

→ Priority queue of pairs :-

• max heap on basis of first (base/default).-

• ~~typedef pair<int, int> pd;~~

int main()

• priority queue<pd> pq;~~priority queue<pd> pq;~~

pq.push({3, 23});

pq.push({10, 23});

pq.push({1, 103});

pq.push({5, 13});

for(;; pq.empty())

while(!pq.empty())

(cout << pq.top().first << " "

<< pq.top().second) << endl;

pq.pop();

3

return 0;

3

Outputs -

10 2

5 1

3 2

1 10

• min heap on basis of first :-

Use syntax of min heap.

typedef pair<int, int> pd;

int main()

```
priority_queue<pd, vector<pd>, greater<pd>> pq;  
pq.push({3, 23});  
pq.push({10, 23});  
pq.push({1, 10});  
pq.push({23, 13});
```

while (!pq.empty()) {

```
cout << pq.top().first << " "  
<< pq.top().second << endl;
```

pq.pop();

}

return 0;

3

Output:-

1 10

3 2

5 1

10 2

• max heap on basis of second :-

Here, idea is to use myamp structure with operator overloading operator to order pairs by it's second element in heap.

typedef pair<int, int> P;

struct myamp

bool operator()(pair<int, int> const& a,
pair<int, int> const& b)

{

return a.second < b.second;

}

};

int main()

priority_queue<P, vector<P>, myamp> pq;

Pq.push({3, 2});

Pq.push({10, 3});

Pq.push({1, 10});

Pq.push({5, 1});

```

while(!pq.empty())
    cout << pq.top().first << " "
    << pq.top().second << endl;
    pq.pop();
}

```

return 0;

}

Output:-

1 10

10 2

3 2

5 1

min heap on basis of second :-

We will use same idea as previous one but will modify comparison structure and overloading method.

typedef pair<int, int> p;

struct mycomp

bool operator()(pair<int, int> const &a,

pair<int, int> const &b)

{

return a.second > b.second;

}

}

int main()

Priority queue(p, vector<P>, myComp> pq;
pq.push({3, 2});
pq.push({10, 2});
pq.push({1, 10});
pq.push({5, 1});

while (!pq.empty()) {

cout << pq.top().first << " "
<< pq.top().second << endl;

}

return 0;

}

Output:-

5 1

3 2

10 2

1 10

→ priority queue of structure:-

we will make priority queue of structure and make priority queue on basis of its any element.

Struct person

int age;

float ht;

person (int a, float b)

{

age = a;

ht = b;

3

};

Struct myComp

bool operator<(person const & p1, person const & p2)

{

return p1.age < p2.age;

3

};

int main()

priority_queue<person, vector<person>, myComp> pq;

pq.push({18, 160});

pq.push({16, 165});

pq.push({20, 200});

```
while(!pq.empty()) {
    cout << pq.top().age << " ";
    cout << pq.top().wt << endl;
    pq.pop();
}
```

Output :-

20	200
16	165
18	160

It is max heap.

→ Sorting array using priority queue -

In daily use we must use `sort()` function to sort ~~the~~ any array, but here, for exercise we will use priority queue method -

```
void sortarr(int arr[], int n) {
    priority_queue<int> pq(arr, arr+n);
    for(int i=n-1; i>=0; i--) {
        arr[i] = pq.top();
        pq.pop();
    }
}
```

```
int main()
{
```

```
    int arr = { 2, 7, 3, 5, 12 };
```

```
    sortarr();

```

```
    int n = sizeof(arr) / sizeof(arr[0]);
```

```
    sortarr(arr, n);
```

```
    return 0;
}
```

→ K largest element in an array :-

In this problem we have given an array and we have print first 'K' largest element.

Ex:-

arr[] = { 10, 5, 30, 2, 13, 9, 14 };

Output:- 10, 13, 30

NOTE:- Order of output is not considered here, it may change.

Method 1:-

In this method it is naive method in which we first sort the array and then print 'K' element of that array.

T.C. will be $O(n \log n + k)$.

$O(n \log n)$ for sorting the array
 $O(k)$ for printing 'K' elements.

Code -

```
void printK(int arr[], int n, int k) {
```

```
    sort(arr, arr+n);
```

```
    for (int i = n - k; i >= 0; i--)
```

```
        cout << arr[i] << " ";
```

```
    }
```

Method 2:- Using max heap-

In this approach we will ~~not~~ make max-heap using priority queue of given array and print first 'K' elements of it.

T.C. will be $O(n + k \log n)$

$O(n)$ for make max heap using PQ

$O(k \log n)$ as K times popping

out of PQ and printing that element.

Code -

```
void printk(int arr[], int n, int k){  
    priority_queue<int> PQ(arr, arr+n);  
  
    while(k){  
        for(int i=0; i < k; i++)  
            cout << PQ.top() << " ";  
        PQ.pop();  
        k--;  
    }  
}
```

Method 3:- Using min heap -

It is more efficient than max-heap.

In this approach we will follow this algo. -

- make min-heap of first 'k' element of array.
- compare element from 'i=k' to 'i=n-1' to top of PQ.

if $arr[i] > pq.top()$ • remove top() and insert $arr[i]$ in pq.

else leave pq as it is.

- After loop ends 'K' element in pq will be largest 'K' element of array.

NOTE:- Beauty of this approach is that, in this approach at any time pq gives 'K' largest element at that point any given point in array.

Here, we have used min-heap as we have to find 'K' largest element if we have to find 'K' smallest element, we can use max-heap with similar algorithm.

T.C of this approach is

$$O(K + (n-K) \log K)$$

$O(K)$ due to making ~~resizing~~ min-heap using pq.

$O((m-k) \log k)$ due to loop on P_2

Code -

```
void printC(int arr[], int n, int k){
```

```
PriorityQueue<int> pq;
PQ.push(arr[0], 0);
```

```
for (int i = k; i < n; i++) {
```

```
if (arr[i] > PQ.top()) {
```

```
PQ.pop();
```

```
PQ.push(arr[i]);
```

```
}
```

```
while (!PQ.empty() == false) {
```

```
cout << PQ.top() << " "
```

```
PQ.pop();
```

```
}
```

```
3
```

→ Purchasing maximum items :-

We have given an array with values of items, and a variable containing total money we have. We have to find the max. no of items we can buy using this money.

Ex -

$$\text{arr}[] = \{1, 30, 7, 5, 12, 39\};$$

$$\text{money} = 20$$

$$\text{result} = 3 \quad (1, 5 \text{ and } 7)$$

method 1:-

In this method we sort the array and compare ~~see~~ ~~total~~ money we have to the ~~i~~th ~~item~~ element of array.

If $(\text{arr}[i] < \text{money})$ we can buy first item $\Rightarrow \text{result}++$ and
 $\text{money} = \text{arr}[i];$

else,

return the result.

Code -

```
int purchaseMax(int arr[], int n, int m){  
    sort(arr, arr+n);  
    int result = 0;  
    for (int i=0; i<n; i++) {  
        if (arr[i] <= m) {  
            result++;  
            m = m - arr[i];  
        }  
        else  
            break;  
    }  
    return result;  
}
```

T.C. will be $O(n + n \log n)$ but we don't need extra space.

method 2:-

In this method we use min-heap, we make min-heap out of array and compare ~~to~~ $PQ.\text{top}()$ with money we have.

```
if(money > PQ.top())
    money = PQ.top();
    PQ.pop();
    res++;
```

~~TIME - O(n log n)~~

T.C. of this app. is

$O(n + \cancel{n} \log n)$

As you can see T.C. looks similar but this T.C. occurs very rarely in min-heap, generally it is

$O(n + \text{res} * \log n)$ and
 $\text{res} \ll n$ in many cases.

Code -

```
int purchaseMax(int arr[], int n, int m) {
```

Priority Queue<int, vector<int>, greater<int>>
pq(arr, arr+n);

```
int res = 0;
```

```
while (m >= 0 && pq.empty() == false &&  
      pq.top() <= m)
```

{

```
m -= pq.top();
```

```
pq.pop();
```

```
res += j;
```

3

```
return res;
```

3

→ Merge k sorted array -

We have given ' k ' sorted array each of size ' n ' in 2-D array form.

$arr[k][n]$:

We have to print the sorted merge out.

Ex -

$k=3, n=4$

$arr[3][4] = \{ \{ 1, 3, 5, 7 \},$

$\{ 2, 4, 6, 8 \},$

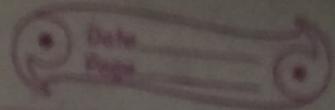
$\{ 0, 9, 10, 11 \} \}$

Output - 0 1 2 3 4 5 6 7 8 9 10 11

Approach 1 -

It is a naive approach we store all input sorted array in one big output array one by one and then sort that output array.

but T.C and S.C will be high.



Approach 2:-

We will use merge sort approach.
We merge those ~~K arrays~~ in half by
first dividing them in half.

1

When devide a given 'K' array ~~rec~~
such that

if only one array is left, return it.

if two array are left merge them

if more than two array are left
call the devide function recursively
one each half.

Code -

```
Void merge_array (int arr[], int sub[], int n,  
                  int m, int out[]){  
    int i=0, j=0, k=0;  
    while (i < n & j < m)  
    {  
        if (arr[i] <= arr[j])  
            out[k++] = arr[i++];  
        else  
            out[k++] = arr[j++];  
    }  
}
```

~~①~~ while ($i < m_1$)

$out[i++] = arr1[i++]$

while ($j < m_2$)

$out[i++] = arr2[j++]$

② 3

3

Void merge_array (int arr1[m], int i, int j,
int arr2[n], int n)

if ($i == j$) { // one array left

for (int p = 0; p < n; p++)

$output[p] = arr1[i][p];$

return;

3

if ($j - i + 1 == 1$) { // two arrays left

merge_array (arr1, arr2, n, n, output)

return;

3

int mid = $(i+j)/2$

~~③ Local part + mid~~

int Out1 [$m * (mid - i + 1)$], Out2 [$n * (j - mid)$]

// two output arrays to store output from array 1

// merge array function,

~~merge k arrays~~

merge-k-array(arr, i, mid, out1);

merge-k-array(arr, mid+1, j, out2);

merge-array(out1, out2, n*(mid-i+1), n*(j-mid))
Output);

}

~~driver code~~

int main()

int arr[2][n] = {{2, 6, 12, 34},
 {1, 9, 20, 1000},
 {23, 34, 90, 2000}};

int x = sizeof(arr)/sizeof(arr[0]);

int output[n*x];

merge-k-array(arr, 0, x-1, output);

3

Approach 3:- $9t$ is best approach to merge k sorted arrays. We use min-heap priority queue to merge k -sorted arrays.

We first make a priority queue of pairs of pairs.

```
typedef pair<int, pair<int, int>> PPI;  
priority_queue<PPI, vector<PPI>, greater<PPI>> PQ;
```

We do so, because we want to store not only value but also in which array that value is and at what index.

SO.

Let

$PPI \text{ curr} = PQ \cdot \text{top}();$ then

$\text{curr}.first \rightarrow \text{Value}$

$\text{curr}.second.first \rightarrow i_n$ which array that value is from.

$\text{curr}.second.second \rightarrow \text{index of element in array.}$

now,

- Create an output array.
- Create a min-heap of size k and insert 1st element in all the array into heap.
- Repeat following step while priority queue is not empty.
 - Remove top/minimum element from heap and store it in output array.
 - Insert next element from the array from which the element is extracted. If the array doesn't have any more elements, then do nothing.

Code -

```
#include<iostream>
// #include<utility>
#include<queue>
#include<vector>
```

```
typedef pair<int, pair<int, int>> ppi;
```

vector<int> merge_k_array(vector<vector<int>> arr) {

vector<int> output;

priority queue<ppi, vector<ppi> greaterppi> pq;

for (int i = 0; i < arr.size(); i++) {

pq.push({arr[i][0], i, 0});

}

while (!pq.empty()) == false) {

ppi cur = pq.top();

pq.pop();

Output.push_back(cur.first);

int j = cur.second.first;

int k = cur.second.second;

if (j + 1 < arr[i].size())

pq.push({arr[i][j + 1], i, j + 1});

}

return Output;

3

int main()

vector<vector<int>> arr {{2, 6, 12},
{1, 9, 30},

{23, 34, 100}, {111, 222}}

vector<int> output = merge-k-array(arr);

for (auto x : output)

cout << x << " "

return 0;