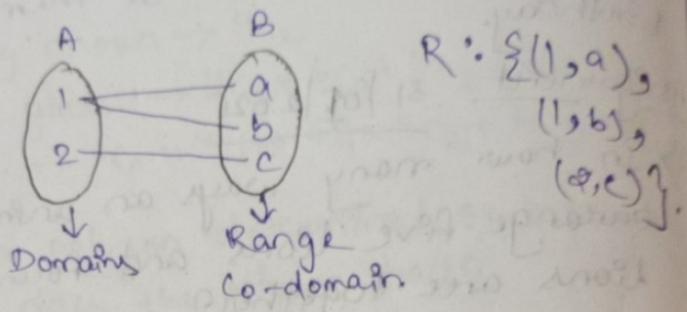


Relations on sets

Let A and B be two sets. A relation R from A to B

$R: A \rightarrow B$ is a subset of $A \times B$.



$$\text{Ex: } A = \{1, 2, 3, 4\}, B = \{1, 2, 3\}$$

$$A \times B = \{(1,1), (1,2), (1,3), (2,1), (2,2), (2,3), (3,1), (3,2), (3,3), (4,1), (4,2), (4,3)\}$$

Define a Relation $R: A \rightarrow B$ by a divisor b .

$$R = \{(1,1), (1,2), (1,3), (2,2), (3,3)\}.$$

$$\text{Domain} = \{1, 2, 3\}$$

$$\text{Range} = \{1, 2, 3\}.$$

Inverse of R :

Let $R: A \rightarrow B$ be a relation then the inverse of R denoted as

$$R^{-1} = \{(y,x) | (x,y) \in R\}$$

$$\Rightarrow R^{-1} = \{(1,1), (2,1), (3,1), (2,2), (3,3)\}$$

Identity Relation:

$$I = \{(a,a) | a \in A\}.$$

Properties of Binary Relation:

A binary relation $R \subseteq A \times A$ is called

(i) Reflexive if $\forall a (a,a) \in R$.

(ii) Symmetric if $\forall (a,y), ((a,y) \in R \Rightarrow (y,a) \in R)$

- (iii) Anti-symmetric if $\forall (x,y)$
- (iv) Transitive if $\forall (x,y)$

Ex:

Q) \leq is

Reflexive

\leq

Q) $=$ is

Yes

Q) $<$ is

No.

Q) If $a \leq$

Q) $\{m,y\}$

Yes.

II. If

→ If A set w/ relation

Ex: If

Q) If and

Composite
Let R_1

(i) Anti-symmetric

If $\forall (x, y) ((x, y) \in R \wedge (y, x) \in R \Rightarrow x = y)$.

(ii) Transitive property

If $\forall (x, y, z) ((x, y) \in R \wedge (y, z) \in R \Rightarrow (x, z) \in R)$

Ex:

Q) \leq is it a reflexive relation?

Reflexive as $=$ is present.

$\leq, =$ reflexive.

Q) \leq is it a symmetric relation?

Yes

Q) \leq is it a symmetric relation?

No.

Q) If $x \leq y$ then it is an anti-symmetric relation.

Q) $\{(x, y) | x+y=3\}$. Is it a symmetric relation?

Yes.

Is it anti-symmetric? No.

→ If A is a set with m elements and B is a set with n elements then the number of possible relations from $A \rightarrow B$ is 2^{mn} .

Ex: If $m=4, n=2$ then the no. of relations

$$2^{4 \times 2} = 2^8.$$

Q) If there are 1024 relations from set A to B and $|B|=2, |A|=?$

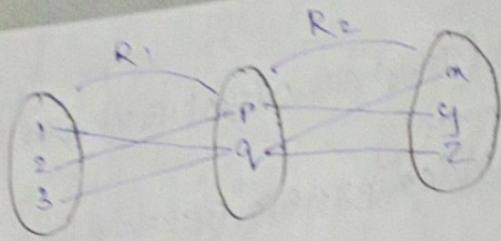
$$|B|=2 \Rightarrow 1024 = 2^{10}$$

$$\Rightarrow |A|=5.$$

Composition of Relations:

Let $R_1 \subseteq A \times B, R_2 \subseteq B \times C$ then

$$R_1 \circ R_2 = \{(x, z) | (x, y) \in R_1 \wedge (y, z) \in R_2\}$$



$$(1, p) \in R_1 \text{ and } (p, m) \in R_2 \Rightarrow (1, m) \in R_1 R_2$$

$$(2, y) = ((2, p) \wedge (p, y)) \in R_1 R_2 \text{ or } R_1 \circ R_2$$

Ex: $A = \{1, 3, 5, 7\}$, $B = \{x, y, z\}$, $C = \{a, b, c\}$

$$R_1: A \rightarrow B, R_1 = \{(1, x), (3, y), (5, x), (7, z)\}$$

$$R_2: B \rightarrow C = \{(x, a), (x, b), (y, a), (z, c)\} \cup \{(x, c), (y, b)\}$$

$$\text{then } R_1 \circ R_2 = \{\ ? \}$$

$$R_2 \circ R_1 = \{\ ? \}$$

$$R_1 \circ R_2 = \{(1, a) \wedge (3, b)\}$$

$$R_2 \circ R_1 = \emptyset$$

Equivalence Relations -

A relation on a set A is called an equivalence relation if it is reflexive, symmetric and transitive.

Ex: $R = \{(a, b) \subseteq \mathbb{Z}^+ \times \mathbb{Z}^+ : a \text{ divides } b\}$

Is it an equivalence relation?

It is not an equivalence relation since symmetric property does not satisfy the above relation.

Ex: Show that the relation $R = \{(a, b) \mid a \equiv b \pmod{m}\}$ is a equivalence relation

Is it a equivalence relation

\rightarrow Reflexive: aRa since m divides $a-a=0$

\rightarrow Symmetric: $(a, b) \in R$

$\Rightarrow a-b = km$ for some integer k

$$\Rightarrow -(a-b) = (-k)m$$

$\rightarrow b$
 $\rightarrow C$
 R is
 suppose a
 $\Rightarrow a-b =$
 Adding
 $a-$
 $\Rightarrow D$
 \Rightarrow

since R
 R is

Reflexive or

Let R be

Then,

$$R_A =$$

$$R \cap R_A$$

$$R \cap R_A^{-1}$$

In other

Transitive

If R

Then

$$R^* =$$

$$\text{Ex: } A =$$

$$R$$

$$\Rightarrow b-a = km$$

$$\Rightarrow (b,a) \in R$$

R is symmetric.

→ Transitive :-

Suppose aRb and bRc

$$\Rightarrow a-b = km \text{ and } b-c = lm \quad (\text{for some integers } k, l)$$

Adding the above relation

$$a-b+b-c = km+lm$$

$$\Rightarrow a-c = (k+l)m$$

↑ integer

$\Rightarrow m$ divides $a-c$

or $(a,c) \in R$ or R is transitive

Since R is reflexive, symmetric & transitive then

R is an equivalence relation.

Reflexive and symmetric closure :-

Let R be a relation on a set A

Then,

$$R_A = \{(a, a) \mid a \in A\}$$

$R \cup R_A$ is Reflexive closure of R.

$R \cup R^{-1}$ is symmetric closure of R.

In other words if $(a, b) \in R$, then

add (b, a) .

Transitive closure :-

If R is a relation on a set A with n elements

then

$$R^* = R \cup R^2 \cup R^3 \cup \dots \cup R^n$$

$$\text{Ex:- } A = \{1, 2, 3\}$$

$$R = \{(1, 2), (2, 3), (3, 1)\}$$

$$R^2 = R \cup R^2 \cup R^3$$

$$R^2 = R \circ R = \{(1,3), (2,3), (3,3)\}$$

$$R^3 = R^2 \cup R = \{(1,3), (2,3), (3,3)\}$$

$$\text{Transitive closure of } R = R \cup R^2 \cup R^3$$

$$= \{(1,2), (1,3), (2,3), (3,3)\}$$

$$\rightarrow R = \{(1,2), (2,3), (3,1)\}$$

$$R = \{(1,2), (2,3), (3,1)\}$$

Reflexive closure:

$$R_A = \{(1,1), (2,2), (3,3)\}$$

$$R \cup R_A = \{(1,2), (2,3), (3,1), (1,1), (2,2), (3,3)\}.$$

Symmetric closure:

$$R^{-1} = \{(2,1), (3,2), (3,3)\}$$

$$R \cup R^{-1} = \{(1,2), (2,3), (3,1), (2,1), (3,2), (3,3)\}.$$

Transitive closure of R° :

$$R^2 = R \circ R = \{(1,3), (2,1), (3,2)\}$$

$$R^3 = R^2 \circ R = \{(1,1), (2,2), (3,3)\}.$$

Transitive closure of R :

$$R \cup R^2 \cup R^3$$

$$= \{(1,1), (3,2), (2,1), (3,3), (2,1), (1,3)\}$$

Posets (Partially Ordered set):

A partially ordered set or a poset is a pair

$P = (X, \leq)$ where X is a non-empty set.
 \leq is a partial order.

i.e. for $x, y \in X$

1. $x \leq x$ (reflexive)

2. $x \leq y, y \leq x$ implies $x = y$ (Anti-symmetric)

3. $x \leq y, y \leq z$ implies $x \leq z$ (transitive).

\hookrightarrow partial ordered set.

\rightarrow A relation
relation is
transitive

$Q(2,3) \rightarrow P$

$a \geq a$

$\therefore Y$

(ii) Anti-Sym

$a \neq b$

$\Rightarrow a$

$\therefore Y$

Posit Trans

a

$\Rightarrow c$

$\therefore Y$

$\therefore \text{This}$

Q) Sha

orde

(i) "

(ii) "

(a)

(b) Re

→ A relation R is said to be partially ordered relation if it is reflexive, anti-symmetric and transitive.

$Q(\mathbb{Z}, \geq)$, partially ordered relation. Check

$$2 \geq 2 \rightarrow$$

(i) Reflexive -

$$\forall a, a \geq a \quad \forall a \in \mathbb{Z}$$

$\therefore \geq$ - reflexive

(ii) Anti-symmetric -

$$a \geq b, b \geq a \Rightarrow a, b \in \mathbb{Z}$$

$$\Rightarrow a = b.$$

$\therefore \geq$ - anti-symmetric.

(iii) Transitive -

$$a \geq b, b \geq c \quad \forall a, b, c \in \mathbb{Z}$$

$$\Rightarrow a \geq c$$

$\therefore \geq$ is transitive

\therefore This relation is partially ordered relation.

Q) Show that the following divides relation are partially ordered on the set A .

(i) ' $m|n$ ', ' m ' divides ' n ' where $A = \mathbb{N}$ (set of all natural numbers)

(ii) ' $m|n$ ', ' m ' divides ' n ' where

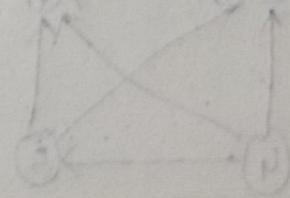
$A =$ the set of +ve integers.

(a)

(i) Reflexive

$$\frac{n}{n} \quad \forall n \in A$$

$\frac{n}{n}$ - Reflexive



(ii) Anti-symmetric

$$\frac{m}{n}, \frac{n}{m} \quad \forall m, n \in \mathbb{Z}$$

$$\Rightarrow m = n$$

$\therefore m|n$ is anti-symmetric.

(iii) Transitive

$$m|n \cdot n|p \rightarrow m|p, \forall m, n, p \in \mathbb{Z}$$

$\therefore m|p$ (transitive).

\therefore The above relation is a partially ordered relation.

all three are

$n|n \rightarrow$ hence the relation is reflexive

\rightarrow If $m|n$ and $n|m$

$\Rightarrow m=n \rightarrow$ anti symmetric

\rightarrow If $m|n$ and $n|p$

$\Rightarrow m|p$ transitive

Step 3% (R)

Step 4%

Ex% - 1

(b) (i) $m|n$, $m|n$ divides n where

$$A = 1, 2, 3, 4, 6, 9, 12, 18, 36.$$

Reflexive

Representation of posets are Hasse diagram% -

A graphical representation of a partial order relation in which all arrows pointing upward is known as Hasse diagram.

Procedure% -

Step - 1% Draw diagram of relation.

Step - 2% Remove self loops.

Step - 3% Remove all transitive edges.

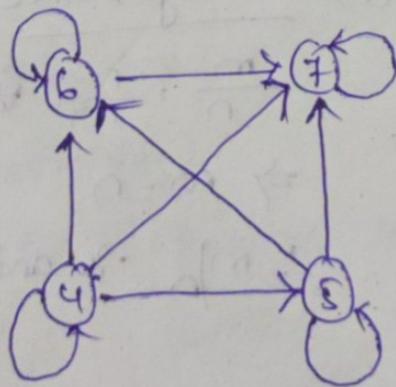
Step - 4% Arrange all edges pointing upward

Step - 5% Replace circles by dots as vertex.

Ex% Considering +ve set $A = \{4, 5, 6, 7\}$. Let R be the relation \leq on A . Draw the Hasse diagram of R .

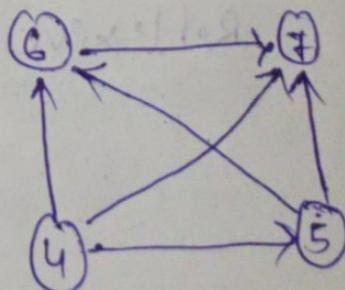
$$R = \{(4, 4), (4, 5), (4, 6), (4, 7), (5, 5), (5, 6), (5, 7), (6, 6), (6, 7)\}$$

Step 1%

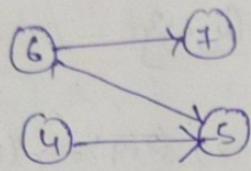


Step 2%

Removing self loop $\rightarrow (4, 4)$
 $(5, 5), (6, 6), (7, 7)$

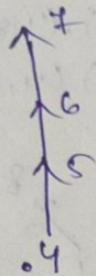


Step 3° (Remove transitive edges)



$$\begin{array}{l} a=b \\ (1,5) \quad (2,6), (5,6)(6,7) \\ (4,6) \quad (4,7) \end{array}$$

Step 4°

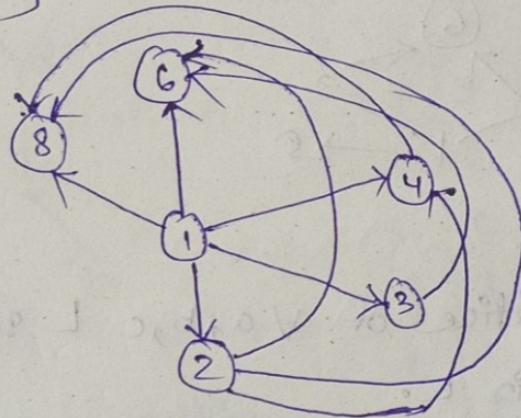


This is own required Hasse diagram.

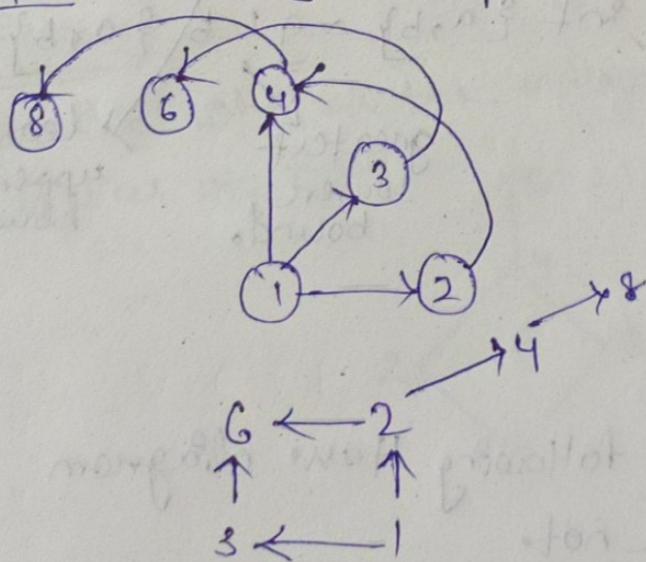
Ex° - Let $A = \{1, 2, 3, 4, 6, 8\}$ be ordered by the relation
'a' divides 'b'.

$$R = [(1,2), (1,3), (1,4), (1,6), (1,8), (2,4), (2,6), (2,8), (3,6), (4,8)]$$

Step 1°



Step 2° Removing self loop



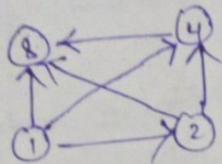
Step 3°

Remove transitive edge

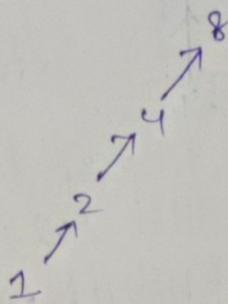
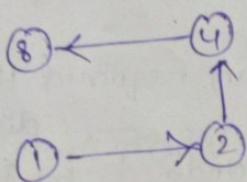
$$\begin{array}{l} (1,2), (2,4) \rightarrow (1,4) X \\ (1,3), (3,6) \rightarrow (1,6) X \\ (1,4), (4,8) \rightarrow (1,8) X \\ (2,4), (4,8) \rightarrow (2,8) X \end{array}$$

Ex:- $D_8 = \{1, 2, 4, 8\}$, The relation is divisibility.

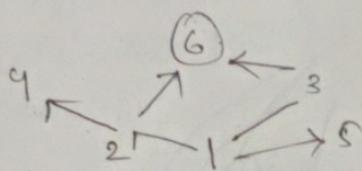
$$R = \{(1,2), (1,4), (1,8), (2,4), (2,8), (4,8)\}$$



$$(1,2) (2,4) \rightarrow (1,4) \\ (1,4), (4,8) \rightarrow (1,8) \\ (2,4), (4,8) \rightarrow (2,8)$$



Ex:- $R = \{1, 2, 3, 4, 5, 6\}$ divisibility.



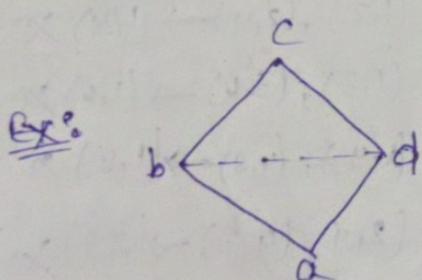
Lattice:-

A poset is a lattice if $\forall a, b, c \in L$ sup{ a, b } and inf{ a, b } exist in L .

$$1. a \vee b = a \text{ join } b = \sup \{a, b\} = \sup \{a, b\}$$

$$2. a \wedge b = a \text{ meet } b = \inf \{a, b\} = g \wedge b. \{a \wedge b\}$$

greatest lower bound. least upper bound



Ex:- Determine whether the following Hasse diagram represents a lattice or not.

Ex 2

Lub

v	a	b	c	d
a	a	b	c	d
b	b	b	c	c
c	c	c	c	c
d	d	c	c	d

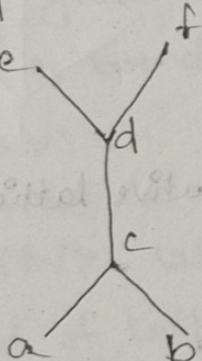
Since each subset of two elements has least upper bound. So, L contain lub and glb.

glb

v	a	b	cd
a	a	a	aa
b	a	b	ba
c	a	b	cd
d	a	a	dd

since each subset of two elements has least upper bound. so L become a lattice and glb.

Bx

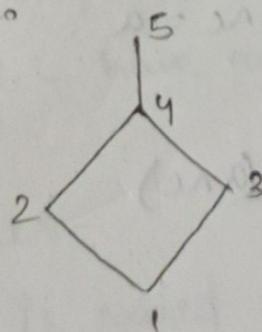


It is not a lattice.

v	a	b	c	d	e	f
a	a	c	c	d	e	f
b	c	b	c	d	e	f
c	c	c	c	d	d	f
d	d	d	d	d	e	ex
e	e	e	e	e	e	f
f	f	f	f	f	f	f x f

Assignment

Determine whether the following diagram represent a lattice or not.



Properties of lattice and distributive lattices

Let (L, \leq) be a lattice & $a, b \in L$

1) Idempotent law, $a \vee a = a$
 $a \wedge a = a$

2) Commutative law, $\forall a, b \in L$

$$a \vee b = b \vee a$$

$$a \wedge b = b \wedge a$$

3) Associative law = $\forall a, b, c \in L$

$$(a \vee b) \vee c = a \vee (b \vee c)$$

$$(a \wedge b) \wedge c = a \wedge (b \wedge c)$$

4) Absorption law for $a, b \in L$

$$(a) a \vee (a \wedge b) = a$$

$$(b) a \wedge (a \vee b) = a$$

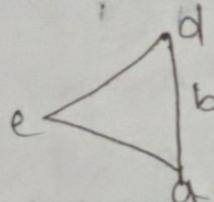
Distributive Lattice

A lattice (L, \leq) is called a distributive lattice if for any $a, b, c \in L$

$$a \vee (b \wedge c) = (a \vee b) \wedge (a \vee c)$$

$$a \wedge (b \vee c) = (a \wedge b) \vee (a \wedge c)$$

Q)



Check whether it is a distributive lattice or not.

$$a \vee (b \wedge c) = a \wedge a = a$$

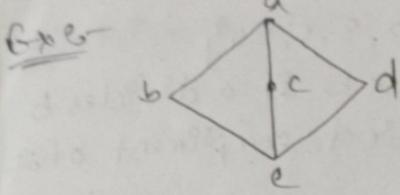
$$(a \vee b) \wedge (a \vee c) = b \wedge c = a$$

$$\therefore LHS = RHS$$

$$a \wedge (b \vee c) = (a \wedge b) \vee (a \wedge c)$$

$$LHS = a \wedge (b \vee c) =$$

$$RHS = (a \wedge b) \vee (a \wedge c) =$$



$b \cap (c \vee d) = (b \cap c) \vee (b \cap d)$
whether it is distributive
lattices or not.

$$\text{LHS} = b \cap (c \vee d) = b \cap a = b$$

$$\text{RHS} = (b \cap c) \vee (b \cap d) = b \cap e = b$$

$$\text{LHS} \neq \text{RHS}$$

Hence, it is a distributive lattice.

Bounded Lattice: A lattice (L, \leq) is called bounded lattice if it has greatest element '1' and least element '0'. Or L has greatest and least element in L .
 $\rightarrow (L, \leq)$ is bounded then for any element $a \in L$

$$a \vee 1 = 1$$

$$a \vee 0 = a$$

$$a \wedge 1 = a$$

$$a \wedge 0 = 0$$

Every finite lattice is bounded.

Complement of Lattice:-

Complement of an element in a lattice (L, \leq) in bounded lattice with least element '0' and greatest element '1'.

$a \in L$, $a' \in L$ is called the complement of a .

$$a \vee a' = 1, a \wedge a' = 0$$

\rightarrow It is not necessary that every element has a complement

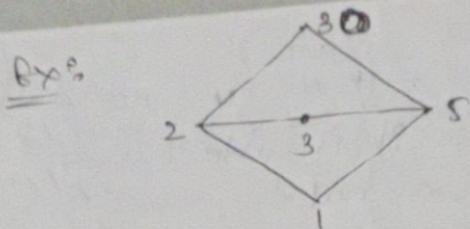
\rightarrow An element $a \in L$ have more than one complement also.

Complemented Lattice:-

A lattice (L, \leq) is called complemented

(i) L is bound

(ii) Every element $a \in L$ has a complement.



$A = \{1, 2, 3, 5, 30\}$
and also if a divides b
then find complement of a.

27 D

$$g \wedge b = 1$$

$$L \cup b = 30$$

$$2 \vee 2' = 30$$

$$2 \wedge 2' = 1$$

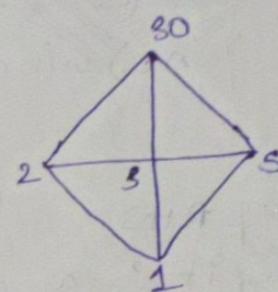
$$2 \vee 3 = 30$$

$$2 \wedge 3 = 1$$

$$2' = 3$$

Complement of 2 is 3.

Complement of 5 is 3.



37 Qd

47 C

If 'a' and 'b' are the elements in a distributive lattice L and if a' complement of a then

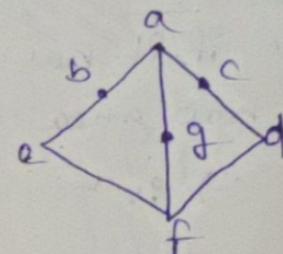
$$(i) (a')' = a$$

$$(ii) (a \vee b)' = a' \wedge b'$$

$$(iii) (a \wedge b)' = a' \vee b'$$

$$(iv) a \vee (a' \wedge b) = a \vee b$$

$$(v) a \wedge (a' \vee b) = a \wedge b$$



some

let

so

o

Boolean Algebra :- A non-empty set B with two binary operation (+) and (.) , a unary operation (') and two distinct element '0' and '1' is called a Boolean algebra denoted by $(B, +, ., ', 0, 1)$ iff the following properties are satisfied.

Axioms of Boolean Algebra

If $a, b, c \in B$, then

1) Commutative Law :-

$$a + b = b + a$$

$$a \cdot b = b \cdot a$$

2) Distributive

$$a + (b \cdot c) = (a + b) \cdot (a + c)$$

$$a \cdot (b + c) = (a \cdot b) + (a \cdot c)$$

3) Identity Law -

$$a + 0 = a$$

$$a \cdot 1 = a$$

4) Complement Law -

$$a + a' = 1$$

$$a \cdot a' = 0$$

Some basic result of Boolean algebra

Let $a, b, c \in B$ then

Idempotent laws

$$a + a = a$$

$$a \cdot a = a$$

$$a = a + 0 \text{ (Identity rules)}$$

$$= a + (aa') \text{ (Complement rules)}$$

$$= (a+a) \cdot (a+a') \text{ (Distributive rules)}$$

$$= 1(a+a) \cdot 1 \text{ (Complement rule)}$$

$$= a+a$$

$$\therefore a+a=a \quad [\text{Hence proved}]$$

$$a \cdot a = a$$

$$a \cdot a = a \cdot a + 0 \quad (\text{Identity law})$$

$$= a \cdot a + a \cdot a' \quad (\text{Complement rule})$$

$$= a(a+a') \quad (\text{Distributive law})$$

$$= a \cdot 1 \quad (\text{Complement rule})$$

$$= a \quad (\text{Idempotent laws})$$

$$\therefore a \cdot a = a$$

2) Boundary laws :-

(a) $a+1 = 1$

(b) $a \cdot 0 = 0$

$$\begin{aligned}
 (a) 1 &= a + a' \quad (\text{complement rule}) \\
 &= a + (a' \cdot 1) \quad (\text{identity law}) \\
 &= (a + a') \cdot (a + 1) \quad (\text{distributive law}) \\
 &= 1 \cdot (a + 1) \quad (\text{complement rule}) \\
 &= (a + 1) \cdot 1 \quad (\text{commutative rule}) \\
 &= a + 1 \quad (\text{identity rule})
 \end{aligned}$$

$$\therefore [a + 1 = 1]$$

$$(b) a \cdot 0 = 0$$

$$\begin{aligned}
 0 &= a \cdot a' \quad (\text{complement rule}) \\
 &= a(a' + 0) \quad (\text{identity law}) \\
 &= (a \cdot a') + (a \cdot 0) \quad (\text{distributive law}) \\
 &= 0 + (a \cdot 0) \quad (\text{complement rule}) \\
 &= (a \cdot 0) + 0 \quad (\text{commutative rule}) \\
 &= a \cdot 0 \quad (\text{identity rule})
 \end{aligned}$$

$$\therefore [a \cdot 0 = 0]$$

Absorption Law :-

Let $a, b \in B$

$$(a) a + (a \cdot b) = a$$

$$(b) a \cdot (a + b) = a$$

$$\begin{aligned}
 (a) a + (a \cdot b) &= (a \cdot 1) + (a \cdot b) \quad (\text{identity law}) \\
 &= a(1 + b) \quad (\text{distributive}) \\
 &= a(b + 1) \quad (\text{commutative law}) \\
 &= a \cdot 1 \quad (b + 1 = 1 \text{ bounded law}) \\
 &= a \quad (\text{identity law})
 \end{aligned}$$

$$(b) a \cdot (a + b) = (a + 0) \cdot (a + b)$$

Augment

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