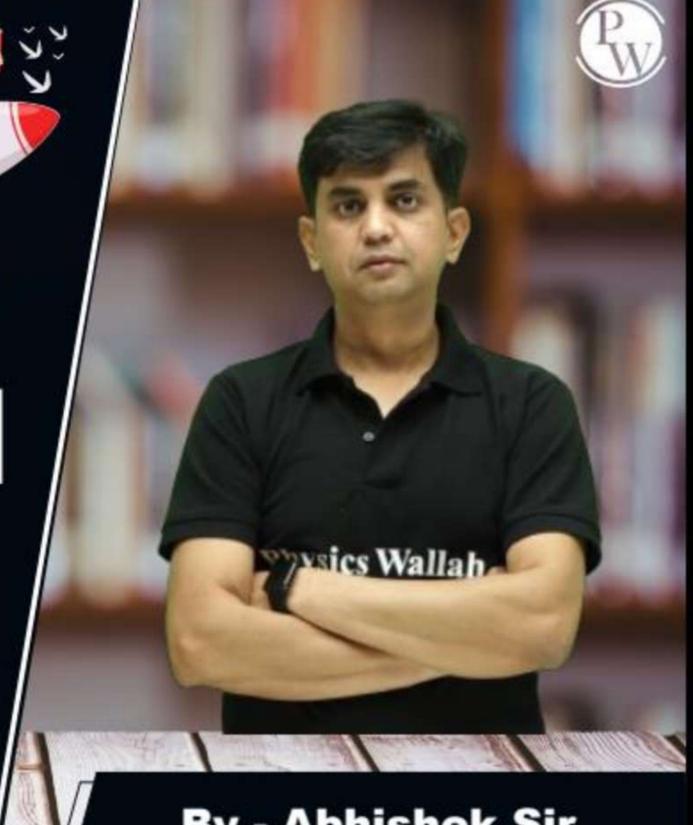
CS&IT ENGNERNG

Computer Network

MAC Layer



By - Abhishek Sir

Lecture No. - 06



Recap of Previous Lecture























Ethernet Bridging



ABOUT ME



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#Q. Consider an Ethernet segment with a transmission speed of 10⁸ bits/sec and a maximum segment length of 500 meters. If the speed of propagation of the signal in the medium is 2×10⁸ meters/sec, then the minimum frame size (*in bits*) required for collision detection is ______.

(SMA/CD tx22tp Tminm Frame Size tx=2tp [GATE-2024, Set-2, 2-Mark]

Ans = 500

Solution:



Minimum frame size

Minimum frame size = 2 * (Distance / Signal Speed) * Bandwidth

= 2 * (500 meters / 2×108 meters/sec) * 108 bits/sec

= 500 bits

#Q. Determine the maximum length of the cable (in km) for transmitting data at a rate of 500 Mbps in an Ethernet LAN with frames of size 10,000 bits. Assume the signal speed in the cable to be 2,00,000 km/s.

(A) 1

tx = 2tp

[GATE 2013]

(B) 2

(C) 2.5

(D) 5

Frame Size = 2 x Distance

Bandwidth = 2 x Signal speed

Ans: B



Solution:

```
Maximum Distance = (Frame size / Bandwidth) * (Signal Speed / 2)

= (10000 bits / 500 Mbps) * (200000 km/sec) / 2

= (10<sup>4</sup> bits / 5 * 10<sup>8</sup> bits/sec) * (2 * 10<sup>5</sup> km/sec) / 2

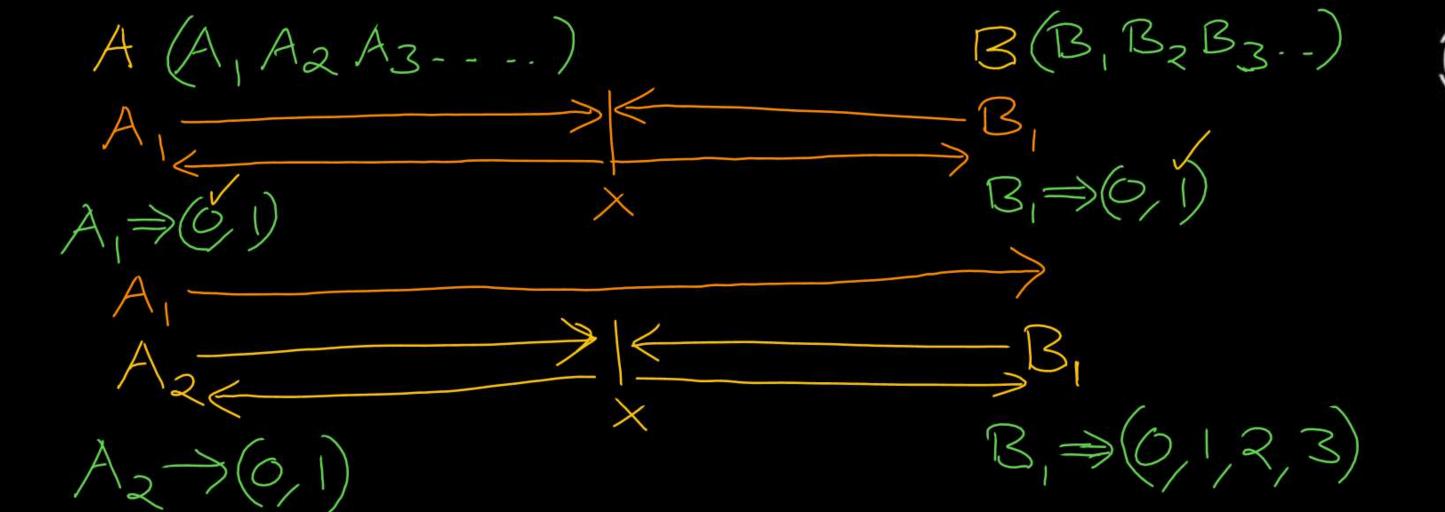
= 2 km
```

#Q. A and B are the only two stations on an Ethernet. Each has a steady queue of frames to send. Both A and B attempt to transmit a frame, collide, and A wins the first backoff race. At the end of this successful transmission by A, both A and B attempt to transmit and collide. The probability that A wins the second backoff race is:

[GATE 2004]

- (A) 0.5
- (B) 0.625
- (C) 0.75
- (D) 1.0

Ans: B

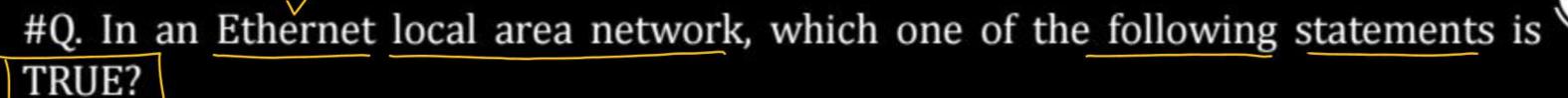








$$Ans = \frac{5}{8} = 0.625$$





- (A) A station stops to sense the channel once it starts transmitting a frame. FALSE
- The purpose of the jamming signal is to pad the frames that are smaller than the minimum frame size.
- (6) A station continues to transmit the packet even after the collision is detected. FALSE
- (D) The exponential backoff mechanism reduces the probability of collision on retransmissions.



Topic: LAN Standards



→ IEEE 802.3 : Ethernet [Bus Topology => CSMA/CD]

→ IEEE 802.4 : Token Bus [Bus Topology => Token]

→ IEEE 802.5 : Token Ring [Ring Topology => Token]

→ IEEE 802.11 : Wireless [Wi-FI] => CSMA/CA]





[GATE 2005]

- (A) IEEE 802.11 wireless LAN runs CSMA/CD protocol. FALSE
- (B) Ethernet is not based on CSMA/CD protocol. FALSE
- (C) CSMA/CD is not suitable for a high propagation delay network like satellite network. TRUE
- There is no contention in a CSMA/CD network. FALSE



Topic: Networking Devices

→ Physical Layer : Layer - 1 Device [Repeater and Hub]

→ Data Link Layer : Layer - 2 Device
[Bridge / Switch]

→ Network Layer : Layer - 3 Device [Router]

→ Application Layer : Layer - 7 Device
[Gateway]







- → Physical Layer [Layer 1 Device]
- → Two port device
- → Signal Regenerator [Regenerate the signal into output link]
- → One bit delay [Bit by bit forwarding]





End-to-end delay for one packet from Host A to Host B

$$= t_x + t_p + [(1 \text{ bit}) / \text{Bandwidth}] + t_p$$

$$= t_x + t_p + t_p + t_p$$

$$= t_x + t_p + t_p + t_p$$

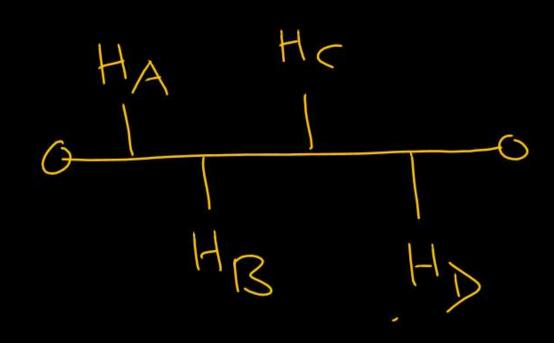


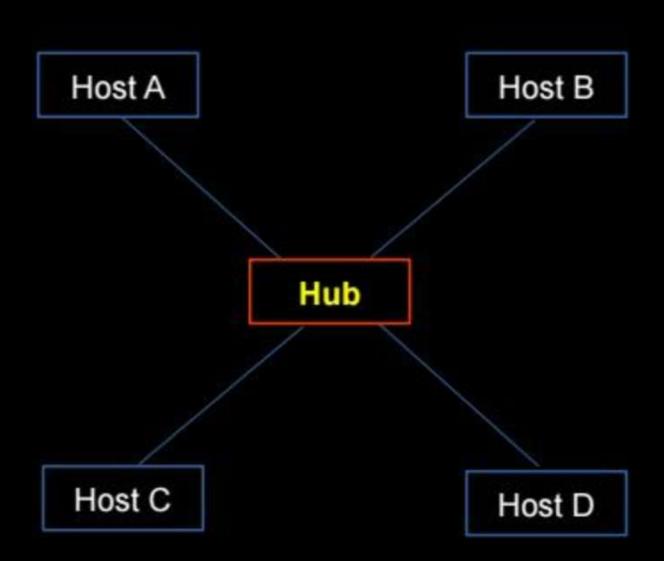


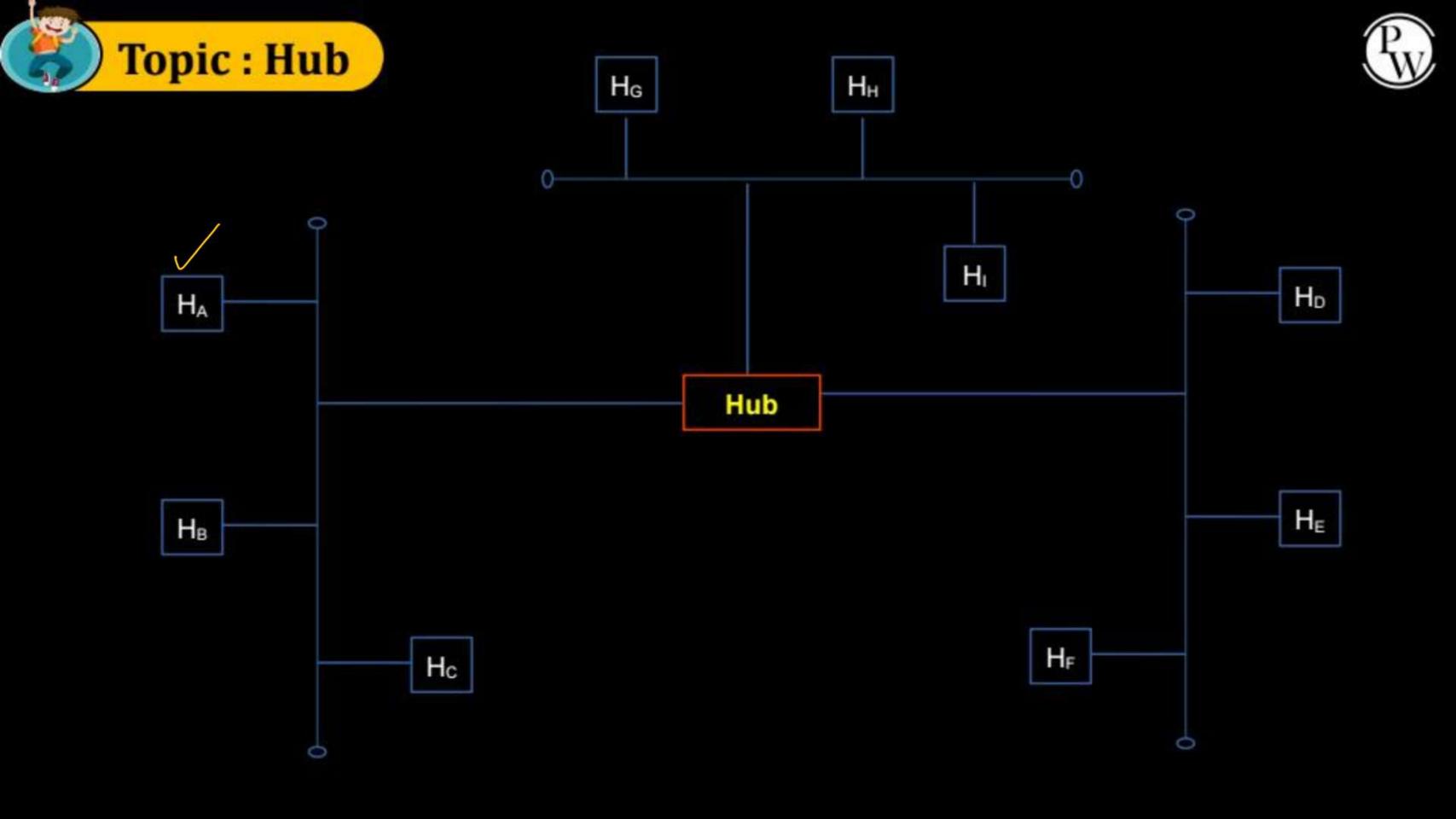
- → Physical Layer [Layer 1 Device]
- → Multiport repeater
- → Regenerate the signal into all connected links, except incoming link [Broadcasting device]
- → One bit delay [Bit by bit forwarding]















CASE I:

Destination MAC Address = Host C MAC Address

CASE II:

Destination MAC Address = Host D MAC Address

CASE III:

Destination MAC Address = Broadcast MAC Address [FF:FF:FF:FF:FF]





- → Entire LAN has same broadcast domain
- → Not collision domain separator [Entire LAN has same collision domain]

Collision domain: Place where collision can occur

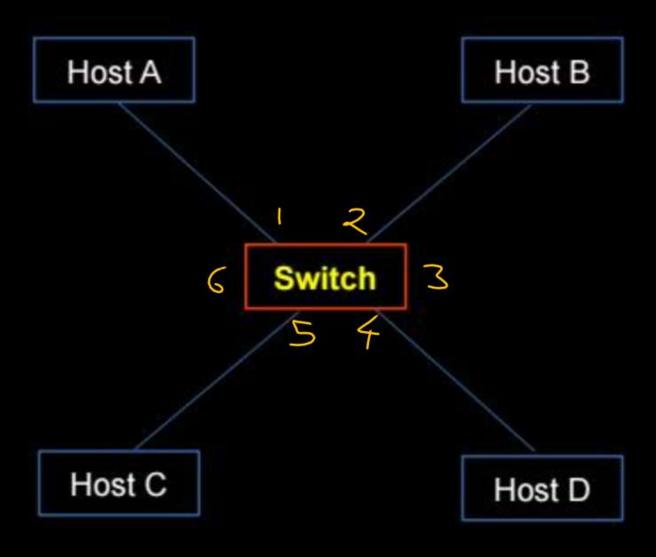




- → Data Link Layer [Layer 2 Device]
- → Old Name : Bridge [Multiport Bridge]
- → Store and Forward [Store, Process and Forward]
- → Switching Table [Forwarding table] [Forwarding based on Destination MAC Address]









Pw

Two types of switch:

- 1. Static
 - -> Non-adaptive
 - -> Manual update
- 2. Dynamic
 - -> Adaptive
 - -> Self learning
 - -> Plug-and-play

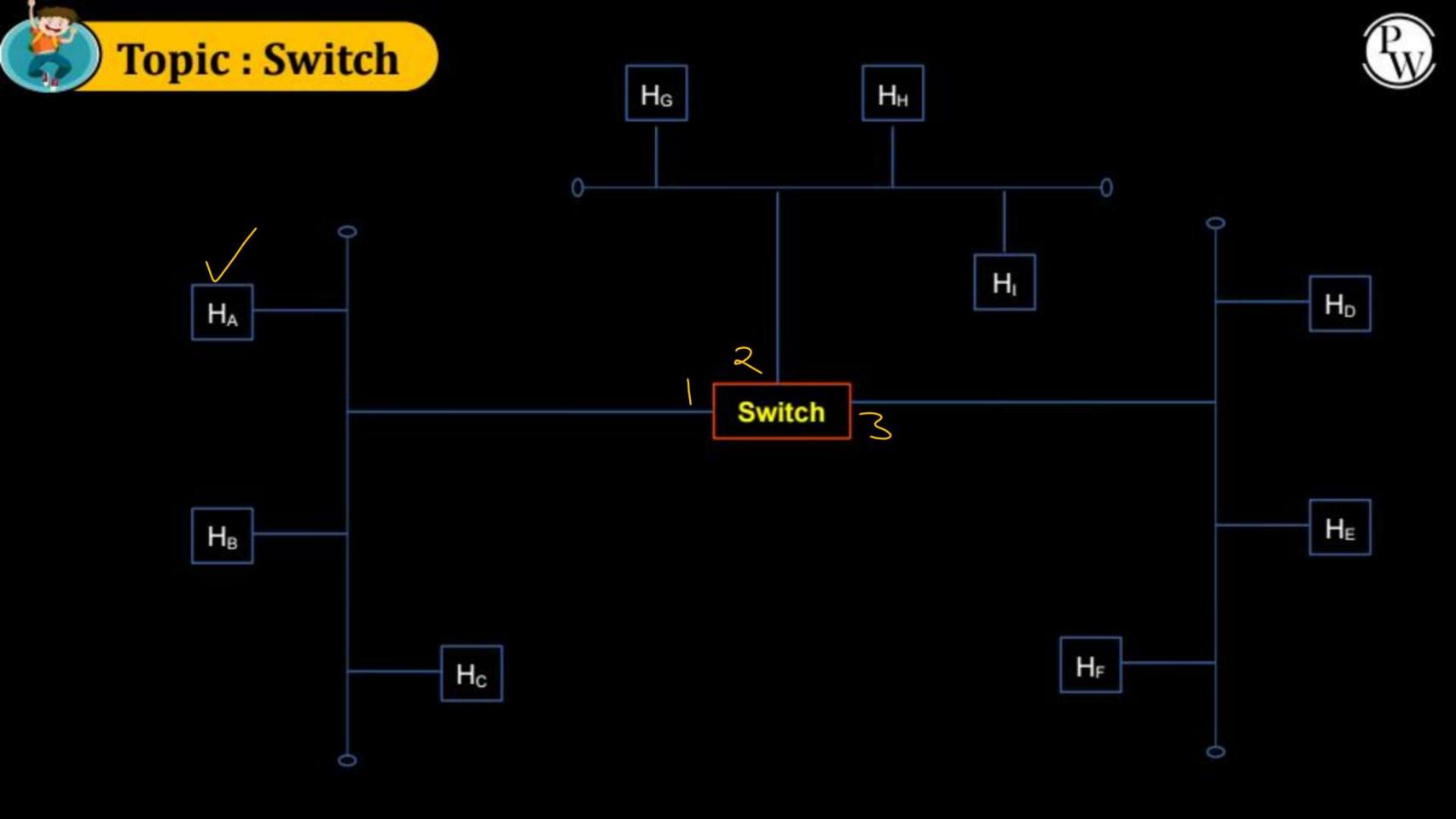


Topic: Switching Table



MAC Address	Interface	TTL
HA (MAC)		
HD (MAC)	4	
FF:FF:FF:FF	Broadcast	Infinite
Default	Broadcast	Infinite

TTL: Time to Live





Source MAC Address = Host A MAC Address

CASE I:

Destination MAC Address = Host C MAC Address

CASE II:

Destination MAC Address = Host D MAC Address

CASE III:

Destination MAC Address = Broadcast MAC Address [FF:FF:FF:FF:FF]





- → Entire LAN has same broadcast domain
- → Collision domain separator [Different LANs have different collision domain]



#Q. Which of the following statements is FALSE regarding a bridge?

[GATE 2005]

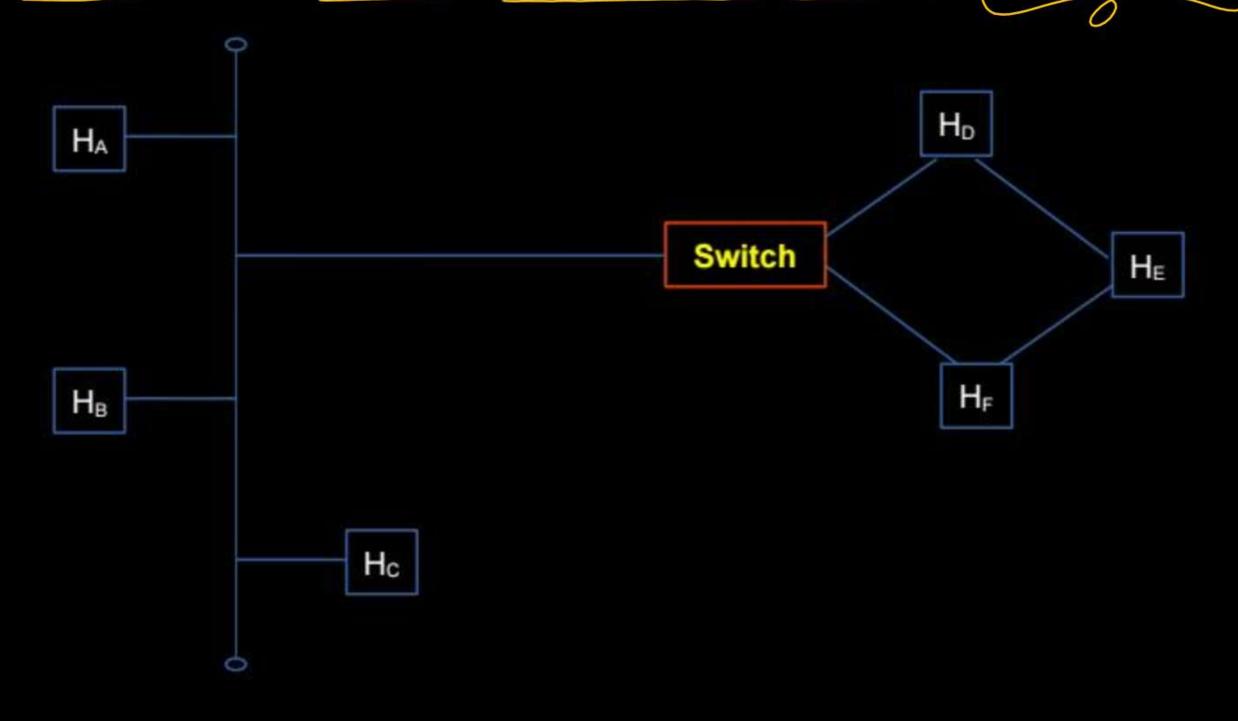
- (A) Bridge is a layer 2 device TRUE
- (B) Bridge reduces collision domain → R∪ ∈
- (C) Bridge is used to connect two or more LAN segments ⊤ ? ∪ ∈
- (D) Bridge reduces broadcast domain FALSE







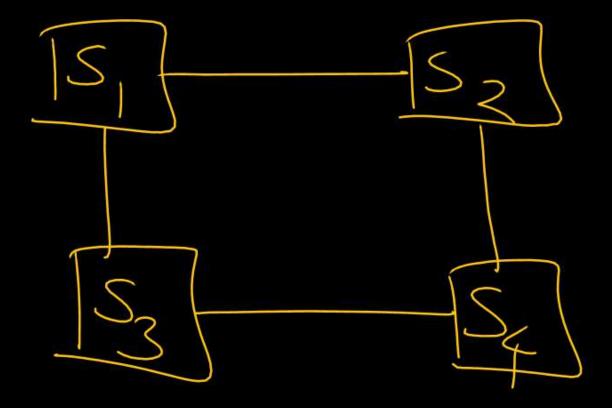
→ To connect different type of LAN segments minimum "Layer-2 device" required.







- → Multiple switches inside a network may form circuit (loop)
- → To avoid indefinite traversing of frames switches uses "Spanning Tree Protocol (STP)"



#Q. In a network of LANs connected by bridges, packets are sent from one LAN to another through intermediate bridges. Since more than one path may exist between two LANs, packets may have to be routed through multiple bridges. Why is the spanning tree algorithm used for bridge-routing?

- (A) For shortest path routing between LANs
- (B) For avoiding loops in the routing paths
- (C) For fault tolerance
- (D) For minimizing collisions



[GATE 2005]









THANK - YOU