

Encryption and decryption of MTP:

File : key_mtp.py

Generates a cryptographically secure key of 2^{20} bits and stores it in a file named otp_secret

File : encrypt_mtp.py

python encrypt_mtp.py plaintext.txt ciphertext.txt

Reads OTP from the otp_secret file. Reads plaintext. The otp_secret is sliced to match the number of bits in the plaintext.

Writes the result of otp_secret xor plaintext to ciphertext

File : decrypt_mtp.py

python decrypt_mtp.py ciphertext.txt viola.txt

Reads OTP from otp_secret file. Reads ciphertext

Writes result of otp_secret xor ciphertext to viola.txt

Problem 2b:

Tries to deduce the location at which space would be present in the plaintext. Once we get the location we can use properties of xor to retrieve a part of the key.

Assuming we have 10 cipher texts encrypted using same key (MTP)

c_1, c_2, \dots, c_{10}

-> Since $c_1 = b_1 \text{ xor key}$, $c_2 = b_2 \text{ xor key}$ etc. If we xor c_1 and c_2 we get a part of plaintext in form of $b_1 \text{ xor } b_2$.

-> Prepare a list of values we could get if we xor english alphabets with ascii value of space.

-> For discovering possible space locations in b_1 ,

- $c_1 \text{ xor } c_2$, $c_1 \text{ xor } c_3$, $c_1 \text{ xor } c_4$ $c_1 \text{ xor } c_{10}$

- Compare the results with the list prepared above.

- If we find a match keep track of it by incrementing the counter of that value. This is done as there can be different combinations that would result in values given in the list.

- Once we have all possible locations, we ignore the one's with low threshold (Some % of number of CTs available). This is done so that other possible combinations can be ignored.

- With the remaining locations, each location of cipher text c_1 is xored with space to iteratively discover a part of key.

-> The above procedure is repeated for all ciphertexts available.