

CSL253 - Theory of Computation

Tutorial 7

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Question 3

Let B be the set of all infinite sequences over $\{0, 1\}$. Show that B is uncountable, using a proof by diagonalization.

Solution

In order to show that B , the set of all infinite sequences over $\{0, 1\}$ is uncountable, we show that no correspondence exists between \mathbb{N} and B . The proof is by contradiction. Suppose that a correspondence f existed between \mathbb{N} and B . We need to show that f fails to work as it is supposed to work. For it to be a correspondence, f must pair all the members of \mathbb{N} with all the members of B uniquely. But we will find an x in B that is not paired with anything in \mathbb{N} , which will be our contradiction.

The way to find this x is to actually construct it. We choose each digit of x to make x different from one of the elements in B that is paired with an element in \mathbb{N} . In the end we are sure that x is different from any element in B that is paired.

We can illustrate this idea by giving an example. Suppose the correspondence f exists. Let $f(1) = 0000\dots$, $f(2) = 1000\dots$, $f(3) = 0100\dots$, and so on, just to make up some values for f . Then f pairs the number 1 with 0000..., the number 2 with 1000..., and so on. The following table shows a few values of a hypothetical correspondence f between \mathbb{N} and B .

n	$f(n)$
1	0000...
2	1000...
3	0100...
4	1101...
\vdots	\vdots

We construct the desired x by giving its binary representation. Our objective is to ensure that $x \neq f(n)$ for any n . To ensure that $x \neq f(1)$, we let the first digit of x be anything different from the first binary bit of $f(1) = 0000\dots$. Hence we let it be 1 as this is the only option. To ensure that $x \neq f(2)$, we let the second binary bit of x be anything different from the second binary bit of $f(2) = 1000\dots$. Hence we let the second binary bit of x be 1. The third binary bit of $f(3) = 0100\dots$ is 0. Hence the third binary bit of x will be 1. The fourth binary bit of $f(4) = 1101\dots$ is 1. Hence the fourth binary bit of x will be 0. Continuing in this way down the diagonal of the table for f , we obtain all the digits of x , as shown in the following table. We know that x is not $f(n)$ for any n because it differs from $f(n)$ in the n th binary bit.

n	$f(n)$
1	<u>0</u> 000...
2	1 <u>0</u> 00...
3	01 <u>0</u> 0...
4	110 <u>1</u> ...
\vdots	\vdots

$$x = 1110...$$

As we can see with $x = 1110...$, $x \neq f(n)$ for every value of $n \in \mathbb{N}$. But as x is a member of B , there should be a pairing between x and some value in \mathbb{N} for the correspondence f to exist. Hence by proof of contradiction no such correspondence between \mathbb{N} and B exists and hence B is uncountable.