

4

16 9

15 7

13 20

17 1

5	[4,5], [4,6], [6,7], [2,3], [1,1]]
7 6	1,1 2,3 3,5 4,6 6,7
8 7	1,1 2,3 3,5 4,6 6,7

20 17	dp = {1, 3, 5, 6, 7}
-------	----------------------

6 18	⑤
------	---

15 5	(1,1) < (2,3) < (4,5) < (6,7) (4,6)
(20, 15)	

list
 $\{1, 2, 7, 4, 5, \}$
 \hookrightarrow 2d \rightarrow sort first one, solve for another
 we call 2d for other two

maximal doll
 w, l

w, l
 h, w

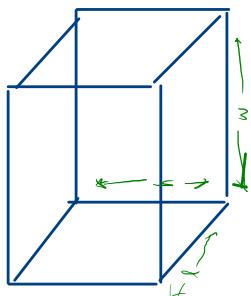
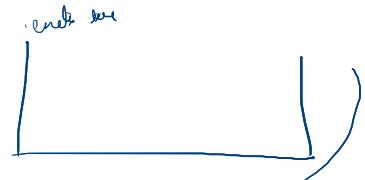
```
public int maxEnvelopes(int[][] envelopes) {
    Arrays.sort(envelopes, (t, o)->{
        return t[0] - o[0];
    });

    int n = envelopes.length;
    ArrayList<Integer> dp = new ArrayList<>();

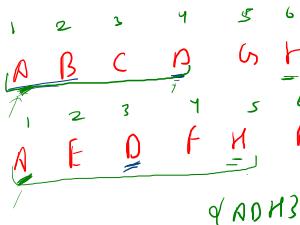
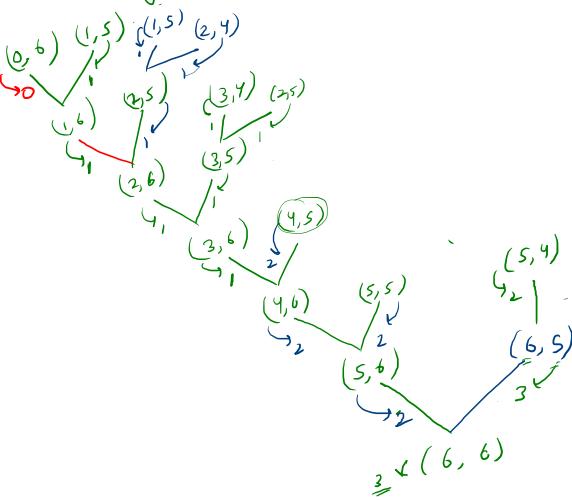
    for(int i=0; i<n; i++){
        int h = envelopes[i][1];

        int pos = findPos(dp, h);
        if(pos==dp.size()){
            dp.add(h);
        } else {
            dp.set(pos,h);
        }
    }

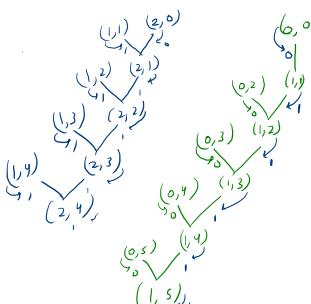
    return dp.size();
}
```



longest common subsequence LCS



LCS

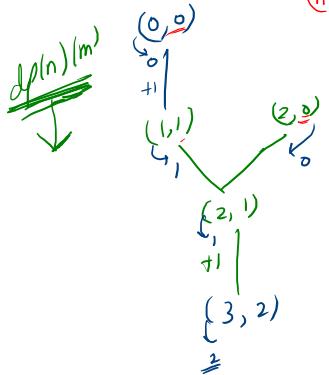


$n=0$ $m=0$

$f_{0,0}(n=0, m=0)$

$(n-1, m-1)$

$(n-1, m)$ $(n, m-1)$



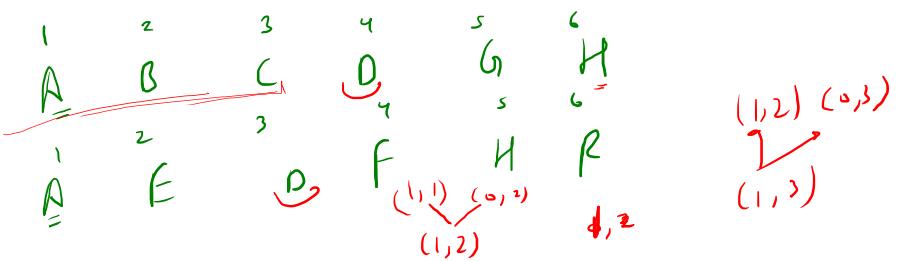
$f_{n,m}$

```
public static int rec_memo(int n, int m, String s1, String s2, int[][] dp){
    if(n==0 || m==0){
        return 0;
    }

    if(dp[n][m]==-1) return dp[n][m];

    if(s1.charAt(n-1) == s2.charAt(m-1)){
        return dp[n][m] = rec_memo(n-1, m-1, s1, s2, dp) + 1;
    } else {
        return dp[n][m] = Math.max(rec_memo(n, m-1, s1, s2, dp), rec_memo(n-1, m, s1, s2, dp));
    }
}
```

$f_{n,m} \Rightarrow$ length of lcs with first n characters from $s1$ and m from $s2$



$s_1 \Rightarrow$	0	1	2	3	4	5	6
$s_2 \Downarrow 0$	0	0	0	0	0	0	0
1	0	1	1	1	1	1	1
2	0	1	1	1	1	1	1
3	0	1	1	1	1	1	1
4	0	1	1	2	2	2	2
5	0	1	1	2	2	2	2
6	0	1	1	2	2	3	3

```

public static int lcs_tab(int N, int M, String s1, String s2, int[][] dp){
    for(int n=0; n<=N; n++){
        for(int m=0; m<=M; m++){
            if(n==0 || m==0){
                dp[n][m]=0;
                continue;
            }

            if(s1.charAt(n-1) == s2.charAt(m-1)){
                dp[n][m] = dp[n-1][m-1] + 1; //rec_memo(n-1, m-1, s1, s2, dp) + 1;
            } else {
                dp[n][m] = Math.max(dp[n][m-1],dp[n-1][m]); //Math.max(rec_memo(n,m-1,s1,s2,dp)
            }
        }
    }
    return dp[N][M];
}

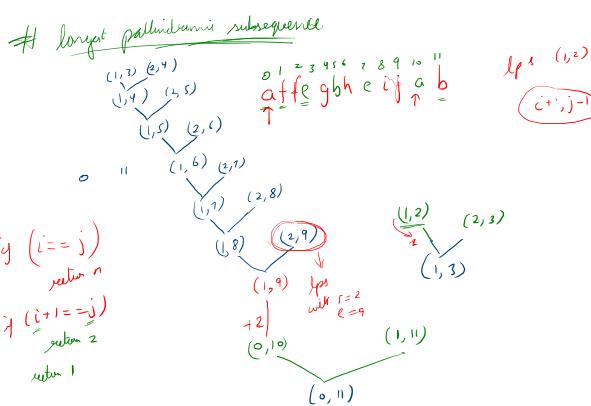
```

break till
||: 6

n-lcs $\leftarrow s \quad e \quad a$  "a"

m-lcs $\leftarrow a \quad t \quad l$

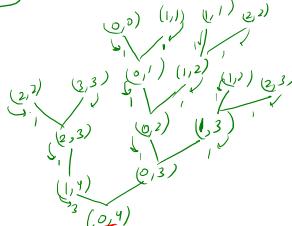
$$\begin{aligned} n - \text{lcs} + m - \text{lcs} \\ \boxed{n+m - 2 * \text{lcs}} \end{aligned}$$



$dp[i][j] = \max$

memo = new int[n][n]

$b \underline{aa} d$



a	b	c	d	e
0	1	2	3	4
1	2	3	4	5

```
public static int lps_rec(String str, int i, int j){
    if(i==j){
        return 1;
    }

    if(str.charAt(i)==str.charAt(j)){
        if(i+1 == j){
            return 2;
        } else {
            return lps_rec(str, i+1, j-1) + 2;
        }
    } else {
        return Math.max(lps_rec(str,i,j-1),lps_rec(str, i+1, j));
    }
}
```

a	b	c	d	e
0	1	2	3	4
1	2	3	4	5

a	b	bb	bb	bb	bb
0	1	2	3	4	5
1	2	3	4	5	6

```
for(int diag=0; diag<n; diag++) {
    for(i=0, j=diag; i<n; i++, j++){
        if(i==j){
            dp[i][j] = 1;
            sdp[i][j] = str.charAt(i) + "";
            continue;
        }

        if(str.charAt(i)==str.charAt(j)){
            if(i+1 == j){
                dp[i][j] = 2;
                sdp[i][j] = str.substring(i, i+2);
            } else {
                dp[i][j] = dp[i-1][j-1] + 2; //lps_memo(str, i+1, j-1, dp) + 2;
                sdp[i][j] = str.charAt(i) + sdp[i+1][j-1] + str.charAt(j);
            }
        } else {
            // dp[i][j] = Math.max(dp[i-1][j-1], dp[i+1][j]); //Math.max(lps_memo(str,i,j-1,dp),lps_
            if(dp[i][j-1] > dp[i+1][j]){
                dp[i][j] = dp[i][j-1];
                sdp[i][j] = sdp[i][j-1];
            } else {
                dp[i][j] = dp[i+1][j];
                sdp[i][j] = sdp[i+1][j];
            }
        }
    }
}
```

min. insertions to make a string pallindrome

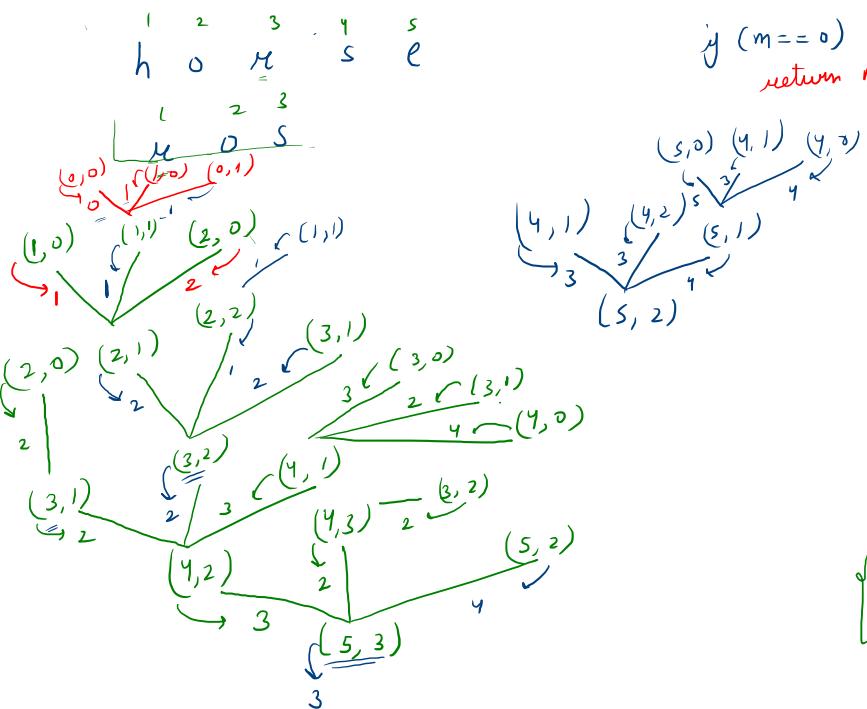
over

\downarrow
n - lps
 \uparrow

$m \underline{a} m \Leftarrow m \underline{b} a \underline{d} m \Rightarrow 2$

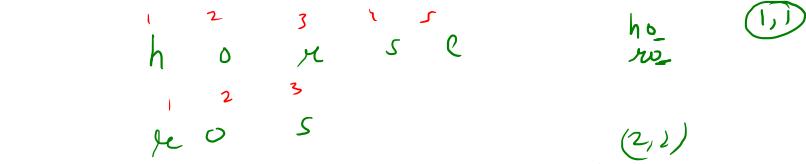
$b c b \Leftarrow \underline{a} \underline{b} c \underline{d} \underline{b} \Rightarrow 2$

$b c b \Leftarrow \underline{\underline{a}} \underline{\underline{b}} \underline{\underline{b}} c \underline{\underline{d}} \underline{\underline{b}} \Rightarrow 3$



$$\left\{ \begin{array}{l} \text{replace} \Rightarrow (n-1, m-1); \rightarrow x \\ \text{delete} \rightarrow (n-1, m) \Rightarrow y \\ \text{insert} \rightarrow \underline{(n, m-1)} \Rightarrow z \end{array} \right.$$

$$\min(x, y, z) + 1$$



break till
10:50

```
int minDistance_tab(int N, int M, String w1, String w2,int[][] dp){
int n=0; n<=N; n++){
for(int m=0; m<=M; m++){
    if(n==0 || m==0){
        dp[n][m] = n==0 ? m : n;
        continue;
    }

    if(w1.charAt(n-1) == w2.charAt(m-1)){
        dp[n][m] = dp[n-1][m-1]; //minDistance_memo(n-1,m-1,w1, w2,dp);
    } else {
        int replace = dp[n-1][m-1]; //minDistance_memo(n-1, m-1, w1, w2,dp);
        int delete = dp[n-1][m]; //minDistance_memo(n-1, m, w1, w2,dp);
        int insert = dp[n][m-1]; //minDistance_memo(n, m-1, w1, w2,dp);

        dp[n][m] = Math.min(replace,Math.min(delete,insert)) + 1;
    }
}

return dp[N][M];
}
```

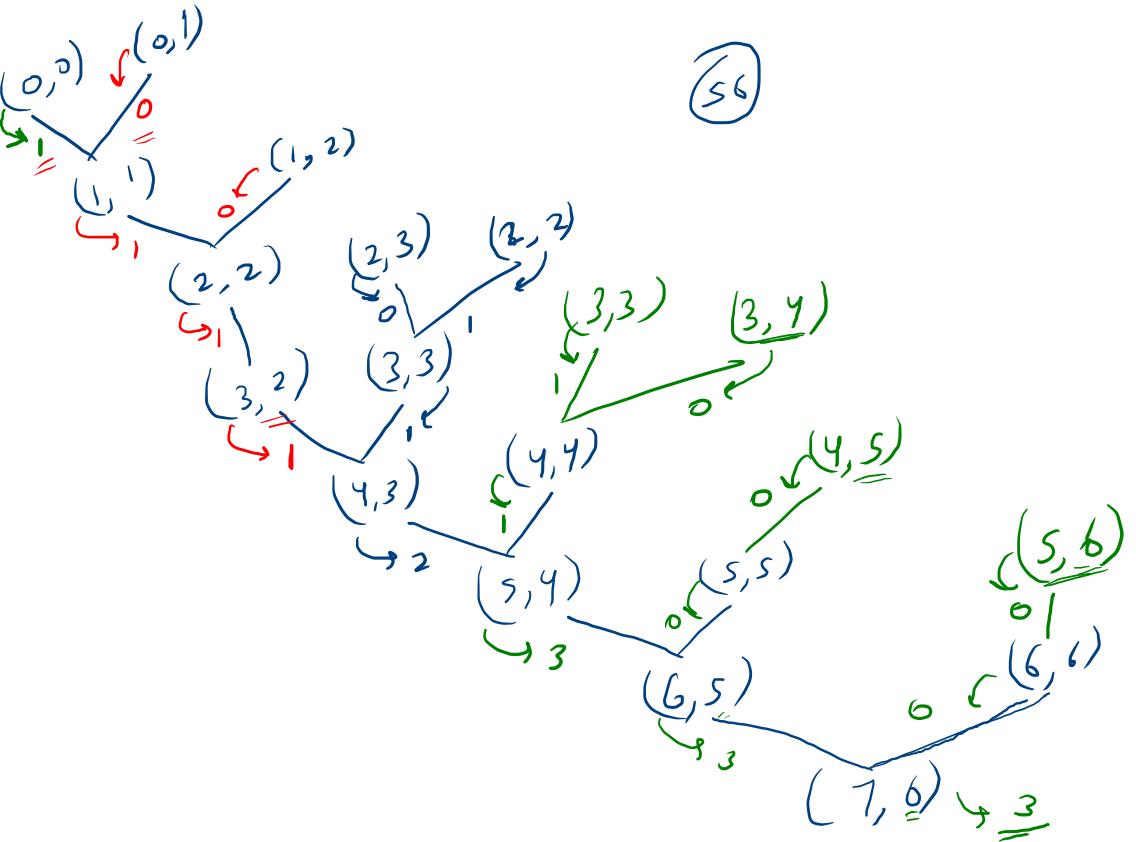
	0	1	2	3
0	0	1	2	3
1	1	1	2	3
2	2	2	1	2
3	3	2	2	2
4	4	3	3	2
5	5	4	4	3

dp(i)(j)

$s \Rightarrow$ "a
 $t \Rightarrow$ "b
 a_1
 a_2
 b_1
 b_2
 b_3
 i_1
 i_2
 t_1
 t_2
 t_3

if ($n < m$)
 return 0;

if ($s(n) = t(m)$)
 $(n-1, m-1)$
 $(n-1, m)$
 del $q_{(n-1, m)}$



(56)