





Integer Range Analysis

Maglione Sandro June 29th, 2020



Context

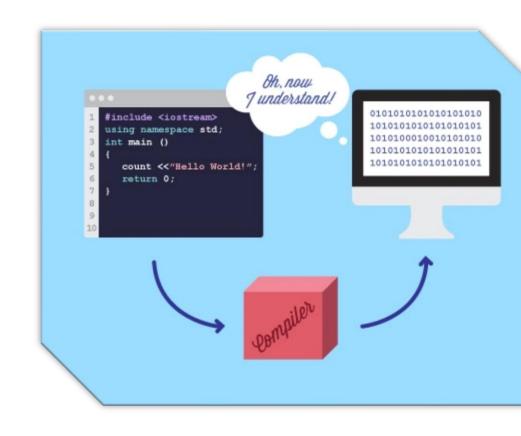
Source codes are made of millions of lines of code, problem of memory and resources!





Context

Optimization of resources using compiler and static analysis







Context





Bitwidth analysis, reduce number of bits allocated for each variable





How to optimize?



Static analysis, reduce memory usage while keeping the code correctness





a {32bit}

b {16bit}

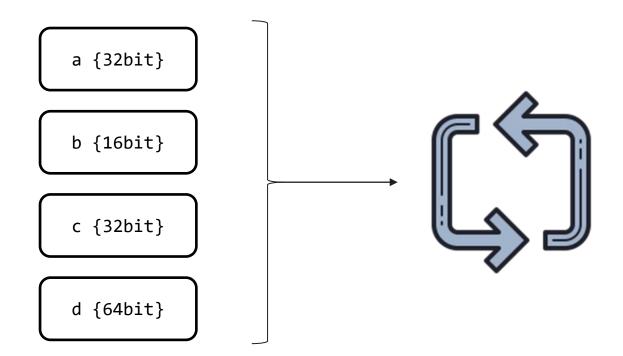
c {32bit}

d {64bit}

144bit

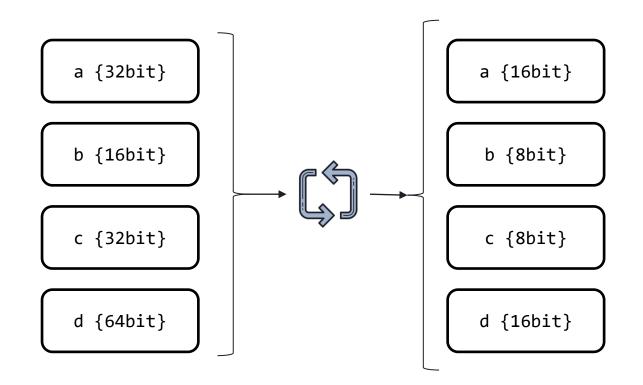






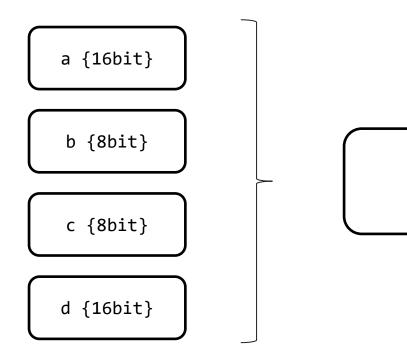










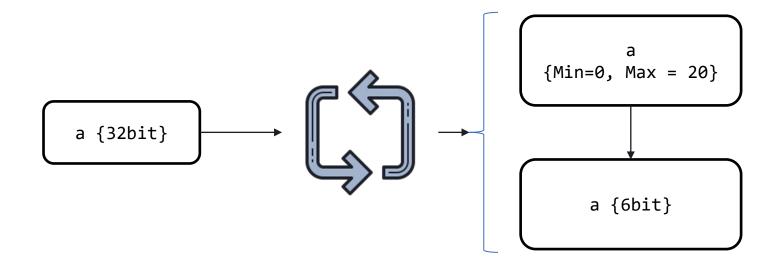


48bit





Purpose



Infer the minimum and maximum value, compute number of bits



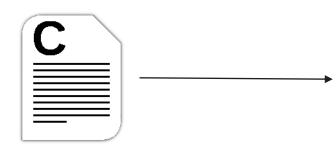


- Open-source compiler infrastructure project
- Intermediate representation (IR)
- LLVM Pass Framework
 - Transformations
 - Optimizations
 - Analysis









```
define dso_local i32 @main() #0 {
entry:
 %retval = alloca i32, align 4
 \%a = alloca i32, align 4
 \%b = alloca i32, align 4
 %c = alloca i32, align 4
 \%d = alloca i32, align 4
  store i32 0, i32 * %retval, align 4
  store i32 0, i32 * \%a, align 4
  store i32 10, i32 * %b, align 4
 \%0 = \text{load i32}, i32 * \%a, align 4
 \%add = add nsw i32 \%0, 3
  store i32 %add, i32 * %c, align 4
 \%1 = \text{load i} 32, i32 * \%b, align 4
 \%add1 = add nsw i32 %1, 10
  store i32 %add1, i32 * %d, align 4
  ret i32 0
```





```
define dso_local i32 @fun() #0 {
entry:
 %j = alloca i32, align 4
 %a = alloca i32, align 4
 store i32 10, i32 * %j, align 4
  store i32 0, i32 * %a, align 4
 br label %for.cond
for . cond:
   ; preds = %for.inc, %entry
 %0 = load i32, i32* %j, align 4
 %cmp = icmp sgt i32 %0, 0
 br il %cmp, label %for.body, label %for.end
for . body:
   ; preds = %for.cond
 %1 = load i32, i32 * %a, align 4
 %sub = sub nsw i32 %1, 5
  store i32 %sub, i32 * %a, align 4
 br label %for.inc
for . inc :
   ; preds = %for.body
 %2 = load i32, i32 * %j, align 4
 \%dec = add nsw i32 %2, -1
 store i32 %dec, i32 * %j, align 4
 br label %for.cond
for . end:
   ; preds = %for.cond
 %3 = load i32, i32 * %a, align 4
 ret i32 %3
```

```
define dso_local i32 @fun() #0 {
entry:
  br label %for.cond
for . cond:
    ; preds = %for.inc, %entry
 %a.0 = phi i32 [ 0, %entry ], [ %sub, %for.inc
 %j.0 = phi i32 [ 10, %entry ], [ %dec, %for.inc ]
 %cmp = icmp sgt i32 %j.0, 0
  br il %cmp, label %for.body, label %for.end
for . body:
    ; preds = %for.cond
 %sub = sub nsw i32 %a.0, 5
  br label %for.inc
for.inc:
    ; preds = %for.body
 \%dec = add nsw i32 %j.0, -1
  br label %for.cond
for . end:
    ; preds = %for.cond
  ret i32 %a.0
```



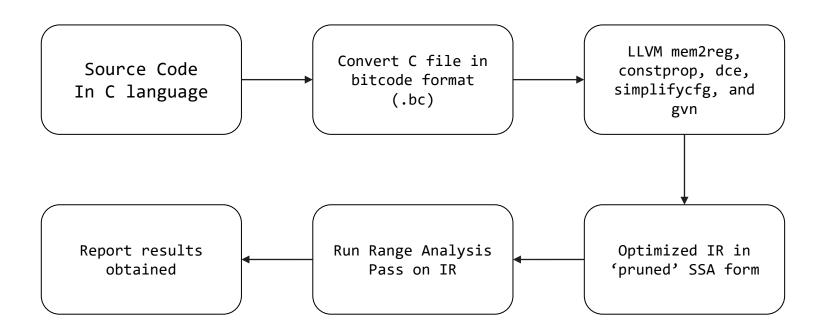


```
define dso_local i32 @fun() #0 {
entry:
 br label %for.cond
for . cond:
   ; preds = %for.inc, %entry
 %a.0 = phi i32 [ 0, %entry ], [ %sub, %for.inc ]
 %j.0 = phi i32 [ 10, %entry ], [ %dec, %for.inc ]
 %cmp = icmp sgt i32 %j.0, 0
 br il %cmp, label %for.body, label %for.end
for . body:
   ; preds = %for.cond
 %sub = sub nsw i32 %a.0, 5
 br label %for.inc
for.inc:
   ; preds = %for.body
 \%dec = add nsw i32 \%j.0, -1
 br label %for.cond
for.end:
   ; preds = %for.cond
  ret i32 %a.0
```





Methodology







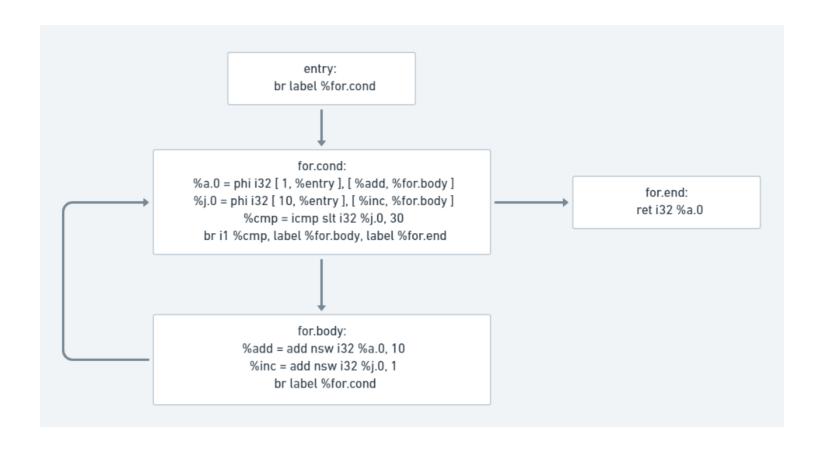
Execution

```
define dso_local i32 @fun() #0 {
 1 int fun()
                                                      entry:
                                                       br label %for.cond
 ^{2}
    {
 3
          int j = 10;
                                                      for.cond:
                                                         ; preds = %for.body, %entry
          int a = 1;
 4
                                                       %a.0 = phi i32 [ 1, %entry ], [ %add, %for.body ]
                                                       %j.0 = phi i32 [ 10, %entry ], [ %inc, %for.body ]
 5
                                                       \%cmp = icmp slt i32 %j.0, 30
          for (; j < 30; ++j)
 6
                                                       br i1 %cmp, label %for.body, label %for.end
                                                      for . body:
                                                         ; preds = %for.cond
                 a = a + 10;
                                                       %add = add nsw i32 %a.0, 10
                                                       \%inc = add nsw i32 \%j.0, 1
 9
                                                       br label %for.cond
10
                                                      for .end:
11
          return a;
                                                         ; preds = %for.cond
                                                       ret i32 %a.0
12 }
```





Execution







Execution

```
1 int fun()
2 {
3
       int j = 10;
       int a = 1;
4
 5
6
       for (; j < 30; ++j)
 7
           a = a + 10;
9
10
11
       return a;
12 }
```

```
---- VALUE-RANGES ----
```

BB: entry

BB: for . cond

$$a.0(1, 201) = 201 \{9 \text{ bit }\}$$

 $j.0(10, 30) = 21 \{6 \text{ bit }\}$

BB: for body
$$\begin{array}{l} j.0(10\,,\ 29) = 20\ \{6\,\mathrm{bit}\,\} \\ \mathrm{add}(-\mathrm{Inf}\,,\ +\mathrm{Inf}) = \mathrm{MAX} \\ \mathrm{inc}\,(11\,,\ 30) = 20\ \{6\,\mathrm{bit}\,\} \end{array}$$

BB: for .end
$$j.0(30, 30) = 1 \{2bit\}$$





```
--- VALUE-RANGES ---
BB: entry

BB: for.cond
    a.0(1, 201) = 201 {9 bit}
    j.0(10, 30) = 21 {6 bit}

BB: for.body
    j.0(10, 29) = 20 {6 bit}
    add(-Inf, +Inf) = MAX
    inc(11, 30) = 20 {6 bit}

BB: for.end
```

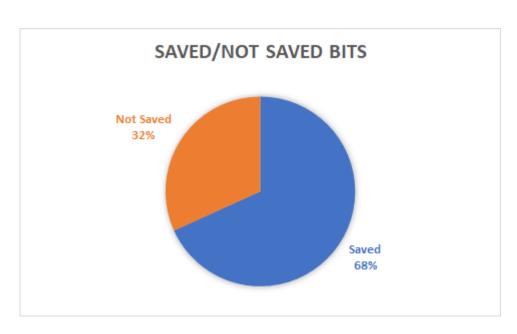
 $j.0(30, 30) = 1 \{2bit\}$



Performance achieved on the example source code.







Performance achieved on the example source code.



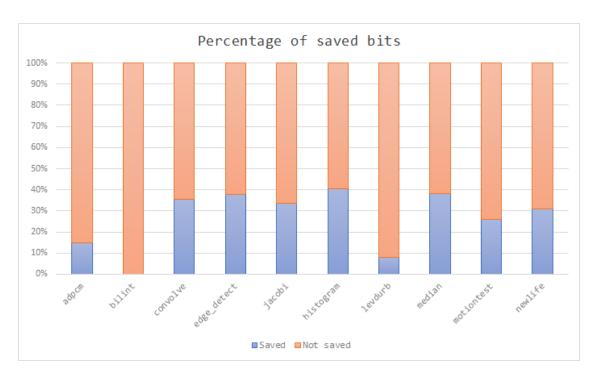


Benchmark	Lines	Performance (%)	Saved bits
adpcm	196	14.72	245
bilint	114	0.00	0
bubblesort	69	34.76	356
convolve	81	35.49	727
edge_detect	182	37.63	289
histogram	110	33.75	594
jacobi	68	40.55	571
levdurb	50	8.06	111
median	99	38.12	427
motiontest	44	26.08	217
newlife	121	30.71	747

Performance achieved on the bitwise benchmark set.



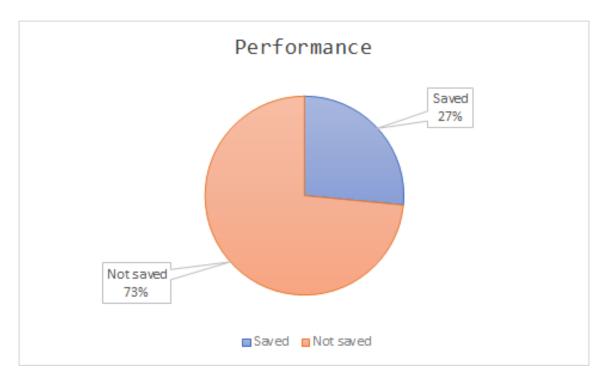




Performance achieved on the bitwise benchmark set.







Performance achieved on the bitwise benchmark set.





Related Works

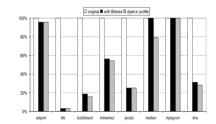
1. A Fast and Low-Overhead Technique to Secure Programs Against Integer Overflows Fernando Magno Quintao Pereira Raphael Ernani Rodrigues Victor Hugo Sperle Campos

Eliminated
24.93% of the
integer overflow
checks

2. The Design and Implementation of a Non-Iterative Range Analysis Algorithm on a Production Compiler Douglas do Couto Teixeira and Fernando Magno Quintao Pereira

Increased 1-bit variables by 1.04x, 8-bit variables by 52.3x and 16-bit variables by 26.6x

3. Bitwidth Analysis with Application to Silicon Compilation Mark Stephenson Jonathan Babb and Saman Amarasinghe









Maglione Sandro

10532096 - 940807 674339166

sandro.maglione@mail.polimi.it



