PROGRAMMING AND DATA STRUCTURES

SORTING ALGORITHMS

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OUTLINE

Sorting problem

Types of sorting solutions

Sorting algorithms

Comparison of sorting algorithms



STUDENT LEARNING OUTCOMES

At the end of this chapter, you should be able to:

- List the types and categories of sorting
- Implement the different sorting algorithms
- Evaluate the complexity of the sorting algorithms
- Compare the sorting algorithms

Sorting problem

- Given a list of data items, arrange the list in an ascending or descending order
- Commonly used in computer science to organize objects based on one specific criterion
- Allows using the efficient binary search algorithm
- ♦ Sorting is more complex than searching

Sorting types

INTERNAL

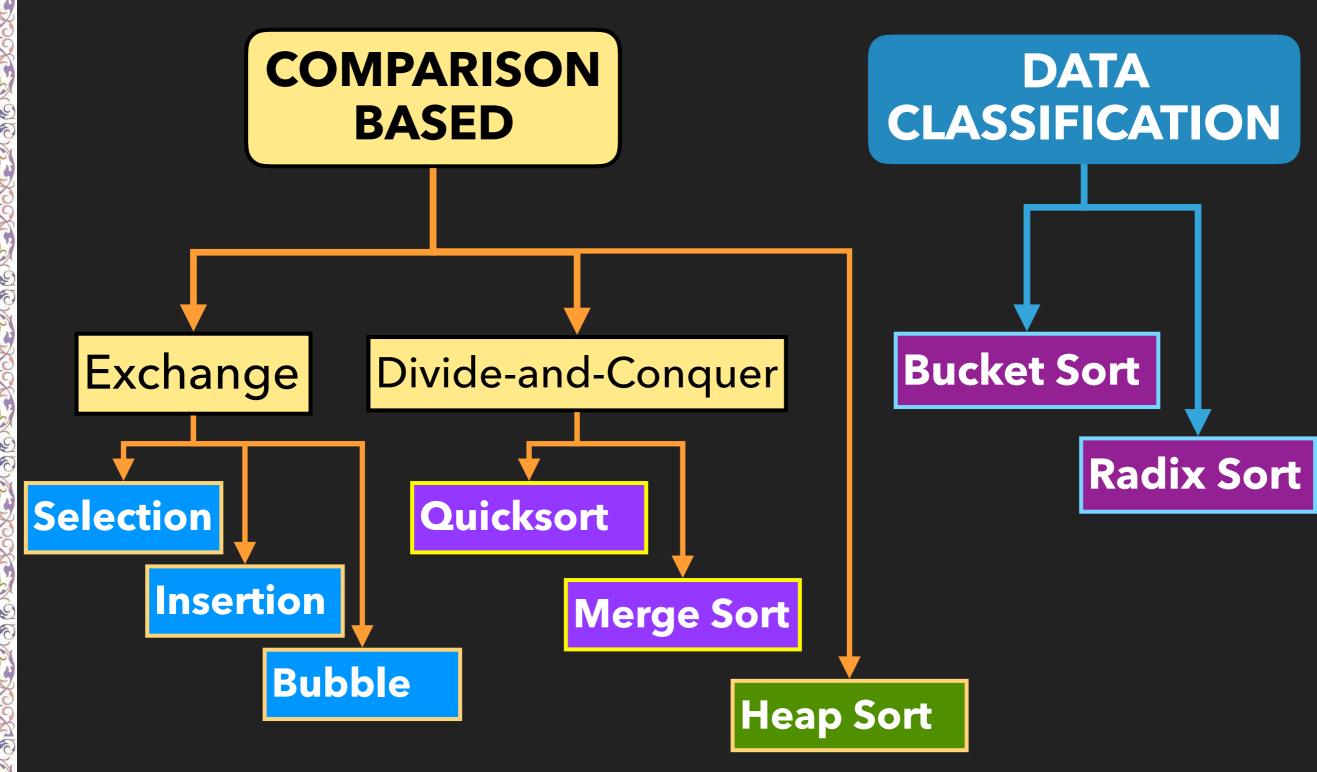
Data is in the main memory (RAM)

EXTERNAL

Data is in a secondary memory (Hard disk: Large Files)

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Sorting categories



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Selection Sort

Selection Sort

```
Algorithm selectionSort

for every element i (N size of the list)

find the smallest element from i to N-1

swap element i with the smallest element

End
```

Selection Sort

```
public static void selectionSort(int[] list) {
    int minIndex;
    for (int i=0; i<list.length-1; i++) {</pre>
       // Find the smallest element from i+1 to N
       int min = list[i];
       minIndex = i;
       for (int j=i+1; j<list.length; j++) {</pre>
          if (list[j] < min) {</pre>
               min = list[j];
               minIndex = j;
       // Swap the smallest element with element i
       swap(list, i, minIndex);
```

Selection Sort

Complexity Analysis

```
Iteration1(outer loop)
        (n)iterations(inner loop)
Iteration 2 (outer loop)
        (n-1) iterations (inner loop)
Iteration k (outer loop)
        (n-k+1) iterations (inner loop)
Iteration n-1 (outer loop)
        1 iteration (inner loop)
        1 + 2 + ... + (n) = n(n+1)/2
```

Time complexity: $O(n^2)$ - Quadratic growth Space complexity: O(1) - no additional space

Insertion Sort

67 33 21 84 49 50 75	Insert 33 at 0
33 67 21 84 49 50 75	Insert 21 at 0
21 33 67 84 49 50 75	Insert 84 at 3
21 33 67 84 49 50 75	Insert 49 at 2
21 33 49 67 84 50 75	Insert 50 at 3
21 33 49 50 67 84 75	Insert 75 at 5
21 33 49 50 67 75 84	

End

Insertion Sort

```
Algorithm insertionSort

for every element i

insert element i in the sorted list
(0 to i-1)

end for
```

Insertion Sort

```
public static void insertionSort(int[] list) {
   for (int i=1; i<list.length; i++) {</pre>
      //Insert element i in the sorted sub-list
      int currentVal = list[i];
      int j = i;
      while (j > 0 \&\& currentVal < (list[j - 1])){
         // Shift element (j-1) into element (j)
         list[j] = list[j - 1];
         j--;
      // Insert currentVal at index j
      list[j] = currentVal;
```

Insertion Sort

Complexity Analysis

```
Iteration1(outer loop)
     (1)iteration(inner loop)
Iteration 2 (outer loop)
     (2) iterations (inner loop)
Iteration k (outer loop)
     (k) iterations (inner loop)
Iteration n-1 (outer loop)
     n-1 iterations (inner loop)
1 + 2 + ... + (n-1) = n(n-1)/2
```

Time complexity: $O(n^2)$ - Quadratic growth Space complexity: O(1)

Bubble Sort

 At each iteration, exchange out of order pairs of elements until all elements are sorted

Pushing the largest element to the end of the list

Bubble Sort

Pass 1

7							
67	33	21	84	49	50	75	Out of order - swap
33	67	21	84	49	50	75	
33	67	21	84	49	50	75	Out of order - swap
33	21	67	84	49	50	75	
33	21	67	84	49	50	75	In order - No swap
33	21	67	84	49	50	75	Out of order - swap
33	21	67	49	84	50	75	
33	21	67	49	84	50	75	Out of order - swap
33	21	67	49	50	84	75	
33	21	67	49	50	84	75	Out of order - swap
22	21	67	40	ΕO	75	0.4	



Bubble Sort

Pass 2

33	21	67	49	50	75	84	Out to order - swap
21	33	67	49	50	75	84	
21	33	67	49	50	75	84	In order - No swap
21	33	67	49	50	75	84	Out of order - swap
21	33	49	67	50	75	84	
21	33	49	67	50	75	84	Out of order - swap
21	33	49	50	67	75	84	
21	33	49	50	67	75	84	In order - No swap



Bubble Sort

Pass 3

```
21 33 49 50 67 75 84 In order - No swap
```

No swaps: the list is sorted

Bubble Sort

```
Algorithm BubbleSort
  sorted = false
  last = N-1 (N size of the array)
 while (sorted is false)
    sorted = true
    for i=0 to last-1
       if(list[i] > list[i+1])
          swap(list[i], list[i+1])
          sorted = false
       end if
   end for
   last = last - 1;;
  end while
End
```

Bubble Sort

```
public static void bubbleSort(int[] list) {
 boolean sorted = false;
  for (int k=1; k < list.length && !sorted; k++) {</pre>
     sorted = true;
     for (int i=0; i<list.length-k; i++) {</pre>
        if (list[i] > list[i+1]) {
           swap(list, i, i+1);
           sorted = false;
```

Bubble Sort

Complexity Analysis

```
Iteration1(outer loop)
    (n-1)iteration(inner loop) to push the max
Iteration 2 (outer loop)
    (n-2) iterations (inner loop)
Iteration k (outer loop)
    (n-k) iterations (inner loop)
Iteration n-1 (outer loop)
    1 iterations (inner loop)
1 + 2 + ... + (n-1) = n(n-1)/2
Time complexity: O(n2) - Quadratic growth
Space complexity: O(1)
```

Comparison (group 1)

Algorithm	Complexity	Performance Analysis
Selection Sort	O(n ²)	Simple Redundant Processing Worst: Same performance all the time
Insertion Sort	O(n ²)	Lower overhead than selection and bubble sort Worst: Reversed list
Bubble Sort	O(n ²)	Complex Inefficient for large sets Worst: Reversed list

Merge Sort

- Recursive sorting algorithm
- Uses divide-and-conquer strategy
- Split the list in halves recursively until obtaining lists with one element
- Merge the lists back in order

Merge Sort

67 33 21 84 49 50 75

67 33 21 84 49 50 75

Split

1. Merge sort the first half

67 33 21

84 49 50 75

Split

67 33 21 84 49 50 75

Split

Merge Sort

67 33 21 84 49 50 75

67 21 33 84 49 50 75

Merge

21 33 67 84 49 50 75

Merge

Merge Sort

2. Merge sort the second half

21 33 67

84 49 50 75

Split

21 33 67 84 49 50 75

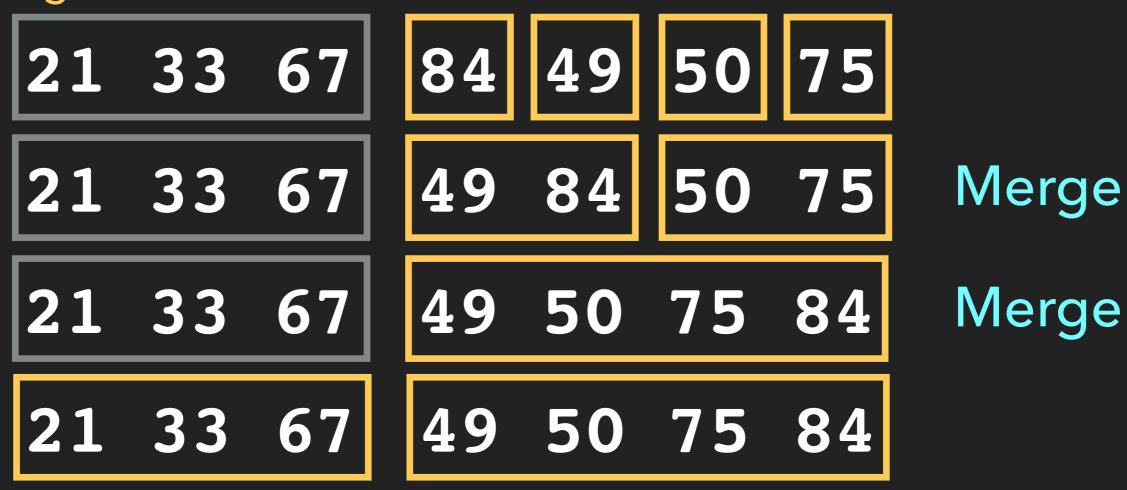
Split

21 33 67 84 49 50 75

Split

Merge Sort

2. Merge sort the second half



3. Merge the two sorted halves



Merge

Merge Sort

```
Algorithm MergeSort (recursive)
```

Split the array in two halves

MergeSort the first half

MergeSort the second half

Merge the two sorted halves

End

Merge Sort

```
public static void mergeSort(int[] list) {
  if (list.length > 1) {
    int[] firstHalf = new int[list.length/2];
    int[] secondHalf = new int[list.length -
                                  list.length/2];
    System.arraycopy(list, 0,
                     firstHalf, 0, list.length/2);
    System.arraycopy(list,list.length/2,
                     secondHalf, 0,
                     list.length-list.length/2);
    mergeSort(firstHalf);
    mergeSort(secondHalf);
    merge(firstHalf, secondHalf, list);
```

Merge Sort

```
public static void merge(int[] list1, int[] list2,
                           int[] list) {
  int list1Index = 0, list2Index = 0, listIndex = 0;
  while (list1Index < list1.length &&
        list2Index < list2.length) {</pre>
    if (list1[list1Index] < list2[list2Index])</pre>
      list[listIndex++] = list1[list1Index++];
    else
      list[listIndex++] = list2[list2Index++];
  // copy the remaining elements from list1 to list if any
  while(list1Index < list1.length)</pre>
    list[listIndex++] = list1[list1Index++];
  // copy the remaining elements from list2 to list if any
  while(list2Index < list2.length)</pre>
    list[listIndex++] = list2[list2Index++];
```

Merge Sort

Complexity Analysis (Time)

```
Splitting the array in halves
   K = (log n) iterations (n/2^k = 1)
Merging halves
   (n) iterations (worst case)
```

Time compelxity: O(n log n) Log Linear

Merge Sort

Complexity Analysis (Space)

```
Splitting the array in halves n/2, n/4, n/8, ...
```

```
Total number of additional space (n/2 + n/4 + n/8 + ...) \approx n
```

Space complexity: O(n)

Quick Sort

- Recursive in-place sorting algorithm
- Uses divide-and-conquer strategy
- Divide the list in two partially sorted parts using a pivot
 - ◆ Part 1: all elements less than the pivot
 - Part 2: all elements greater than the pivot
- Repeat quickSort recursively on each part

Quick Sort

50 33 21 49 (67) 84 75

Find position of the pivot

Part1

Part2

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Quick Sort

67) 33 21 84 49 50 75 pivot = 67

50 33 21 49 67 84 75 List1 List2

49 21 33 50 67 75 84 List3 List4

pivot (List1)= 50pivot (List2) = 84

Quick Sort

67) 33 21 84 49 50 75

50 33 21 49 67 84 75 List1 List2

49 33 21 50 67 75 84 List3

21 33 49 50 67 75 84 pivot (List3)= 49

List5

Quick Sort

21 33 49 50 67 75 84 Sorted list

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```
Algorithm QuickSort (recursive)
  Select a pivot
 Partition the array in two parts
 Part1: Elements less than the pivot
 Part2: Elements greater than the pivot
 QuickSort (part1)
```

End

QuickSort (part2)

Quick Sort

```
public static void quickSort(int[] list) {
   quickSort(list, 0, list.length-1);
// Recursive Helper Method
public static void quickSort(int[] list,
                             int first, int last) {
  if (last > first) {
    int pivotIndex = partition(list, first, last);
    quickSort(list, first, pivotIndex-1);
    quickSort(list, pivotIndex+1, last);
```

Quick Sort

```
public static int partition(int list[],
                              int first, int last) {
   int pivot;
   int index, pivotIndex;
   pivot = list[first];// the first element
   pivotIndex = first;
   for (index = first + 1; index <= last; index++) {</pre>
      if (list[index] < pivot) {</pre>
         pivotIndex++;
         swap(list, pivotIndex, index);
   swap(list, first, pivotIndex);
   return pivotIndex;
```

Complexity Analysis

```
Partitioning the array in ~ halves (log n) iterations (average)
```

```
Arranging elements around the pivot (n) iterations - worst case
```

```
Time complexity: average case O(n log n) - Log Linear
```

Complexity Analysis

```
Partitioning the array not in halves (n) iterations - worst case
```

```
Arranging elements around the pivot (n) iterations - worst case
```

```
Time complexity: O(n^2) - worst case Space complexity: O(1) - No additional space
```

How to select the pivot?

Original data is random,pivot = first element

Original data partially sorted, use "middle of the three" rule (median of first, middle, and last)

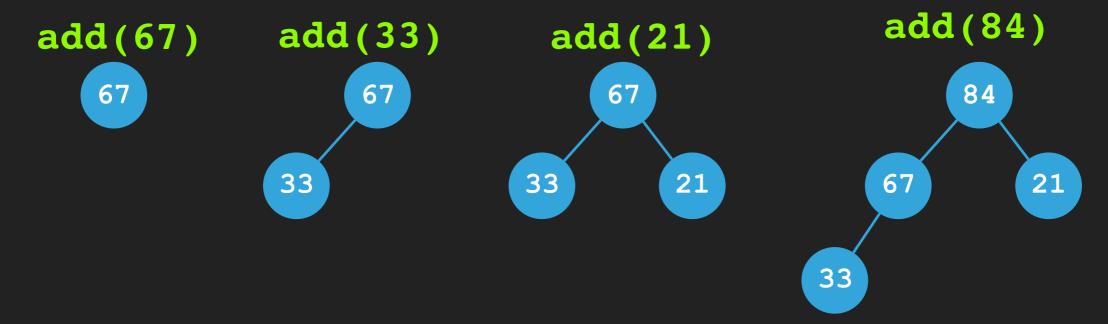
Heap Sort

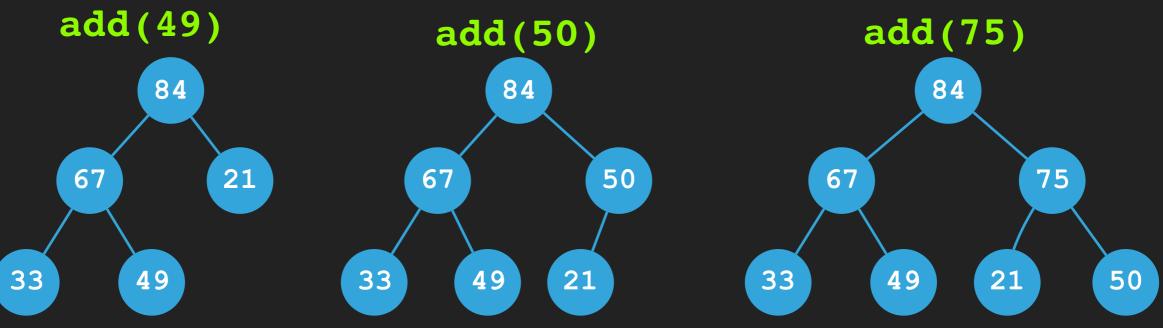
- Sorting using the heap data structure
 - Add the data to be sorted to the heap using add() method

Remove the elements from the heap using remove() method (elements are removed in ascending/descending order)

Heap Sort

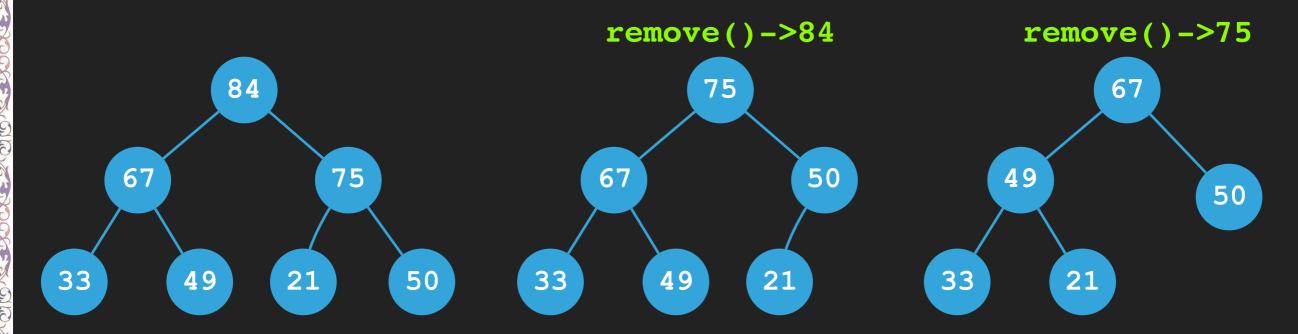
67 33 21 84 49 50 75

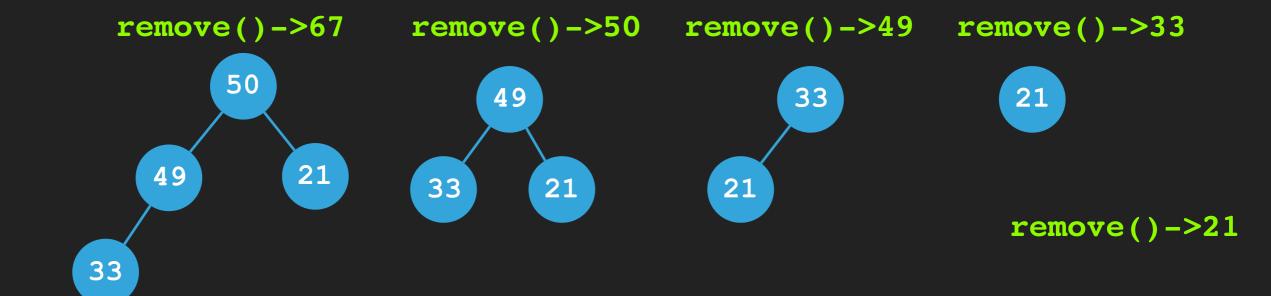




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Heap Sort





84 75 67 50 49 33 21

Heap Sort

```
Algorithm HeapSort
 create a heap h
 for each element i (0 to N-1) in list{
     h.add(element i)
 for each index i (N-1 to 0) in list{
  list[i] = h.remove()
End
```

Heap Sort

```
public static void heapSort(int[] list) {
    // Create a max heap
    Heap<Integer> heap = new Heap<>();
    // Add the elements of list to the heap
    for(int i=0; i<list.length; i++) {</pre>
       heap.add(list[i]);
    // Move the data from the heap back to list
    for (int i=list.length-1; i>=0; i--) {
       list[i] = heap.remove();
```

Heap Sort

Complexity Analysis (Time)

```
add() method O(log n)
    One path from a leaf to the root
 remove() method O(log n)
    One path from the root to a leaf
add() and remove() are called n times
                                    O(n)
Heap Sort: Worst case
    O(n log n) - Log Linear
```

Heap Sort

Complexity Analysis (Space)

Heap with **n** nodes required to sort the array list

Heap Sort space complexity: O(n)

Summary (group 2)

	Merge Sort	Quick Sort	Heap Sort	
Туре	Divide and Conquer Recursive	Divide and Conquer Recursive	Complete Binary Tree	
Main task	Merging O(n)	Partitioning <i>O(n)</i>	Adding/removing nodes to/from heap	
Time Complexity	Worst case: O(n logn)	Average: O(n logn) Worst case: O(n²)	Worst case: O(n logn)	
Space Complexity	Temporary arrays O(n)	In-place (No additional space)	Heap data structure O(n)	

- Sorting algorithms are general work for any type of data
- Sorting criterion defined in the method compareTo() or compare()
- Comparison based sorting cannot perform better than O(n log n)
- How can we sort without comparing? Data
 Classification Algorithms



Bucket Sort

- For integers only
- Range of values to be sorted [0, t]
- ◆ Use (t+1) buckets
- ◆ An element equal to v is put in bucket v
- A bucket holds all the elements with the same value

Bucket Sort

33

21 84 49

50

75

```
Bucket Bucket
               Bucket Bucket
                                     Bucket Bucket
                                                      Bucket
         21
                                                        84
                 33
                        49
                               50
                                       67
                                              75
                               50
                                       67
                                                        84
```

Bucket Sort

```
Algorithm bucketSort(list)
   create t+1 buckets
       (t is the maximum value in list)
   for each value in list (n)
     Assign value to bucket(value)
   for each bucket i (0 to t)
    Assign the elements of bucket i to list
```

Bucket Sort

```
public static void bucketSort(int[] list) {
   int t = max(list);
   ArrayList<ArrayList<Integer>> buckets;
   buckets = new ArrayList<>(t+1);
   // create t+1 buckets
   for(int i=0; i<t+1; i++)
     buckets.add(new ArrayList<>());
   //Distribute the data on the buckets
   for(int i=0; i<list.length; i++) {</pre>
     ArrayList<Integer> bucket = buckets.get(list[i]);
     bucket.add(list[i]);
   // Move the data from the buckets back to the list
   int k = 0;
   for(int i=0; i<buckets.size(); i++) {</pre>
     ArrayList<Integer> bucket = buckets.get(i);
     for(int j=0; j<bucket.size(); j++)</pre>
       list[k++] = bucket.get(j);
```

Bucket Sort

Complexity Analysis

```
Create the buckets O(t)
Distribute data on the buckets O(n)
Move data from the buckets to list O(t)
```

```
Bucket Sort:
```

```
Time Complexity: O(n+t) - Linear
Space Complexity: O(n+t)
```

Not practical for large t



Radix Sort

- Bucket Sort with fixed number of buckets
- Radix-10: 10 buckets (decimal numbers)
- Divide data into subgroups based on their radix positions
- ♦ Bucket Sort is applied on each radix position

Radix Sort

67

33 21 84 49 50

75





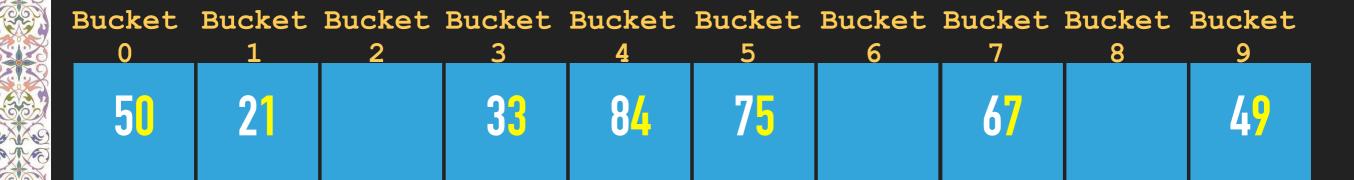
Radix Sort

Bucket sort using the ones position (10°)

33 21 84

50

75



50

21

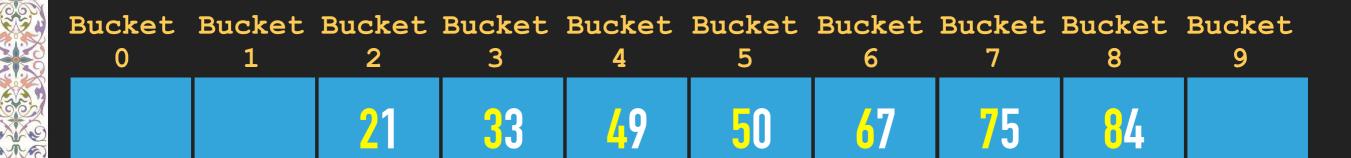
33

84 75 67

Radix Sort

Bucket sort using the tens position (10^{1})

50 21 33 84 75 67



33

49 50 67

75

End

Radix Sort

```
Algorithm radixSort(list)
   create 10 buckets
   max = number of digits of the largest value in list
   for each digit (0 to max-1){
    for each value in list
     Assign value to the bucket with number
                  (value % 10(digit+1) / 10digit)
    for each bucket i (0 to 9)
     Assign the elements of bucket i to list
    clear all buckets
```

```
public static void radixSort(int[] list) {
        ArrayList<ArrayList<Integer>> buckets;
        buckets = new ArrayList<>();
        Integer maxValue = max(list);
        int digits = maxValue.toString().length();
        for(int d=0; d<digits; d++) {</pre>
            // create 10 buckets
            for(int j=0; j<10; j++) {
                buckets.add(new ArrayList<>());
            //Distribute the data on the buckets
            for(int j=0; j<list.length; j++){</pre>
                // find the index of the bucket where list[j] should be placed
                int bucketIndex = (list[j] % (int)(Math.pow(10, d+1))) /
                                               (int)(Math.pow(10,d));
                ArrayList<Integer> bucket = buckets.get(bucketIndex);
                bucket.add(list[j]);
            // Move the data from the buckets back to the list
            int k=0;
            for(int j=0; j<10; j++) {
                ArrayList<Integer> bucket = buckets.get(j);
                for(int 1=0; 1<bucket.size(); 1++)</pre>
                    list[k++] = bucket.get(1);
            // clear all the buckets for the next iteration
            buckets.clear();
```

Radix Sort

Complexity Analysis

```
Classifying the data into buckets O(n)

Classifying for each position O(d)

Radix Sort:

Time Complexity: O(d.n)

Space complexity: O(n)

d: maximum number of radix positions
```

Testing the sorting algorithms

```
public class Test{
   public static void main(String[] args){
       int[] list = {67, 33, 21, 84, 49, 50, 75};
       print(list);
       Sort.selectionSort(list);
       System.out.print("Selection Sort: ");
       print(list);
       shuffle(list);
       print(list);
       Sort.insertionSort(list);
       System.out.print("Insertion Sort: ");
       print(list);
       shuffle(list);
       print(list);
       Sort.bubbleSort(list);
       System.out.print("Bubble Sort: ");
       print(list);
       shuffle(list);
       print(list);
       Sort.mergeSort(list);
       System.out.print("Merge Sort: ");
       print(list);
       shuffle(list);
       print(list);
       Sort.quickSort(list);
       System.out.print("Quick Sort: ");
       print(list);
       shuffle(list);
       print(list);
       Sort.heapSort(list);
       System.out.print("Heap Sort: ");
       print(list);
       shuffle(list);
       print(list);
       Sort.bucketSort(list);
       System.out.print("Bucket Sort: ");
       print(list);
       shuffle(list);
       print(list);
       Sort.radixSort(list);
       System.out.print("Radix Sort: ");
       print(list);
```

Summary

W-2/0		Quadratic	Log Linear			Line	ear	
6) 2 36 0 C 2 C	Sorting Algorithm	Selection Insertion Bubble Sort	Merge Sort	Quick Sort	Heap Sort	Bucket Sort (t buckets)	Radix Sort (d digits)	External Merge Sort
(6) 25 3/6 CM = 2/6	Туре	Exchange	Divide and Conquer		Binary Tree	Data Classification		Divide and Conquer
16 15 16 CK	Time	O(n²)	O(n logn)	O(n logn) to O(n²)	O(n logn)	O(n+t)	O(d.n)	O(n logn)
360X	Space	No additional space	Require temporary array O(n)		Heap O(n)	Require buckets		Require temporary files O(n)