

Mid Sem 2024 (old) Netsec



what is end to end encryption. explain how can it be used to secure app based chat

What is End-to-End Encryption (E2EE)?

End-to-End Encryption (E2EE) is a security method where:

- Data is encrypted **on the sender's device**
- It remains encrypted while traveling through servers/network
- It is decrypted **only on the recipient's device**

 No one in between — not ISPs, hackers, or even the service provider — can read the message.

Example Apps That Use E2EE

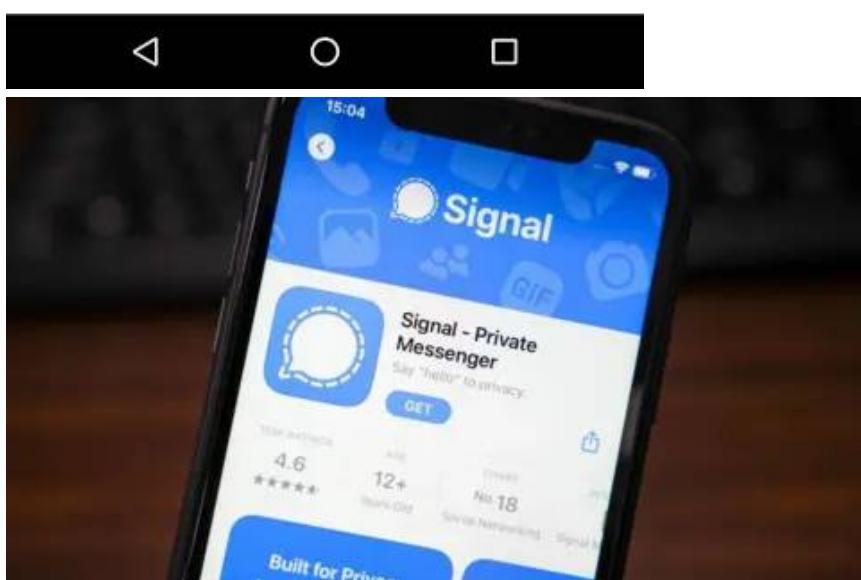


When possible, the messages you send and your calls are secured with end-to-end encryption, which means WhatsApp and third parties can't read or listen to them.

[Learn more about WhatsApp security.](#)

Show security notifications

Turn on this setting to receive notifications when a contact's security code has changed. Your calls and the messages you send are encrypted regardless of this setting, when possible.



**Bas**

last seen at 22:25



You have invited Bas to join a secret chat.

Secret chats:

- 🔒 Use end-to-end encryption
- 🔒 Leave no trace on our servers
- 🔒 Have a self-destruct timer
- 🔒 Do not allow forwarding



Waiting for Bas to get online...



Apps like:

- WhatsApp
- Signal
- Telegram (only Secret Chats)
- iMessage

use E2EE to protect conversations.

🔍 How E2EE Works (Conceptually)

Let's say Alice sends a message to Bob.

Step 1: Key Generation

- Bob generates:
 - **Public key** → shared with everyone
 - **Private key** → kept secret

Step 2: Encrypting the Message

- Alice encrypts the message using Bob's **public key**
- Ciphertext is sent over the internet

Step 3: Decryption

- Bob uses his **private key** to decrypt the message

Even if someone intercepts the data:

```
pgsql
```

```
HELLO → (encrypted) → 8F3A9X!QZ
```

They cannot read it without Bob's private key.

🛡 How To Use E2EE in an App-Based Chat System

Now let's design it like a Computer Science student would.

1 Use Public-Key Cryptography for Key Exchange

Example:

- **Diffie-Hellman**
- **Elliptic Curve Diffie-Hellman**

Purpose:

- Securely establish a shared symmetric key between users
-

2 Use Symmetric Encryption for Messages

Once shared key is created:

- Use **AES**
- Or secure modes like:
 - AES-GCM
 - ChaCha20-Poly1305

Symmetric encryption is:

- ✓ Fast
 - ✓ Efficient
 - ✓ Suitable for real-time chat
-

3 Use Authentication (Integrity + No Tampering)

Add:

- **Message Authentication Code**
- Or authenticated encryption (AES-GCM)

This ensures:

- Message is not modified
 - Sender is authentic
-

4 Implement Perfect Forward Secrecy

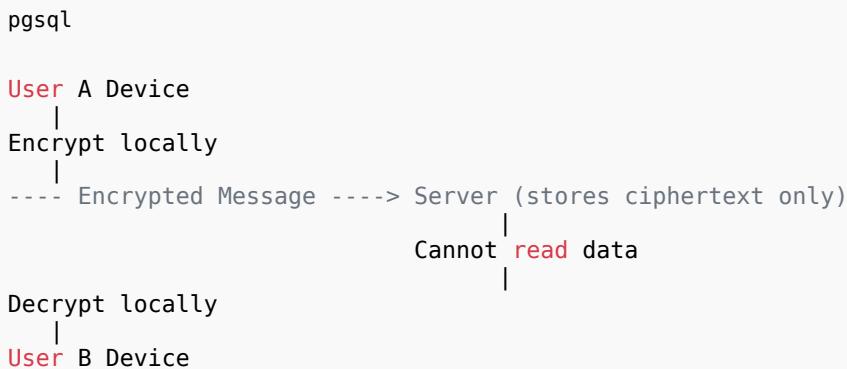
Use:

- Ephemeral keys (new keys per session/message)
- Like in Signal Protocol

Even if long-term key is compromised:

- Old chats remain secure

Architecture of Secure Chat with E2EE



The server only stores encrypted blobs.

Without End-to-End Encryption

If encryption is only between user ↔ server:



- Server can read messages
- Insider attacks possible
- Government or breach can expose chats

That is NOT true E2EE.

How Apps Like Signal Secure Chat

Signal Protocol uses:

- Double Ratchet Algorithm
- Prekeys
- X3DH key agreement
- Forward secrecy
- Post-compromise security

Used by:

- **Signal**
- **WhatsApp**

Security Properties Achieved

Property	How E2EE Provides It
Confidentiality	Only sender & receiver decrypt
Integrity	MAC / AEAD modes
Authentication	Public key identity
Forward Secrecy	Ephemeral keys
Post-compromise security	Key ratcheting

If You Were Building One

Since you're into backend + crypto systems:

You'd need:

- Client-side key generation
- Secure key storage
- Secure random number generator
- Use battle-tested libraries (libsodium, OpenSSL)
- Never write custom crypto
- End-to-end key verification (QR code fingerprint)

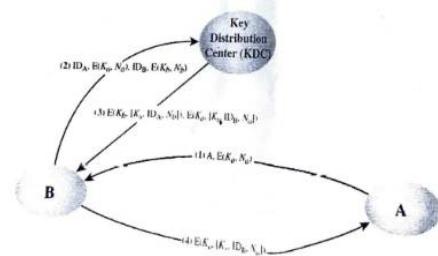
If you want, I can next explain:

-  How Signal's Double Ratchet works (step-by-step)
-  How to implement E2EE in your own Node.js app
-  Common E2EE implementation mistakes
-  How E2EE differs from TLS

Tell me which direction you want 



Q3: The following figure shows that one local area network vendor provides a key distribution facility.



(i) Explain the key distribution scheme

The figure represents a **Key Distribution Center (KDC) based symmetric key distribution scheme**, similar to the **Needham-Schroeder Symmetric Key Protocol**.

There are three parties:

- **A** → Initiator
- **B** → Receiver
- **KDC** → Trusted Key Distribution Center

Each user shares a **long-term secret key** with the KDC:

- A shares K_A with KDC
- B shares K_B with KDC

The goal is to establish a **session key K_s** between A and B securely.

Step-by-Step Explanation of the Scheme

Step (1): A → KDC

A sends a request to the KDC:

$$ID_A, ID_B, N_A$$

- ID_A = Identity of A
- ID_B = Identity of B
- N_A = Nonce generated by A (to prevent replay attacks)

Purpose:

- A tells KDC that it wants to communicate with B.
- The nonce ensures freshness.

Step (2): KDC → A

KDC generates a **session key** K_s and sends:

$$E_{K_A}(K_s, ID_B, N_A, E_{K_B}(K_s, ID_A))$$

Encrypted using A's long-term key K_A .

Inside this message:

1. $K_s \rightarrow$ session key
2. $ID_B \rightarrow$ confirms intended receiver
3. $N_A \rightarrow$ proves freshness
4. Ticket for B $\rightarrow E_{K_B}(K_s, ID_A)$

Purpose:

- Only A can decrypt because it's encrypted with K_A .
- A gets:
 - The session key
 - A "ticket" meant for B.

Step (3): A → B

A forwards B's ticket:

$$E_{K_B}(K_s, ID_A)$$

Since only B knows K_B , only B can decrypt this.

After decryption, B learns:

- Session key K_s
- Identity of A

Now both A and B share K_s .

Step (4): B → A (Mutual Authentication)

B sends:

$$E_{K_s}(N_B)$$

A decrypts it and returns:

$$E_{K_s}(f(N_B))$$

This proves:

- Both parties know the session key
- Mutual authentication is achieved

What This Scheme Achieves

Security Property	How It Is Achieved
Confidentiality	Session key encrypted with long-term keys
Authentication	Ticket contains identity
Freshness	Nonces N_A, N_B
Secure key delivery	KDC generates and distributes K_s

Why KDC is Needed

Without KDC:

- Every pair of users must share a key.
- For n users $\rightarrow \frac{n(n-1)}{2}$ keys required.

With KDC:

- Each user only shares **one long-term key** with KDC.
- KDC dynamically creates session keys.

Important Points (Exam-Oriented)

- It is a **symmetric key distribution protocol**
- KDC is a **trusted third party**
- Uses **nonces** to prevent replay attacks
- Establishes a **temporary session key**
- Reduces key management complexity

If you want, I can also explain:

- The replay attack vulnerability in original Needham-Schroeder
- How Kerberos improves this scheme
- Compare KDC vs Diffie-Hellman key exchange



(ii) Highlight the advantages and disadvantages of key distribution scheme

The question refers to the **KDC-based symmetric key distribution scheme** (like Needham-Schroeder / Kerberos style).

Here are the **advantages and disadvantages** clearly for exam writing:

✓ Advantages of KDC-Based Key Distribution

1 Reduces Key Management Complexity

- Without KDC:
 $\frac{n(n-1)}{2}$ keys required for n users.
- With KDC:
 - Each user shares only **one long-term key** with KDC.
 - Total keys required = **n**.

✓ Scalable for large networks.

2 Session Keys for Every Communication

- KDC generates fresh **session key** K_s for each session.
- Limits exposure if a session key is compromised.

✓ Improves confidentiality.

3 Provides Mutual Authentication

- Use of nonces ensures:
 - Freshness
 - No replay attacks (partially)
- Both A and B prove knowledge of K_s .

4 Centralized Control

- Easy to:
 - Add new users
 - Revoke compromised users
 - Monitor authentication

- ✓ Better administrative control.

5 Efficient (Uses Symmetric Cryptography)

- Uses symmetric encryption (faster than public key crypto).
- Suitable for LAN environments.

✗ Disadvantages of KDC-Based Key Distribution

1 Single Point of Failure

- If KDC fails → entire system stops.
- If KDC is compromised → attacker can:
 - Generate fake session keys
 - Decrypt all communications

⚠ Major security risk.

2 Key Storage Risk at KDC

- KDC stores all long-term secret keys.
- If compromised, full network security collapses.

3 Scalability Bottleneck

- All key requests must pass through KDC.
- In large systems → performance bottleneck.

4 Vulnerable to Replay (Original Version)

- Basic Needham-Schroeder protocol had replay attacks.
- Requires timestamps or additional protection.

(Kerberos fixed this using timestamps.)

5 Trust Dependency

- Entire system security depends on trusting the KDC completely.

Final Exam Summary (Short Form)

Advantages:

- Reduces number of keys required
- Provides fresh session keys
- Enables mutual authentication
- Easy centralized management
- Efficient (symmetric cryptography)

Disadvantages:

- Single point of failure
- High-value attack target
- Scalability issues
- Possible replay attacks
- Requires full trust in KDC

If you want, I can also provide a **5-mark perfect exam answer format** (compact writing style).