

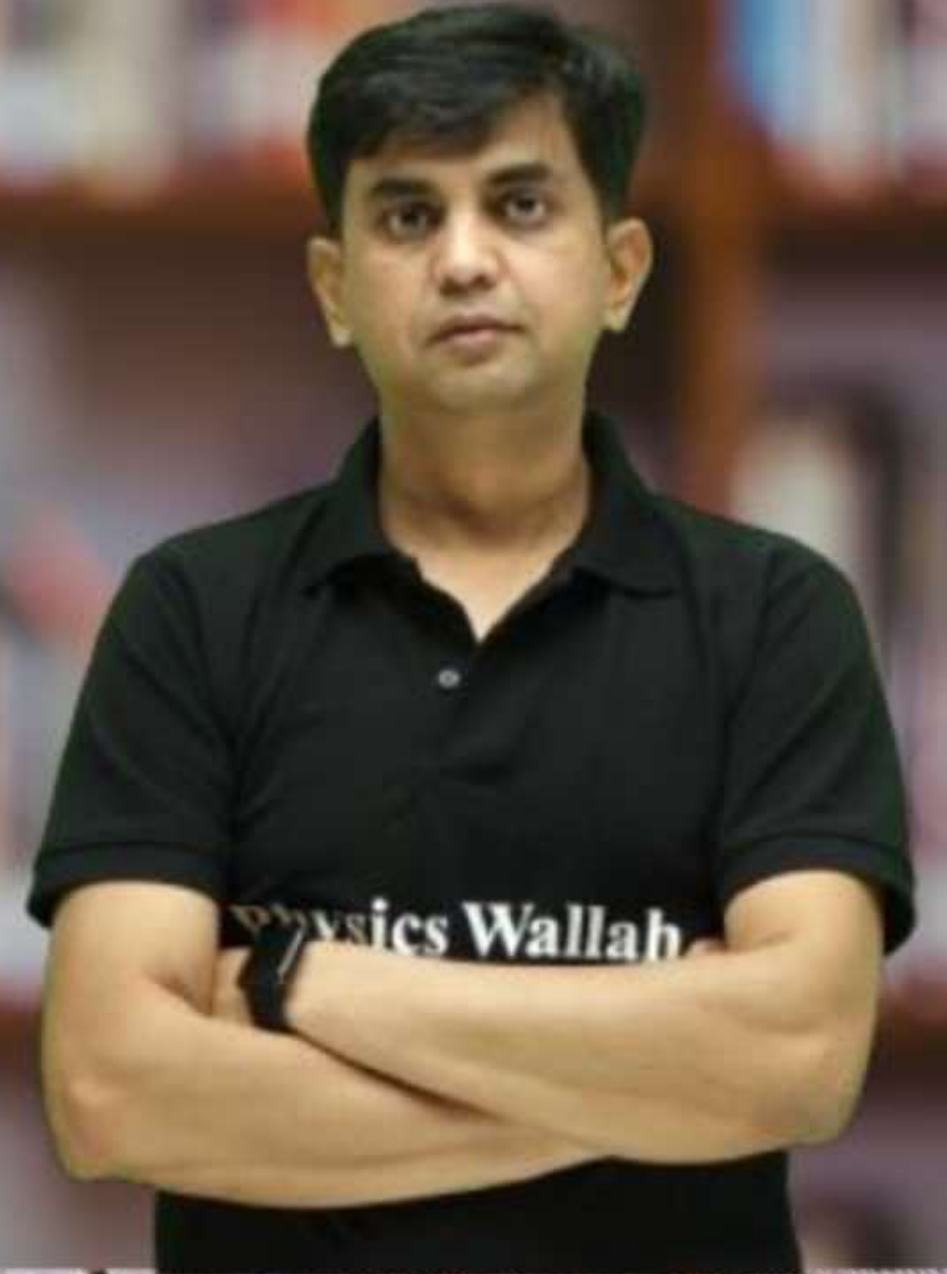
CS & IT ENGINEERING



Computer Network

Error Control

Lecture No. - 02



By - Abhishek Sir



Recap of Previous Lecture



Topic

Error Control

Topic

One-bit Parity



Topics to be Covered



Topic

Block Code

Topic

CRC

ABOUT ME



Hello, I'm **Abhishek**

- GATE CS AIR - 96
- M.Tech (CS) - IIT Kharagpur
- 12 years of GATE CS teaching experience

Telegram Link : https://t.me/abhisheksirCS_PW





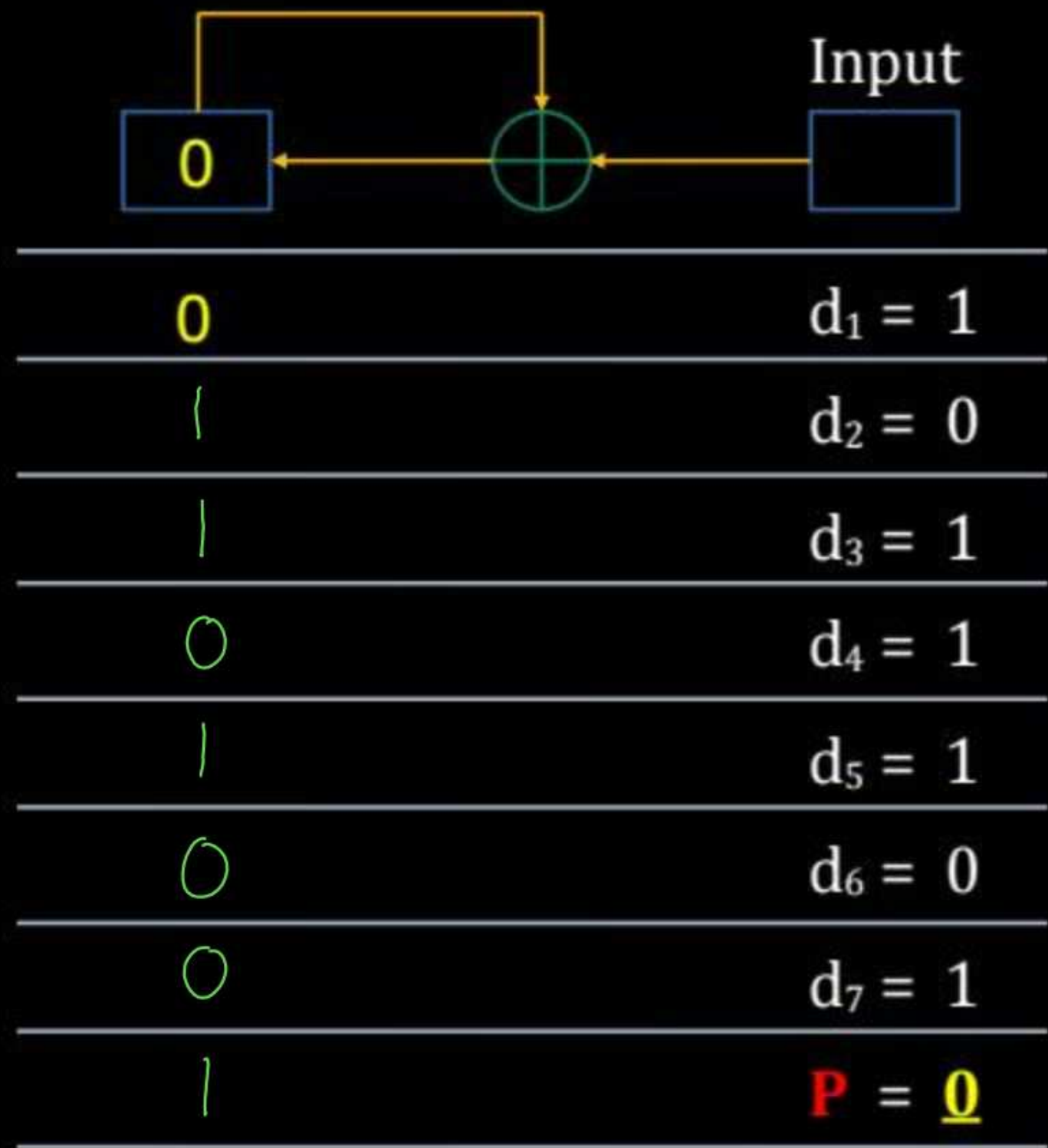
DATA = " 1 0 1 1 1 0 1"
 d₁ d₂ d₃ d₄ d₅ d₆ d₇

Transmitted Data = 1 0 1 1 1 0 1 **1**
P

$$a \oplus 1 = \bar{a}$$

$$a \oplus 0 = a$$

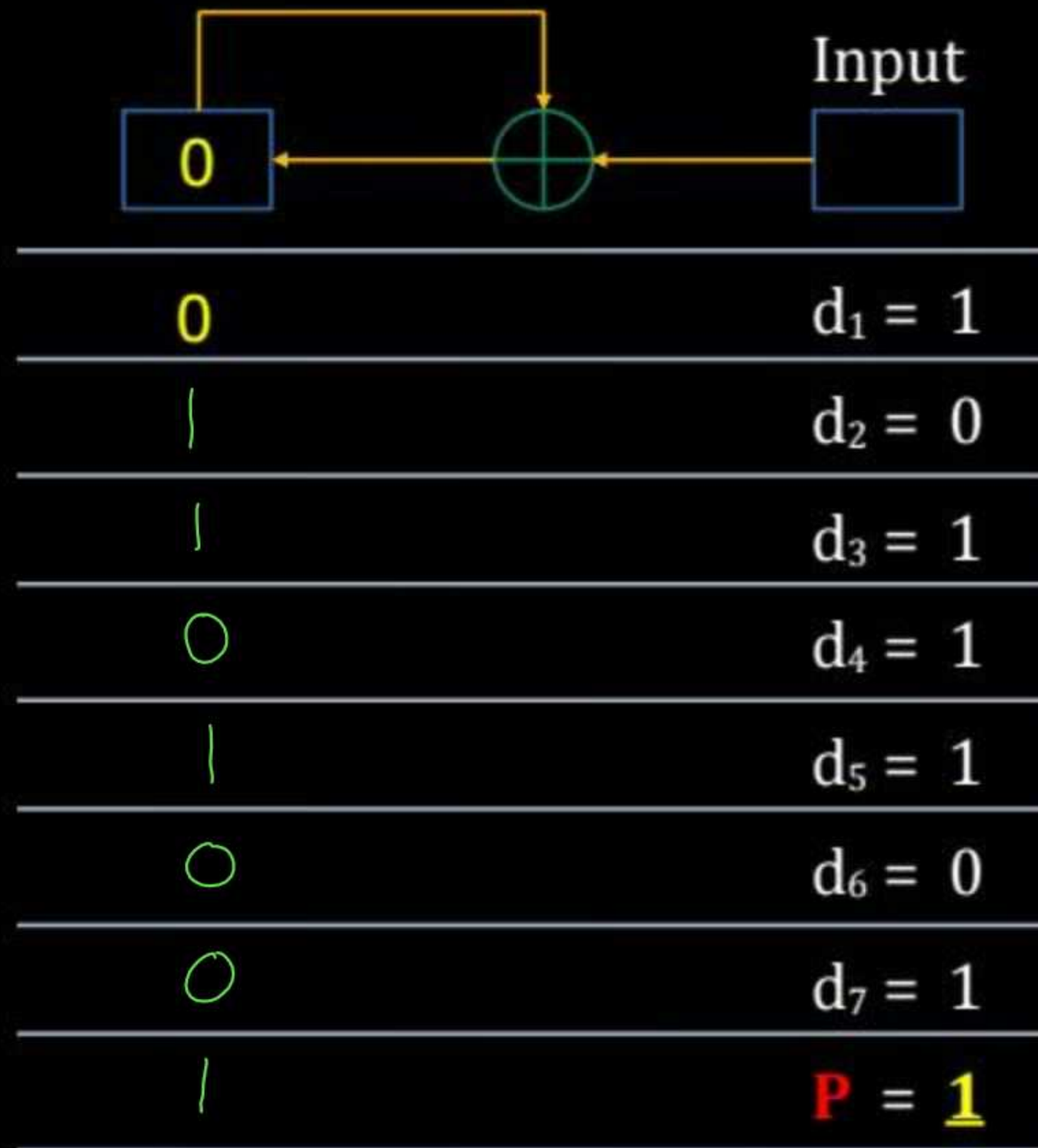




P = 1

AT Sender (Transmitter)

Input = " 1 0 1 1 1 0 1 0 "
 d_1 d_2 d_3 d_4 d_5 d_6 d_7 **P**



Result = 0

AT Receiver

Input = "1 0 1 1 1 0 1 1"
 d_1 d_2 d_3 d_4 d_5 d_6 d_7 **P**

if **Result** == **ZERO** :
 then Receiver concluded
 "**No any error detected**"
 else
 Receiver concluded
 "**Error detected**"



Topic : Block Code



k-bits input

→

n-bits output

k = input data bits

n = code length

→ 2^k codewords of length n.

No. of parity bit

in each codeword = $(n - k)$

Codeword = [data with parity]



Topic : Block Code



One-bit parity (with even parity) and 3 data bits

Data	-->	Codeword
<u>$d_1 d_2 d_3$</u>	-->	<u>$d_1 d_2 d_3$</u> <u>P</u>
<hr/>		
000	-->	0000
001	-->	0011
010	-->	0101
011	-->	0110
100	-->	1001
101	-->	1010
110	-->	1100
111	-->	1111

Set of ($2^k = 8$)
Codewords

Block Code
+
Linear Code
+
Cyclic Code



Topic : Linear Code

→ Linear combination of codewords

if C_i and C_j are codewords in set
then C_k is also be a codeword in that set
where $C_k = C_i \oplus C_j$

$$\begin{array}{r} 0110 \\ 1010 \\ \hline 1100 \end{array}$$

Binary arithmetic 

Modulo₂ Arithmetic :-

$$0+0=0$$

$$0+1=1$$

$$1+0=1$$

$$1+1=0$$



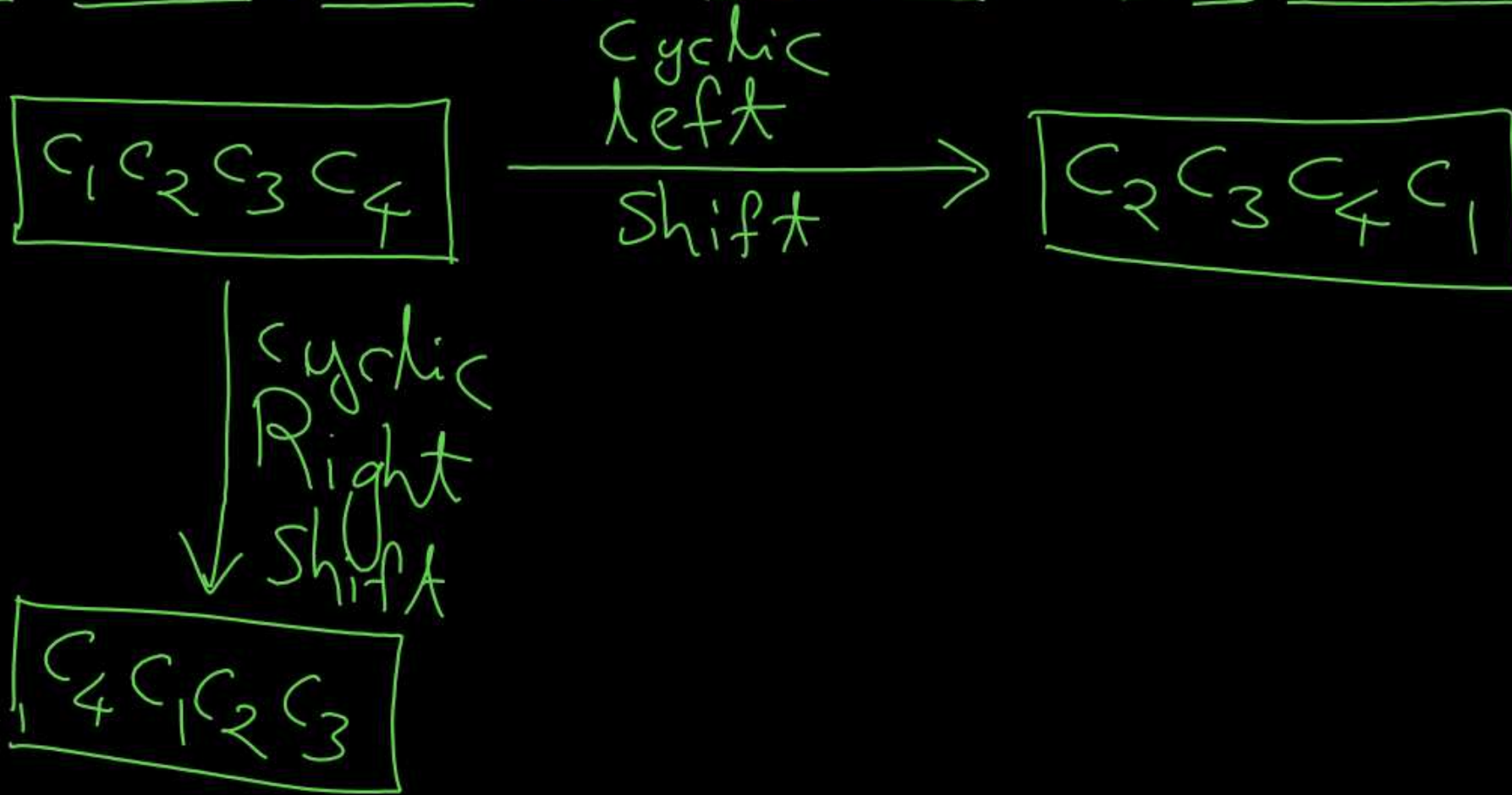
Topic : Cyclic Code



→ Cyclic combination of linear codewords

if C is codeword $[c_1 c_2 c_3 c_4]$

then bit-wise cyclic left or right shift on C is also be a codeword





Topic : Valid Codewords vs Invalid Codewords

→ Valid codewords :

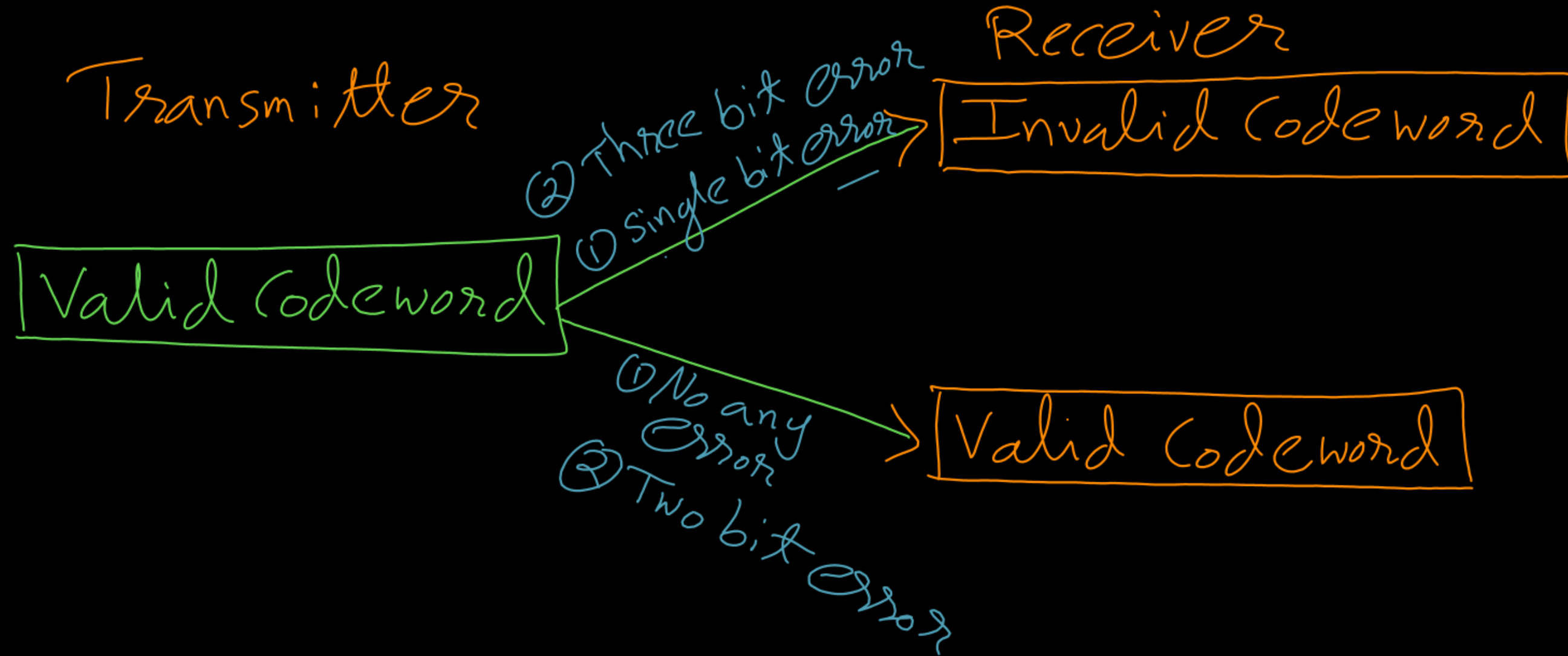
0000, 0011, 0101, 0110, 1001, 1010, 1100, 1111

→ Invalid codewords :

0001, 0010, 0100, 0111, 1000, 1011, 1101, 1110

→ Transmitter always transmit valid code word in the channel.

- If receiver receives a codeword from channel, which belongs to valid code set, then receiver conclude "no error detected".
- If receiver receives a codeword from channel, which belongs to Invalid Codeword set then receiver conclude "Error Detected".





Topic : CRC



- Cyclic Redundancy Check (CRC)
- Frame Check Sequence (FCS)



Topic : Polynomial Function



→ Polynomial function, where coefficients are either zero or one.

Example :

$$\text{Function} = \underline{X^5 + X^2 + 1}$$

$$= 1 \cdot X^5 + 0 \cdot X^4 + 0 \cdot X^3 + 1 \cdot X^2 + 0 \cdot X^1 + 1 \cdot X^0$$



Topic : Generator Polynomial

$G(X)$: Generator Polynomial function

→ $(n+1)$ terms [X^n to X^0]

→ $\text{Degree}[G(X)] = n$ \Rightarrow coefficient of the term x^n should be one.

→ Coefficient of term X^0 should be "one"

[$G(X)$ should not be ^{completely} divisible by X]

→ Both transmitter and receiver must agree on same $G(X)$.

$$\underline{G(X)} = X^n + \dots + 1$$



Topic : Divisor



$$\underline{G(X)} = X^n + \dots + 1$$

Divisor : binary string, (n+1) bits [1 1]

Example :

$$G(X) = X^3 + X^2 + 1$$

$$= 1 * X^3 + 1 * X^2 + 0 * X^1 + 1 * X^0$$

$$\text{Divisor} = 1\ 1\ 0\ 1$$



Topic : Message Polynomial

M(X) : Message Polynomial function

→ m terms, [$X^{(m-1)}$ to X^0]

→ coefficients are either zero or one

DATA (Message) : binary string (m - bits)



Topic : Message Polynomial



DATA (Message) : binary string (m - bits)

Example :-

$$\underline{M(X)} = \underline{X^7} + X^4 + X^3 + X$$

$$= 1 \cdot X^7 + 0 \cdot X^6 + 0 \cdot X^5 + 1 \cdot X^4 + 1 \cdot X^3 + 0 \cdot X^2 + 1 \cdot X^1 + 0 \cdot X^0$$

$$\text{DATA} = \underline{1} \underline{0} \underline{0} \underline{1} \underline{1} \underline{0} \underline{1} \underline{0}$$



Topic : CRC



Transmitter protocol :

$$\underbrace{[M(X)]}_{f_1(x)} \underbrace{* X^n}_{f_2(x)} \text{ [Modulo-2 Division] } \underbrace{[G(X)]}_{f_3(x)}$$

$[f_1(x) * f_2(x)] \text{ division } [f_3(x)]$



2 mins Summary



Topic

Block Code

Topic

CRC



THANK - YOU