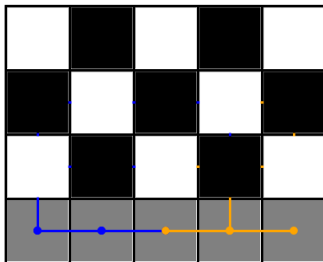


# The Floortile Problem

DD2380 Artificial Intelligence

Group Project



October 17th, 2017

KTH Royal Institute of Technology

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# Plan

Introduction

PDDL

Optimisation

Discussion

Conclusion

# Introduction

## International Planning Competition

### Floortile domain

N robots

Deterministic environment

Paint\_up and Paint\_down restriction

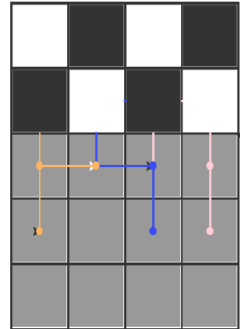
Only paint with Black and White

Robots can not stand on a painted tile

## How we solved the problem?

PDDL

Optimisation



# Plan

Introduction

**PDDL**

Optimisation

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Conclusion



# Formulation

EFFECT:  $\text{RobotAt}(r_1, t_y) \neg \text{RobotAt}(r_1, t_x) \wedge \text{Clear}(t_x) \neg \text{Clear}(t_y) \neg \text{Painted}(t_y)$

Action(Down)

PRECOND:  $\text{RobotAt}(r_1, t_x) \wedge \text{Down}(t_y, t_x) \wedge \text{Clear}(t_y)$

EFFECT:  $\text{RobotAt}(r_1, t_y) \neg \text{RobotAt}(r_1, t_x) \wedge \text{Clear}(t_x) \neg \text{Clear}(t_y) \neg \text{Painted}(t_y)$

Action(Right)

PRECOND:  $\text{RobotAt}(2_1, t_x) \wedge \text{Right}(t_y, t_x) \wedge \text{Clear}(t_y)$

EFFECT:  $\text{RobotAt}(r_1, t_y) \neg \text{RobotAt}(r_1, t_x) \wedge \text{Clear}(t_x) \neg \text{Clear}(t_y) \text{notPainted}(t_y)$

Action(Left)

PRECOND:  $\text{RobotAt}(r_1, t_x) \wedge \text{Left}(t_y, t_x) \wedge \text{Clear}(t_y)$

EFFECT:  $\text{RobotAt}(r_1, t_y) \neg \text{RobotAt}(r_1, t_x) \wedge \text{Clear}(t_x) \neg \text{Clear}(t_y) \neg \text{Painted}(t_y)$

# Interpretation

1	0	3
0	0	2
3	0	1
0	0	0

FIGURE – Robot representation

Value	Use
0	The Cell is clear
1	Cell has been painted white
2	Cell has been painted black
3	Robot is in top of the cell

TABLE – Meaning of the values in the matrix

# Interpretation

1	1	2	100	100
---	---	---	-----	-----

FIGURE – Board representation

Position	Use
1	X coordinate
2	Y coordinate
3	Current paint
4	Remaining black paint
5	Remaining white paint

TABLE – Meaning of the values in the vector



# Implementation

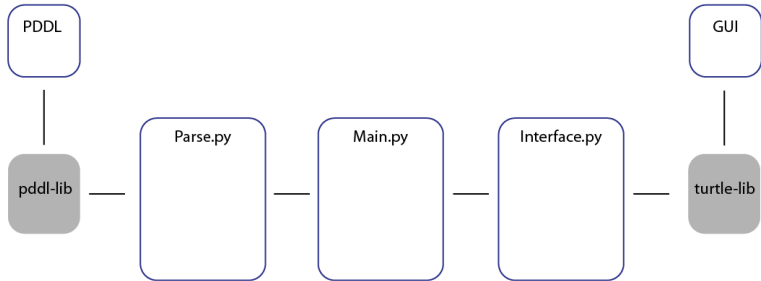


FIGURE – Python code schematic

# Implementation

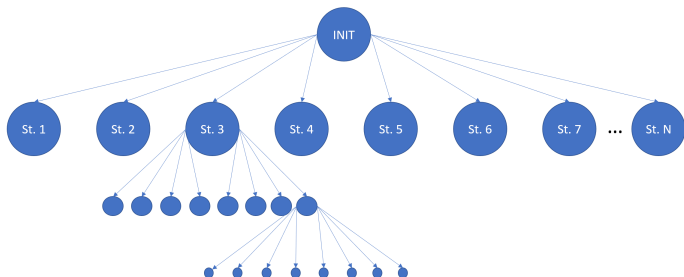


FIGURE – States traversal

# Implementation

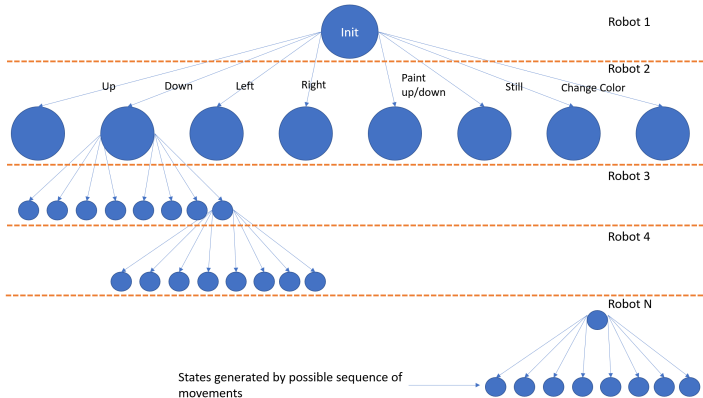


FIGURE – Obtaining possible states

# Implementation

---

**Algorithm 1** Planning algorithm

---

```
1: procedure PLANNING(ROBOTS,INITIALSTATE,TARGETSTATE)
2:   state its defined as [robots,initialState,sequenceMoves]
3:   q  $\leftarrow$  queue of states
4:   memory  $\leftarrow$  dictionary of states
5:   q  $\leftarrow$  push [robots,initialState,[] ]
6:   while:
7:     if q is empty then
8:       break
9:     current  $\leftarrow$  top(q)
10:    pop(q)
11:    if memory[current] == true then
12:      continue
13:    if current[1] == targetState then
14:      return current
15:    memory[current] = true
16:    newPossibleStates  $\leftarrow$  getSequence(current[0], 0, current[1], targetState).
17:    For possible in newPossibleStates:
18:      q  $\leftarrow$  push [possible[0], possible[1], current[2].append(possible[2])
19:    goto For
20:  goto while.
21:  return null.
```

---

FIGURE – Planning algorithm

# Implementation

---

**Algorithm 2** getSequence algorithm

---

```
1: procedure GETSEQUENCE(ROBOTS,INDEX,STATE,TARGETSTATE,MOVEMENTS)
2:   if index == len(robots) then
3:     resStates.append([robots,index,movements])
4:     return
5:   nextMovements  $\leftarrow$  getPossiblesFOC(robot[index],state,targetState)
6:   For next in nextMovements:
7:     getPossibles(robots,index+1,next[0],targetState,movements.append(next[1])
8:   goto For
9:   if index == 0 then
10:    return resStates
```

---

FIGURE – getSequence

# Implementation

*(painted tile(3,1) black) < (painted tile(2,1) white) < (painted tile(1,1) black);*  
*(painted tile(3,2) white) < (painted tile(2,2) black) < (painted tile(1,2) white);*  
*(painted tile(3,3) black) < (painted tile(2,3) white) < (painted tile(1,3) black).*

## FIGURE – Forced Ordering Constraints

**Algorithm 3** getPossiblesFOC algorithm

---

```
1: procedure GETPOSSIBLES(ROBOT, STATE, TARGET STATE)
2:   res = []
3:   movements = {up, down, left, right}
4:   For mov in movements:
5:     aux = tryMovement(mov, robot, state)
6:     if aux == NULL then
7:       res.append(aux)
8:   goto For
9:   if tryMovement(up, robot, state) and paintedColumn(robot, state) and state[robot[0:1]] == robot[2] then
10:    res.append(paintUp(robot, state))
11:   if tryMovement(down, robot, state) and paintedColumn(robot, state) and state[robot[0:1]] == robot[2] then
12:    res.append(paintDown(robot, state))
13:   res.append(changeColor(robot, state))
14:   res.append(still(robot, state))
15:   return res
```

---

## FIGURE – getPossiblesFOC

# Demo

Demonstration

# Case studies

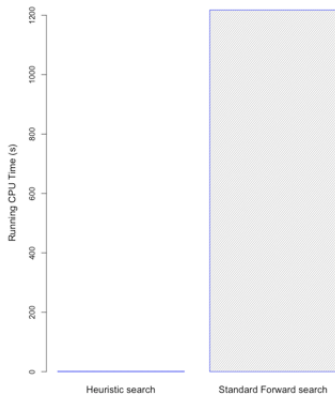
Case 0 : Algorithms	<ul style="list-style-type: none"> <li>Algorithm without using FOC</li> <li>Algorithm using FOC</li> </ul>
Case 1 : Number of robots	<ul style="list-style-type: none"> <li>2 robots</li> <li>3 robots</li> <li>4 robots</li> </ul>
Case 2: Pattern	<ul style="list-style-type: none"> <li>Target: all tiles painted</li> <li>Target: some tiles painted</li> </ul>
Case 3: Size	<ul style="list-style-type: none"> <li>(3,3) (3,4) (5,4) (5,5) (5,6)</li> <li>(4,3) (4,4) (4,5) (6,5) (6,6)</li> </ul>

FIGURE – Different case studies in PDDL

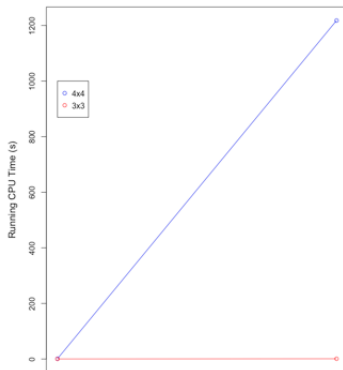


# Case 0

Case 0a: Standard versus Heuristic

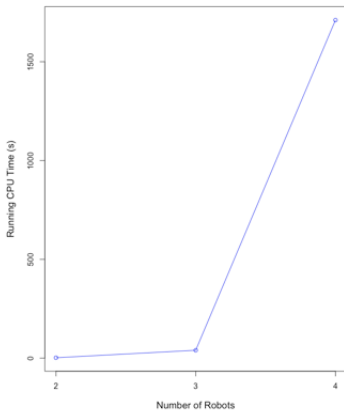


Case 0b: Different Search: Heuristic vs Forward

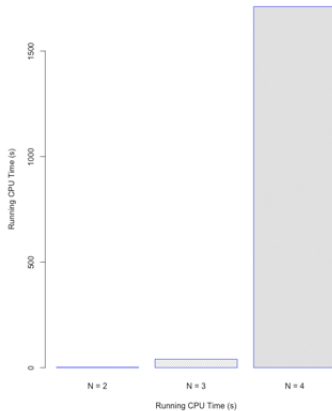


# Case 1

**Case 1a: Different Number of Robots**

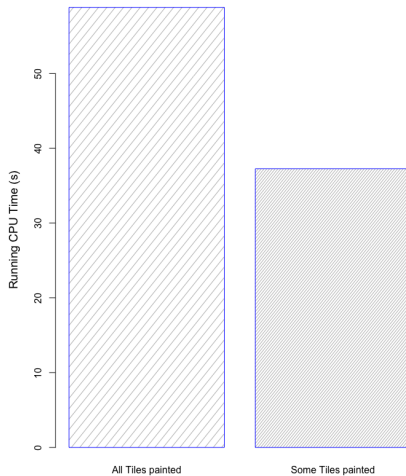


**Case 1b: Different Number of Robots**



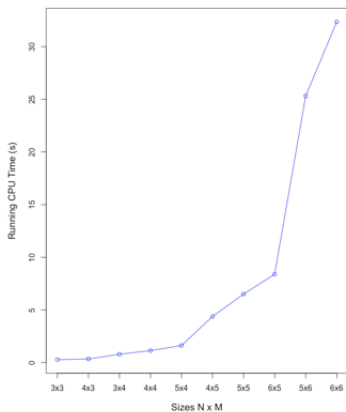
# Case 2

**Case 2: Different Patterns**

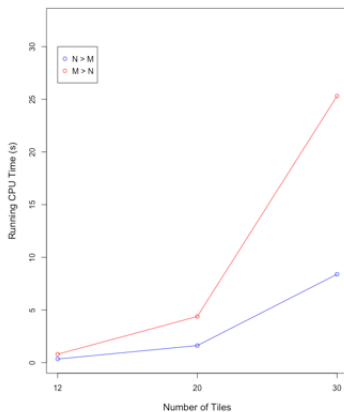


# Case 3

Case 3a: Different Sizes



Case 3b: Different Dimensions



# Plan

Introduction

PDDL

**Optimisation**

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# Formulation

- $t^*$  the time needed for complete board
- $\text{cell}(i, j, t)$  the state of cell  $(i, j)$  at time  $t$ . 0 = not painted, 1 = painted
- $\text{painting}(i, j, t, r) = 1$  if robot  $r$  is painting cell  $(i, j)$  at time  $t$
- State of the robots
  - ▶  $y(t, r)$  the vertical position of robot  $r$  at time  $t$
  - ▶  $x(t, r)$  the horizontal position of robot  $r$  at time  $t$
  - ▶  $\text{color}(t, r)$  the current color of robot  $r$  at time  $t$ . 1 = white, 0 = black
  - ▶  $\text{stock0}(t, r)$  the current stock of black paint of robot  $r$  at time  $t$
  - ▶  $\text{stock1}(t, r)$  the current stock of white paint of robot  $r$  at time  $t$
- Actions of the robot
  - ▶  $\text{paint}(t, r) = 1$  if robot  $r$  is painting at time  $t$
  - ▶  $\text{move}(t, r) = 1$  if robot  $r$  is moving at time  $t$
  - ▶  $\text{switch}(t, r) = 1$  if robot  $r$  is switching color at time  $t$

# Formulation

min  $t^*$

such that  $\sum_{i,j} \text{cell}(i,j,t^*) = nm$

$\text{paint}(t,r) + \text{move}(t,r) + \text{switch}(t,r) \leq 1 \quad \forall t,r$

$\text{cell}(x(t,r),y(t,r),t) = 0 \quad \forall t,r$

$x(t,r) = x(t,r') \Rightarrow y(t,r) \neq y(t,r') \quad \forall r,r' \neq r$

$y(t,r) = y(t,r') \Rightarrow x(t,r) \neq x(t,r') \quad \forall r,r' \neq r$

$\text{paint}(t,r) = 1 \Rightarrow \text{painting}(x(t,r),y(t,r) + 1,t,r) \\ + \text{painting}(x(t,r),y(t,r) - 1,t,r) = 1 \quad \forall t,r$

$\text{paint}(t,r) = 1 \Rightarrow \sum_{i,j} \text{painting}(i,j,t,r) = 1 \quad \forall t,r$

$\text{paint}(t,r) = 0 \Rightarrow \text{painting}(i,j,t,r) = 0 \quad \forall i,j,t,r$

$\text{cell}(i,j,t) = 0 \text{ and } \sum_r \text{painting}(i,j,t,r) \geq 1 \Rightarrow \text{cell}(i,j,t+1) = 1 \quad \forall i,j,t$

$\text{cell}(i,j,t) = 1 \Rightarrow \text{cell}(i,j,t+1) = 1 \quad \forall i,j,t$

Complete board

One action at a time

Not stand on paint

One robot per cell – 1

One robot per cell – 2

Painting update

Painting only one cell

Not painting

Cells update

Stay painted



# Formulation

$\text{cell}(i, j, t) = 0 \text{ and } \sum_r \text{painting}(i, j, t, r) = 0 \Rightarrow \text{cell}(i, j, t + 1) = 0 \quad \forall i, j, t$	Not painted
$\text{move}(t, r) = 1 \Rightarrow  y(t + 1, r) - y(t, r)  +  x(t + 1, r) - x(t, r)  = 1 \quad \forall t, r$	Moving
$\text{move}(t, r) = 0 \Rightarrow  y(t + 1, r) - y(t, r)  +  x(t + 1, r) - x(t, r)  = 0 \quad \forall t, r$	Not moving
$\text{switch}(t, r) = 1 \Rightarrow  \text{color}(t + 1, r) - \text{color}(t, r)  = 1 \quad \forall t, r$	Switching colors
$\text{switch}(t, r) = 0 \Rightarrow \text{color}(t + 1, r) = \text{color}(t, r) \quad \forall t, r$	Not switching
$\text{paint}(t, r) = 1 \text{ and } \text{color}(t, r) = 0 \Rightarrow \text{stock0}(t + 1, r) = \text{stock0}(t, r) - 1$ $\text{and } \text{stock1}(t + 1, r) = \text{stock1}(t, r) \quad \forall t, r$	Decrement black stock
$\text{paint}(t, r) = 1 \text{ and } \text{color}(t, r) = 1 \Rightarrow \text{stock1}(t + 1, r) = \text{stock1}(t, r) - 1$ $\text{and } \text{stock0}(t + 1, r) = \text{stock0}(t, r) \quad \forall t, r$	Decrement white stock
$\text{paint}(t, r) = 0 \Rightarrow \text{stock0}(t + 1, r) = \text{stock0}(t, r)$ $\text{and } \text{stock0}(t + 1, r) = \text{stock0}(t, r) \quad \forall t, r$	Constant stocks
$\text{paint}(t, r) = 1 \text{ and } \text{color}(t, r) = 1 \Rightarrow \text{pattern}(x(t, r) + 1, y(t, r))$ $= \text{color}(t, r) \quad \forall t, r$	Respect pattern
$1 \leq y(t, r) \leq n, \quad 1 \leq x(t, r) \leq m \quad \forall t, r$	Stay inside the board
$t^*, y(t, r), x(t, r), \text{stock0}(t, r), \text{stock1}(t, r) \in \mathbb{Z}^+ \quad \forall t, r$	Integer variables
$\text{cell}(i, j, t), \text{painting}(i, j, t, r), \text{color}(t, r), \text{paint}(t, r), \text{move}(t, r),$ $\text{switch}(t, r) \in \{0, 1\} \quad \forall i, j, t, r$	Binary variables



# Solver



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```
graph LR; MODEL{MODEL} --> DATA{DATA}; DATA --> SOLVE{SOLVE}; SOLVE --> DEPLOY[DEPLOY]; ANALYZE{ANALYZE} --> MODEL; ANALYZE --> SOLVE;
```

*“IBM ILOG CPLEX CP Optimizer is a necessary and important complement to the optimization specialists’ toolbox for solving real-world operational planning and scheduling problems – without a significant investment in R&D.”*

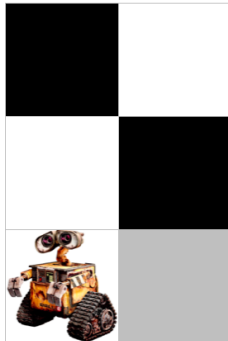
# Demo

Demonstration

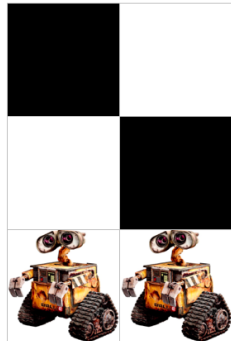
## Case studies



(a)



(b)



(c)

(a) 1.63608    (b) 36.309    (c) 117.117

TABLE – CPU times in seconds

# Plan

Introduction

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Optimisation

Discussion

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# Discussion

PDDL

Optimisation

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# Conclusion

Applications

Future work

