

Number Theory

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10:59 AM

① Euclid's algorithm:

$$\gcd(a_0, a_1) = \gcd(a_0 \bmod a_1, a_1)$$

$$= \gcd(a_1, a_0 \bmod a_1)$$

② Extended Euclidean algorithm:

$$\gcd(a_0, a_1) = s a_0 + t a_1$$

↓
by Euclid's algorithm

Idea: compute regular EA

how?

$$\begin{aligned} \gcd(a_0, a_1): \quad a_0 &= q_1 a_1 + a_2 & a_2 &= s_2 a_0 + t_2 a_1 \\ \gcd(a_1, a_2): \quad a_1 &= q_2 a_2 + a_3 & a_3 &= s_3 a_0 + t_3 a_1 \\ &\vdots & &\vdots \\ \gcd(a_{l-1}, a_l): \quad a_{l-1} &= q_l a_l + a_{l+1} & a_{l+1} &= s_{l+1} a_0 + t_{l+1} a_1 = \gcd(a_0, a_1) \\ & & a_{l+1} &= q_l a_l + 0 \end{aligned}$$

∴ s and t are
s_l, t_l here. (d)

how?

eg: $\gcd(973, 301) = s \cdot 973 + t \cdot 301 = 7$

i	a_0	a_1	a_2
2	973	=	3 · 301 + 70

$$a_2 = 70 = [1] 973 + [-3] 301$$

3	301	=	4 · 70 + 21
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$$a_3 = 21 = 301 - 4 \cdot 70$$

we need a multiple mth terms of 973 & 301
sepⁿ → substitute from previous iteration

$$a_3 = 21 = 301 - 4 [973 - 3 \cdot 301]$$

$$a_3 = [-4] a_2 + [13] a_1$$

$$4: \quad 70 = 3 \cdot 21 + 7$$

$$a_4 = 70 - 3 \cdot 21 \quad \text{FROM 2 LINES CITATIONS BACK}$$

$$7 = [1] a_2 + [-3] a_1 - 3 [-4] a_2 + [13] a_1$$

$$7 = [13] a_2 + [-42] a_1$$

in general:

$$s_{i-2} = s_{i-2} a_0 + t_{i-2} a_1$$

$$s_{i-1} = s_{i-1} a_0 + t_{i-1} a_1$$

next iteration:

$$\text{E.A: } s_{i-2} = q_{i-1} s_{i-1} + s_i$$

$$s_i = s_{i-2} - q_{i-1} s_{i-1}$$

substitute back

$$s_i = [s_{i-2} - q_{i-1} s_{i-1}] a_0 + [t_{i-2} - q_{i-1} t_{i-1}] a_1$$

$$s_i = s_i a_0 + t_i a_1$$

Recursive Formulae

$$s_i = s_{i-2} - q_{i-1} s_{i-1}, \quad i \geq 2$$

$$t_i = t_{i-2} - q_{i-1} t_{i-1}, \quad i \geq 2$$

$$\text{where } \begin{matrix} s_0 = 1 & t_0 = 0 \\ s_1 = 0 & t_1 = 1 \end{matrix}$$

MAIN APPLICATION OF EEA is

COMPUTING OF INVERSES Mod n

problem: $a^{-1} \equiv ? \pmod{n}$

$$a^{-1} a \equiv 1 \pmod{n} \quad (\text{by def}^n)$$

$$\gcd(n, a) = 1 \quad \left[\begin{array}{l} \text{if } \exists a^{-1}, \text{ then} \\ \gcd(n, a) \text{ must be } 1 \end{array} \right]$$

$$= s \cdot n + t \cdot a \quad (\text{by EEA})$$

$$1 = s \cdot n + t \cdot a$$

take mod n both sides

$$1 = s \cdot n + t \cdot a$$

take mod n both sides

$$1 \bmod n = (s \cdot n + t \cdot a) \bmod n$$

$$1 = 0 + t \cdot a \bmod n$$

$$\Rightarrow ta \equiv 1 \bmod n$$

t is actually a^{-1}
 \hookrightarrow parameter of the smaller number

SOME THEOREMS:

(1) EULER'S PHI FUNCTION:

$$\mathbb{Z}_m = \{0, 1, \dots, m-1\}$$

$$\begin{aligned} \gcd(0, m) &= m \\ (1, m) &= \\ (2, m) &= \\ \vdots & \\ (m-1, m) &= \end{aligned} \quad \left. \begin{array}{l} \text{how} \\ \text{many are} \\ \text{coprime to} \\ m? \end{array} \right\}$$

eg: $m = 6$

$$\mathbb{Z}_6 = \{0, 1, 2, 3, 4, 5\}$$

$$\begin{aligned} \gcd(0, 6) &= 6 \\ (1, 6) &= 1 \quad \checkmark \\ (2, 6) &= 2 \\ (3, 6) &= 3 \\ (4, 6) &= 2 \\ (5, 6) &= 1 \quad \checkmark \end{aligned} \quad \phi(6) = 2$$

but this BRUTE FORCE
 is not viable for
 higher m .

egⁿ:

$$m = p_1^{e_1} \cdot p_2^{e_2} \cdot \dots \cdot p_n^{e_n}$$

$p_i \rightarrow$ distinct prime numbers
 $e_i \rightarrow$ +ve integers

$$\phi(m) = \prod_{i=1}^n (p_i^{e_i} - p_i^{e_i-1})$$

eg: $m = 240$
 $\phi(240) = ?$

$$\begin{aligned} m &= 16 \cdot 15 \\ &= 2^4 \cdot 3 \cdot 5 \end{aligned}$$

$$\begin{aligned}
 \phi(240) &= \prod_{i=1}^3 (p_i^{e_i} - p_i^{e_i-1}) \\
 &= (2^4 - 2^3)(3^1 - 3^0)(5^1 - 5^0) \\
 &= 8 \cdot 2 \cdot 4 \\
 &= 64
 \end{aligned}$$

(2)

FERMAT'S LITTLE THEOREM

$a \rightarrow \text{integer}$, $p \rightarrow \text{prime}$

$$a^p \equiv a \pmod{p}$$

EULER'S THEOREM

$a, m \rightarrow \text{integers}$ s.t. $\gcd(a, m) = 1$
then:

$$a^{\phi(m)} \equiv 1 \pmod{m}$$