

any doubts?

$$\frac{P_N C}{2}$$

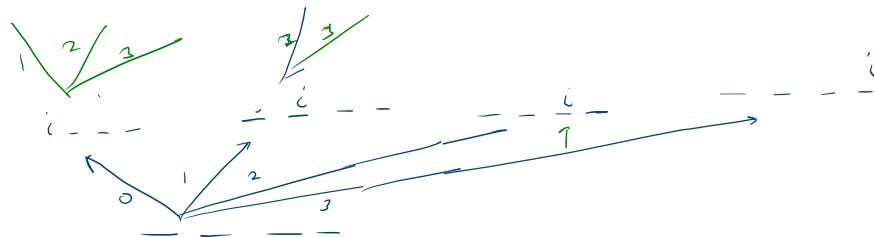
what  
how  
why

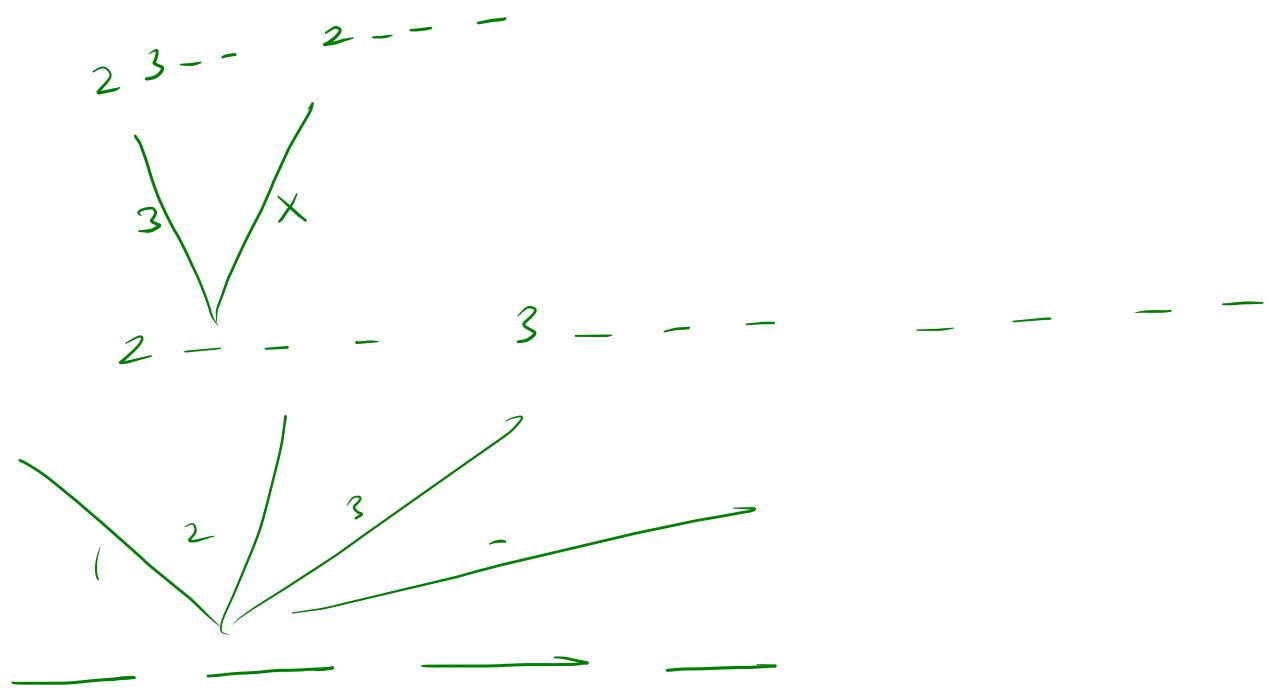
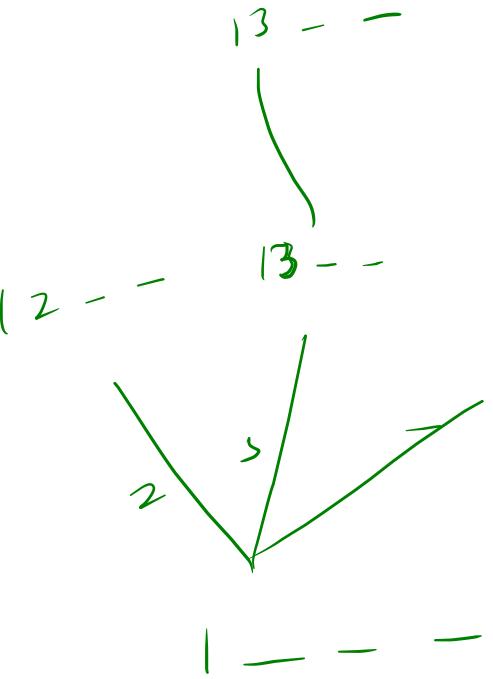
dry run  
watch the  
dry run

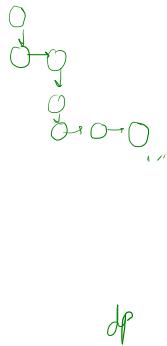
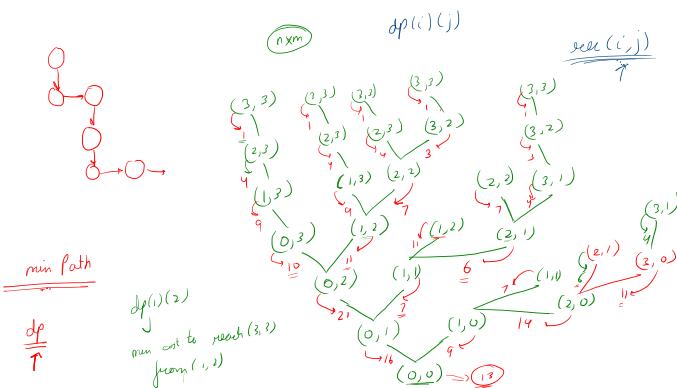
```

public static void combinations(boolean[] boxes, int ci, int ti, int lb){
    if(ci > ti){
        for(int i=0; i<boxes.length; i++){
            if(boxes[i]==true){
                System.out.print(s: "i");
            } else {
                System.out.print(s: "-");
            }
            System.out.println();
        }
        return;
    }
    for(int i = lb +1; i<boxes.length; i++){
        if(boxes[i]==false){
            boxes[i]=true;
            combinations(boxes, ci+1, ti, i);
            boxes[i]=false;
        }
    }
}

```







dp & d

0	1	2	3	
0	13	16	21	10
1	9	7	11	9
2	14	6	7	4
3	11	4	3	1

dp

```
public int minPathTab(int[][] grid, int i, int j, int[][] dp){
    for(int grid.length-1; i>0; i--){
        for(j=grid[0].length-1; j>0; j--){
            if(i==grid.length-1 && j==grid[0].length-1){
                dp[i][j] = grid[i][j];
                continue;
            }

            int right = Integer.MAX_VALUE;
            int down = Integer.MAX_VALUE;

            if(j+1 < grid[0].length){
                right = dp[i][j+1]; //minPath_memo(grid, i, j+1, dp);
            }

            if(i+1 < grid.length){
                down = dp[i+1][j]; //minPath_memo(grid, i+1, j, dp);
            }

            int ans = grid[i][j];

            if(right < down){
                ans += right;
                dp[i][j] = "r" + dp[i][j+1];
            } else {
                ans += down;
                dp[i][j] = "d" + dp[i+1][j];
            }

            dp[i][j] = ans;
        }
    }

    return dp[0][0];
}
```

```
public int minPath_tabString(int[][] grid, int i, int j, int[][] dp){
    int n = grid.length;
    int m = grid[0].length;

    String[][] sdp = new String[n][m];

    for(i=grid.length-1; i>0; i--){
        for(j=grid[0].length-1; j>0; j--){
            if(i==grid.length-1 && j==grid[0].length-1){
                sdp[i][j] = grid[i][j];
                continue;
            }

            int right = Integer.MAX_VALUE;
            int down = Integer.MAX_VALUE;

            if(j+1 < grid[0].length){
                right = dp[i][j+1]; //minPath_memo(grid, i, j+1, dp);
            }

            if(i+1 < grid.length){
                down = dp[i+1][j]; //minPath_memo(grid, i+1, j, dp);
            }

            int ans = grid[i][j];

            if(right < down){
                ans += right;
                sdp[i][j] = "r" + sdp[i][j+1];
            } else {
                ans += down;
                sdp[i][j] = "d" + sdp[i+1][j];
            }

            dp[i][j] = ans;
        }
    }

    System.out.println(sdp[0][0]);
    return dp[0][0];
}
```

0	1	2	3	
0	13	9	11	1
1	2	1	4	5
2	8	2	4	3
3	7	0	2	1

0	1	2	3	
0	ddd	dddr	rrdd	rrr
1	rddr	ddrr	rrdr	rrr
2	rdr	dr	r	rr
3	rrr	rr	r	rr

min path from (i,j) to end

$sdp(i,j)$   
 $\rightarrow$   
 $min\ path\ from(i,j)$   
 $to\ (n-1,m-1)$

0	1	2	3	
0	13	16	21	10
1	9	7	11	9
2	14	6	7	4
3	11	4	3	1

$\downarrow$   
 $min\ path\ from(i,j)$   
 $to\ end$

$\mathcal{H}=1$   
 $d=4$   
 $ans=2+4 \Rightarrow 6$

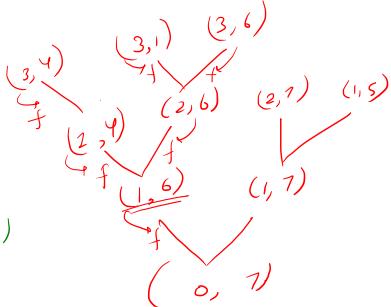
0	1	2	3	
0	ddd	dddr	rrdd	rrr
1	rddr	ddrr	rrdr	rrr
2	rdr	dr	r	rr
3	rrr	rr	r	rr

$\downarrow$   
 $min\ path\ from(i,j)$   
 $to\ end$

{1, 2, 5}

0 1 2

t<sub>target</sub> = 7



$(dp(idx) \neq \text{null}) \Rightarrow \text{subset in } dp(idx)$

(1, 6)  
f

true, false

idx = {0, n}  
target = 0, target  
memo

```
public static boolean isSubsetSum_memo(int idx, int[] arr, int target, boolean[][] dp){
    if(target == 0){
        return dp[idx][target] = true;
    }

    if(idx == arr.length){
        return dp[idx][target] = false;
    }

    if(dp[idx][target] != false) return dp[idx][target];
    boolean pick = false;

    if(target - arr[idx] >= 0)
        pick = isSubsetSum_memo(idx+1, arr, target-arr[idx], dp);

    boolean notPick = isSubsetSum_memo(idx+1, arr, target, dp);

    return dp[idx][target] = pick || notPick;
}

static Boolean isSubsetSum(int N, int arr[], int target){
    boolean[][] dp = new boolean[N+1][target+1];
    // return isSubsetSum_rec(0, arr, target);
    return isSubsetSum_memo(0, arr, target, dp);
}
```

b, int	0	1	2	3	4	5	6	7
0								
1								0
2								.
3							.	.

```
public static int isSubsetSum_memo(int idx, int[] arr, int target, int[][] dp){
    if(target == 0){
        return dp[idx][target] = 1;
    }

    if(idx == arr.length){
        return dp[idx][target] = 0;
    }

    if(dp[idx][target] != -1) return dp[idx][target];
    int pick = 0;

    if(target - arr[idx] >= 0)
        pick = isSubsetSum_memo(idx+1, arr, target-arr[idx], dp);

    int notPick = isSubsetSum_memo(idx+1, arr, target, dp);

    return dp[idx][target] = pick | notPick;
}

static Boolean isSubsetSum(int N, int arr[], int target){
    int[][] dp = new int[N+1][target+1];
    for(int[] d:dp){
        Arrays.fill(d,-1);
    }
    // return isSubsetSum_rec(0, arr, target);
    int ans = isSubsetSum_memo(0, arr, target, dp);
    return ans == 1 ? true : false;
}
```

(1, 6)  
0

(1, 6)  
false

false = 0  
true = 1

$$0, 1, 2, 3, 4, 5, 6, 7, 8, 9$$

3, 34, (4), 12, 5, 2, (6,3) (6,5)

$\text{tar} = 9$

$dp(i \text{ dn}) (\text{ target})$

$\overbrace{dp(5)(2)}^{\leftarrow} \overbrace{dp(6)(\cdot)}^{\rightarrow}$

0	1	(2)	3	4	5	6	7	8	9
✓	X	X	✓	✓	✓	✓	✓	✓	✓
✓	X	✓	X	✓	✓	✓	X	X	X
✓	X	✓	X	✓	✓	✓	X	X	X
✓	X	✓	X	X	✓	X	X	X	X
✓	X	✓	X	X	✓	X	X	X	X
✓	X	✓	X	X	✓	X	X	X	X
✓	X	✓	X	X	✓	X	X	X	X
✓	X	✓	X	X	✓	X	X	X	X
✓	X	✓	X	X	✓	X	X	X	X

$$(2)(4) = (3)(0) \mid (3)(2)$$

$(2)(6) \Rightarrow (3)(2) \mid (3)6$

```

public static boolean isSubsetSum_tab(int idx, int[] arr, int Target, boolean[][] dp){
    for(idx = arr.length; idx>=0; idx--){
        for(int target = 0; target<=Target; target++){
            if(target == 0){
                dp[idx][target] = true;
                continue;
            }

            if(idx == arr.length){
                dp[idx][target] = false;
                continue;
            }

            boolean pick = false;

            if(target - arr[idx] >=0)
                pick = dp[idx+1][target-arr[idx]]; //isSubsetSum_memo(idx+1, arr, target-arr[idx], dp);

            boolean notPick = dp[idx+1][target]; //isSubsetSum_memo(idx+1, arr, target, dp);

            dp[idx][target] = pick | notPick ;
        }
    }

    return dp[0][Target];
}

```

$dp(5)(5) = \overbrace{dp(6)(3)}^{\leftarrow} \overbrace{dp(6)(2)}^{\rightarrow}$

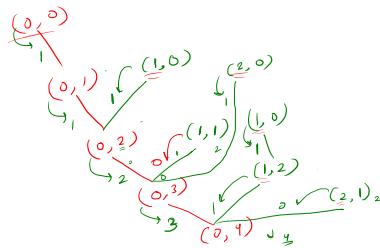
1,1,1  
1,1,2  
1,1,3  
2,2

(1, 2, 3)

✓ yes  
✗ no

$tar=0$

for (int j=idx;j<coins.length();j++)  
    ans+=combination(coins,j,tar-coins[j]);

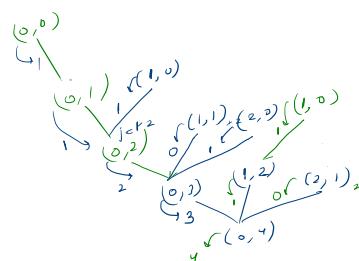


```
public long combination_rec(int[] coins, int idx, int tar){  
    if(tar==0)  
        return 1;  
  
    long ans = 0;  
    for(int j=idx;j<coins.length();j++){  
        if(tar-coins[j]>=0){  
            ans += combination_rec(coins, j, tar - coins[j]);  
        }  
    }  
  
    return ans;  
}
```

(1, 2, 3)

$tar=4$

$dp(0)[tar]$



(0, 1, 2, 3)

	0	1	2	3	4	
0	1	1	-	2	3	4
1	1	+	0	-	1	
2	1	0				

```
public long combination_memo(int[] coins, int idx, int tar, long[][] dp){  
    if(tar==0)  
        return dp[idx][tar]=1;  
  
    if(dp[idx][tar]!=-1) return dp[idx][tar];  
  
    long ans = 0;  
    for(int j=idx;j<coins.length();j++){  
        if(tar-coins[j]>=0){  
            ans += combination_memo(coins, j, tar - coins[j],dp);  
        }  
    }  
  
    return dp[idx][tar]=ans;  
}
```

No two consecutive houses have same color

k colors

seed    b    g  
0th    14    2    11

1st    11    (14)    5

2nd    (14)    (3)    10

color    1st house with red

$$\min(b, g) = \min(11, 10) = 10$$

	red	blue	green
0	14	2	11
1	13	25	7
2	21	(10)	11

```

public int solve(int cost[][], int n) {
    int[][] dp = new int[n][3];

    for(int house=0; house<n; house++){
        if(house==0){
            dp[house][0]=cost[house][0];
            dp[house][1]=cost[house][1];
            dp[house][2]=cost[house][2];
        } else {
            dp[house][0] = cost[house][0] + Math.min(dp[house-1][1],dp[house-1][2]);
            dp[house][1] = cost[house][1] + Math.min(dp[house-1][0],dp[house-1][2]);
            dp[house][2] = cost[house][2] + Math.min(dp[house-1][0],dp[house-1][1]);
        }
    }

    int ans = Integer.MAX_VALUE;
    for(int color=0; color<3; color++){
        ans = Math.min(ans,dp[n-1][color]);
    }

    return ans;
}
  
```

	0	1	2	3	4
0	0	1	-	-1	1
1	0	1	1	1	0
2	1	1	1	1	1

# LIS  
longest increasing subsequence.

0 1 2 3 4 5 6 7

dp =	<table border="1" style="width: 100%; border-collapse: collapse;"> <tr><td style="width: 25px; height: 25px;"></td><td style="width: 25px; height: 25px;"></td><td style="width: 25px; height: 25px;"></td><td style="width: 25px; height: 25px;"></td><td style="width: 25px; height: 25px; color: red;">2</td><td style="width: 25px; height: 25px; color: red;">2</td><td style="width: 25px; height: 25px;"></td><td style="width: 25px; height: 25px; color: red;">4</td><td style="width: 25px; height: 25px;"></td><td style="width: 25px; height: 25px; color: red;">4</td></tr> </table>					2	2		4		4
				2	2		4		4		
	$\{2,3\}$ $\{2,5\}$ $\{2,3\}$ $\{2,3,7\}$										

$dp[i] \Rightarrow$  length of increasing subsequence ending with  $arr[i]$

$$dp[5] = 3 \\ \Rightarrow \max(dp[1], dp[2]+1) \\ \downarrow \\ 3 \quad 1+1$$

$$dp[4] = 2 \\ dp[5] = \max(dp[4], dp[4]+1)$$

at the  $i^{\text{th}}$  point,  $(i)$  can be added

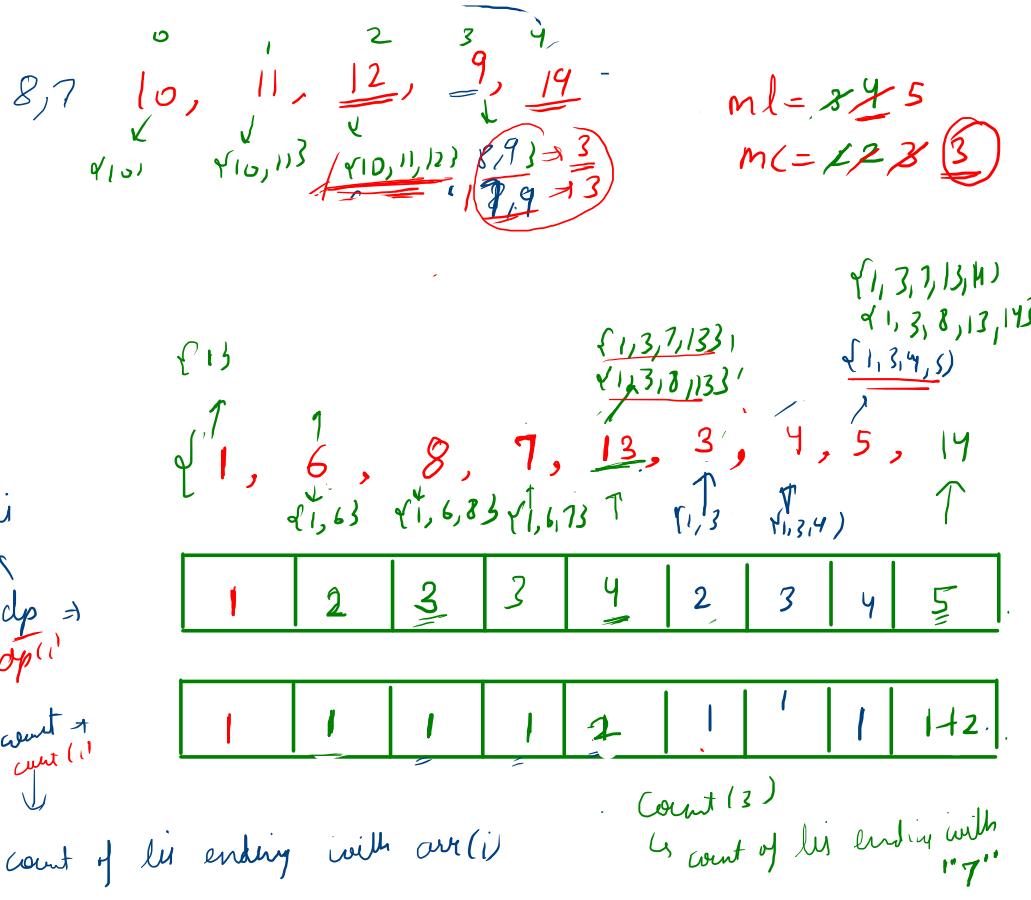
Q2

Q | 0, 9, 11, 12, 1, 2 3

0	1	2	3	4	5
			3		2

$dp[i] \rightarrow$  max sum of increasing subsequence ending with  $arr[i]$

```
public int lengthOfLIS(int[] nums) {
    int n = nums.length;
    int[] dp = new int[n];
    int ans = 0;
    for(int i=0; i<n; i++){
        dp[i]=1;
        for(int j=i-1; j>=0; j--){
            if(nums[i] > nums[j]){
                dp[i] = Math.max(dp[i], dp[j]+1);
            }
        }
        ans = Math.max(ans, dp[i]);
    }
    return ans;
}
```



```

public int findNumberOfLIS(int[] nums) {
    int n = nums.length;
    int[] dp = new int[n];
    int[] count = new int[n];
    int maxLength = 0; → max length
    int maxCount = 0; → max count

    for(int i=0; i<n; i++){
        dp[i] = 1;
        count[i] = 1;
    }

    for(int j=i-1; j>=0; j--){
        if(nums[i] > nums[j]){
            if(dp[i] < dp[j]+1){
                dp[i] = dp[j]+1;
                count[i] = count[j];
            } else if(dp[i] == dp[j]+1){
                count[i] += count[j];
            }
        }
    }

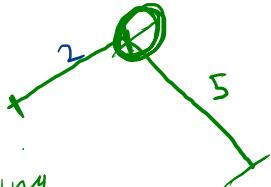
    if(dp[i] > maxLength){
        maxLength = dp[i];
        maxCount = count[i];
    } else if(dp[i] == maxLength){
        maxCount += count[i];
    }
}

return maxCount; → max count
  
```

$b=+6$

$\{1, 11\} \uparrow \quad \{11, 10\} \downarrow \quad 9, 2, 13$

increasing, decreasing



list ending  
with arr(i)  $\leftarrow lis \rightarrow$

[1, 11, 2, 10, 4, 5, 2, 1]

1	2	2	3	3	4	2	1
---	---	---	---	---	---	---	---

$\{1, 2, 4\} \uparrow \quad \{1, 2, 4, 13\} \uparrow \quad \{1, 2\}, \{13\}$

lds  
starting with  
arr(i)  $\leftarrow lds$

1	5	2	4	3	3	2	1
$\{1, 3\}$	$\{2, 11\}$	$\{10, 4, 2, 13\}$	$\{9, 2, 13\}$	$\{5, 2, 1\}$	$\{2, 1\}$	$\{1\}$	

$\{11, 10, 9, 2, 1\}$

bil.onic  $\Rightarrow$  lis ending with arr(i) + lds starting with arr(i) - 1

min

$\Rightarrow \underline{\underline{2}}, \underline{\underline{3}}, \underline{\underline{5}}, \cancel{1}, \cancel{4}, \cancel{7}, \cancel{\frac{6}{1}} \cancel{\frac{6}{1}}$

(3)

$\Rightarrow \underline{\underline{\{2, 3, 5, 6, 6\}}}$

$n - \text{lis}$

already sorted things

We can sort above array in 3 insertion steps as shown below,

1 before array value 2

4 before array value 5

6 before array value 7

$n^2$

[10, 9, 2, 5, 3, 7, 101, 18]

$n \log n$

dp: {2, 3, 7, 18}

sorted {increasing}

lis

Binary Search

Create lis

```
// LIS Better ======  
public int findPos(ArrayList<Integer> dp, int ele){  
    int left = 0;  
    int right = dp.size() - 1;  
  
    while(left <= right){  
        int mid = (left+right)/2;  
  
        if(dp.get(mid) < ele){  
            left = mid + 1;  
        } else {  
            right = mid - 1;  
        }  
    }  
  
    return left;  
}  
  
public int lengthOfLIS(int[] nums) {  
    int n = nums.length;  
  
    ArrayList<Integer> dp = new ArrayList<>();  
  
    for(int i=0; i<n; i++){  
        int pos = findPos(dp,nums[i]);  
        if(pos==dp.size()){  
            dp.add(nums[i]);  
        } else {  
            dp.set(pos,nums[i]);  
        }  
    }  
  
    return dp.size();  
}  
Run | Debug
```

$\beta S$

$\log n$

6      4      2      1  
 7      3      6      5

```

public static int maxNonOverlappingBridges(int[][] bridges){
    Arrays.sort(bridges,(int[] t, int[] o)->{
        if(t[0] == o[0]){
            return t[1] - o[1];
        }
        return t[0] - o[0];
    });

    int n = bridges.length;
    int[] dp = new int[n];
    int ans=0;

    for(int i=0; i<n; i++){
        dp[i] = 1;
        for(int j=i-1; j>=0; j--){
            if(bridges[i][1] > bridges[j][1]){
                dp[i] = Math.max(dp[i],dp[j] + 1);
            }
        }
        ans = Math.max(ans,dp[i]);
    }
    return ans;
}
  
```

$n \log n$

ed       $\{ \{6, 2\}, \{4, 3\}, \{2, 6\}, \{1, 5\} \}$

$\{ \underline{\underline{6}}, \underline{\underline{2}} \}$        $\{ \underline{\underline{1}}, \underline{\underline{5}} \}$   
 $\{ \underline{\underline{4}}, \underline{\underline{3}} \}$        $\{ \underline{\underline{2}}, \underline{\underline{6}} \}$   
 $\{ \underline{\underline{2}}, \underline{\underline{6}} \}$        $\{ \underline{\underline{4}}, \underline{\underline{3}} \}$   
 $\{ \underline{\underline{1}}, \underline{\underline{5}} \}$        $\{ \underline{\underline{6}}, \underline{\underline{2}} \}$

