


Time comp briefly

while loop
for loop

method



every jump takes 1 sec
Total = 100 sec \Rightarrow N secs

$$\frac{N}{2} \text{ secs}$$

int a=5;

a = a+1;

int b=9; b = a+1;

-
- 1) Take an array a .
 - 2) Create an array of size n .
 - 3) Increase value of every element of arr by 1.

lets suppose, every line execution takes k unit of time.

$\left\{ \begin{array}{l} \text{for (int } i=0; i < n; i++) \\ \text{arr[i]}++; \end{array} \right.$

order of n

$$\boxed{\text{Total time} = n * k}$$

$O(n)$

$\text{arr}(0)++;$
 $\text{arr}(1)++;$
 $\text{arr}(2)++;$
 $\text{arr}(3)++;$
.....
 n

```

int a=7; k           int n= k
int b=7;
for (int i=0; i<n; i++) {
    for (int j=0; j<n; j++) {
        a++;
    }
}

```

$$a = +2 -3 +4 -5 +6 -7 +8 +9 +10 +11 \\ +12 -13 +14 -15 +16 +17 -18 +19 -20 +21 \\ -22 +23 -24 +25 -26$$

int n=5; → k constant term

$$\lim_{a \rightarrow 5} k$$

for \leftarrow for $\rightarrow n^2k$

$O(n^2)$

$$n^2k + k + k + k \dots$$

$$(n^2 + i)k = n^2 k$$

for (int i=0; i<n; i=i+2) {
 System.out.print(bis[ow(i)]);
}

worst case scenario

$$\begin{array}{l}
 \text{arr(0)} \\
 \text{arr(2)} \\
 \text{arr(4)} \\
 \text{arr(6)} \\
 \vdots \\
 n \\
 \hline
 O\left(\frac{n}{2}\right) \Rightarrow O(n) \\
 O(2n) \Rightarrow O(n) \\
 O(3n) \Rightarrow O(n) \\
 O(kn) \Rightarrow O(n) \\
 \downarrow \\
 k \ll\ll n
 \end{array}$$

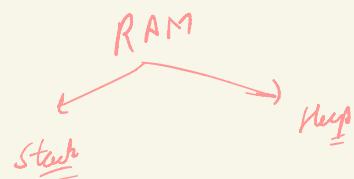
0	1	2	3	4	5	
1	2	3	4	5	6	

tar = 11

```
public static int search(int arr[], int n, int tar){
    for(int i=0; i<n; i++){
        if(arr[i]==tar){
            return i;
        }
    }
    // if my compiler at this line, i was not able to return any index
    return -1;
}
```

Space Complexity

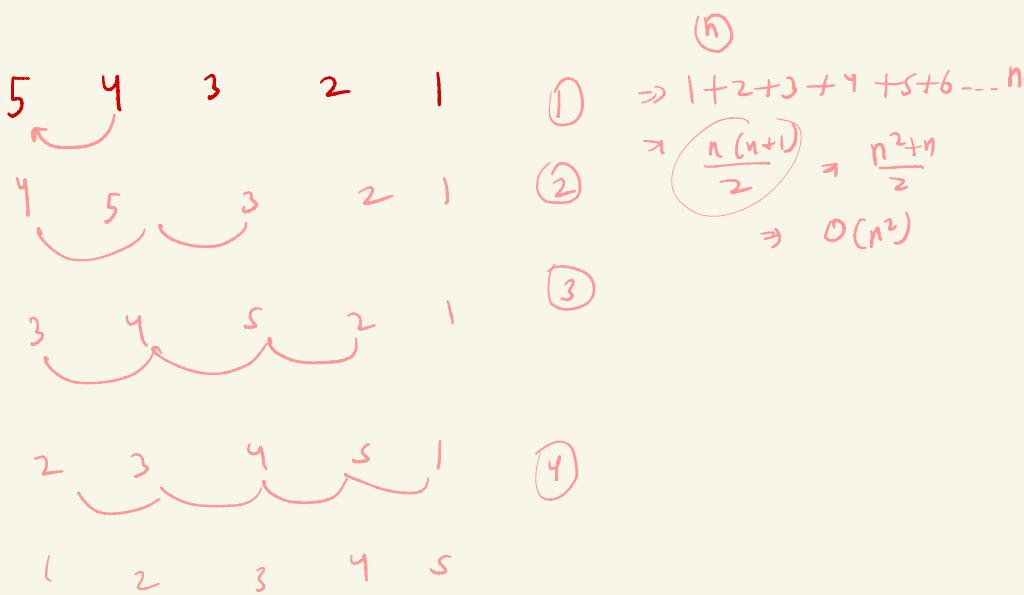
$\text{int } a = 5;$
 $\text{int } b = 7;$
 $\text{int } n = \text{scn.nextInt();}$
 $\text{int } [] arr = \text{new int}[2];$
 $\text{int } [] arr = \text{new int}(n);$
 $O(n) \rightarrow \text{space}$



1 second $\rightarrow 10^9$ operations

$(10^8 - 10^9)$

$$n \Rightarrow 10^3$$



1, 2, 3, 4, 5, 6, 7

$$a=1 \quad d=1$$

$$1+2+3+4+\dots+5+\dots+n \\ \Rightarrow \frac{n(n+1)}{2} \Rightarrow \frac{n^2+n}{2} \Rightarrow O(n^2)$$

What is the time, space complexity of following code :

```
int a = 0, b = 0;
for (i = 0; i < N; i++) {
    a = a + rand();
}
for (j = 0; j < M; j++) {
    b = b + rand();
}
```

$\rightarrow O(N)$

$O(M)$

time
 $\rightarrow O(N+M)$

Assume that `rand()` is $O(1)$ time, $O(1)$ space function.

constant space $\Rightarrow O(1)$ space

pseudo code

What is the time, space complexity of following code :

```
int a = 0, b = 0;
```

```
for (i = 0; i < N; i++) {
```

```
    for (j = 0; j < N; j++) {
```

```
        a = a + j;
```

```
}
```

```
}
```

```
for (k = 0; k < N; k++) {
```

```
    b = b + k;
```

```
}
```

constant space

$O(1)$

$O(N^2)$

$O(N^2 + N)$

$O(N^2)$

What is the time complexity of the following code :

```
int a = 0, i = N;
```

```
while (i > 0) {
```

```
    a += i;
```

```
    i /= 2;
```

```
}
```

$a = 190 \ 180 \ 170 \ 160 \ 150 \ 140 \ 130 \ 120 \ 110 \ 100 \ 90 \ 80 \ 70 \ 60 \ 50 \ 40 \ 30 \ 20 \ 10 \ 0$

number of terms = c

$N, \frac{N}{2}, \frac{N}{4}, \frac{N}{8}, \frac{N}{16}, \dots$

$\left[\frac{N}{2^0}, \frac{N}{2^1}, \frac{N}{2^2}, \frac{N}{2^3}, \frac{N}{2^4}, \dots \right] \frac{N}{2^c}$

$$\frac{N}{2^x} = 1$$

$$\Rightarrow N = 2^x$$
$$\Rightarrow \log_2 N = \log_2 2^x$$

$$\begin{aligned} &\log_a a^b \\ &\Rightarrow b \log_a a \\ &\Rightarrow b \end{aligned}$$

$$\Rightarrow \log_2^N = n \log_2^x$$

$$\Rightarrow x = \log_2^N$$

$$N = 100$$

What is the time complexity of the following code :

```
int i, j, k = 0;
for (i = n/2; i <= n; i++) {
    for (j = 2; j <= n; j = j * 2) {
        k = k + n/2;
    }
}
```

number of loops = α

$i = s_0, s_1, s_2, \dots$

$$i = s_0 - - - - - \quad 10^0$$

$$i = s_0 \log n \quad \log n$$

$$i = s_1 \log n \quad \log n$$

$$i = s_2 \log n \quad \log n$$

$$\left. \right\} \rightarrow \frac{N}{2} \log n$$

$$\Rightarrow O(n \log n)$$

$j = 2, 4, 8, 16, 32 - - - N$

$2, 2^2, 2^3, 2^4, 2^5 - - - 2^\alpha$

$$i = 100 \log n$$

$$2^\alpha = N$$

$$\log_2 2^\alpha = \log_2 N$$

$$\alpha \log_2 2 = \log_2 N$$

$$\alpha = \log_2 N$$

$$\log_a^b$$

$$\Rightarrow b \log a$$

$$\log_a^a = 1$$

$$i+j$$

$$(i = i+j)$$

$$i = i+2$$

$$(i+2)$$

$$a \leftarrow a \times 2$$

$$(a \times 2)$$

C) $\text{for}(i = 1; i < n; i *= 2) \rightarrow \log n$

$$n=100 \\ i=1+2+4+8+16+32+64$$

D) $\text{for}(i = n; i > -1; i /= 2) \rightarrow \text{infinite}$

$$i=x-48+16-\dots-n \Rightarrow \log n$$

$$2^0, 2^1, 2^2, \dots, 2^x$$

$$2^x=n \rightarrow x=\log n$$

d)

$$n=100$$

$$i=100 \rightarrow 25+12+6+3+2+1+0+0+0+0+0+0+0$$

(Worst-case scenario)

$$\textcircled{3} f_1 = 2^n$$

$$\textcircled{2} f_2 = \sqrt{n^3}$$

$$\textcircled{1} f_3 = n \log n$$

$$\textcircled{4} f_4 = n^{\log n}$$

$$f(2) \Rightarrow \sqrt{n^3}$$

$$f(3) \Rightarrow n \log n$$

↓
larger value

$$\text{int } a=1;$$

$$\Rightarrow \underline{n\sqrt{n}}$$

for ($\text{int } i=1; i \leq n; i++$) {
 for ($\text{int } j=2; j*j \leq n; j++$) {
 $a++$;
 }
}

$$\sqrt{n} \leftarrow \}$$

$$n=100$$

$$j^2 \leq n$$

$$j \leq \sqrt{n}$$

$$j=2+3+4+5+6+7+8+9+10$$

$$j=9 \rightarrow \sqrt{n}$$

$$\textcircled{9} \Rightarrow 10$$

\downarrow
 $10-1$

$$2 \times 2 < 100$$

$$3 \times 3 < 100$$

$$4 \times 4 < 100$$

$$5 \times 5 < 100$$

$$6 \times 6 < 100$$

$$7 \times 7 < 100$$

$$8 \times 8 < 100$$

$$9 \times 9 < 100$$

$$10 \times 10 > 100$$

