## Assignment 5 OS.

## A) False.

Like we saw during the lectures, in the claim 5.17:

Peterson's algorithm is 1-bounded but is not 0-bounded.

That means that if the process 0 wait in the entry of the <CS> before the process 1, it's possible that the process 1 actually enter in the <CS> before the process 0.

We can show this with the next scenario:

- 1) The process 1 sets the turn\_to\_wait variable to 1.
- 2) The process 0 sets the turn\_to\_wait variable to 0 and enters wait area.
- 3) The process 1 enter in the <CS> and so bypass process 0.

```
B)
shared: boolean flag for p = False
Algorithm for q:
//Entry code
1. while (flag ==True);
2. Peterson Enter 0(q, r)
3. // useless line
<CS>
//Exit code
1. flag = True
2. Peterson Leave 0(q, r)
Algorithm for r:
//Entry code
1. while (flag == True);
2. Peterson_Enter_1(r, q)
3. // useless line
\langle CS \rangle
//Exit code
1. flag = True
2. Peterson_Leave_1(r, q)
Algorithm for p:
//Entry code
1. while (flag == False);
<CS>
//Exit code
1. flag = False
```