CSE-3103: Microprocessor and Microcontroller

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Unconditional Jumps

Jump (JMP) allows programmer →
- skip sections of program,
- branch to any part of memory for next instruction.

3 types of unconditional jump instructions →
- (i) Short jump →
  - jumps or branches to memory locations within +127 and -128 bytes.
- (ii) Near jump →
  - branch or jump within ±32K bytes or
  - anywhere in current code segment.
- (iii) Far jump:
  - jump to any memory location within real memory system.

Short and near jumps = intrasegment jumps,
Far jumps = intersegment jumps.

Unconditional Jumps

Short Jump →

Short jumps = relative jumps.

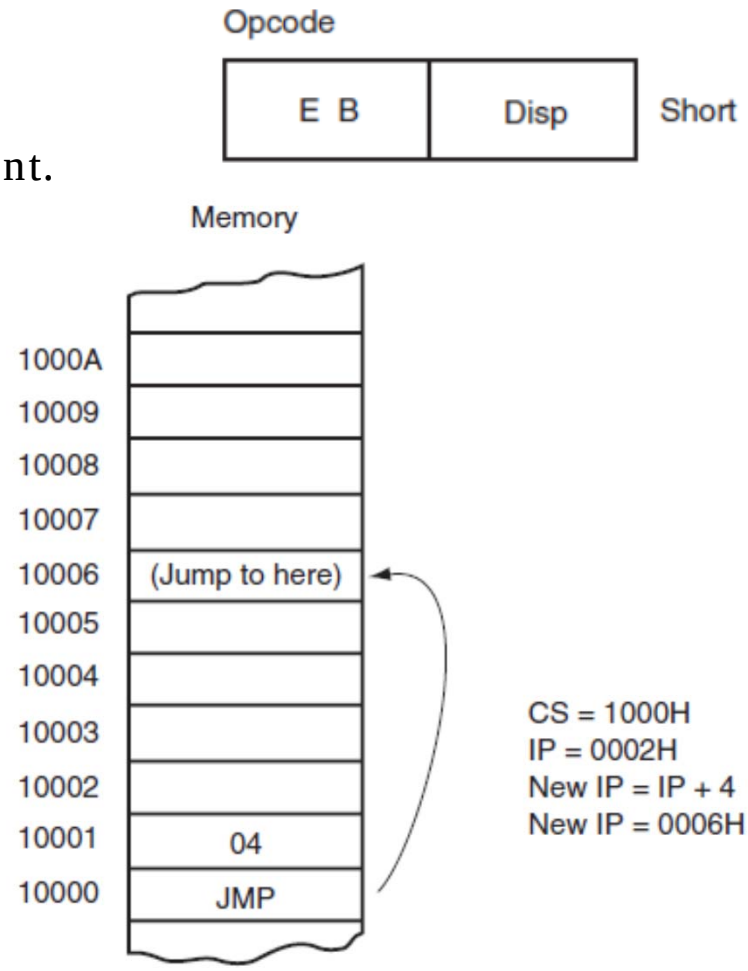
Can be moved to any location in current code segment.

Jump address is not stored with opcode.

Distance or displacement follows opcode.

Displacement value = between +127 and -128.

Jump address = displacement is sign extended + IP.



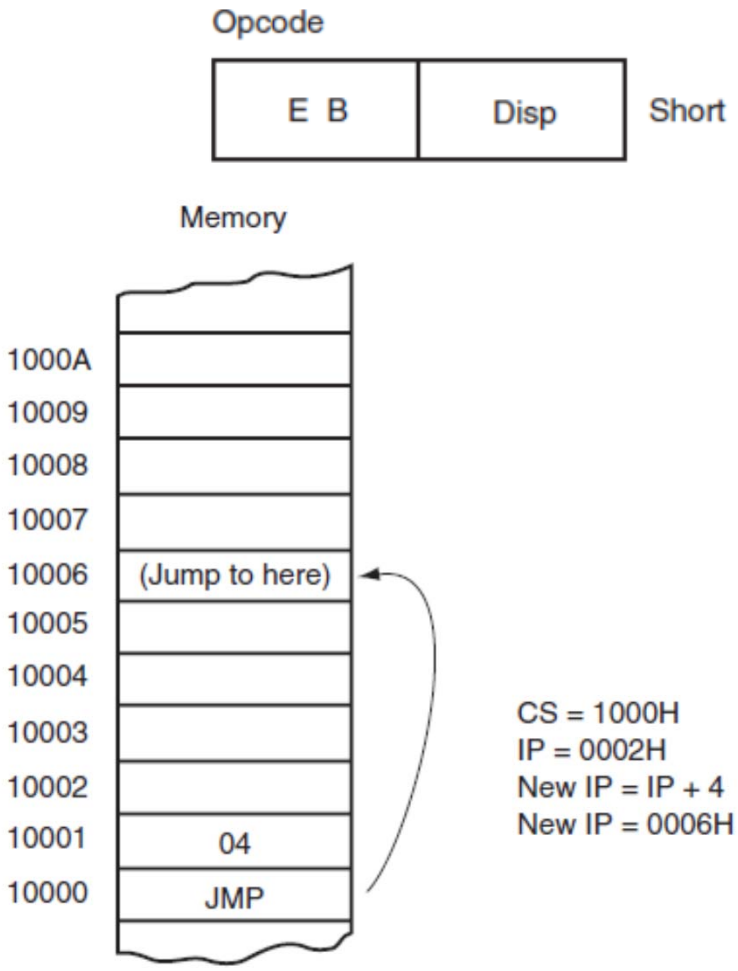
Unconditional Jumps

Short Jump →

```
                XOR     BX, BX
START:          MOV     AX, 1
                ADD     AX, BX
                JMP     SHORT  NEXT

; <skipped memory locations>

NEXT:           MOV     BX, AX
                JMP     START
```



Unconditional Jumps

Near Jump →

passes control to instruction in current code segment located within ±32K bytes.

3-byte instruction =

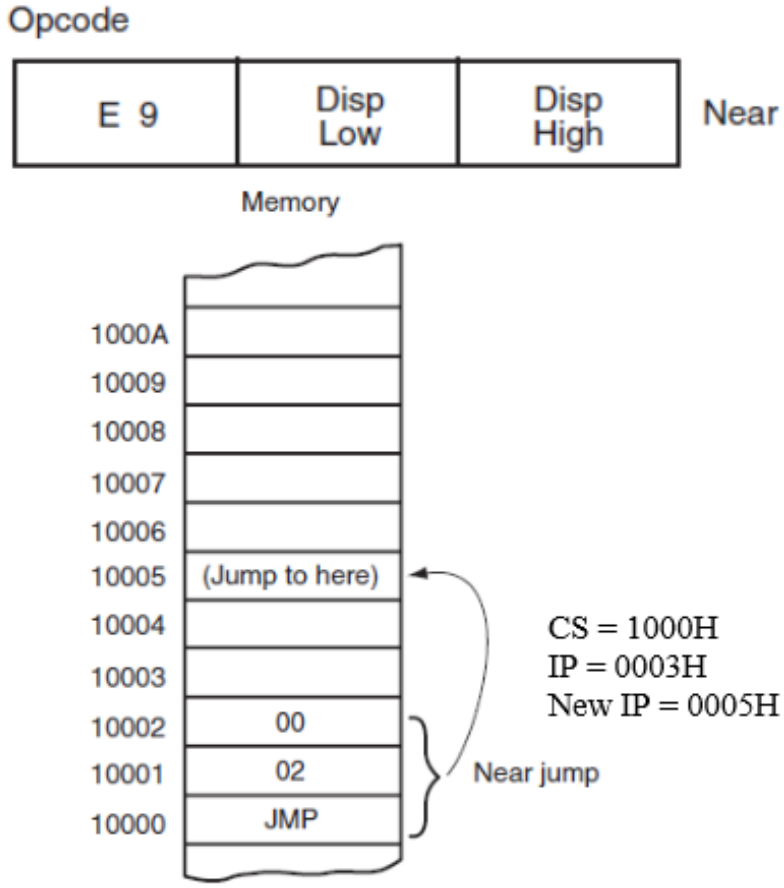
opcode + signed 16-bit displacement.

signed displacement + IP = jump address.

can jump to any memory location within current code segment.

relocatable and relative jump →

if code segment moves to new location, distance between jump instruction and operand address remains same.



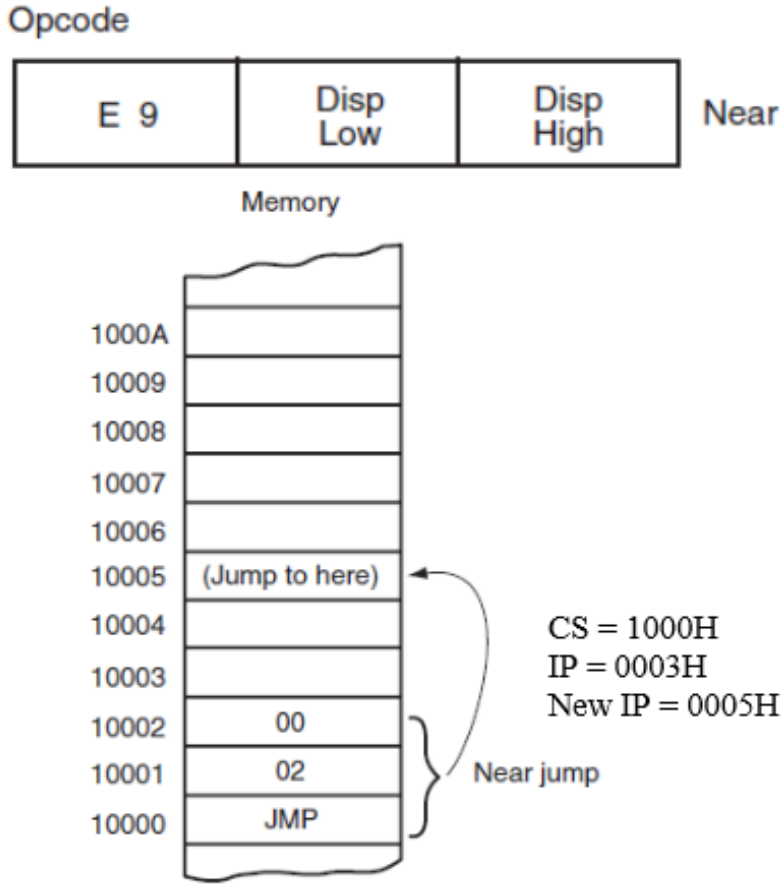
Unconditional Jumps

Near Jump →

```
                XOR     BX, BX
START:          MOV     AX, 1
                ADD     AX, BX
                JMP     NEXT

; <skipped memory locations>

NEXT:           MOV     BX, AX
                JMP     START
```



Unconditional Jumps

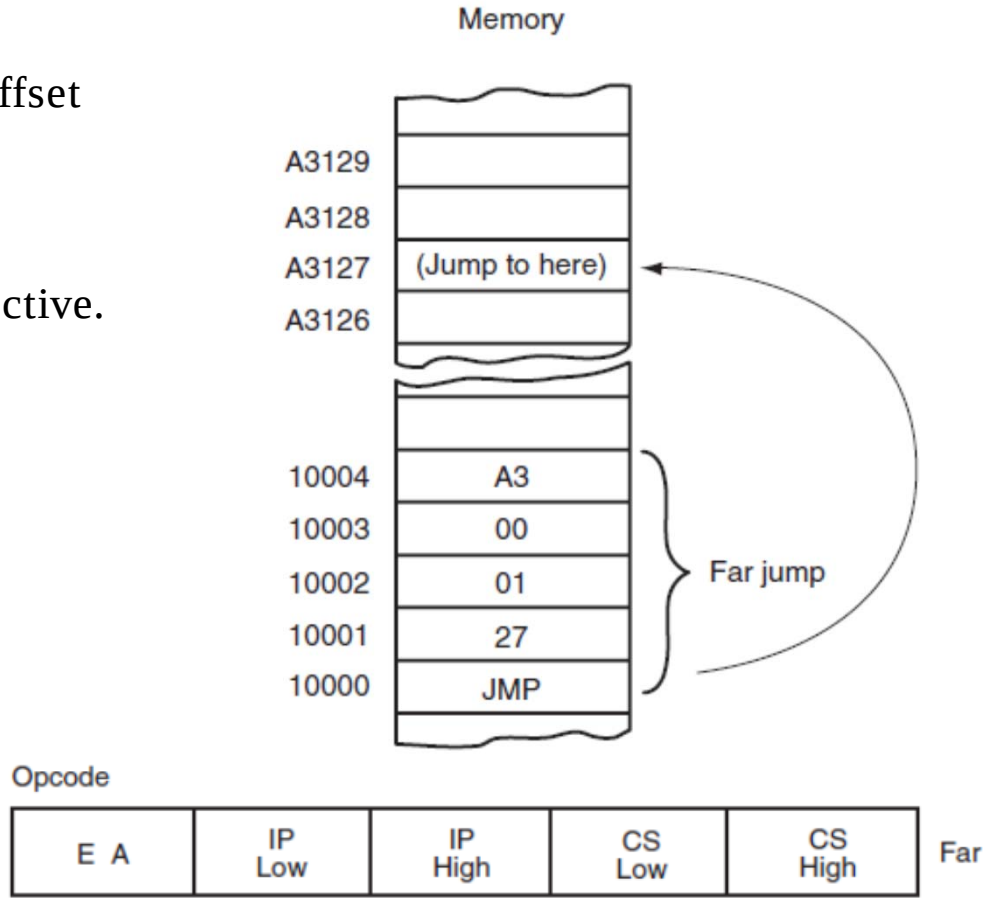
Far Jump →

instruction obtains new segment and offset address to accomplish jump.
bytes 2 and 3 = new offset address;
bytes 4 and 5 = new segment address.
instruction appears with FAR PTR directive.
obtained by defining label as far label.

```
            EXTRN   UP:FAR
            XOR     BX, BX
START:      ADD     AX, 1
            JMP     NEXT

; <skipped memory locations>

NEXT:       MOV     BX, AX
            JMP     FAR  PTR  START
            JMP     UP
```



Conditional Jumps

Conditional jump instructions = short jumps.
Range of jump = +127 bytes to -128 bytes.

Conditional jump instructions test flag bits →
sign (S), zero (Z), carry (C), parity (P), overflow (O).

Test condition is true =
branch to label associated with jump occurs.
Test condition is false =
next sequential step in program executes.

2 sets of conditional jump instructions for comparison.
When signed numbers are compared,
JG, JL, JGE, JLE, JE, JNE instructions.

When unsigned numbers are compared,
JA, JB, JAE, JBE, JE, JNE instructions.

Conditional Jumps

Assembly Language	Tested Condition	Operation
JA	Z=0 and C=0	Jump if above
JAE	C=0	Jump if above or equal
JB	C=1	Jump if below
JBE	Z=1 OR C=1	Jump if below or equal
JC	C=1	Jump if carry
JE or JZ	Z=1	Jump if equal or jump if zero
JG	Z=0 and S=0	Jump if greater than
JGE	S=0	Jump if greater than or equal
JL	S != 0	Jump if less than
JLE	Z=1 or S != 0	Jump if less than or equal
JNC	C=0	Jump if no carry
JNE or JNZ	Z=0	Jump if not equal or jump if not zero
JNO	O=0	Jump if no overflow
JNS	S=0	Jump if no sign (positive)
JNP or JPO	P=0	Jump if no parity or jump if parity odd
JO	O=1	Jump if overflow
JP or JPE	P=1	Jump if parity or jump if parity even
JS	S=1	Jump if sign (negative)
JCXZ	CX=0	Jump if CX is zero
JECXZ	ECX=0	Jump if ECX equals zero
JRCXZ	RCX=0	Jump if RCX equals zero (64-bit mode)

Conditional Jumps

```
; SCASB searches a table of 100 bytes for 0AH
; address of TABLE is assumed to be in ES:DI

        MOV     CX, 100             ; load counter
        MOV     AL, 0AH             ; load AL with 0AH
        CLD                         ; auto-increment
        REPNE  SCASB                ; search for 0AH
                                    ; repeat SCASB until CX=0
        JCXZ    NOT_FOUND           ; if not found
        STC                         ; set carry if found
NOT_FOUND
```

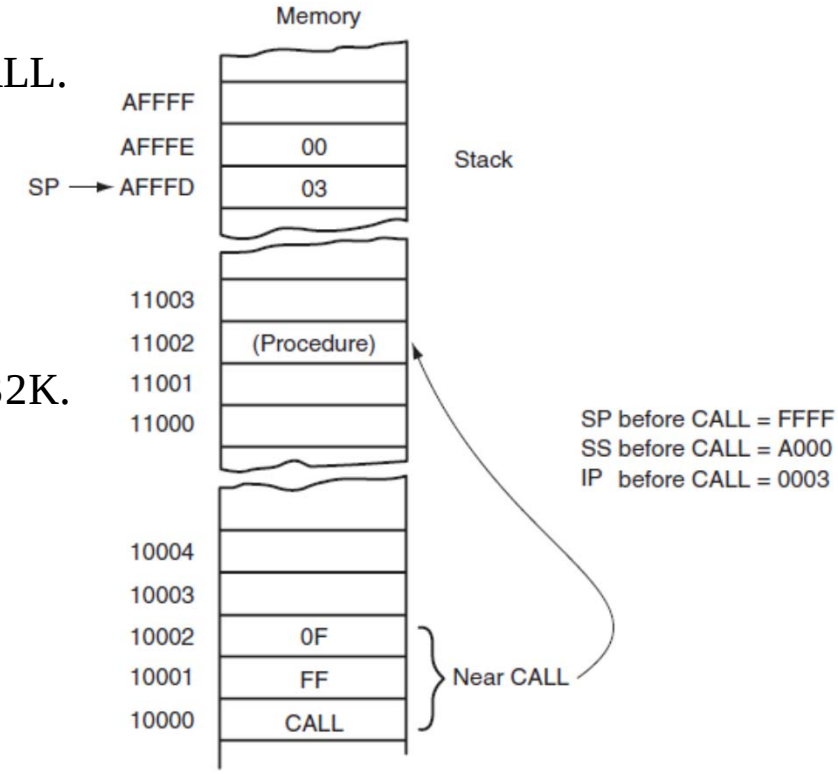
CALL Instruction

CALL transfers flow of program to procedure.
CALL saves return address on stack.
Return address = instruction that immediately follows CALL.
RET = returns control to Return address.

Near CALL →

Near CALL = 3 bytes long;
1st byte = opcode,
2nd + 3rd bytes = displacement, or distance of ±32K.

When near CALL executes,
i) pushes offset address of next instruction onto stack.
ii) IP + displacement →
transfer control to procedure.
iii) RET →
pops offset address from stack (saved in step-i).
program control passes to instruction following CALL.



CALL Instruction

Far CALL →
Call procedure stored in any memory location.
Far CALL = 5-byte instruction =
Opcode +
Bytes 2 and 3 = new contents of IP +
Bytes 4 and 5 = new contents for CS.

Far CALL pushes contents of both IP and CS on stack.
It then jumps to address indicated by bytes 2 through 5.