

4

16 9

15 7

13 20

17 1

5		[4,5],
7	6	[4,6],
7	6	[6,7],
8	7	[2,3],
8	7	[1,1]]

$$\begin{array}{r} 1, 1 \\ 2, 3 \\ \underline{\underline{+}} \quad \Rightarrow \\ 4, 5 \\ 4, 6 \\ 6, 7 \end{array}$$

20 17  $d_p = \{1, 3, 5, 6, 13\}$

5

$$\begin{matrix} 15 & 5 \\ (20, 15) \end{matrix} \quad (1,1) < (2,3) < (4,5) < (6,7) \\ (4,6)$$

lisi  
q = {1, 2, 7, 4, 5, }  
2d → sort first one, solve for another  
↑  
one call 2d for other two

~~W L~~ ~~must im Doll~~

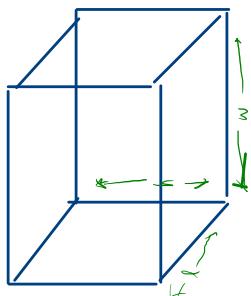
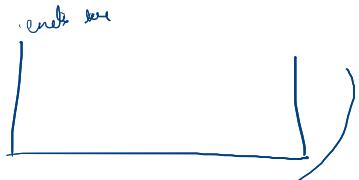
$$\begin{array}{l} \cancel{\text{h}} \\ \cancel{\text{w}} \\ \cancel{\text{l}} \\ \cancel{\text{h}} \\ \cancel{\text{w}} \end{array}$$

```
public int maxEnvelopes(int[][] envelopes) {
    Arrays.sort(envelopes, (t, o)->{
        return t[0] - o[0];
    });

    int n = envelopes.length;
    ArrayList<Integer> dp = new ArrayList<>();

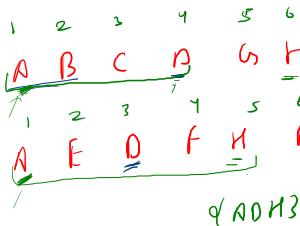
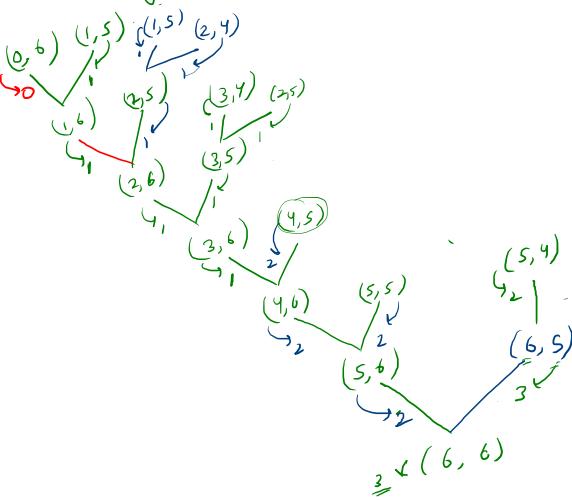
    for(int i=0; i<n; i++){
        int h = envelopes[i][1];
        int pos = findPos(dp, h);
        if(pos==dp.size()){
            dp.add(h);
        } else {
            dp.set(pos,h);
        }
    }

    return dp.size();
}
```

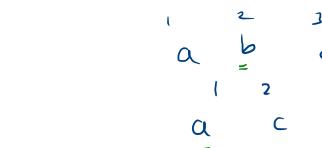
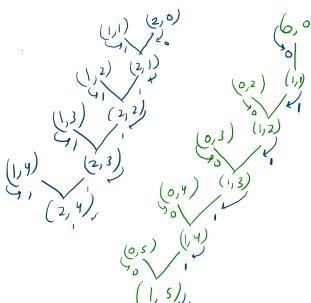




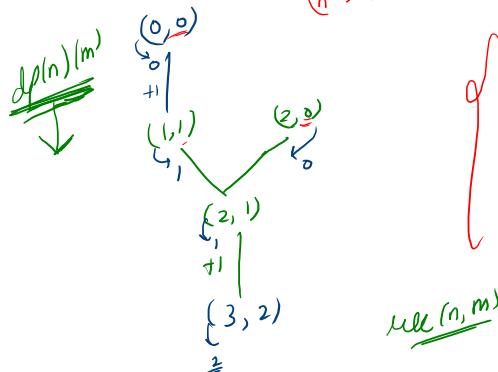
# longest common subsequence LCS



LCS



$f_{0k}(n=0, m=0)$



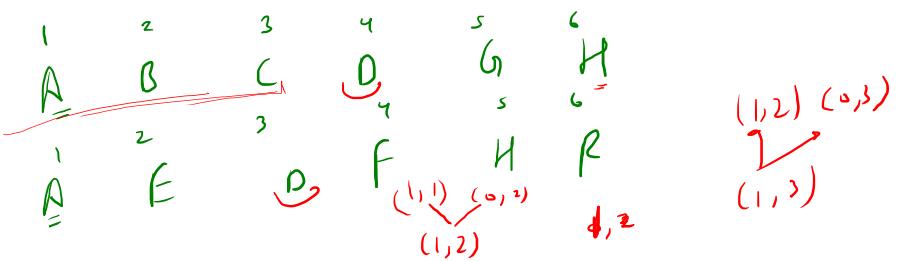
lcs(n, m)

```
public static int rec_memo(int n, int m, String s1, String s2, int[][] dp){
    if(n==0 || m==0){
        return 0;
    }

    if(dp[n][m]==-1) return dp[n][m];

    if(s1.charAt(n-1) == s2.charAt(m-1)){
        return dp[n][m] = rec_memo(n-1, m-1, s1, s2, dp) + 1;
    } else {
        return dp[n][m] = Math.max(rec_memo(n, m-1, s1, s2, dp), rec_memo(n-1, m, s1, s2, dp));
    }
}
```

$dp(n, m) \Rightarrow$  length of lcs with first  $n$  characters from  $s1$  and  $m$  from  $s2$



$s_1 \Rightarrow$	0	1	2	3	4	5	6
$s_2 \Downarrow 0$	0	0	0	0	0	0	0
1	0	1	1	1	1	1	1
2	0	1	1	1	1	1	1
3	0	1	1	1	1	1	1
4	0	1	1	2	2	2	2
5	0	1	1	2	2	2	2
6	0	1	1	2	2	3	3

```

public static int lcs_tab(int N, int M, String s1, String s2, int[][] dp){
    for(int n=0; n<=N; n++){
        for(int m=0; m<=M; m++){
            if(n==0 || m==0){
                dp[n][m]=0;
                continue;
            }

            if(s1.charAt(n-1) == s2.charAt(m-1)){
                dp[n][m] = dp[n-1][m-1] + 1; //rec_memo(n-1, m-1, s1, s2, dp) + 1;
            } else {
                dp[n][m] = Math.max(dp[n][m-1],dp[n-1][m]); //Math.max(rec_memo(n,m-1,s1,s2,dp)
            }
        }
    }
    return dp[N][M];
}

```

break till  
||: 6

$n$ -les       $\leftarrow s \quad e \quad a$        "a"

$m$ -les       $\leftarrow a \quad t \quad l$

$$\begin{aligned} n\text{-les} + m\text{-les} \\ [n+m - 2*les] \end{aligned}$$



min. insertions to make a string pallindrome

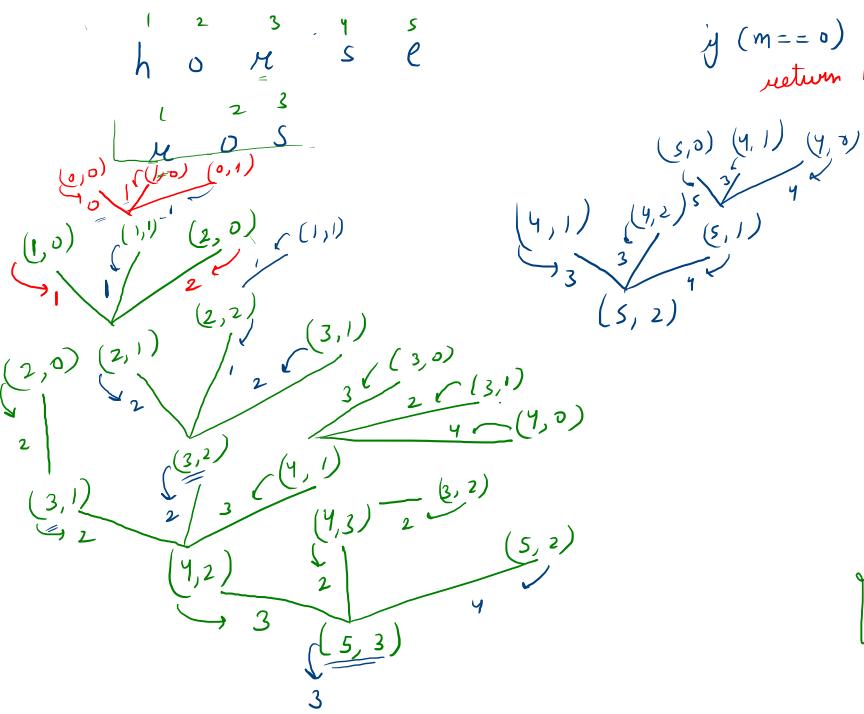
over

$\downarrow$   
n - lps  
 $\uparrow$

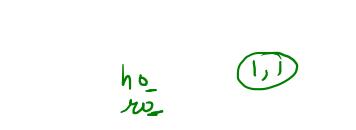
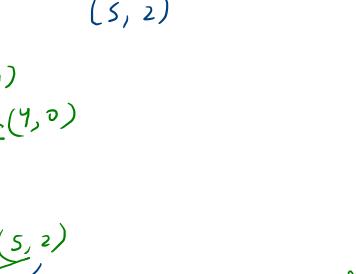
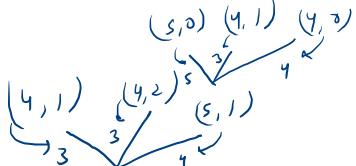
$m \underline{a} m \Leftarrow m \underline{b} a \underline{d} m \Rightarrow 2$

$b c b \Leftarrow \underline{a} \underline{b} c \underline{d} \underline{b} \Rightarrow 2$

$b c b \Leftarrow \underline{\underline{a}} \underline{\underline{b}} \underline{\underline{b}} c \underline{\underline{d}} \underline{\underline{b}} \Rightarrow 3$



$y(m==0)$  return n;  
 $y(n==0)$  return m;



$\begin{cases} \text{replace} \rightarrow (n-1, m-1); \\ \text{delete} \rightarrow (n-1, m); \\ \text{insert} \rightarrow (n, m-1) \end{cases}$

$\min(x, y, z) + 1;$

h o u s e  
 h o u se  
 h o u se

	0	1	2	3
0	0	1	2	3
1	1	1	2	3
2	2	2	1	2
3	3	2	2	2
4	4	3	3	2
5	5	4	4	3

$dp(i)(j)$

h o  
h o u

(1,1)  
(2,2)

break till  
10:50

```

public int minDistance_tab(int N, int M, String w1, String w2,int[][] dp){
    for(int n=0; n<=N; n++){
        for(int m=0; m<=M; m++){
            if(n==0 || m==0){
                dp[n][m] = n==0 ? m : n;
                continue;
            }

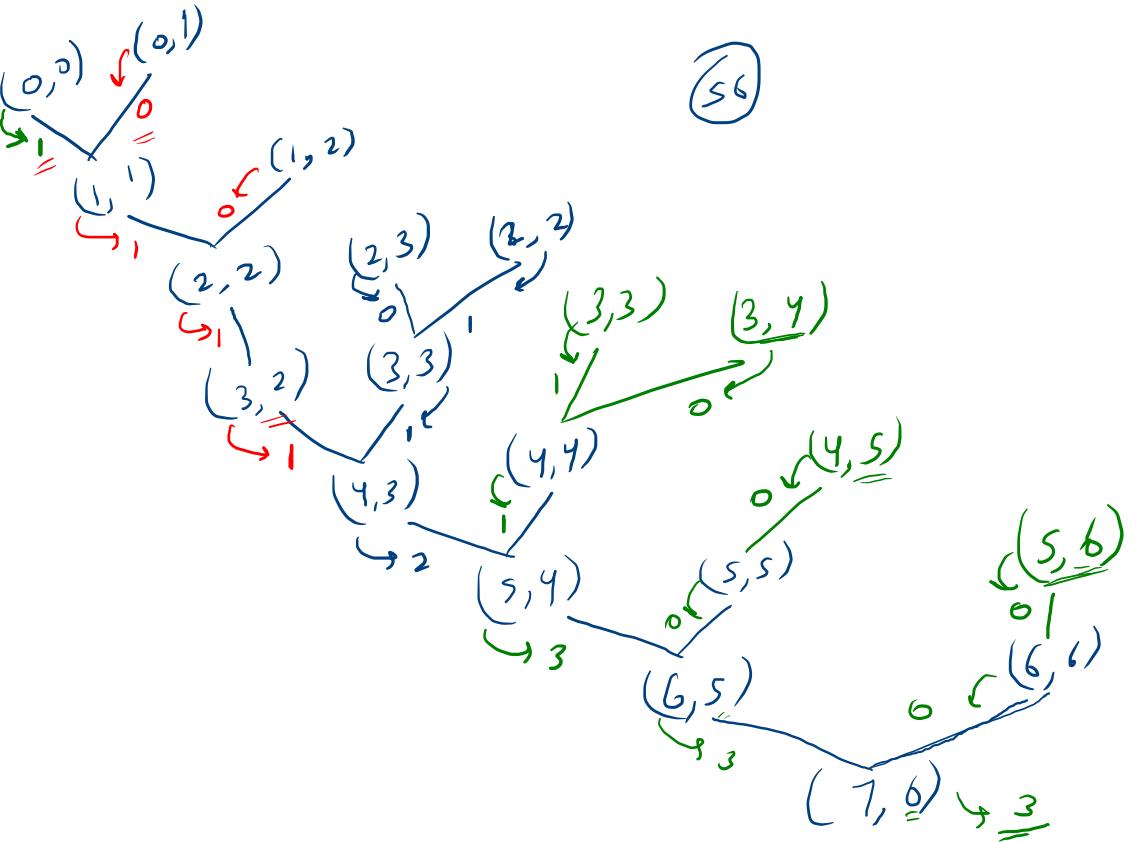
            if(w1.charAt(n-1) == w2.charAt(m-1)){
                dp[n][m] = dp[n-1][m-1]; //minDistance_memo(n-1,m-1,w1, w2,dp);
            } else {
                int replace = dp[n-1][m-1]; //minDistance_memo(n-1, m-1, w1, w2,dp);
                int delete = dp[n-1][m]; //minDistance_memo(n-1, m, w1, w2,dp);
                int insert = dp[n][m-1]; //minDistance_memo(n, m-1, w1, w2,dp);

                dp[n][m] = Math.min(replace,Math.min(delete,insert)) + 1;
            }
        }
    }
    return dp[N][M];
}
  
```

$s \Rightarrow$  "a  
 $t \Rightarrow$  "b  
 $a_1$   
 $a_2$   
 $b_1$   
 $b_2$   
 $b_3$   
 $i_1$   
 $i_2$   
 $t_1$   
 $t_2$   
 $t_3$

if ( $n < m$ )  
 return 0;

if ( $s(n) = t(m)$ )  
 $(n-1, m-1)$   
 $(n-1, m)$   
 else  $d(n-1, m)$



$$\begin{aligned} \cancel{n} \cancel{s} cb &= \cancel{e^a} + \cancel{e^b} + \cancel{e^b} \\ n \cancel{s} cb &\Rightarrow \underline{\underline{2e^b}} + \underline{\underline{e^a}} \end{aligned}$$

~~b → add after every "a"~~

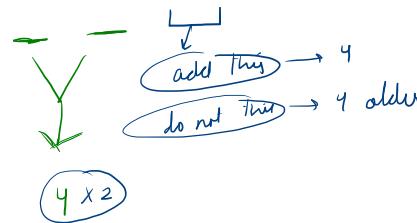
~~add after every "b", do not add~~

$\hookrightarrow$  added after every  $\Rightarrow$   
 $\hookrightarrow$  add / not add  $\Leftarrow$

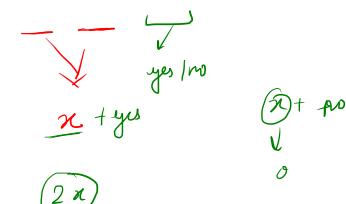
$$\begin{aligned}
 & \Rightarrow 39 \\
 & \begin{array}{l} \text{C} \\ \text{C} \\ \uparrow \end{array} \quad \begin{array}{l} ea = 5 \cancel{1} - 7 \\ eb = \cancel{5} + 8 \cancel{13} \\ ec = \cancel{5} 13 \end{array} \\
 & \cancel{ec} = 13 + 2 \times 13 \Rightarrow 39 \\
 & \boxed{ea = 2 * \cancel{ea} + 1} \\
 & \boxed{eb = ea + 2 * b} \\
 & \boxed{eb = 2 * eb + ea} \\
 & \boxed{ec = 2 * ec + eb} \\
 & \text{a'block} \\
 & \begin{array}{c} aabc \\ abbc \\ abc \\ abc \end{array} \\
 & \begin{array}{c} aabb \\ abbb \\ abbc \\ abbc \end{array} \\
 & \begin{array}{c} aabb \\ abbb \\ abbc \\ abbc \end{array}
 \end{aligned}$$

a → get added after 'a', not get added after  
create a new ss.

b → after every a,  
after every b, not



$b$	$b$
$b$	$c$
$c$	$c$
$a$	$d$



\* → any number of characters

$\rightarrow \text{O}_2, 3$

? → one character

$$\begin{array}{ccccccccc} & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 \\ \text{S} \Rightarrow & a & a & a & b & \textcircled{5} & d & a & a \\ & 1 & 2 & 3 & 4 & 5 & & & \\ \text{P} \Rightarrow & \textcircled{2} & & & d & a & a & . & \end{array}$$

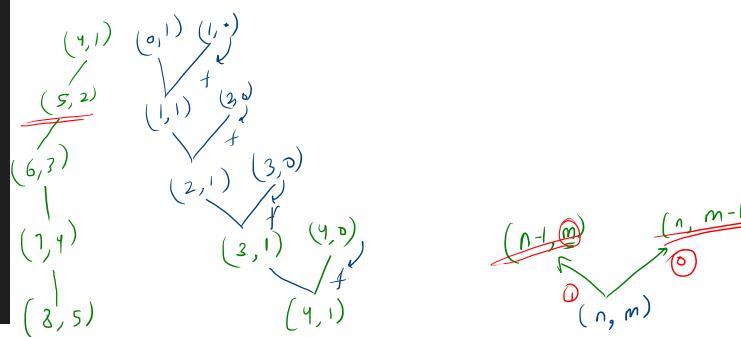
  $y(m = 0)$   
 return false  
  $(n = 0 \text{ & } m = 0)$   
 return true  
  $(n = 0 \text{ & } m = 1)$   
 $\cdot p(m) = 1$   
 return true

```

public boolean wildCard_rec(int n, int m, String s, String p){
    if(n==0 || m==0){
        if(n==0 && m==0){
            return true;
        } else if(m==1 && p.charAt(m-1)=='*'){
            return true;
        } else {
            return false;
        }
    }

    if(p.charAt(m-1)=='?'){
        return wildCard_rec(n-1, m-1, s, p);
    } else if(p.charAt(m-1)=='*'){
        return wildCard_rec(n-1, m, s, p) || wildCard_rec(n, m-1, s, p);
    } else if(s.charAt(n-1) == p.charAt(m-1)){
        return wildCard_rec(n-1, m-1, s, p);
    } else {
        return false;
    }
}

```



$s = "a b c d e f g h"$

$p = "x w x * x x * x * x x x n * p o"$

groups of step = \*\*

$\text{o} = \text{false}$   
 $\text{i} \Rightarrow \text{true}$

$$\text{dy} \cdot (i)(l) = ! = -1$$

$\stackrel{0}{\equiv} \rightarrow \text{false}$

$\downarrow \text{not true}$

$y (dp(i)(j)) = \text{false}$

$\text{dt}(3, 4)$

$(3, 4) \Rightarrow \text{false}$

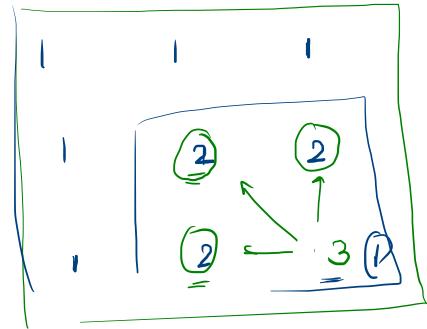
~~if~~

$\ell m = 1$

1 len square

1  
1  
 $\begin{array}{c} 1 \\ \swarrow \\ 1 \end{array}$   
 $\begin{array}{c} 2 \\ \searrow \\ 2 \end{array}$

0	1	1	0	1
1	-	2	1	1
0	-	2	2	2
1	0	1	2	3
0	0	0	1	2



2 len

2 len

2 len

$m(n+1)$