fit, Next fit , Best fit Worst fit

First fit \rightarrow Allocate the first hole that is big enough

Next fit \rightarrow same as first fit but start search always from last allocated hole.

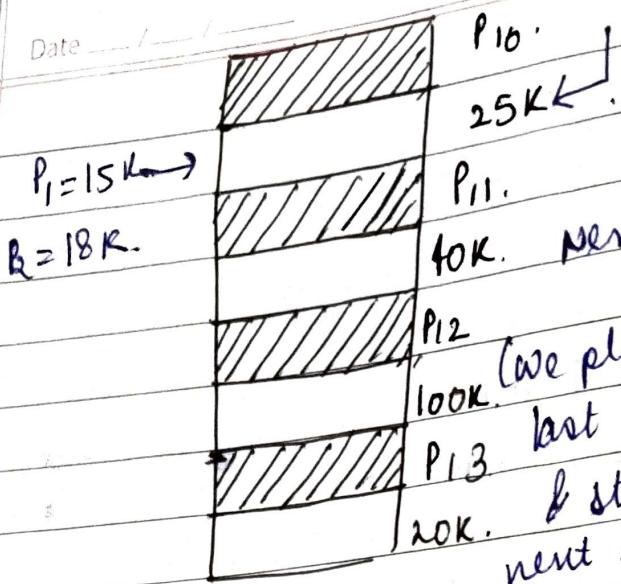
Best fit \rightarrow Allocate the smallest hole that is big enough

Worst fit \rightarrow Allocate the largest hole.



first fit
(find the best
possible large loc^n
which could
accommodate process
& place it in the
RAM)

Date / /



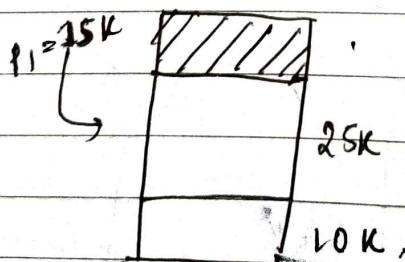
(we place a pointer at
10K. last allocated first fi
20K. & start searching fo
next empty hole. we
don't need to start
from begining.)



Best fit
It would search the entire
list & then return the
hole where there is
min^m fragmentation
It is slow

~~Interval
fragmentation
is least~~

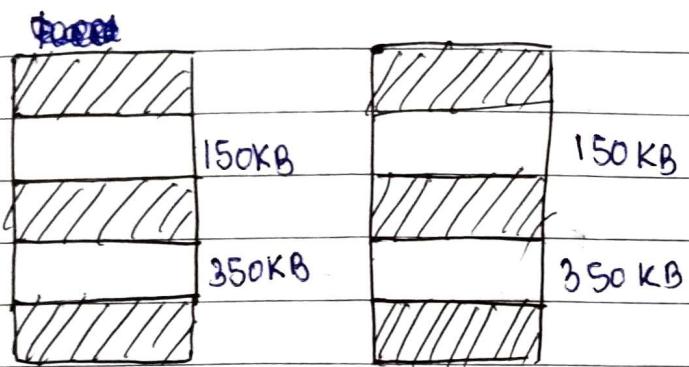
worst fit - It would search the entire list
but would search the biggest hole among
every partitions. (Slow due to searching of
entire list).



5.6 Create Question On ~~fit~~ fits

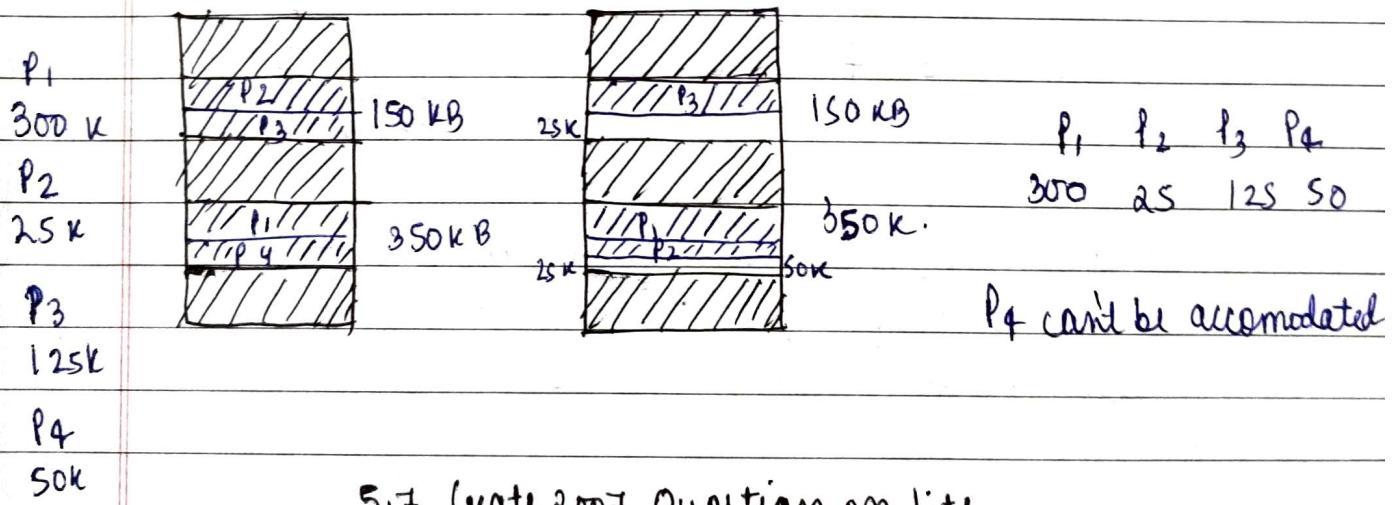
Q. Request from processes are 300K, 25K, 125K, 50K respectively. The above req. could be satisfy with.

- a) best fit but not first fit
- b) first fit but not best fit
- c) Both
- d) None



For first fit

For best fit



5.7 Create 2007 Question on fits

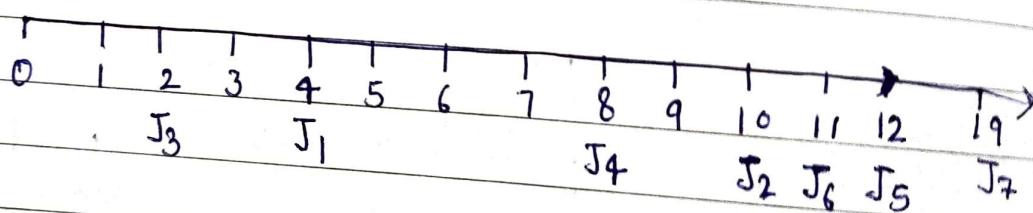
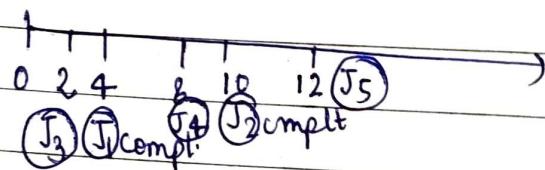
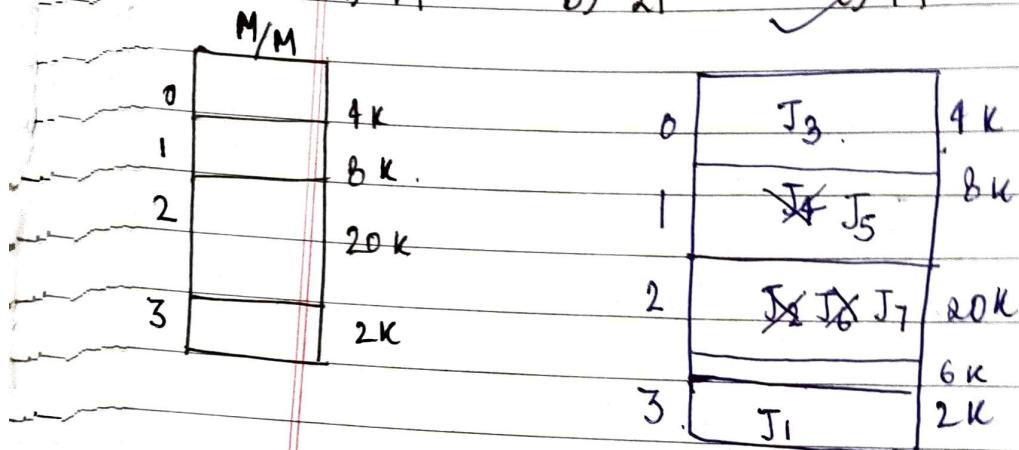
Date _____

Best fit

Req. No	J ₁	J ₂	J ₃	J ₄	J ₅	J ₆	J ₇	J ₈
Request size	2K	19K	3K	6K	6K	10K	7K	20K
Usagetime	4	10	2	8	4	1	8	6.

Calculate the time at which J₇ will be completed

- a) 17 b) 21 c) 19 d) 20

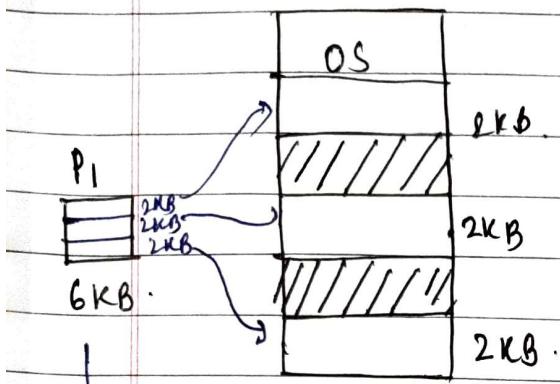


In case of J₅ & J₆ we would take out the one which is completed first. And then put it & calc. the final time by initial + usage-time.



5.8 Need of Paging (Non contiguous Memory Allocation)

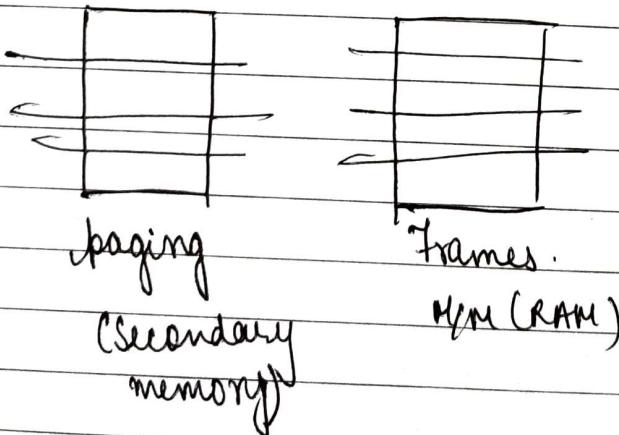
A process can be spanned & put up in fragments in different memory location.



process is divided accⁿ. to the hole size but the hole size changes on runtime so it's a bit time consuming (division of process)

process is divided & put M/M in diff. memory locations. So a process is divided before it gets memory in pages.

When we do partition of process in secondary memory it's called paging while when we do partition of process in RAM it is called framing.



THANK YOU