

→ Previous greater element :-

In this problem we have given an array, we have to print closest (position wise) greater element than given element in it's left, if there is no such element, print -1.

This problem is much like previous STOCK SPAN problem, approach is same, we make

we make a stack, with one initial value at arr[0];

traverse array, and for given element, remove all smaller element than it from stack, then top element give required element, and if stack is empty print -1.

push current element in stack.

Code :-

```

void nextgrater(int arr[], int n)
{
    stack<int> s;
    s.push(arr[0]);
    cout << "-1 ";
    for (int i = 1; i < n; i++)
    {
        while (s.empty() || arr[i] >= s.top())
            s.pop();
        int pg = (s.empty()) ? -1 : s.top();
        cout << pg << " ";
        cout << endl;
        s.push(arr[i]);
    }
}
  
```

→ Next greater element :-

In this problem we have given an array, and we have to return an vector in which we have to store next greater element ~~of each~~ for each value of that array.

next greater element means, number greater than given number and nearer on basis of distance in right side of it.



Ex -

arr = {7, 3, 6, 15, 10, 9, 30, 12, 20, 3}

Output -

vec = {8, 15, 15, 30, 30, 30, -1, 20, -1}

-1 for those element, for which, no greater element in right side of it.

It is similar problem to previous
greater element, we just have to
~~not~~ traverse the array from backward
and ~~not~~ instead of printing the element
we store them in vector and rotate
the vector before returning.

If we ~~not~~ ^{not} rotate it, the result will
be wrong, as we are traversing the
array in reverse.

Code :-

vector<int> greater(~~const~~ int arr[], int n)

{

stack<int> s; vector<int> v;

s.push(~~arr[0]~~);

~~const~~ reverse(v); v.push_back(-1);

for(int i = n - 2; i >= 0; i--)

{

while((!s.empty()) && arr[i] >= s.top())

{

s.pop();

}

int ans = (s.empty()) ? -1 : s.top();

v.push_back(ans);

s.push(~~arr[i]~~);

}

reverse(v.begin(), v.end());

return v;

}

→ Delete middle element of a stack :-

approach 1:- Using extra space :-

We have a stack and we have to remove middle element of it then print it.

If we use extra space, we will use another stack in which, we pop top element from given stack and push it in temp stack until we reach middle elem. We pop ~~it~~ that ele. from stack but not push it in temp stack.

After that we ~~can~~ POP element one by one from temp ~~array~~ stack and push ~~it~~ them in given array.

Ex. -

Stack (int) S = { 1, 3, 7, 8, 12 }

Result :-

{ 12 , 8 , 3 , 1 }
↑
Top

Result is not reversed but while printing stack, it prints from top to bottom, so.

Code:-

void stack(Stk &S, int size)

{

Stack <int> t;

int m = ~~size~~ cell (size/2);

// n is element number which we have to

// remove, if size = 5, m = cell (5) = cell(3)

// which is 3, i.e. have to remove 3rd elem.

for (int i = 1; i <= m; i++)

{

if (i == n)

{ S.pop();

} else

int x = S.top();

t.push(x);

S.pop();

}

}

while (!empty())

{

int x = t.top();

S.push(x);

t.pop();

}

}

Approach 2:- without using extra space

In this approach we will make recursive calls to avoid ~~the~~ using extra stack. As we know, recursion stack calls are made so, we will use that.

Code :-

```
Void deleteuntil(stack<int>&S, int n, int cu)
```

{

```
if(cu == n/2)
```

{

```
S.pop();
```

```
return;
```

}

```
int x = S.pop();
```

```
S.pop();
```

```
cu++;
```

```
deleteuntil(S, n, cu);
```

```
S.push(x);
```

3

delete mid
void[↑](stack<int> &s, int size)
{

 delete until(s, size, 0);

}

int main()

{

 stack<int> s;

 s.push(7);

 s.push(10);

 s.push(2);

 s.push(12);

 s.push(5);

 delete mid(s, s.size());

 for(!s.empty()) {

 cout << s.top() << " ";

 s.pop();

}

 return 0;

}

Output:-

5 12 10 7

~~strb~~ = abc
~~str~~ = abc

→ Removing consecutive duplicates from string:-

We have given string, in which we have to ~~remove~~ modify the string such that it will remove consecutive duplicate characters.

Ex. -

string s = "aaaaabbccaaaddd";

After modification s = "abcad";

So, what we do is we make another string say (st) of same size as S and a stack (say t)

We will push str[0] in stack and run a loop from i=1 to i < s.length().

Check every str[i], if it's equal to t.top() if yes continue if no

Add ~~top~~ top of t (t.top()) in second

~~string~~ string str then push that str[i] in stack.

return str or copy it in s.

Code (1) same approach:-

String removeduplicate(string s)

{
String str = " ";

Stack<char> t;
t.push(s[0]);

for (int i = 1; i < s.length(); i++)
{

if (s[i] == t.top())
 continue;

else {

str = str + t.top();

t.push(s[i]);

}

3

str += t.top();

return str;

3

Code (2) same approach:-

String removeduplicate(string s)

{
String str = " ";

Stack<char> t;

str = str + s[0];

```

t.push(s[0]);
for(int i=1; i < s.length(); i++) {
    if(s[i] == t.top())
        continue;
    else {
        st.push(s[i]);
        t.push(s[i]);
    }
}
return st;

```

input:- a a a b b c c d

Output - abcd

→ Remove Consecutive elements :-

We have given an array of integers in which we have also given two integers x and y . When we encounter two x or two y consecutive remove them from array.

Ex -

vector $\langle \text{NA} \rangle = \{2, 1, 2, 2, 1, 5\}$ $x=2, y=1$

Output: vector $v_1 = \{2, 5\}$

Qs, first 2,2 in middle will be removed then v becomes $\{2, 1, 1, 5\}$ then 5,1,3 in the middle will be removed so final vector will be $\{2, 5\}$

Approach:-

We make second vector to store final result and return it, make stack for operations

~~we store~~ $v[0]$ is stack, ~~and~~ and run loop from $i=1$ to $i=v.size;$

If $v[i] == x$ or y then check if $v[i]$ is $= t.top()$ if yes $t.pop()$ if not $t.push(v[i]);$

If $v[i]$ is not x or y then also $t.push(v[i]);$

After that copy stack to vector result, and return that.

We can store values in result vector while iterating, so we don't need to copy stack to vector.

Code 1:-

```
vector<int> remove (vector<int> v, int x, int y) {
    stack<int> t;
    t.push(v[0]);
```

```
for (int i = 1; i < v.size(); i++) {
    if (v[i] == x || v[i] == y)
```

```
        if (v[i] == t.top())
            t.pop();
```

```
    else
```

```
        t.push(v[i]);
```

```
    }
```

```
    t.push(v[i]);
```

```
}
```

```
vector<int> v1;
```

while (!t.empty())

{

int x = t.top();

v1.push_back(x);

t.pop();

}

return v1;

}

Code 2:-

vector<int> remove(vector<int> v, int x, int y) {

vector<int> v1;

stack<int> s;

s.push(v[0]); v1.push_back(vc[0]);

for (int i = 1; i < v.size(); i++) {

if (vc[i] == x || vc[i] == y) {

if (vc[i] == t.top()) {

t.pop();

v1.pop_back();

}

t.push(vc[i]);

t.push_back(vc[i]);

}

else {

t.push(vc[i]);

v1.push(vc[i]);

}

3

return vi;

→ Get minimum element from stack:-

In this problem we have to make changes in push and pop such that after and make another method getmin() such that when we call getmin() it will return minimum element from stack in O(1) time; also pop and push operation also need to work in O(1).

ex:-

Stack<int> S;

S.push(7);

S.push(2);

S.push(5);

S.getmin(); // returns '2'

S.pop()

S.getmin(); // returns '2'

S.push(1);

S.getmin(); // returns '1'.

5
2
7

1
2
7

Approach 1:-

In this approach we will use ~~some~~ extra space in form of stack. We have our main stack and an auxiliary stack. We use main stack for push and pop and use auxiliary array for getmin.

Whenever push check if pushed element is less than or equal to top of stack, push that element also ^{aux} to stack. Else only push element to main stack.

while popping if popped element is greater than top of auxiliary stack only pop from main stack else also pop from aux stack.

When asked for min element return top of auxiliary stack.

Code:-

```
void push(stack<int> &S, stack<int> &Q) {
```

Class stack-example

Stack<int> s;

Stack<int> a;

Public:

```
void push(int x);  
int pop();  
int getmin();
```

}

Void stack-example::push(int x){

if(s.empty()) {

s.push(x);

a.push(x);

}

else if(x <= a.top())

{

s.push(x);

a.push(x);

}

else {

s.push(x);

}

}

Int stack-example :: getmin()

if(s.empty()) return -1;

else return a.top();

}

```
int stackExample :: pop() {
    if (s.empty()) return -1;
    else if (s.top() > q.top())
        return s.top();
```

~~return s.top();~~

```
int x = s.top();
s.pop();
return x;
```

else

```
int x = s.top();
s.pop();
q.pop();
return x;
```

int main()

StackExample obj;

obj.push(3);

obj.push(2);

obj.push(10);

obj.push(1);

~~obj~~

cout << "current min is " << obj.getMin();

obj.pop();

cout << "current min is " << obj.getMin();

return 0;

~~3~~

Output:-

current min is 1
current min is 2

Approach 2:-

In this approach we will not use another auxiliary stack. We use a variable to store current minimum. We modify the main stack such that it will store element which is pushed and previous minimum element.

when getmin() is called return minEle element.

when pushing if stack is empty push 'x' to the stack and assign minEle. $\text{minEle} = x$.

if x is less than minEle, ~~then~~ push $2x - \text{minEle}$ to stack and assign x to minEle.

else

just push 'x' to the stack.

while popping

if stack is empty return -1:

else if ($s.top() \geq minEle$) no change
in ~~minEle~~; return $s.top()$ and
Pop it.

else

Store current minEle in variable say (y).

~~POP > STOP~~

Change minEle as $2 * minEle - s.top$;

Pop s.top;

return y;

Code:-

Class StackExample

Stack<int> s;

int minEle;

public:

int getmin();

if(s.empty()) return -1;

return minEle;

}

void push(int x){

if(s.empty())

s.push(x);

minEle = x;

}

else if($x < \minEle$)
 s.push($2 * x - \minEle$);
 minEle = x ;

{

else

 s.push(x);

{

int pop() {

if(s.empty()) return -1;

else if(s.top() >= minEle) {

int x = s.top();

s.pop();

return x;

{

else

int y = minEle;

 minEle = $2 * \minEle - s.top()$;

s.pop();

return y;

{

{

};

int main()

~~Stack~~ stack example obj;

obj.push(2);

obj.push(12);

obj.push(7);

obj.push(1);

cout <> obj.getmin(); endl;

obj.pop();

cout <> obj.getmin(); endl;

return 0;

3

Output:-

1

2