

Lecture 11

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Arthur Molnar

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Quick overview

Recursion. Computational complexity

Lect. PhD. Arthur Molnar

Babes-Bolyai University

Overview

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Quick overview

Circular definition

In order to understand recursion, one must first understand recursion.

What is recursion?

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Quick overview

- A recursive definition is used to define an object in terms of itself.
- A recursive definition of a function defines values of the functions for some inputs in terms of the values of the same function for other inputs.
- Recursion can be:
 - **Direct** - a function **p** calls itself
 - **Indirect** - a function **p** calls another function, but it will be called again in turn

Demo

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Recursion

Examine the source code in **ex25_recursion.py**

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Quick overview

■ Main idea

- Base case: simplest possible solution
- Inductive step: break the problem into a simpler version of the same problem plus some other steps

■ How recursion works

- On each method invocation a new symbol table is created. The symbol table contains all the parameters and the local variables defined in the function
- The symbol tables are stored in a stack, when a function is returning the current symbol table is removed from the stack

■ Recursion and stack memory

- Stack memory size is allocated by the compiler/runtime environment
- Compilers can optimize recursive computation (e.g. see Ackermann's function)

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Quick overview

■ Advantages

- + Clarity
- + Simplified code

■ Disadvantages

- Large recursion depth might run out of stack memory
- Large memory consumption in the case of branched recursive calls (for each recursion a new symbol table is created - see Ackermann's function)

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Quick overview

What is it?

Studying algorithm efficiency mathematically

- **We study algorithms with respect to**
 - Run time required to solve the problem
 - Extra memory required for temporary data
- **What affects runtime for a given algorithm**
 - Size and structure of the input data
 - Hardware
 - Changes from a run to another due to hardware and software environment

Running time example

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Quick overview

As a first example, let's take a well-understood function:
computing the n^{th} term of the Fibonacci sequence

- What is so special about it?
 - Easy to write in most programming languages
 - Iterative and recursive implementation comes naturally
 - Different run-time complexity!

Demo

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Computational complexity

Examine the source code in **ex26_complexity.py**¹

¹To run the example, install the texttable component from
<https://github.com/foutaise/texttable>

Overcalculation in recursive Fibonacci

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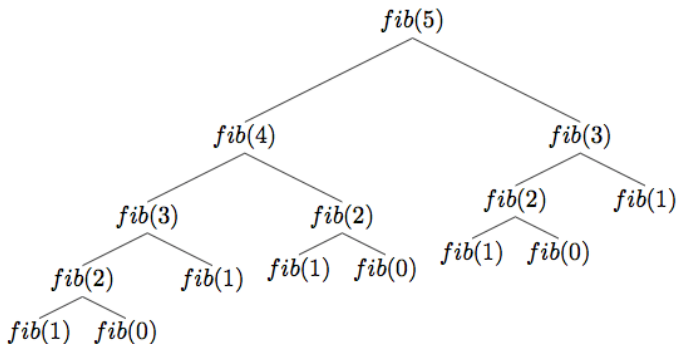
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Discussion

How can overcalculation be eliminated?

Memoization

Examine the source code in **ex27_complexityOptimized.py**²

²To run the example, install the texttable component from
<https://github.com/foutaise/texttable>

Efficiency of a function

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What is function efficiency?

The amount of resources they use, usually measured in either the space or time used.

■ Measuring efficiency

- **Empirical analysis** - determines exact running times for a sample of specific inputs, but cannot predict algorithm performance on all inputs.
- **Asymptotic analysis** - mathematical analysis that captures efficiency aspects for all possible inputs but cannot provide execution times.

■ Function run time is studied in direct relation to data input size

- We focus on asymptotic analysis, and illustrate it using empirical data.

Complexity

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Quick overview

- **Best case (BC)**, for the data set leading to minimum running time $BC(A) = \min_{I \in D} E(I)$
- **Worst case (WC)**, for the data set leading to maximum running time $WC(A) = \max_{I \in D} E(I)$
- **Average case (AC)**, average running time of the algorithm $AC(A) = \sum_{I \in D} P(I)E(I)$

Legend

A - algorithm; **D** - domain of algorithm; **E(I)** - number of operations performed for input **I**; **P(I)** the probability of having **I** as input data

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Observation

Due to the presence of the **P(I)** parameter, calculating average complexity might be challenging

Run time complexity

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The essence

- How the running time of an algorithm increases with the size of the input at the limit: **if** $n \rightarrow \infty$, **then** $3n^2 \approx n^2$
- We compare algorithms using the magnitude order of the number of operations they make

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Quick overview

- Running time is not a fixed number, but rather a function of the input data size n , denoted $T(n)$.
- Measure basic steps that the algorithm makes (e.g. number of statements executed).
- + It gets us within a small constant factor of the true runtime most of the time.
- + Allows us to predict run time for different input data
- Does not exactly predict true runtime

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Quick overview

■ Example:

$$T(n) = 13 * n^3 + 42 * n^2 + 2 * n * \log_2 n + 3 * \sqrt{n}$$

- Because $0 < \log_2 n < n, \forall n > 1$, and $\sqrt{n} < n, \forall n > 1$, we conclude that the n^3 term dominates for large n .
- Therefore, we say that the running time $T(n)$ grows "*roughly on the order of n^3* ", and we write it as $T(n) \in O(n^3)$.
- Informally, the statement above means that "*when you ignore constant multiplicative factors, and consider the leading term, you get n^3* ".

"Big-O" notation

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Quick overview

We denote function $f : \mathbb{N} \rightarrow \mathbb{R}$, and by T the function that gives the number of operations performed by an algorithm, $T : \mathbb{N} \rightarrow \mathbb{N}$.

Definition, "Big-oh" notation

We say that $T(n) \in O(f(n))$ if there exist c and n_0 positive constants independent of n such that $0 \leq T(n) \leq c * f(n), \forall n \geq n_0$.

"Big-O" notation

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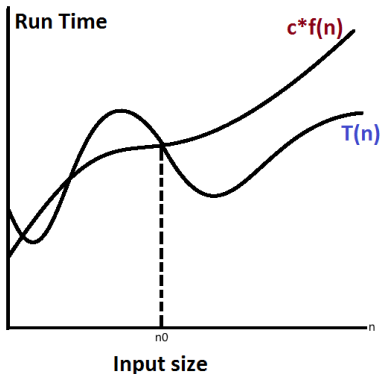
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- In other words, $O(n)$ notation provides the asymptotic upper bound.

"Big-O" notation

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Alternative definition, "Big-oh" notation

We say that $T(n) \in O(f(n))$ if $\lim_{n \rightarrow \infty} \frac{T(n)}{f(n)}$ is 0 or a constant, but not ∞ .

- If $T(n) = 13 * n^3 + 42 * n^2 + 2 * n * \log_2 n + 3 * \sqrt{n}$, and $f(n) = n^3$, then $\lim_{n \rightarrow \infty} \frac{T(n)}{f(n)} = 13$. So, we say that $T(n) \in O(n^3)$.
- The O notation is good for putting an upper bound on a function. We notice that if $T(n) \in O(n^3)$, it is also $O(n^4)$, $O(n^5)$, since the limit will then go to 0.
- To be more precise, we also introduce a lower bound on complexity.

"Big-omega" notation

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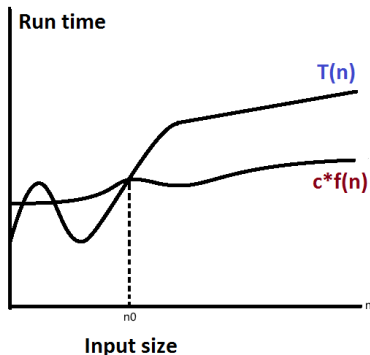
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Definition, "Big-omega" notation

We say that $T(n) \in \Omega(f(n))$ if there exist c and n_0 positive constants independent of n such that $0 \leq c * f(n) \leq T(n), \forall n \geq n_0$.



"Big-omega" notation

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Alternative definition, "Big-omega" notation

We say that $T(n) \in \Omega(f(n))$ if $\lim_{n \rightarrow \infty} \frac{T(n)}{f(n)}$ is a constant or ∞ , but not 0.

- If $T(n) = 13 * n^3 + 42 * n^2 + 2 * n * \log_2 n + 3 * \sqrt{n}$ and $f(n) = n^3$, then $\lim_{n \rightarrow \infty} \frac{T(n)}{f(n)} = 13$. So, we say that $T(n) \in \Omega(n^3)$.
- The Ω notation is used for establishing a lower bound on an algorithm's complexity.

"Big-theta" notation

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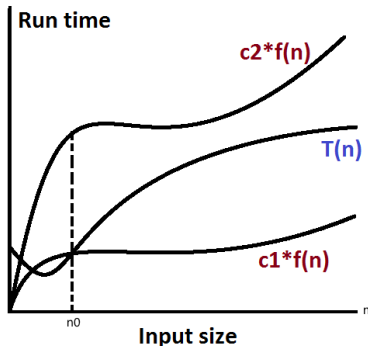
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Quick overview

Definition, "Big-theta" notation

We say that $T(n) \in \Theta(f(n))$ if $T(n) \in O(f(n))$ and $T(n) \in \Omega(f(n))$, i.e. there exist c_1, c_2 and n_0 positive constants, independent of n such that $c_1 * f(n) \leq T(n) \leq c_2 * f(n), \forall n \geq n_0$.



"Big-theta" notation

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Alternative definition, "Big-theta" notation

We say that $T(n) \in \Theta(f(n))$ if $\lim_{n \rightarrow \infty} \frac{T(n)}{f(n)}$ is a constant (but not 0 or ∞).

- If $T(n) = 13 * n^3 + 42 * n^2 + 2 * n * \log_2 n + 3 * \sqrt{n}$ and $f(n) = n^3$, then $\lim_{n \rightarrow \infty} \frac{T(n)}{f(n)} = 13$. So, we say that $T(n) \in \Theta(n^3)$. This can also be deduced from $T(n) \in O(n^3)$ and $T(n) \in \Omega(n^3)$
- The run time of an algorithm is $\Theta(f(n))$ if and only if its worst case run time is $O(f(n))$ and best case run time is $\Omega(f(n))$.

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Quick overview

```
for i in data_list:  
    # do something here...
```

- Assuming that the loop body takes $f(i)$ time to run, the total running time is given by the summation

$$T(n) = \sum_{i=1}^n f(i)$$

Observation

Nested loops naturally lead to nested sums.

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Quick overview

Solving summations breaks down into two basic steps

- Simplify the summation as much as possible - remove constant terms and separate individual terms into separate summations.
- Solve each of the remaining simplified sums.

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Quick overview

```
def f(n):  
    s = 0  
    for i in range(1, n+1):  
        s=s+1  
    return s
```

- $T(n) = \sum_{i=1}^n 1 = n \Rightarrow T(n) \in \Theta(n)$
- BC/AC/WC complexity is the same

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Quick overview

```
def f(n):  
    i = 0  
    while i <= n:  
        # do something here ...  
        i += 1
```

- $T(n) = \sum_{i=1}^n 1 = n \Rightarrow T(n) \in \Theta(n)$
- BC/AC/WC complexity is the same

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Quick overview

```
def f(l):  
    '''  
    Return True if list contains an even number  
    '''  
    poz = 0  
    while poz < len(l) and l[poz]%2 !=0:  
        poz += 1  
    return poz<len(l)
```

- BC - first element is even number, $T(n) = 1, T(n) \in \Theta(1)$
- WC - no even number in list, $T(n) = n, T(n) \in \Theta(n)$

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```
def f(l):  
    '''  
    Return True if list contains an even number  
    '''  
    poz = 0  
    while poz < len(l) and l[poz]%2 !=0:  
        poz += 1  
    return poz<len(l)
```

- AC - the **while** can be executed $1, 2, \dots, n$ times, with same probability (lacking additional information). The number of steps is then the average number of iterations:

$$T(n) = \frac{1+2+\dots+n}{n} = \frac{n+1}{2} \Rightarrow T(n) \in \Theta(n)$$

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```
def f(n):  
    for i in range(1, 2*n-1):  
        for j in range(i+2, 2*n+1):  
            # do something...
```

$$\blacksquare T(n) = \sum_{i=1}^{2n-2} \sum_{j=i+2}^{2n} 1 = \sum_{i=1}^{2n-2} (2n - i - 1)$$

$$\blacksquare T(n) = \sum_{i=1}^{2n-2} 2n - \sum_{i=1}^{2n-2} i - \sum_{i=1}^{2n-2} 1$$

$$\blacksquare T(n) = 2n * \sum_{i=1}^{2n-2} 1 - \frac{(2n-2)(2n-1)}{2} - (2n-2)$$

$$\blacksquare T(n) = 2 * n^2 - 3 * n + 1 \in \Theta(n^2).$$

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```
def f():  
    for i in range(1, 2*n-1):  
        j = i+1  
        cond = True  
        while j < 2*n and cond:  
            # do something ...  
            if someCondition:  
                cond = False
```

- Best Case - while executed once,

$$T(n) = \sum_{i=1}^{2n-2} 1 = 2n - 2 \in \Theta(n)$$

- Worst Case - while executed $2n - i - 1$ times,

$$T(n) = \sum_{i=1}^{2n-2} (2n - i - 1) = \dots = 2n^2 - 3n + 1 \in \Theta(n^2)$$

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```
def f():  
    for i in range(1, 2*n-1):  
        j = i+1  
        cond = True  
        while j < 2*n and cond:  
            # do something ...  
            if someCondition:  
                cond = False
```

- Average Case - for a given i the "while" loop can be executed $1, 2, \dots, 2n - i - 1$ times, average steps:

$$c_i = \frac{1+2+\dots+2n-i-1}{2n-i-1} = \dots = 2n - i$$

- $$T(n) = \sum_{i=1}^{2n-2} c_i = \sum_{i=1}^{2n-2} 2n - i = \dots \in \Theta(n^2)$$

- Overall complexity is therefore $\Theta(n^2)$

Summation - important sums

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Quick overview

- Constant series $\sum_{i=1}^n 1 = n$
- Arithmetic series $\sum_{i=1}^n i = \frac{n(n+1)}{2}$
- Quadratic series $\sum_{i=1}^n i^2 = \frac{n(n+1)(2n+1)}{2}$
- Harmonic series $\sum_{i=1}^n \frac{1}{i} = \ln(n) + O(1)$
- Geometric series $\sum_{i=1}^n c^i = \frac{c^{n+1} - 1}{c - 1}, c \neq 1$

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Quick overview

- **Constant time:** $T(n) \in O(1)$. It means that run time does not depend on size of the input. It is very good complexity.
- $T(n) \in O(\log_2 \log_2 n)$. This is also a very fast time, it is practically as fast as constant time.
- **Logarithmic time:** $T(n) \in O(\log_2 n)$. It is the run time of binary search and height of balanced binary trees. About the best that can be achieved for data structures using binary trees. Note that $\log_2 1000 \approx 10$, $\log_2 1000^2 \approx 20$.

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Quick overview

- **Polylogarithmic time:** $T(n) \in O((\log_2 n)^k)$.
- **Linear time:** $T(n) \in O(n)$. It means that run time scales linearly with the size of input data.
- $T(n) \in O(n * \log_2 n)$. This is encountered for fast sort algorithms, such as merge-sort and quick-sort.

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Quick overview

- **Quadratic time:** $T(n) \in O(n^2)$. Empirically, ok with n in the hundreds but not with n in the millions.
- **Polynomial time:** $T(n) \in O(n^k)$. Empirically practical when k is not too large.
- **Exponential time:** $T(n) \in O(2^n), O(n!)$. Empirically usable only for small values of input.

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Quick overview

What is a recurrence?

A recurrence is a mathematical formula defined recursively.

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Quick overview

- A **recurrence** is a mathematical formula that is defined recursively.
- For example, let us consider the problem of determining the number $N(h)$ of nodes of a complete 3-ary tree of height h . We can observe that $N(h)$ can be described using the following recurrence:

$$\begin{cases} N(0) = 1 \\ N(h) = 3 * N(h - 1) + 1, h \geq 1 \end{cases}$$

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Quick overview

The explanation is given below:

- The number of nodes of a complete 3-ary tree of height 0 is 1.
- A complete 3-ary tree of height h , $h > 0$ consists of a root node and 3 copies of a 3-ary tree of height $h - 1$. If we solve the above recurrence, we obtain that:

$$N(h) = 3^h * N(0) + (1 + 3^1 + 3^2 + \dots + 3^{h-1}) = \sum_{i=0}^h 3^i.$$

Example II - Recursive list summation

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Quick overview

```
def recursiveSum(l):  
    '''  
    Compute the sum of numbers in a list  
    l – input list  
    return int, the sum of the numbers  
    '''  
    # base case  
    if l == []:  
        return 0  
    # recursion step  
    return l[0] + recursiveSum(l[1:])
```

- n represents list length
- In this case, the recurrence is:

$$T(n) = \begin{cases} 1, & n = 0 \\ T(n-1) + 1, & n > 0 \end{cases}$$

Example II - Recursive list summation

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Quick overview

- Solving the recurrence:

$$T(n) = \begin{cases} 1, & n = 0 \\ T(n-1) + 1, & n > 0 \end{cases}$$

- $T(n) = T(n-1) + 1$
- $T(n-1) = T(n-2) + 1$
- $T(n-2) = T(n-3) + 1 \Rightarrow T(n) = n + 1 \in \Theta(n)$

Example III - Tower of Hanoi

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Quick overview

Legend says there is an Indian temple containing a large room with three posts and surrounded by 64 golden discs. Brahmin priests, acting out an ancient prophecy, are moving these discs since time immemorial, according to the rules of the Brahma. According to the legend, when the last move is completed, **the world will end.**

Example III - Tower of Hanoi

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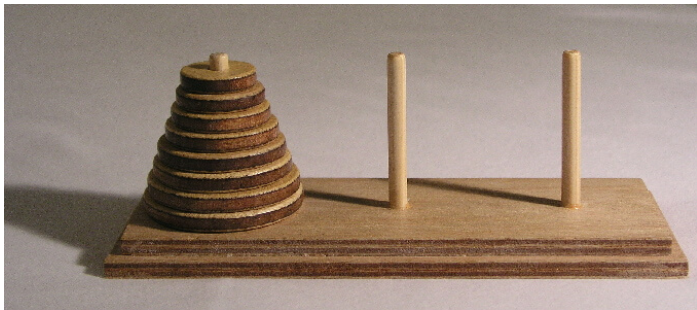
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Quick overview

A mathematical game. Starts with three rods and a number of discs of increasing radius placed on one of them. The objective of the game is to move all the discs to another rod, observing the following rules:

- You can only move one disk at a time
- You can only move the uppermost disc from a rod
- You cannot place a larger disc on a smaller one

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Quick overview

So ... are we safe (for now)? Let's study this:

- **Mathematically**
- **Empirically**

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Quick overview

The idea of the algorithm (for n discs):

- Move $n-1$ discs from source to intermediate stick
- Move the last disc to the destination stick
- Solve problem for $n-1$ discs

Example III - Tower of Hanoi

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Quick overview

```
def hanoi(n, x, y, z):  
    '''  
    n — number of disks on the x stick  
    x — source stick  
    y — destination stick  
    z — intermediate stick  
    '''  
    if n==1:  
        print("disk 1 from ",x, " to ",y)  
        return  
    hanoi(n-1, x, z, y)  
    print("disk ",n, " from ",x," to ",y)  
    hanoi(n-1, z, y, x)
```

- The recurrence is:

$$T(n) = \begin{cases} 1, n = 1 \\ 2T(n-1) + 1, n > 1 \end{cases}$$

Example III - Tower of Hanoi

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Quick overview

- Solving the recurrence:

$$T(n) = \begin{cases} 1, n = 1 \\ 2T(n-1) + 1, n > 1 \end{cases}$$

- $T(n) = 2T(n-1) + 1$, $T(n-1) = 2T(n-2) + 1$,
 $T(n-2) = 2T(n-3) + 1, \dots$, $T(1) = T(0) + 1$
- $T(n) = 2T(n-1) + 1$, $2T(n-1) = 2^2T(n-2) + 2$,
 $2^2T(n-2) = 2^3T(n-3) + 2^2, \dots, 2^{n-2}T(2) =$
 $2^{n-1}T(1) + 2^{n-2}$
- We have $T(n) = 2^{n-1} + 2^0 + 2^1 + 2^2 + \dots + 2^{n-2}$
- Therefore $T(n) = 2^n - 1 \in \Theta(2^n)$

Example III - Tower of Hanoi

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Quick overview

So ... are we safe for now? Let's study this:

- **Mathematically**
- **Empirically**

Demo

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Examine the source code in **ex28_hanoi.py**

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Quick overview

What is the space complexity of an algorithm?

The space complexity estimates the quantity of memory required by the algorithm to store the input data, the final results and the intermediate results. As the time complexity, the space complexity is also estimated using "O" and "Omega" notation.

- All the remarks from related to the asymptotic notations used in running time complexity analysis are valid for the space complexity, also.

Example I - List summation

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Quick overview

```
def iterative_sum(l):  
    '''  
    Compute the sum of numbers in a list  
    l - input list  
    return int, the sum of the numbers  
    '''  
    res = 0  
    for nr in l:  
        rez += nr  
    return rez
```

- We need memory to store the numbers, so
 $T(n) = n \in \Theta(n)$.

Example I - List summation

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Quick overview

```
def recursive_sum(l):  
    '''  
    Compute the sum of numbers in a list  
    l - input list  
    return int, the sum of the numbers  
    '''  
  
    # base case  
    if l == []:  
        return 0  
  
    # recursion step  
    return l[0] + recursive_sum(l[1:])
```

- The recurrence is:

$$T(n) = \begin{cases} 0, n = 1 \\ T(n-1) + n - 1, n > 1 \end{cases}$$

Complexity overview

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Quick overview

1 If there is Best/Worst case

- Describe Best case
- Compute complexity for Best Case
- Describe Worst Case
- Compute complexity for Worst case
- Compute average complexity (if possible)
- Compute overall complexity (if possible)

2 If Best = Worst = Average

- Compute complexity