

1) Drop in PSP

$\text{pow}(a, n)$

$\text{pow}(a, n, p)$

TC for recursive codes

SC for recursive codes

Q1) Given  $a, n$  find  $a^n$  using recursion

$n > 0$

$a$	$n$	$a^n$
2	5	$2^5 = 32$
3	4	$3^4 = 81$
5	3	$5^3 = 125$

Assumption: Given  $a, n$  calc & return  $a^n$

$$A^n = A \times A^{n-1}$$

$$\text{pow}(A, N) = \text{pow}(A, N-1) \times A$$

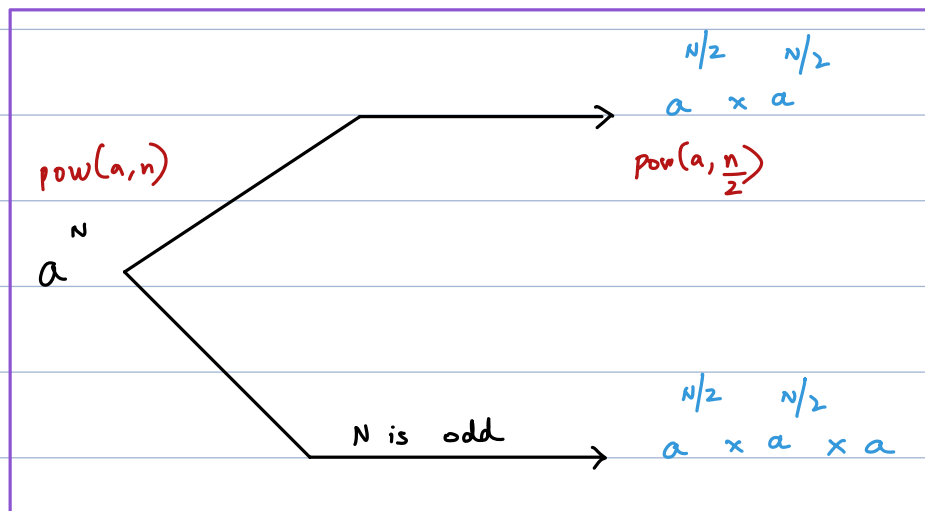
```
int pow(a, n) {  
    Base case: if (n == 1) { return a } // if (n == 0) { return 1 }  
               return pow(a, n-1) * a  
Main logic: }  
✓ ✓
```

$\text{pow}(a, 5) \rightarrow \text{pow}(a, 4) \rightarrow \text{pow}(a, 3) \rightarrow \text{pow}(a, 2) \rightarrow \text{pow}(a, 1)$

$\text{pow}(a, N)$

TC:  $O(N)$

$$\begin{array}{lcl}
 \begin{array}{c} 8 \\ a \end{array} & = & \begin{array}{c} 7 \quad 1 \\ a \times a \end{array} \quad \longrightarrow \quad \begin{array}{c} 8 \\ a \end{array} = \begin{array}{c} 4 \quad 4 \\ a \times a \end{array} \\
 & & \begin{array}{c} 10 \\ a \end{array} = \begin{array}{c} 5 \quad 5 \\ a \times a \end{array} \\
 & & \begin{array}{c} 13 \\ a \end{array} = \begin{array}{c} 6 \quad 6 \\ a \times a \end{array} \times a \\
 & & \begin{array}{c} 15 \\ a \end{array} = \begin{array}{c} 7 \quad 7 \\ a \times a \end{array} \times a
 \end{array}$$



$$\frac{N-1}{2}$$

N is odd

$$N = 13$$

$$\frac{N-1}{2} = 6$$

$$\frac{N}{2} = 6$$

Fast exponentiation

math.pow(a, n)

Main logic:

```
int pow(a, n) {  
    if (n == 1) { return a }  
    ha = pow(a, n/2)  
    if (N & 1 == 1) { # N is odd  
        | return ha * ha * a  
    }  
    else {  
        | return ha * ha  
    }  
}
```

$a \times \times 1000$

$a^{1000}$

```

int pow(a, n) {  n=23  a=2
1 if (n==1) { return a; }
2 ha = pow(a, n/2)  2048
3 if (N&1==1) { # N is odd
   return ha*ha*a
}
4 else {
   return ha*ha
}
}

```

2048 = 2048 \* 2 → output

```

int pow(a, n) {  n=11  a=2
1 if (n==1) { return a; }
2 ha = pow(a, n/2)  32
3 if (N&1==1) { # N is odd
   return ha*ha*a  32*32*2
}
4 else {
   return ha*ha
}
}

```

```

int pow(a, n) {  n=2  a=2
1 if (n==1) { return a; }
2 ha = pow(a, n/2)  a
3 if (N&1==1) { # N is odd
   return ha*ha*a
}
4 else {
   return ha*ha
}
}

```

```

int pow(a, n) {  n=5  a=2
1 if (n==1) { return a; }
2 ha = pow(a, n/2)  a^2=4
3 if (N&1==1) { # N is odd
   return ha*ha*a  4*4*2=32
}
4 else {
   return ha*ha
}
}

```

```

int pow(a, n) {  n=1  a=2
1 if (n==1) { return a; }
2 ha = pow(a, n/2)
3 if (N&1==1) { # N is odd
   return ha*ha*a
}
4 else {
   return ha*ha
}
}

```

$$N \rightarrow \frac{N}{2} \rightarrow \frac{N}{4} \dots 1$$

$$TC: O(\log N)$$

Q3) Given  $a, n, m$  calculate  $a^n \% m$   
 $\hookrightarrow [0, m-1]$

$$1 \leq a \leq 10^9$$

$$1 \leq n \leq 10^9$$

$$1 \leq m \leq 10^9$$

HW: Update this method to take  
care of overflow

10<sup>1000</sup>

```
int pow(a, n) {  
    if (n == 1) return a;  
    long ha = pow(a, n/2);  
    long hp = ((ha % m) * ha % m) % m;  
    if (n & 1 == 1) { # N is odd  
        return (hp % m * a % m) % m;  
    }  
    else {  
        return hp;  
    }  
}
```

1) Modular arithmetic class revise

9:53 - 10:08

1) Do a dry run

Substitution method for TC calculation

# TC for recursive codes using recursive solution

$N=1$

$f(1)$

```
int sum(N) {
```

```
    if (N == 1) { return 1 }
```

```
    return sum(N-1) + N
```

```
}
```

TC  $\rightarrow f(N)$

$$f(N) = O(1) + f(N-1)$$

- 1) ✓ Find recursive time equation
- 2) ✓ Find general solution after R substitution
- 3) Equate unknown with base condition and find TC

$$f(N) = f(N-1) + 1 \quad \text{Recursive time equation}$$

$$f(N) = f(N-1) + 1$$

substitute  $N \rightarrow N-1$

$$f(N) = f(N-1) + 1$$

$$f(N-1) = f(N-2) + 1$$

$$f(N) = (f(N-2) + 1) + 1$$

$$f(N) = f(N-2) + 2$$

$N \rightarrow N-2$

$$f(N) = f(N-1) + 1$$

$$f(N-2) = f(N-3) + 1$$



$$f(N) = (f(N-3) + 1) + 2$$

$$f(N) = f(N-3) + 3$$

⋮

$$f(N) = f(N-k) + k$$

↓ final solution

↓ unknown term

$$f(N) = f(1) + N - 1$$

$$f(N) = 1 + N - 1$$

$$f(N) = N$$

$$f(1) = 1$$

$$f(N-k) = f(1)$$

$$N - k = 1$$

$$k = N - 1$$

$$f(1) = 1$$

int pow(a, n) {

if (n == 1) { return a }

ha = pow(a, n/2)

if (N & 1 == 1) { # N is odd

return ha \* ha \* a

}

else {

return ha \* ha

}

}

step 1)

TC  $\rightarrow f(N)$

$$f(N) = O(1) + f(N/2)$$

$$f(N) = f(N/2) + 1$$

$$f(N) = f(N/2) + 1$$

$$N \rightarrow N/2$$

$$f(N/2) = f(N/4) + 1$$

$$f(N) = (f(N/4) + 1) + 1 = f(N/4) + 2$$

$$f(N) = f(N/4) + 2$$

$$f(N) = f(N/2) + 1$$

$$N \rightarrow N/4$$

$$f(N/4) = f(N/8) + 1$$

$$f(N) = f\left(\frac{N}{8}\right) + 3$$

$$1) f(N) = f\left(\frac{N}{2}\right) + 1$$

$$2) f(N) = f\left(\frac{N}{4}\right) + 2$$

$$3) f(N) = f\left(\frac{N}{8}\right) + 3$$

$$\log x^{x^n} = n$$

$$f(N) = f\left(\frac{N}{2^k}\right) + k$$

$$f(N) = f(1) + \log_2 N$$

$$f(N) = 1 + \log_2 N$$

$$f\left(\frac{N}{2^k}\right) = f(1)$$

$$\frac{N}{2^k} = 1$$

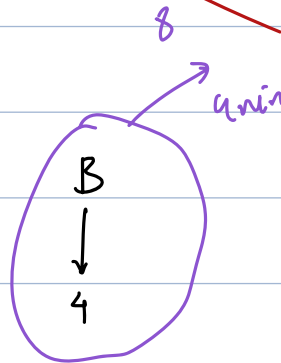
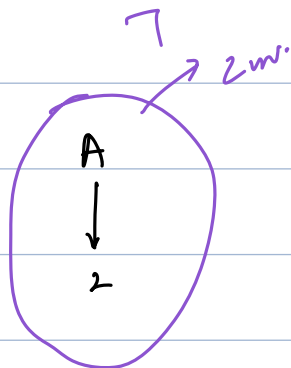
$$N = 2^k$$

Taking  $\log_2$  both sides

$$\log_2 N = \log_2 2^k$$

$$\log_2 N = k$$

$$f(N) = O(\log N)$$



$A \neq B$

4, 6, 8

Time equation:

```
int pow(a, n) {
```

```
    if (n == 1) { return a }
```

```
    ha = pow(a, n/2) * pow(a, n/2)
```

```
    if (N & 1 == 1) { # N is odd
```

```
        | return ha * a
```

```
    }
```

```
    else {
```

```
        | return a
```

```
    }
```

```
}
```

$$f(N) = 1 + f\left(\frac{N}{2}\right) + f\left(\frac{N}{2}\right)$$

$$f(N) = 2f\left(\frac{N}{2}\right) + 1$$

$$f(N) = 2f\left(\frac{N}{2}\right) + 1$$

$N \rightarrow N/2$

$$f\left(\frac{N}{2}\right) = 2f\left(\frac{N}{4}\right) + 1$$

$$f(N) = 2 \left[ 2f\left(\frac{N}{4}\right) + 1 \right] + 1$$

$$f(N) = 4f\left(\frac{N}{4}\right) + 3$$

$$f(N) = 2f\left(\frac{N}{2}\right) + 1$$

$N \rightarrow N/4$

$$f\left(\frac{N}{4}\right) = 2f\left(\frac{N}{8}\right) + 1$$

$$f(N) = 4 \left[ 2f\left(\frac{N}{8}\right) + 1 \right] + 3$$

$$f(N) = 8f\left(\frac{N}{8}\right) + 7$$

$$f(N) = 2f\left(\frac{N}{2}\right) + (2-1)$$

$$f(N) = 4f\left(\frac{N}{4}\right) + (4-1)$$

$$f(N) = 8f\left(\frac{N}{8}\right) + (8-1)$$

$$f(N) = 16f\left(\frac{N}{16}\right) + (16-1)$$

$$f(N) = 2^k f\left(\frac{N}{2^k}\right) + (2^k - 1)$$

$$f(1) = O(1)$$

$$f\left(\frac{N}{2^k}\right) = 1$$

$$N = 2^k$$

$$f(N) = Nf(1) + N - 1$$

$$= 2N - 1$$

$$f(N) = O(N)$$

Hw

$$f(N) = 4f\left(\frac{N}{2}\right) + 1$$

$$f(1) = N$$

$$f(N) = f(N-1) + N$$

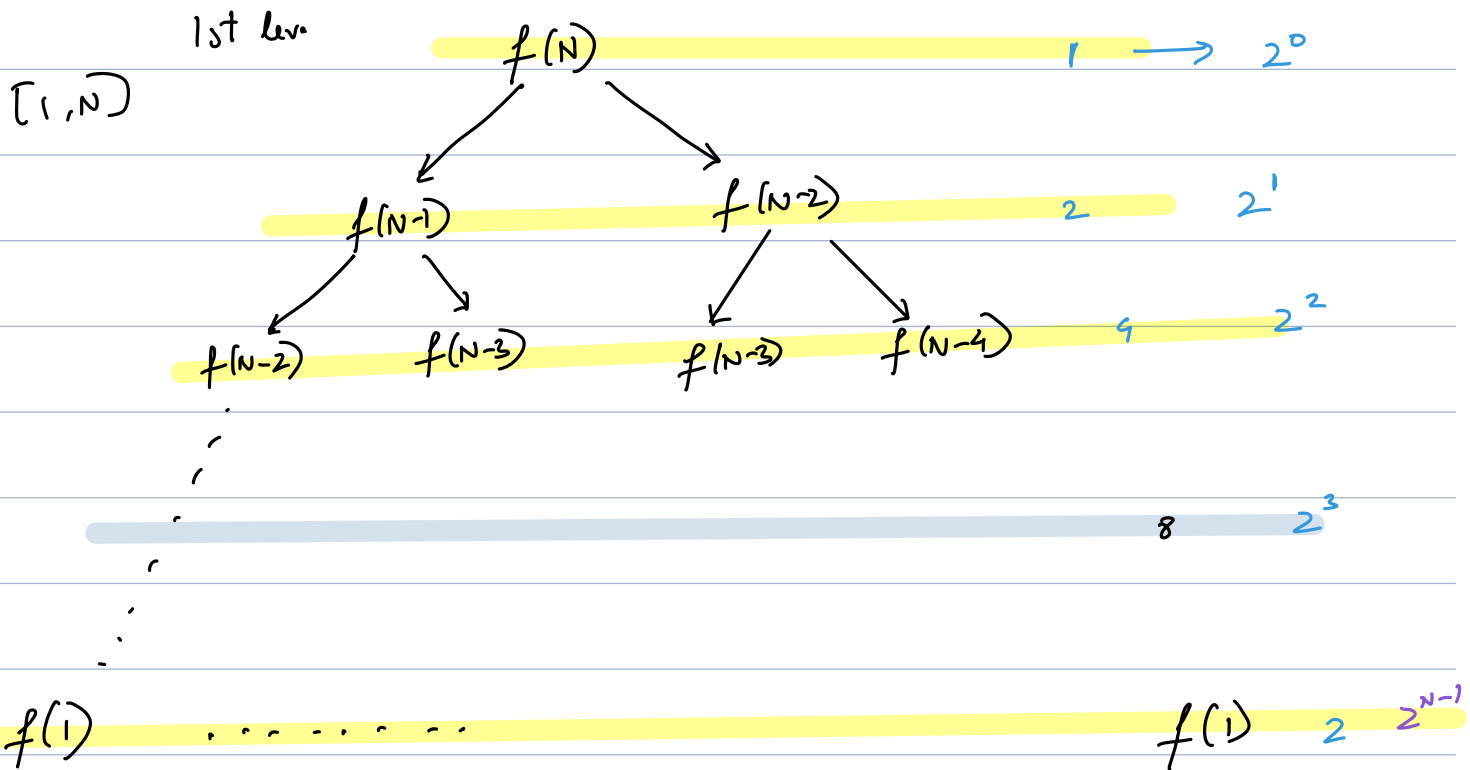
$$f(0) = 1$$

$$f(N) = f(N-1) + f(N-2) + 1$$

Right side more than one recursive term  
substitution fails

```
int fib(N) {  
    if (N <= 1) { return N }  
    return fib(N-1) + fib(N-2)    for (i=0; i < N; i++)  
}
```

Time equation:  $f(N) = f(N-1) + f(N-2) + 1$



N levels

TC: (Total No. of function calls)  $\times$  Time taken for every function call

$$1 + 2 + 4 + 8 \dots 2^{N-1}$$

sum of GP

$$\frac{a(r^n - 1)}{r - 1}$$

$$\boxed{TC: O(2^N)}$$

$$r = 2 \quad n = N$$

$$a = 1$$

Space complexity

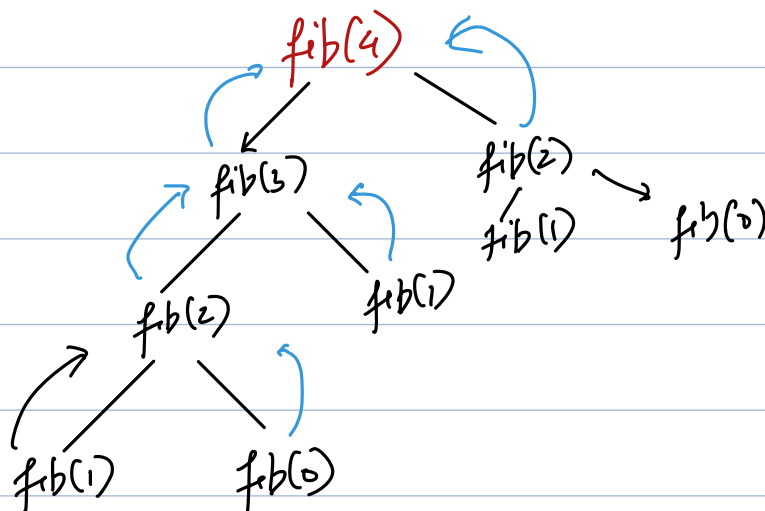
Dry run

Max no. of entries in function call stack

$\text{fib}(4)$

TC:  $O(2^N)$

SC:  $O(2^N)$

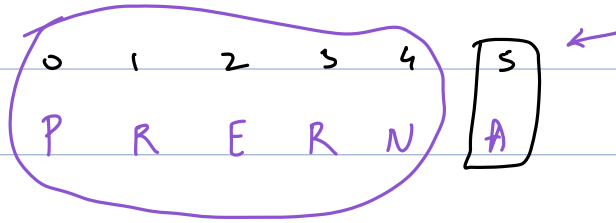


size = 4

SC:  $O(N)$



Print string in reverse order using recursion



Assumption: rev will print the string in rev order [0:i]

```
void rev(s, i) {  
    if (i == -1) { return }  
    print(s[i])  
    rev(s, i-1)  
}
```

$i = \text{len} - 1$