1. Consider three threads (T1, T2, and T3) that are created from the following code.

```
semaphore s1 = 1, s2 = 1, s3 = 1;
void T1()
                        void T2()
                                               void T3()
    semWait(s3);
                           semWait(s2);
                                                  semWait(s1);
    semWait(s1);
                           semWait(s3);
                                                  semWait(s2);
    semSignal(s3);
                         semSignal(s2);
                                                   semSignal(s1);
    semSignal(s1);
                           semSignal(s3);
                                                   semSignal(s2);
}
void main()
{
     parbegin (T1, T2, T3);
}
```

- a) Show a case in which deadlock will occur.
- b) Draw a resource allocation graph that shows the deadlock state.
- c) Will deadlock occur if the order of resource requests of T1 is changed as follows?

- 2. Suppose that there is a resource deadlock in a system. Is it possible for a process which is NOT in the circular chain in the corresponding resource allocation graph to be deadlocked?
- 3. Consider a system with a total of 150 units of memory, allocated to three processes as shown:

Process	Claim	Allocation
1	70	45
2	60	40
3	60	15

Apply the *banker's algorithm* to determine whether it would be safe to grant each of the following requests. If yes, indicate a sequence of terminations that could be guaranteed possible.

- a) A fourth process arrives, with a maximum memory need of 60 and an initial need of 25 units.
- b) A fourth process arrives, with a maximum memory need of 60 and an initial need of 35 units.

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1. A. B. C. D.	A condition of policy that must be present for a deadlock to be possible is Mutual exclusion Hold and wait No preemption All of the above
A. r B. r C. c	A closed chain of processes exists, such that each process holds at least one resource ed by the next process in the chain is the condition of no preemption nutual exclusion circular wait nold and wait
	The condition can be prevented by requiring that a process request all of quired resources at one time and blocking the process until all requests can be granted ltaneously. mutual exclusion hold and wait circular wait no preemption
4. types A. B. C. D.	The condition can be prevented by defining a linear ordering of resource hold and wait no preemption mutual exclusion circular wait
5. A. B. C. D.	In the banker's algorithm, a safe state is defined as one in which At least one potential process sequence does not result in a deadlock All potential process sequences do not result in a deadlock: Several potential process sequences do not result in a deadlock: None of the above
6.	A strategy for dealing with deadlocks that allows the presence of deadlock is called
A. B. C. D.	Deadlock Prevention Deadlock Avoidance Deadlock Detection None of the above