

# CS & IT ENGINEERING



**Computer Network**

Switching & Routing

**Lecture No. - 03**



**By - Abhishek Sir**



# Recap of Previous Lecture



Topic

Packet Switching

Topic

Virtual Circuit Switching

Topic

Routing





# Topics to be Covered



Topic

Routing

Topic

Link State Routing

Topic

Distance Vector Routing

# ABOUT ME



Hello, I'm **Abhishek**

- GATE CS AIR - 96
- M.Tech (CS) - IIT Kharagpur
- 12 years of GATE CS teaching experience

Telegram Link : [https://t.me/abhisheksirCS\\_PW](https://t.me/abhisheksirCS_PW)





## Topic : Routing



→ Identify optimal (best) path between Source and Destination.

→ Types of Routing :

1. Static Routing

→ Non-adaptive Routing

2. Dynamic Routing

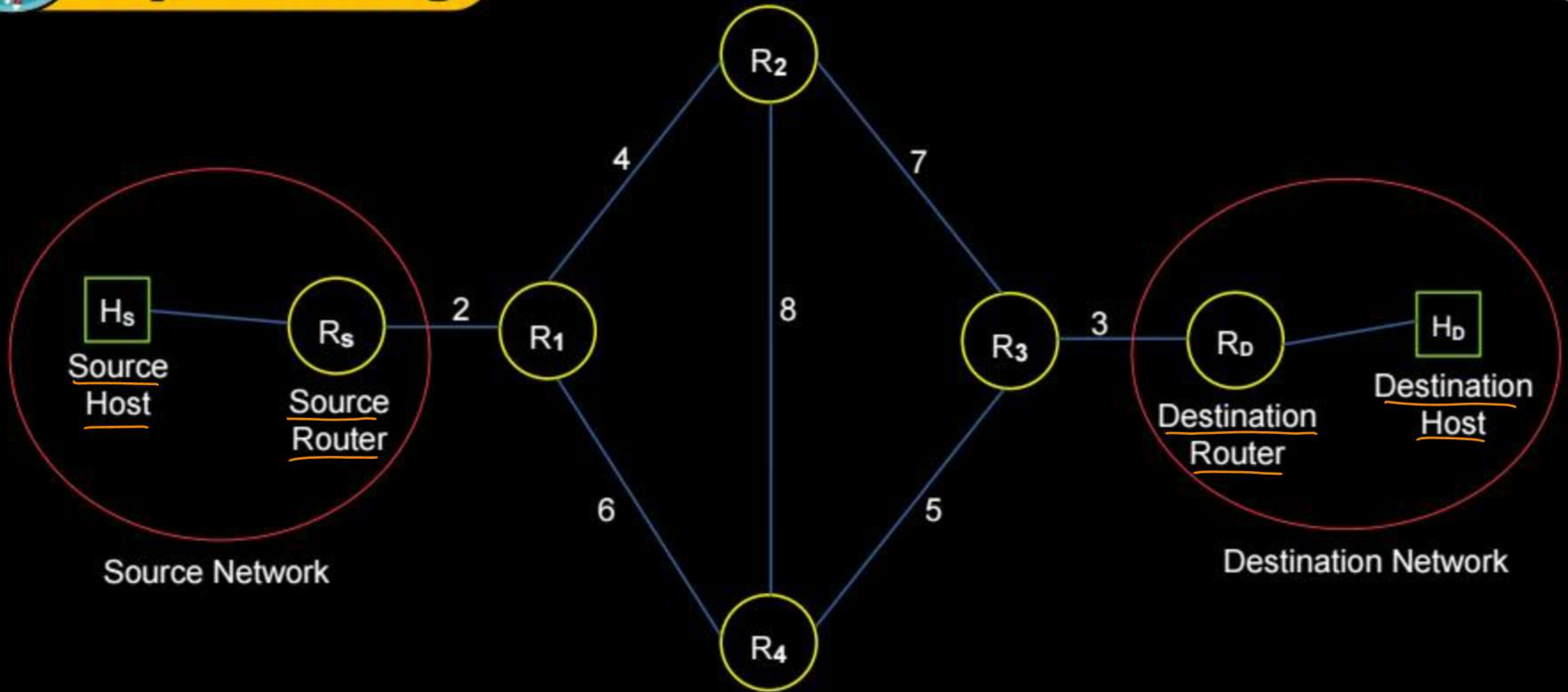
→ Adaptive Routing

→ Load sensitive or Load insensitive





# Topic : Routing





## Topic : Link Cost



→ Metric to compare link :

1. Distance (Hop Count)
2. Delay
3. Bandwidth

Cost  $\propto$  Distance

Cost  $\propto$  Delay

Cost  $\propto \frac{1}{\text{Bandwidth}}$



## Topic : Static Routing



→ Non-adaptive Routing ✓

→ Types of Static Routing :

1. Shortest Path Algorithm ✓ ⇒ Input

2. Flooding

→ Broadcast

(1) Graph [Matrix/List]

(2) Source





## Topic : Dynamic Routing



→ Adaptive Routing

→ Types of Routing Algorithm:

1. Dijkstra's Link State Routing Algorithm
2. Bellman Ford Distance Vector Routing Algorithm



# Topic : Routing Protocols



Two categories :

1. Interior Gateway Protocol (IGP) [With in ISP (A.S.)]

1.1 Routing Information Protocol (RIP)

→ Distance Vector Routing (RIP-DV)

1.2 Open Shortest Path First (OSPF)

→ Link State Routing [OSPF-LS]

2. Exterior Gateway Protocol (EGP) [Between ISPs (A.S.)]

2.1 Border Gateway Protocol (BGP)

→ Path Vector Routing



#Q. Which one of the following is TRUE about interior Gateway routing protocols - Routing Information Protocol (RIP) and Open Shortest Path First (OSPF) ?

[GATE-2014, Set-2, 1-Mark]

- RIP-DV                      OSPF-LS
- ✓ (A) RIP uses distance vector routing and OSPF uses link state routing
  - ✗ (B) OSPF uses distance vector routing and RIP uses link state routing
  - ✗ (C) Both RIP and OSPF use link state routing
  - ✗ (D) Both RIP and OSPF use distance vector routing

Ans: A





## Topic : Routing



→ Routing Algorithm Classification :

1. Decentralized Routing : [Distributed Routing]
2. Global Routing : [Centralized Routing]



## Topic : Decentralized Routing

- Distributed Routing Algorithm
- Iterative process of computation
- Exchange of information with neighbors
- Routers initially only know link costs to attached neighbors
- "Distance Vector" Routing Algorithm



## Topic : Global Routing



- Centralized Routing Algorithm
- All routers have complete topology link cost information
- “Link State” Routing Algorithm





## Topic : Link State Routing



- Centralized Routing Algorithm
- Link State Routing is divided into three steps :

1. Maintain Link State information ✓
2. Link State broadcast ✓
3. Share updated Link State information } → make algo. adaptive



## Topic : Link State Routing

### 1. Maintain Link State information

- Every router maintain separate "Link State" information
- Maintain information about adjacent link only  
[Adjacent (neighbor) routers only]

OSPF protocol in every Router

- i) Link state info construct
- ii) Pass this LS info to IP

Source IP Add. = Own IP Add.

Dest. IP Add. = Class D Special IP Add. (Broadcast)

OSPF → IP





## Topic : Link State Routing



### 2. Link State broadcast

- Every router flood (broadcast) its "Link State" information  
[To all other routers in the network]
- Every router have "Link State" information of all other routers  
[Complete information about the network topology]
- Every router construct "Adjacency Matrix" or "Adjacency List" ✓
- Every router execute "Dijkstra's Algorithm" locally ✓  
[To find optimal distance and path to all other routers]





## Topic : Link State Routing

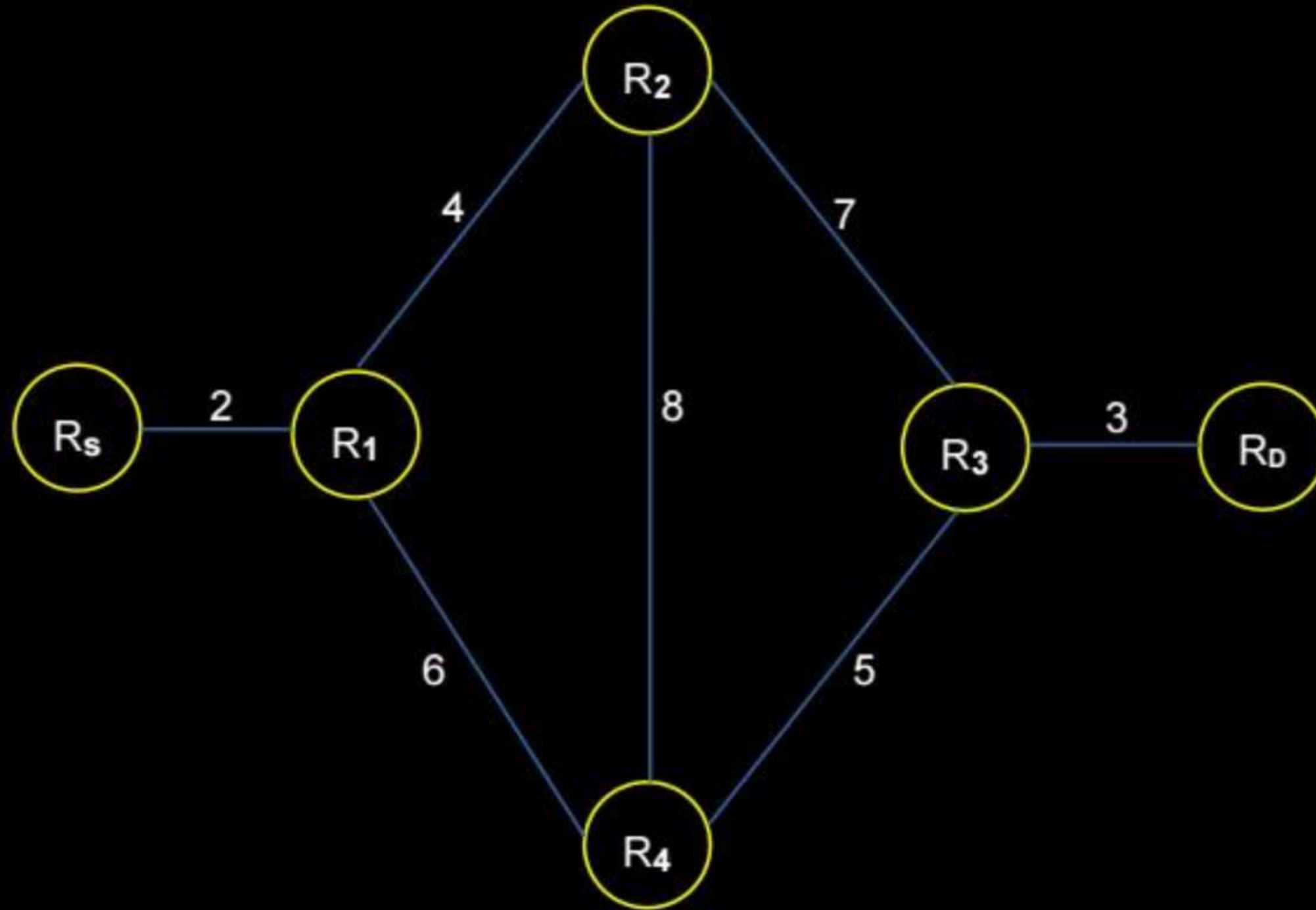


### 3. Share updated Link State information

- Whenever changes occur in the network topology,  
[Relevant routers update their “Link State” information accordingly]
- Only those routers broadcast their updated “Link State” information  
[To all other routers in the network]
- Whenever a router receive any updated “Link State” information,  
it construct new “Adjacency Matrix / List” accordingly and  
re-execute “Dijkstra’s Algorithm” locally



## Topic : Link State Routing





# Topic : Link State Routing



R <sub>s</sub> Link State Info	
R <sub>1</sub>	2

R <sub>D</sub> Link State Info	
R <sub>3</sub>	3

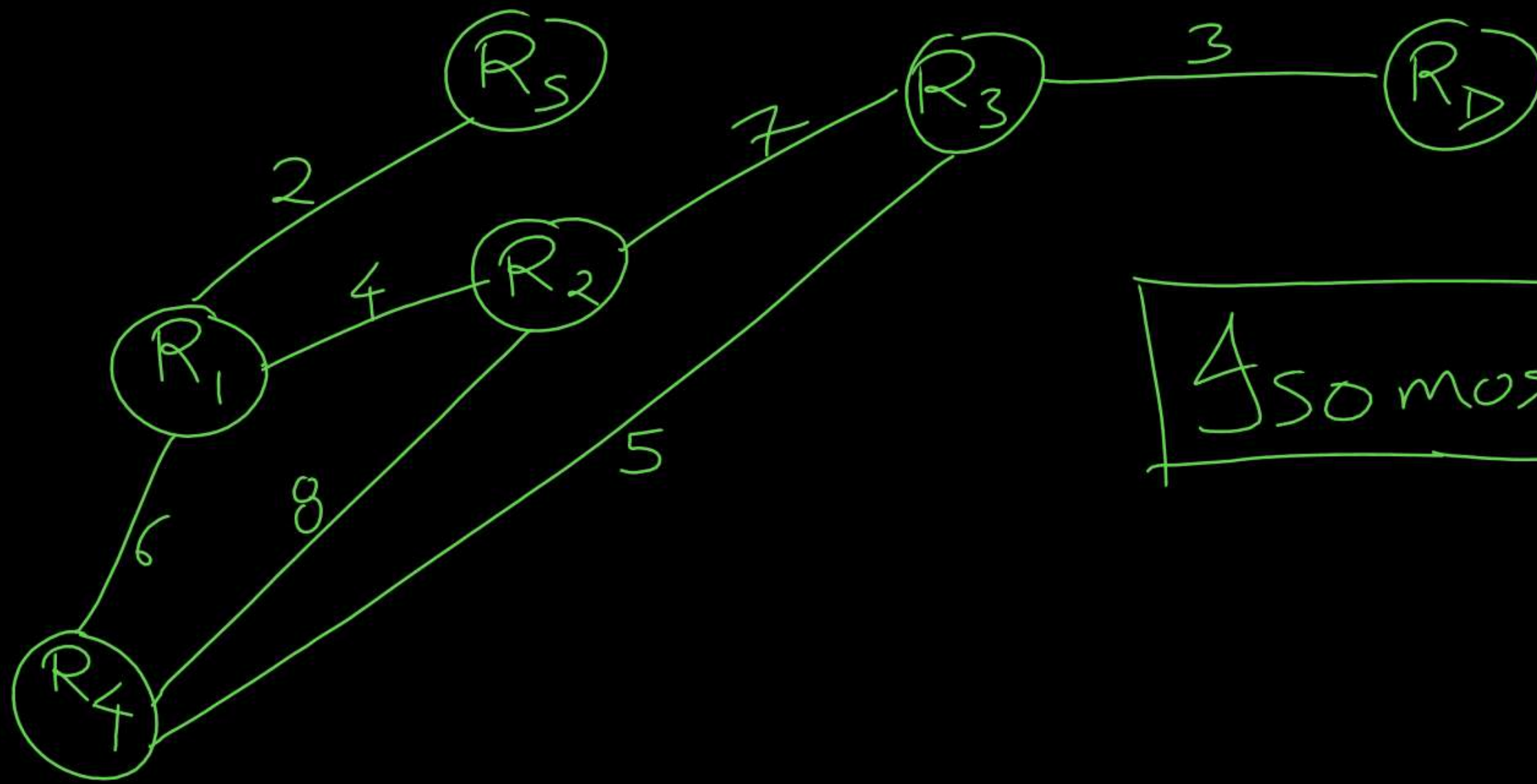
R <sub>1</sub> Link State Info	
R <sub>s</sub>	2
R <sub>2</sub>	4
R <sub>4</sub>	6

R <sub>4</sub> Link State Info	
R <sub>1</sub>	6
R <sub>2</sub>	8
R <sub>3</sub>	5

R <sub>2</sub> Link State Info	
R <sub>1</sub>	4
R <sub>3</sub>	7
R <sub>4</sub>	8

R <sub>3</sub> Link State Info	
R <sub>D</sub>	3
R <sub>2</sub>	7
R <sub>4</sub>	5





Isomorphic



## Topic : Dijkstra's Algorithm



- Single Source Shortest Path Algorithm
- Dijkstra's Algorithm is divided into two steps :
  1. Initialization
  2. Iterative distance calculation



# Topic : Dijkstra's Algorithm



Forwarding Table

Dest.	Next Hop
-------	----------

Source =  $R_1$

Dest.	Cost	Next Hop
$R_1$	0	$R_1$
$R_2$		
$\vdots$		
$R_5$		
$\vdots$		
$R_6$		





## Topic : Dijkstra's Algorithm



Dijkstra's Algorithm time complexity (per node/router)

- Number of nodes = n
- Total n iteration
- Total number of comparisons =  $O(n^2)$



## Topic : Link State Routing



### Message Complexity :

- Number of nodes =  $n$
- Each router must broadcast its Link State information to all other router
- Each router's message crosses  $O(n)$  links
- Overall message complexity =  $O(n^2)$



## Topic : Link State Routing



- High number of packets per network (Disadv.)
- Adapt network topology very fast (Adv.)





## Topic : Distance Vector Routing



→ Based on Bellman-Ford equation  
[Dynamic programming]

→ Bellman-Ford equation

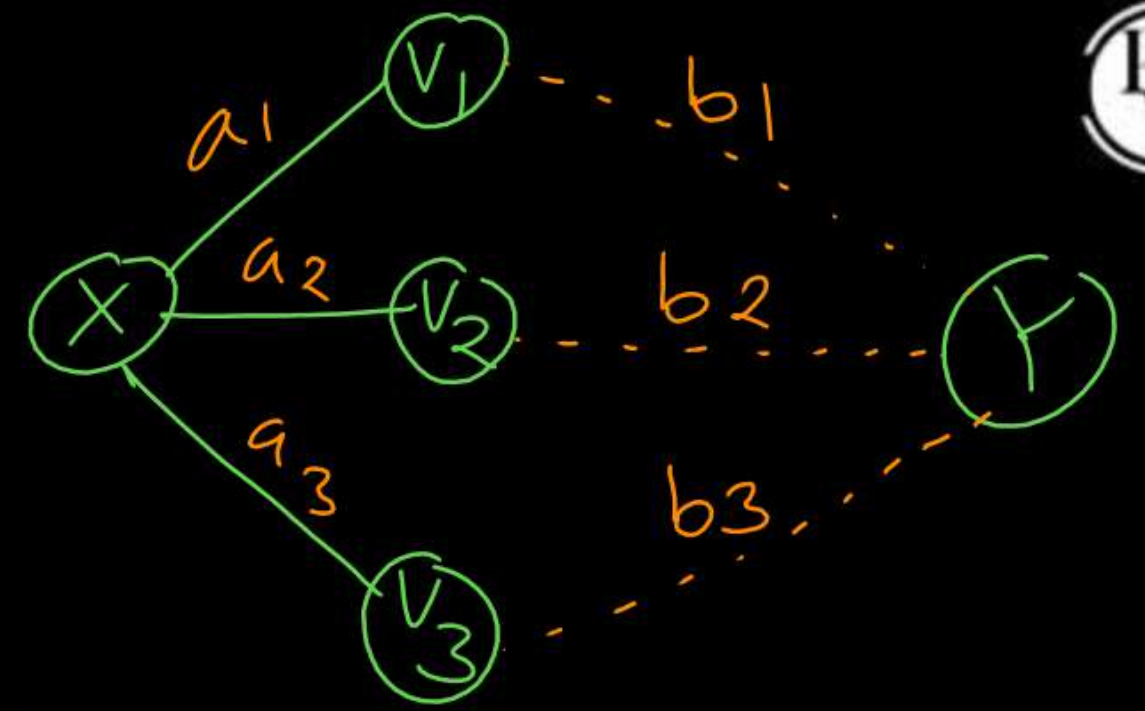
Let  $D_x(Y)$  : Cost of the least-cost path from X to Y

$$D_x(y) = \min_v \{ C(X, V) + D_v(Y) \}$$

→ Minimum taken over all neighbors  $V$  of  $X$

$C(X, V)$  : Direct cost of link from X to V

$D_v(Y)$  : V's estimated least-cost path cost to Y



$$\begin{aligned} D_x(Y) &= \\ &\min \left[ \begin{aligned} &\{ C(X, V_1) + D_{V_1}(Y) \} \\ &\{ C(X, V_2) + D_{V_2}(Y) \} \\ &\{ C(X, V_3) + D_{V_3}(Y) \} \end{aligned} \right] \\ &= \min [a_1 + b_1, a_2 + b_2, a_3 + b_3] \end{aligned}$$





## Topic : Distance Vector Routing

- Each router maintain separate “Distance Vector” estimate  
[Best known minimum distance to all other routers]
- Each router sends their own “Distance Vector” estimate to  
their neighbor routers only
- When a router receives new “Distance Vector” estimate from any neighbor,  
It update its own “Distance Vector” using Bellman-Ford equation



## 2 mins Summary



**Topic**

**Routing**

**Topic**

**Link State Routing** ✓

**Topic**

**Distance Vector Routing**





**THANK - YOU**