

Exam 6

Question 1-

a. Mutual exclusion:

Either if process 0 or 1 are interested to enter in the CS, he must set the variable $Q[i]$ as true. Then, even if one of them enter in the CS, the second will not enter by examining the code:
 $Read_i(Q[j] = T) \rightarrow$ blocked in the entry code.

b. Deadlock free:

The algorithm doesn't provide deadlock freedom as the next scenario shows:

Beginning with the process 0:

$Write_0(Q[0] = T) \rightarrow Write_0(R[0] = 0)$

Then, process 1 gets running time:

$Write_1(Q[1] = T)$

Then, process 0 gets running time:

$Read_0(Q[1] = T \ \&\& \ R[0] = R[1] = 0) \rightarrow$ blocking by the until statement, the process must go to line 1.

Then, process 1 gets running time:

$Write_1(R[1] = (0 + 1) \% 2 = 1) \rightarrow Read_1(Q[0] == T \ \&\& \ R[1] = (R[0] + 1) \% 2 = 1) \rightarrow$ blocking by the until statement, the process must go to line 1.

repeating this scenario...

c. Deadlock free:

The algorithm doesn't provide deadlock freedom as the next scenario shows:

Beginning with the process 0:

$Write_0(Q[0] = T) \rightarrow Write_0(Choose[0] = T) \rightarrow Write_0(R[0] == 0)$

Then, process 1 gets running time:

$Write_1(Q[1] = T)$

Then, process 0 gets running time:

$Write_0(Choose[0] = F) \rightarrow Read_0(Q[1] = T \ \&\& \ Choose[1] = F \ \&\& \ R[0] = R[1] = 0) \rightarrow$ blocking by the until statement, the process must go to line 1.

Then, process 1 gets running time:

$Write_1(Choose[1] = T) \rightarrow Write_1(R[1] = (0 + 1) \% 2 = 1) \rightarrow Write_1(Choose[1] = F) \rightarrow Read_1(Q[0] = T \ \&\& \ Choose[0] = F \ \&\& \ R[1] = (R[0] + 1) \% 2 = 1) \rightarrow$ blocking by the until statement, the process must go to line 1.

repeating this scenario...

d.

The algorithm doesn't provide deadlock freedom as the next scenario shows:

Beginning with the process 0:

$Write_0(Q[0] = T) \rightarrow Write_0(R[0] == 0) \rightarrow Write_0(Choose[0] = T)$

Then, process 1 gets running time:

$Write_1(Q[1] = T)$

Then, process 0 gets running time:

$Read_0(Q[1] = T \ \&\& \ Choose[1] = F \ \&\& \ R[0] = R[1] = 0) \rightarrow$ blocking by the until statement, the process must go to line 1.

From now only process 0 gets running time.

repeating this scenario...

Question 2-

b.

Shared :

```
semaphore deleter_thread_bc_of_seach = 1
```

```
semaphore deleter_thread_bc_of_insert = 1
```

```
int count_searcher = 0
```

code for Searcher thread:

```
<entry code>
```

```
count_searcher++
```

```
if (count_searcher == 1)
```

```
    down(deleter_thread_bc_of_seach)
```

```
<CS>
```

```
<exit code>
```

```
count_searcher--
```

```
if(count_searcher == 0)
```

```
    up(deleter_thread_bc_of_seach)
```

code for Inserted thread:

```
<entry code>
```

```
down(deleter_thread_bc_of_insert)
```

```
<CS>
```

```
<exit code>
```

```
up(inserted_thread_bc_of_insert )
```

code for Deleter thread:

```
<entry code>
```

```
down(deleter_thread_bc_of_insert)
```

```
down(deleter_thread_bc_of_search)
```

```
<CS>
```

```
<exit code>
```

```
up(inserted_thread_bc_of_search)
```

```
up(inserted_thread_bc_of_insert )
```