

# ADS Homework 7

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## 1 Problem 7.1

### 1.1 a)

Please see the attached file counting\_sort.cpp

### 1.2 b)

Please see the attached file bucket\_sort.cpp

### 1.3 c)

The required pseudo-code is as follows:

```
1 type algorithm(arr[], a, b){
2     m = max(arr);
3     count = [0]*(m+1);
4     for(i:range of arr){
5         count[arr[i]]+=1;
6     }
7     for(i:range of arr){
8         count[i] += count[i-1];
9     }
10    return count[b]-count[a-1];
11 }
12 }
```

### 1.4 d)

Please see the attached file trei\_sort.cpp.

### 1.5 e)

When all the entries go in a single bucket then, the time complexity for bucket sort is worst. In this condition the time complexity of the algorithm depends upon the sub-sorting algorithm used, which in our case is insertion sort. The worst-case complexity of insertion-sort is  $O(n^2)$ , so the worst-case complexity for bucket-sort is also  $O(n^2)$ .

Example:

Array = (0.9, 0.5, 0.1, 0.2, 0.6)

The elements are divide into following buckets:

Bucket[0] = Null

Bucket[1] = Null

Bucket[2] = Null

Bucket[3] = Null

Bucket[4] = (0.9, 0.5, 0.1, 0.2, 0.6)

Bucket[5] = Null

Bucket[6] = Null

Bucket[7] = Null

Bucket[8] = Null

Bucket[9] = Null

Here, the time-complexity depends upon the sorting algorithm for Bucket[4] which is insertion sort in our case.

## 1.6 f)

The required pseudo-code is as follows:

```
1
2 type distance(arr[])
3     return square root of (arr[x]^2 + arr[y]^2)
4
5 type algorithm(arr[[]], size){
6     for(i: 0 to size){
7         flag=0;
8         for(j: 0 to size-i-1){
9             if(distance(A[j]) > distance(A[j+1])){
10                 swap(A[j], A[j+1]);
11                 flag=1;
12             }
13         }
14         if(flag == 0)
15             break;
16     }
17 }
```