Evolution of CSMA Protocols for the IEEE 802.11 Standard

Michael Shell School of Electrical and Computer Engineering Georgia Institute of Technology Atlanta, Georgia 30332-0250

Email: http://www.michaelshell.org/contact.html

Homer Simpson Twentieth Century Fox Springfield, USA

James Kirk and Montgomery Scott Starfleet Academy Email: homer@thesimpsons.com San Francisco, California 96678-2391 Telephone: (800) 555-1212

Fax: (888) 555-1212

Abstract—In this paper we present the requirements of candidate protocols to replace the pervasive CSMA/CA medium access control. We discuss the possibility of further preventing collisions and provide an overview of the related work. We specify protocols that are candidates of replacing CSMA/CA in pseudocode and use simulation to assess performance metrics such as throughput, fairness and collision probability.

I. Introduction

This demo file is intended to serve as a "starter file" for IEEE conference papers produced under LATEX using IEEEtran.cls version 1.7 and later. I wish you the best of success.

A candidate to replace CSMA/CA should

- Provide performance advantages, either in the form of throughput or short term fairness.
- Be backward compatible with current implementation.
- Be simple a simple evolution implementation to ease the transition and reduce time to market (Optional but desirable).

II. RELATED WORK

Since the popularization of IEEE 802.11, several papers have proposed modifications to the contention protocol that is used for sharing the medium. They can be categorized in three groups regarding the approach they use. The first one prevents that the contention window is reset to its minimum value after successful transmissions. Examples of this first group include [1], [2]. This solution improves throughput in saturation conditions at the price of lowering short term fairness.

The second groups involves the accurate estimation of the number of contenders to adjust the contention parameters. Two examples of this group are [3], [4]. This approach offers some throughput and fairness gains at the expense of increased implementation complexity. As the number of contenders is estimated relying on the number of collisions, the presence of channel errors further complicates the estimation. Furthermore, there is a fundamental trade-off between the accuracy and the reaction time of the estimation.

The aforementioned solutions are not able to fairly share the medium with legacy devices. In fact, these proposals are, generally speaking, less aggressive than the currently implemented protocol. Consequently, in a hypothetical mixed network in which the new and old protocols coexist, the new stations would receive a smaller share of the available bandwidth in a scenario.

A more important limitation of the solutions exposed so far is that the throughput is bounded by that of CSMA/CA with optimal configuration [3], [5]. In the present paper we focus on a third group of solutions that delivers throughput above the maximum attainable by CSMA/CA.

This third group of solutions uses a deterministic backoff after successful transmissions to further reduce the chances of collisions. Under certain conditions, collision-free operation is reached. It was introduced in [6] and a more detailed analysis that includes both saturated and non-saturated conditions is presented in [7]. A more in-depth study is carried out in [8], including realistic elements such as clear channel assessment errors. Different aspects of fairness are addressed in [8]–[10]. The performance in realistic channels taking into account the Auto Rate Fallback mechanism is evaluated [11], [12].

Even though initial research efforts where focused on the WLAN collision problem, more recent papers try to extend the idea to multi-hop networks. In [13], the multi-hop slotted case explored. The more realistic situation in multi-hop networks in which the time is not slotted is studied in [14].

The same principles that we exploit to prevent collisions in WLANs can be used in other areas of radio resource management in wireless area networks [15]-[18].

In all the previous work on collision-free operation in WLANs mentioned so far, there is the limitation that the number of contenders should not exceed the value of the deterministic backoff used after successful transmissions. If this value is exceeded, it is no longer possible to achieve collision-free operation. A first solution to solve this problem is presented in [19], but it requires the presence of a central entity (typically the access point) that instructs the other nodes to adjust the value of their deterministic backoff.

In the present paper, we study a completely distributed solution to accommodate a large number of contenders in a fair collision-free fashion.

```
1 b \leftarrow \mathcal{U}[0, CW_{\min} - 1];
                                                                       1 b \leftarrow \mathcal{U}[0, \mathrm{CW}_{\min} - 1];
2 while there is a packet to transmit do
                                                                          /★ Hysteresis: The backoff stage is
       a \leftarrow 0;
                                                                               reset only when a node joins the
       while a < A do
                                                                               contention or the queue is empty. */
           while b > 0 do
5
                                                                       a \leftarrow 0:
               wait 1 slot;
6
                                                                       3 while there is a packet to transmit do
               b \leftarrow b - 1;
                                                                              while a < A do
                                                                                  while b > 0 do
           Attempt transmission;
                                                                                      wait 1 slot;
           if success then
10
                                                                                      b \leftarrow b - 1;
               /* Random backoff.
                                                                                  end
                                                                       8
11
               b \leftarrow \mathcal{U}[0, \text{CW}_{\text{min}} - 1];
                                                                                  Attempt transmission;
               break;
12
                                                                                  if success then
           else
13
                                                                                      /* Deterministic backoff.
               /* Random backoff.
                                                                                      b \leftarrow \text{CW}_{\text{min}}/2;
                                                                       11
               a \leftarrow a + 1:
14
                                                                       12
                                                                                      break;
               b \leftarrow \mathcal{U}[0, \min(CW_{\min} * 2^{a}, CW_{\max}) - 1];
15
                                                                       13
           end
16
                                                                                      a \leftarrow a + 1;
                                                                       14
17
       end
                                                                                      /* fall to random backoff.
                                                                                      b \leftarrow \mathcal{U}[0, \min(CW_{\min} * 2^{a}, CW_{\max}) - 1];
18 end
                                                                       15
                   Algorithm 1: CSMA/CA
                                                                       16
                                                                                      if a > m then
                                                                                          Discard packet;
                                                                       17
                                                                                          break;
                                                                       18
1 b \leftarrow \mathcal{U}[0, CW_{\min} - 1];
                                                                       19
                                                                                      end
2 while there is a packet to transmit do
                                                                                  end
                                                                       20
       a \leftarrow 0:
                                                                              end
                                                                       21
       while a < A do
4
                                                                       22 end
           while b > 0 do
                                                                                 Algorithm 3: CSMA/ECA with hysteresis
               wait 1 slot;
6
               b \leftarrow b - 1;
7
           end
8
                                                                        1 b \leftarrow \mathcal{U}[0, \mathrm{CW}_{\min} - 1];
           Attempt transmission;
                                                                          /★ Hysteresis: The backoff stage is
           if success then
                                                                               reset only when a node joins the
10
               /* Deterministic backoff.
                                                               * /
                                                                               contention or the queue is empty.
11
               b \leftarrow \text{CW}_{\text{min}}/2;
                                                                       a \leftarrow 0;
               break;
12
                                                                       3 while there is a packet to transmit do
           else
                                                                              while a < A do
13
               a \leftarrow a + 1 ;
14
                                                                                  while b > 0 do
               /* fall to random backoff.
                                                                       6
                                                                                      wait 1 slot;
15
               b \leftarrow \mathcal{U}[0, \min(\mathrm{CW}_{\min} * 2^{\mathrm{a}}, \mathrm{CW}_{\max}) - 1];
                                                                       7
                                                                                      b \leftarrow b - 1;
           end
                                                                                  end
16
       end
17
                                                                                  /* Fair-share: 2^a packets are
18 end
                                                                                       transmitted.
                  Algorithm 2: CSMA/ECA
                                                                                  Attempt aggregate transmission of 2^a packets;
                                                                       9
                                                                                  if success then
                                                                       10
                                                                                      /* Deterministic backoff.
                                                                                      b \leftarrow \text{CW}_{\text{min}}/2;
                    III. ENHANCED CSMA
                                                                       11
                                                                                      break;
                                                                       12
               IV. PERFORMANCE EVALUATION
                                                                       13
                                                                                  else
                                                                                      a \leftarrow a + 1;
                        V. Conclusion
                                                                                      /* fall to random backoff.
                                                                                      b \leftarrow \mathcal{U}[0, \min(\mathrm{CW_{\min}} * 2^{\mathrm{a}}, \mathrm{CW_{\max}}) - 1];
                                                                       15
   The conclusion goes here.
                                                                       16
                                                                                  end
                                                                              end
                      ACKNOWLEDGMENT
                                                                       18 end
                                                                         Algorithm 4: CSMA/ECA with hysteresis and fair-share
   The authors would like to thank...
```

REFERENCES

- [1] V. Bharghavan, A. Demers, S. Shenker, and L. Zhang, "MACAW: a media access protocol for wireless LAN's," in ACM SIGCOMM Computer Communication Review, vol. 24, no. 4. ACM, 1994, pp. 212–225.
- [2] C. Wang, B. Li, and L. Li, "A new collision resolution mechanism to enhance the performance of IEEE 802.11 DCF," *IEEE Transactions on Vehicular Technology*, vol. 53, no. 4, pp. 1235–1246, 2004.
- [3] F. Cali, M. Conti, E. Gregori, and P. Aleph, "Dynamic Tuning of the IEEE 802.11 Protocol to Achieve a Theoretical Throughput Limit," IEEE/ACM Trans. Netw., vol. 8, no. 6, pp. 785–799, 2000.
- [4] A. Lopez-Toledo, T. Vercauteren, and X. Wang, "Adaptive Optimization of IEEE 802.11 DCF Based on Bayesian Estimation of the Number of Competing Terminals," *IEEE Trans. Mobile Comput.*, vol. 5, no. 9, p. 1283, 2006.
- [5] G. Bianchi, "Performance Analysis of the IEEE 802.11 Distributed Coordination Function," *IEEE J. Sel. Areas Commun.*, vol. 18, no. 3, pp. 535–547, 2000.
- [6] J. Barcelo, B. Bellalta, C. Cano, and M. Oliver, "Learning-BEB: Avoiding Collisions in WLAN," in *Eunice*, 2008.
- [7] J. Barcelo, B. Bellalta, A. Sfairopoulou, C. Cano, and M. Oliver, "CSMA with Enhanced Collision Avoidance: a Performance Assessment," in *IEEE VTC Spring*, 2009.
- [8] Y. He, R. Yuan, J. Sun, and W. Gong, "Semi-Random Backoff: Towards Resource Reservation for Channel Access in Wireless LANs," in *IEEE ICNP*, 2009, pp. 21–30.
- [9] J. Barcelo, A. Lopez-Toledo, C. Cano, and M. Oliver, "Fairness and Convergence of CSMA with Enhanced Collision Avoidance," in *IEEE ICC*, 2010.
- [10] M. Fang, D. Malone, K. Duffy, and D. Leith, "Decentralised Learning MACs for Collision-free Access in WLANs," Wireless Networks, to appear, preprint available in arXiv.
- [11] G. Martorell, F. Riera, G. Femenias, J. Barcelo, and B. Bellalta, "On the performance evaluation of CSMA/E2CA protocol with open loop ARFbased adaptive modulation and coding," in *European Wireless*, 2012.
- [12] G. Martorell, F. Riera-Palou, and G. Femenias, "Tuning Fast Link Adaptation Algorithms for CSMA/CA-and CSMA/E2CA-Based WLANs," in *ICCCN*. IEEE, 2012, pp. 1–7.
- [13] K. Hui, T. Li, D. Guo, and R. Berry, "Exploiting peer-to-peer state exchange for distributed medium access control," in *Information Theory Proceedings (ISIT)*, 2011 IEEE International Symposium on. IEEE, 2011, pp. 2368–2372.
- [14] J. Barcelo, B. Bellalta, C. Cano, A. Sfairopoulou, and M. Oliver, "On the Distributed Construction of a Collision-Free Schedule in Multi-Hop Packet Radio Networks," in *Telecommunication Systems*, to appear. Available in ArXiV.
- [15] K. R. Duffy, C. Bordenave, and D. J. Leith, "Decentralized Constraint Satisfaction," CoRR, vol. abs/1103.3240, 2011.
- [16] A. Checco, R. Razavi, D. Leith, and H. Claussen, "Self-Configuration of Scrambling Codes for WCDMA Small Cell Networks," in *PIMRC*, 2012.
- [17] A. Checco and D. Leith, "Learning-Based Constraint Satisfaction With Sensing Restrictions," arXiv preprint arXiv:1210.7156, 2012.
- [18] Z. Khan, J. Lehtomaki, L. DaSilva, and M. Latva-aho, "Autonomous Sensing Order Selection Strategies Exploiting Channel Access Information," *Mobile Computing, IEEE Transactions on*, to appear. Available in IEEE Xplore.
- [19] J. Barcelo, B. Bellalta, C. Cano, A. Sfairopoulou, and M. Oliver, "Towards a Collision-Free WLAN: Dynamic Parameter Adjustment in CSMA/E2CA," in EURASIP Journal on Wireless Communications and Networking, 2011.