

INF 212

ANALYSIS OF PROG. LANGS

FUNCTION COMPOSITION

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Topics

- Recursion
- Higher-order functions
- Continuation-Passing Style
- Monads (take 1)
 - ▣ Identity Monad
 - ▣ Maybe Monad



Recursion

Prototypical Example

```
fact(n):  
    if (n <= 1) then 1  
    else n * fact(n-1)
```

Thinking Recursively

- Add numbers in a list
- Print a list of numbers
- Check if a number is in a list
- (Live coding)

Tail Recursion

(first-order case)

slide 6

- Function g makes a **tail call** to function f if return value of function f is return value of g

- Example

tail call

not a tail call

$\text{fun } g(x) = \text{if } x > 0 \text{ then } f(x) \text{ else } f(x) * 2$

- Optimization: can pop current activation record on a tail call
 - ▣ Especially useful for recursive tail call because next activation record has exactly same form

Example of Tail Recursion

slide 7

Calculate least power of 2 greater than y

$f(1,3)$

control		↑
return val		
x	1	↑
y	3	

```
fun f(x,y) = if x>y  
  then x  
  else f(2*x, y);  
f(1,3) + 7;
```

control		↑
return val		
x	1	↑
y	3	

control		↑
return val		
x	2	↑
y	3	

control		↑
return val		
x	4	↑
y	3	

Tail Recursion Elimination

slide 8

$f(1,3)$

control		↑
return val		↑
x	1	
y	3	

$f(2,3)$

control		↑
return val		↑
x	2	
y	3	

$f(4,3)$

control		↑
return val		↑
x	4	
y	3	

```
fun f(x,y) = if x>y  
  then x  
  else f(2*x, y);  
f(1,3) + 7;
```

Optimization

- pop followed by push - reuse activation record in place
- Tail recursive function is equivalent to iterative loop

Tail Recursion and Iteration

slide 9

$f(1,3)$

control		↑
return val		↑
x	1	
y	3	

$f(2,3)$

control		↑
return val		↑
x	2	
y	3	

$f(4,3)$

control		↑
return val		↑
x	4	
y	3	

```
fun f(x,y) = if x > y  
  then x  
  else f(2*x, y);  
f(1,y);
```

test

loop body

initial value

```
function g(y) {  
  var x = 1;  
  while (!x > y)  
    x = 2*x;  
  return x;  
}
```



Higher-order functions

Higher-Order Functions

slide 11

- Function passed as argument
 - ▣ Need pointer to activation record “higher up” in stack
- Function returned as the result of function call
 - ▣ Need to keep activation record of the returning function
- Functions that take function(s) as input and return functions as output are known as functionals

Return Function as Result

slide 12

- Language feature (e.g., Python, ML, ...)
- Functions that return “new” functions
 - ▣ Example: `fun compose(f,g) = (fn x => g(f x));`
 - ▣ Function is “created” dynamically
 - Expression with free variables; values determined at run-time
 - ▣ Function value is closure = $\langle \text{env}, \text{code} \rangle$
 - ▣ Code not compiled dynamically (in most languages)
 - ▣ Need to maintain environment of the creating function

Closures

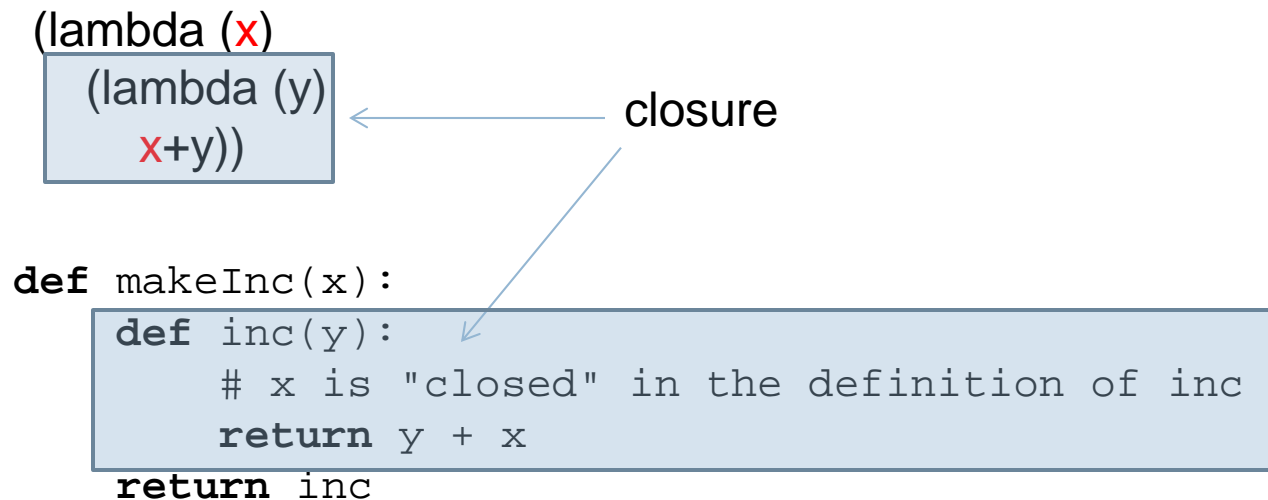


slide 13

- Function value is pair **closure** = $\langle \text{env}, \text{code} \rangle$
 - ▣ Statically scoped function must carry a link to its static environment with it
 - ▣ Only needed if function is defined in a nested block
- When a function represented by a closure is called...
 - ▣ Allocate activation record for call (as always)
 - ▣ Set the access link in the activation record using the environment pointer from the closure

Closures

- Function with free variables that are bound to values in the enclosing environment



Note to self: illustrate closures in Python and C# (my examples)

What are closures good for?



- For changing your mind later!
 - ▣ Replaces constants and variables with functions
 - ▣ Replaces conditionals
 - ▣ ...



Return Function with Private State

slide 16

```
fun mk_counter (init : int) =  
  let val count = ref init  
      fun counter(inc:int) =  
        (count := !count + inc; !count)  
      in  
        counter  
      end;  
val c = mk_counter(1);  
c(2) + c(2);
```

- Function to “make counter” returns a closure
- How is correct value of `count` determined in `c(2)` ?

Implementing Closures

slide 17

- Closures as used to maintain static environment of functions as they are passed around
- May need to keep activation records after function returns
 - ▣ Stack (last-in-first-out) order fails! (why?)
- Possible “stack” implementation:
 - ▣ Put activation records on heap
 - ▣ Instead of explicit deallocation, invoke garbage collector as needed



Continuations

Continuations

- Representation of the control state of a program
 - ▣ Data structure available to the programmer instead of hidden
 - ▣ Contains the current stack and point in the computation
- Can be later used to return to that point

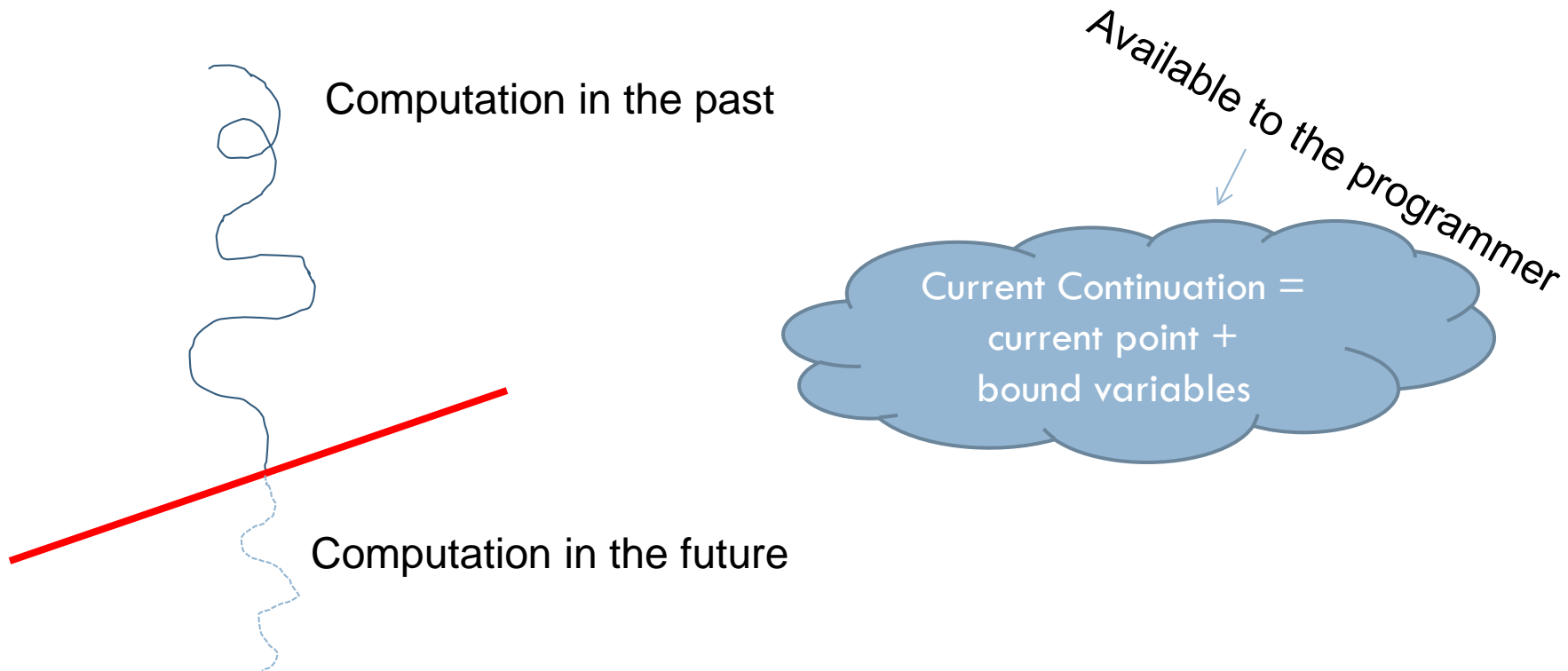
Remember Goto

```
A: blah  
    blah  
    if something GOTO A else GOTO B  
B: ...
```

Flow control via textual labels mixes computation (beta-reductions) and representation (the text of the program)

Continuations continued

- Elegant concept for arbitrary flow control



Note to self: illustrate continuations in Scheme (Wikipedia)

What are continuations good for?

- Everything control flow!
 - ▣ Co-routines
 - ▣ Exceptions
 - ▣ Preserving flow
in non-blocking I/O
 - (rhymes!)

The continuation nature of exceptions

```
function fact (n) {  
  if (n < 0)  
    throw "n < 0" ;  
  else if (n == 0)  
    return 1 ;  
  else  
    return n * fact(n-1) ;  
}
```

```
function total_fact (n) {  
  try {  
    return fact(n) ;  
  } catch (ex) {  
    return false ;  
  }  
}
```

```
document.write("total_fact(10): " + total_fact(10)) ;  
document.write("total_fact(-1): " + total_fact(-1)) ;
```

The continuation nature of exceptions – desugaring the previous slide

```
function fact (n,r,t) {  
  if (n < 0)  
    t ("n < 0")  
  else if (n == 0)  
    r(1)  
  else  
    fact(n-1,  
        function (t0) {  
          r (n*t0) ;  
        },  
        t)  
}
```

```
function total_fact (n,ret) {  
  fact (n,ret,  
    function (ex) {  
      ret(false) ;  
    }) ;  
}
```

```
total_fact(10, function (res) {  
  document.write("total_fact(10): " + res)  
}) ;  
  
total_fact(-1, function (res) {  
  document.write("total_fact(-1): " + res)  
}) ;
```


I/O and continuations

Blocking (I/O in most systems)

```
contents = fs.ReadFile(path);  
with contents do  
    blah
```

*Blocks here until
we have the result*

Non-blocking

```
contents = fs.ReadFileAsync(path);  
with contents do  
    blah
```

*Uh-oh, we still don't
have it*

How to solve this?

I/O and continuations

Non-blocking

```
fs.readFileAsync(path, lambda(contents)
{
    with contents do
        blah
} ) ;
```

It's a callback!

It's the “current continuation” of the blocking form

JavaScript is FULL of this, so are jquery and node.js



Monads

Monads – what is the problem?

- The problem: how to affect the world
- Problem is more prevalent in pure functional programming style
 - ▣ No side-effects
 - ▣ That's right: no side-effects!
- But you've all seen it too!

No side effects?! Why?

- Easier to test: idempotent functions
- Easier to parallelize

- But the world is ALL about side-effects, right?
 - ▣ Storage, network, UI, ...
 - ▣ Programs affect and control objects and activities in the real world

Example – a Tracing monad

```
def hypotenuse(x, y):  
    return math.sqrt(math.pow(x, 2) + math.pow(y, 2))
```

Now we want to trace it, or affect the world in it:

```
def hypotenuse(x, y):  
    h = math.sqrt(math.pow(x, 2) + math.pow(y, 2))  
    print "In hypotenuse " + h  
    return h
```



Example – a Tracing monad

```
def hypotenuse(x, y):  
    h = math.sqrt(math.pow(x, 2) + math.pow(y, 2))  
    return h, "In hypotenuse" + h
```

Signature was float, float -> float

Signature now is float, float -> float, string

```
> math.pow(hypotenuse(6, 16), 4);
```



What is a monad?

- It's a container
- An active container... it has behavior to:
 - ▣ Wrap itself around a [typed] value
 - ▣ Bind *functions* together

What is a monad?

- [A type constructor, m]
- A function that builds values of that type
 $a \rightarrow m\ a$ (what you'd normally call a constructor in OOP)
- A function (bind) that combines values [of that type] with computations that produce values [of that type]
 $m\ a \rightarrow (a \rightarrow m\ b) \rightarrow m\ b$
- An unwrap function that shows “what’s inside”