

Tutorial-2

Ques-1) What is the time Complexity of below code & how?

```
void fun(int n)
```

```
{
    int j=1, i=0;
```

```
    while (i < n)
```

```
    {
        i += j;
```

```
        j++;
```

```
    }
```

```
}
```

j=1	i=1	} m-level
j=2	i=1+2	
j=3	i=1+2+3	

for(i)

$$\therefore 1+2+3+ \dots < n$$

$$1+2+3+m < n$$

$$\frac{n(n+1)}{2} < n$$

$$m = \sqrt{n}$$

By Summation method

$$\sum_{i=1}^m 1 \Rightarrow 1+1+ \dots + \sqrt{n} \text{ times}$$

$$\boxed{T(n) = \sqrt{n}} \text{ Ans.}$$

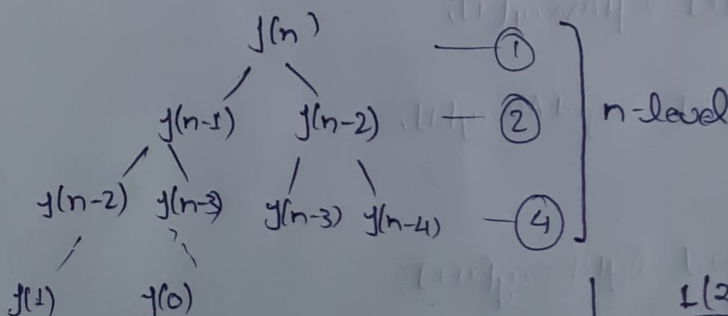
Ques-2) What recurrence relation for function that prints Fib. series. Solve it to get the time Complexity. What will be Space Complexity & why For Fibonacci Series

$$f(n) = f(n-1) + f(n-2)$$

$$f(0) = 0$$

By forming a tree,

$$f(1) = 1$$



$$1, 2, 4, 8, \dots, n$$

$$a=1, x=2$$

$$\frac{1(2^n - 1)}{2 - 1} = 2^n - 1$$

$$\boxed{TC = O(2^n)}$$

Space Complexity

Recursive :- $T(n) = O(n)$

Iterative :- $T(n) = O(1)$

Ques. 3) Write programs which have time complexity $n \log n$, n^3 , $\log(\log n)$

① $n \log n \rightarrow$ Merge Sort

```
void mergeSort(int arr[], const int low, const int high)
```

```
{
    if (low >= high) return;
    int mid = low + (high - low) / 2;
    mergeSort(arr, low, mid);
    mergeSort(arr, mid + 1, high);
    merge(arr, low, mid, high);
}
```

```
void merge(int arr[], const int low, int mid, int high)
```

```
{
    int i, j, k;
    int n1 = mid - low + 1;
    int n2 = high - mid;
    int leftArray[n1], rightArray[n2];
    for (int d = 0; d < n1; d++)
        leftArray[d] = arr[low + d];
```

```
    for (int d = 0; d < n2; d++)
        rightArray[d] = arr[mid + 1 + d];
    i = 0, j = 0, k = low;
```

```
    while (i < n1 && j < n2)
```

```
    {
        if (leftArray[i] <= rightArray[j])
            arr[k] = leftArray[i];
        else
```

```
            arr[k] = rightArray[j];
```

```
    }
```

```
    while (i < n1)
```

```
        arr[k++] = leftArray[i++];
```

```
    while (j < n2)
```

```
        arr[k++] = rightArray[j++];
}
```

(ii) $n^3 \rightarrow$ Multiplication of two matrices.

for ($i=0; i < n1; i++$)

for ($j=0; j < n2; j++$)

for ($k=0; k < n1; k++$)

{

res[i][j] += a[i][k] * b[k][j];

}

(iii) $\log(\log(n)) -$

for ($i=2; i < n; i=i*i$)

count++;

Ques-4) Solve the following recurrence relation:-

At level

$$0 \rightarrow cn^2$$

$$1 \rightarrow \frac{n^2}{4^2} + \frac{n^2}{2^2} = \frac{5n^2}{16}$$

$$2 \rightarrow \frac{n^2}{8^2} + \frac{n^2}{16^2} + \frac{n^2}{4^2} + \frac{n^2}{8^2} = \left(\frac{5}{16}\right)^2 n^2 c^2$$

!

$$\text{max level} = \frac{n}{2^k} = 1$$

$$= k = \log_2 n$$

$$T(n) = c \left(n^2 + \frac{5n^2}{16} + \frac{5^2 n^2}{16^2} + \dots + \left(\frac{5}{16}\right)^{\log n} n^2 \right)$$

$$T(n) = cn^2 \left[1 + \frac{5}{16} + \frac{5^2}{16^2} + \dots + \frac{5^{\log n}}{16^{\log n}} \right]$$

$$T(n) = cn^2 \times 1 \times \left(\frac{1 - \left(\frac{5}{16}\right)^{\log n}}{1 - \left(\frac{5}{16}\right)} \right)$$

$$T(n) = cn^2 \times \frac{11}{5} \times \left(1 - \left(\frac{5}{16}\right)^{\log n} \right)$$

$$T(n) = O(n^2 c)$$

$$\boxed{O(n^2 c)} \text{ Ans}$$

Ques 5) What is time complexity of following func()?

int func(int n) {

for(int i=1; i<n; i++) {

for(int j=1; j<n; j+=i) {

// Some O(1) work

}}}

for i j $j = (n-1)/i$ times

1 1

2 1+3+5

3 1+4+7

!

n

$$\sum_{i=1}^n \frac{(n-1)}{i}$$

$$T(n) = \frac{n-1}{1} + \frac{(n-1)}{2} + \frac{(n-1)}{3} + \dots + \frac{n-1}{n}$$

$$T(n) = n \left[1 + \frac{1}{2} + \frac{1}{3} + \dots + \frac{1}{n} \right] - 1 \left[1 + \frac{1}{2} + \frac{1}{3} + \dots + \frac{1}{n} \right]$$

$$= n \log n - \log n$$

$$T(n) = O(n \log n) \text{ Ans}$$

Ques 6) What should be time complexity of

for(int i=2; i<n; i=pow(i,k)) $K \rightarrow \text{Constant}$

// Some O(1)

for i

2^1

2^k

2^{k^2}

2^{k^3}

!

2^{k^m}

where

$$2^{k^m} \leq n$$

$$k^m = \log_2 n$$

$$m = \log_k \log_2 n$$

$$\therefore \sum_{i=1}^m 1$$

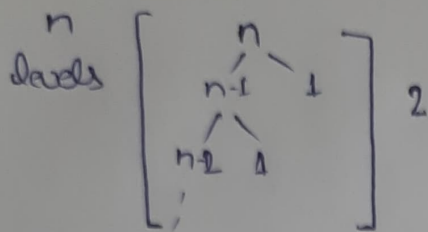
1+1+1 ... m times

$$T(n) = O(\log_k \log n) \text{ --- Ans}$$

Ques. 7)

Sol:- Given Algorithm divides array in 99% and 1% part

$$\therefore T(n) = T(n-1) + O(1)$$



(n) work is done at each level

$$T(n) = (T(n-1) + T(n-2) + \dots + T(1) + O(1)) \times n$$

$$= n \times n$$

$$\boxed{T(n) = O(n^2)}$$

Lowest height = 2

highest height = n

$$\boxed{\therefore \text{difference} = n-2} \quad n > 1$$

The given algorithm produces linear result.

Ques. 8) Arrange the following in increasing order of growth:

a) $n, n^2, \log n, \log \log n, \sqrt{n}, \log(n!), n \log n, \log^2(n), 2^n, 2^{2^n}, 4^n, n!, 100$
 $100 < \log \log n < \log n < (\log n)^2 < \sqrt{n} < n < n \log n < \log(n!) < n^2 < 2^n < 4^n < 2^{2^n}$

b) $2(2^n), 4n, 2n, 1, \log(n), \log(\log(n)), \sqrt{\log(n)}, \log 2n, 2 \log(n), n \log(n!), n!, n^2$
 $1 < \log \log n < \sqrt{\log n} < \log n < \log 2n < 2 \log n < n < n \log n < 2n^2 < 4n < \log(n!) < n^2 < n! < 2^{2^n}$

c) $8^{2^n}, \log_2 n, n \log_6(n), n \log_2(n), \log n!, n!, \log_5(n), 96, 8n^2, 7n^3, 5n$
 $96 < \log_2 n < \log 2n < 5n < n \log_6 n < n \log_2 n < \log(n!) < 8n^2 < 7n^3 < n! < 8^{2^n}$