

Parallelization of Solving Influence Maximization Problem in Social Network

Zeyu Fu

Model of Influence

Independent Cascade Model (IC)

For a network $G = (V, E)$, state of node $v_i \in V$ will only be active (1) or inactive (0)

- (1) Initial state, all nodes are inactive
- (2) At $t = 0$, an initial seed set X_0 are activated ($0 \rightarrow 1$)
- (3) At $t = n$, the set of newly activated nodes at time $n - 1$, $v \in X_{n-1}$, will infect (activate) the inactive outgoing neighbors u with a probability $P_{u,v}$
- (4) Propagation ends when no newly active node exists

Linear Threshold Model (LT)

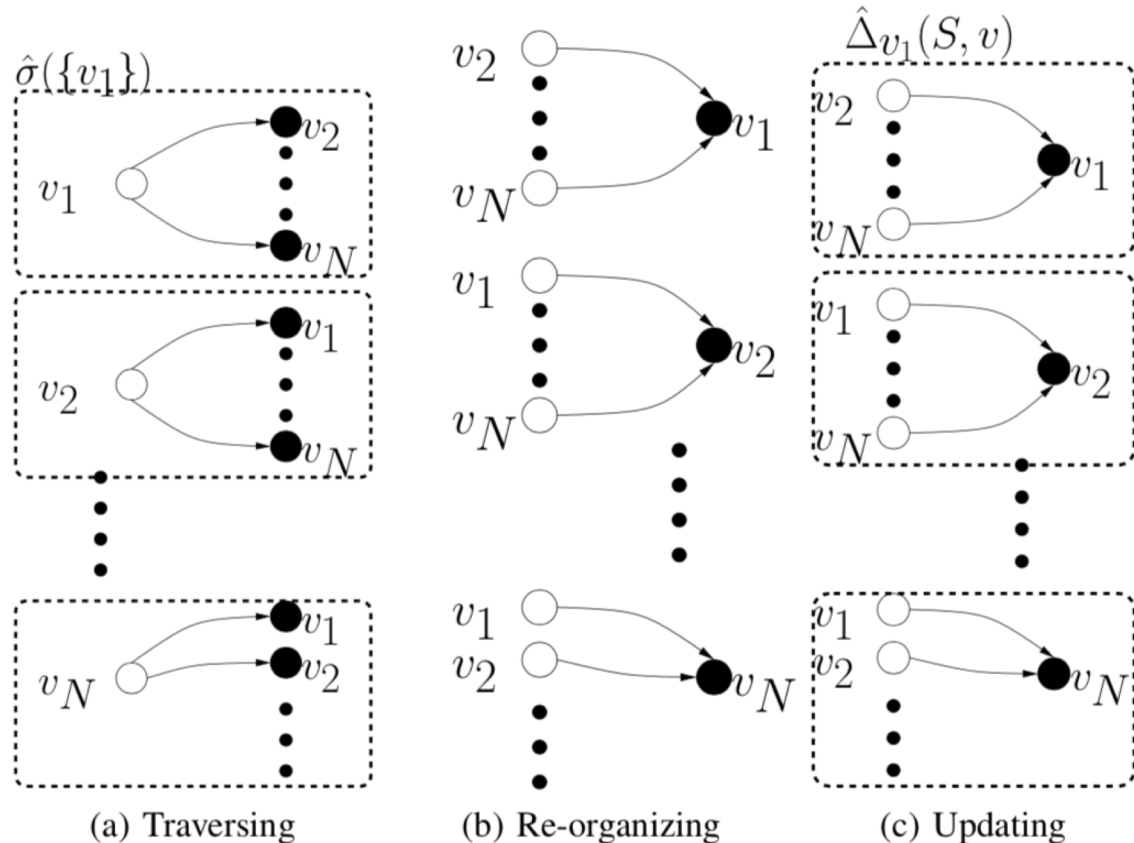
For a network $G = (V, E)$, node $v_i \in V$ will have an activation threshold $\theta_v \in [0, 1]$ uniformly at random

- (1) Same as IC model
- (2) Same as IC model
- (3) At $t = n$, node v become active if $\sum_{u \in \eta^{in}(v) \cap H_{t-1}} b(u, v) \geq \theta_v$, where $b(u, v) = 1 / L_{in}(v)$, $L_{in}(v)$ is the incoming degree of node v
- (4) Same as IC model

Influence Maximization Problem

Given a network $G(V, E)$ and a size k , under a specific influence model, we want to find a seed set S of k nodes such that the expected spread of influence is maximized

Independent Path Algorithm (IPA) : A parallel way to solve this problem



Project goal:

1. Practice IPA algorithm (maybe traversing step of IPA)
2. Compare the performance with serial code

Ref: Kim, J., Kim, S. K., & Yu, H. (2013, April). Scalable and parallelizable processing of influence maximization for large-scale social networks?. In *2013 IEEE 29th International Conference on Data Engineering (ICDE)* (pp. 266-277). IEEE.

IPA Pseudo Code

Algorithm 2 traversing(G, θ)

Require: G : graph, θ : threshold

Ensure: R : set of $\hat{\sigma}_I(\{v\})$

```
1:  $R \leftarrow \phi$ 
2: parallel for  $v \in V$  do
3:    $\hat{\sigma}_I(\{v\}) \leftarrow 1$ 
4:   generate  $P_{v \rightarrow V}$  by traversing  $G$  with  $\theta$ 
5:   for  $u \in O_v$  do
6:      $\hat{\sigma}_I(\{v\}) \leftarrow \hat{\sigma}_I(\{v\}) + \hat{\sigma}_I^u(\{v\})$ 
7:   end for
8:    $R \leftarrow R \cup \{\hat{\sigma}_I(\{v\})\}$ 
9: end for
10: return  $R$ 
```

Algorithm 3 updating(S, v)

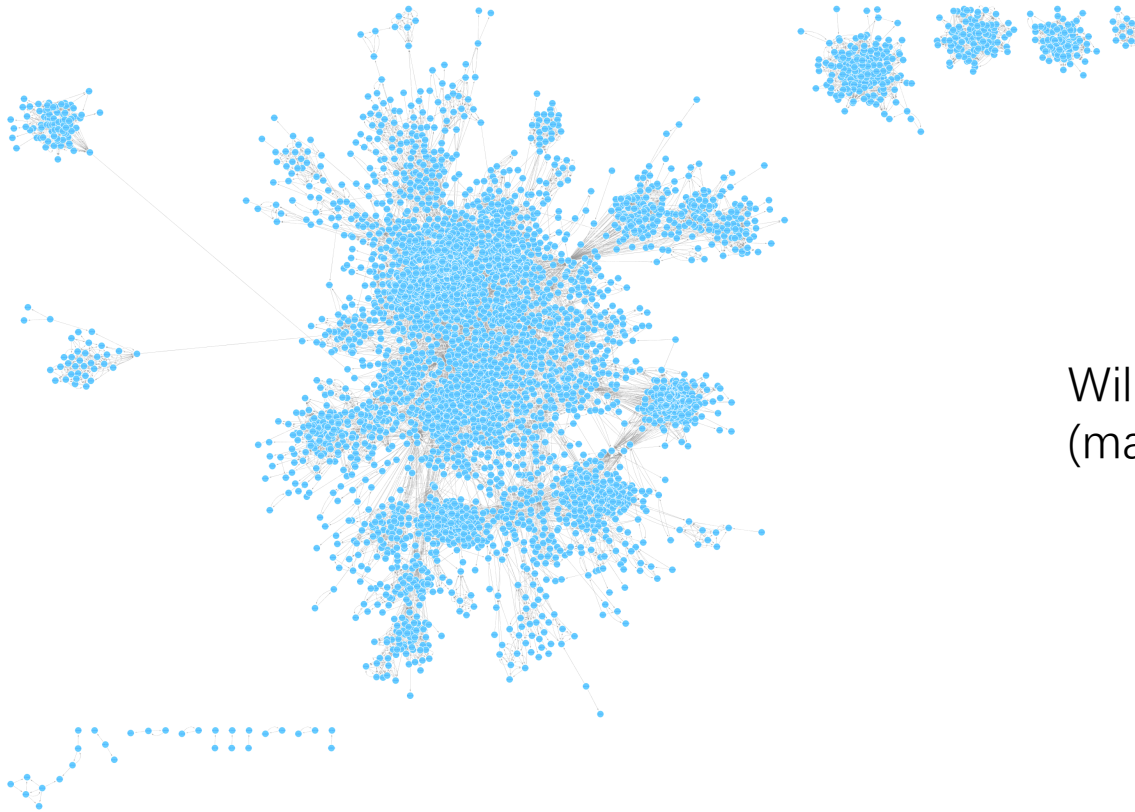
Require: S : current seed nodes, v : new seed node

Ensure: $\hat{\Delta}(S, v)$

```
1:  $ret \leftarrow 1$ 
2: parallel for  $u \in O_v \cup \{v\}$  do
3:    $new \leftarrow 1; old \leftarrow 1$ 
4:   for  $p \in P_{v \rightarrow u}$  and  $p \cap S = \phi$  do
5:      $new \leftarrow new \times (1 - ipp(p))$ 
6:   end for
7:   for  $s \in S$  do
8:     for  $p \in P_{s \rightarrow u}$  and  $p \cap S \subseteq \{s, v\}$  do
9:        $old \leftarrow old \times (1 - ipp(p))$ 
10:      if  $p \cap S = \{s\}$  then
11:         $new \leftarrow new \times (1 - ipp(p))$ 
12:      end if
13:    end for
14:  end for
15:   $ret \leftarrow ret + (1 - new) - (1 - old)$ 
16: end for
17: return  $ret$ 
```

Dataset: Stanford Large Network Dataset

Twitter users relationship data (directed)



Will first work on a very small sub-graph
(maybe tens of nodes)

Partial network with 4514 nodes and 79600 edges