

Elliptic Curve Cryptography

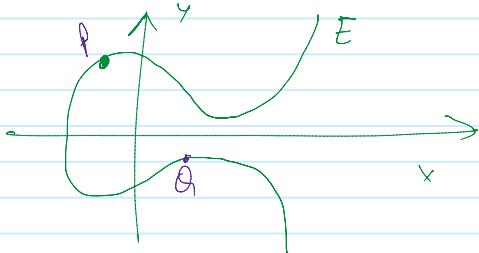
28 November 2019 17:32

Can we find another cyclic group in which the DLP is difficult? Ideally more difficult than in \mathbb{Z}_p ?

for use in crypto we need to consider polynomials over \mathbb{Z}_p .

DEFINITION → The EC over \mathbb{Z}_p , $p > 3$ is the set of all pairs $(x, y) \in \mathbb{Z}_p^2$:
 $y^2 = x^3 + ax + b \pmod{p}$
 together w/ an imaginary point at infinity \mathcal{O} , where $a, b \in \mathbb{Z}_p$
 for $4a^3 + 27b^2 \neq 0 \pmod{p}$

e.g.: $y^2 = x^3 - 3x + 3$ over \mathbb{F}_R



Note: symmetry wrt to x-axis

$$y^2 = x^3 + ax + b$$

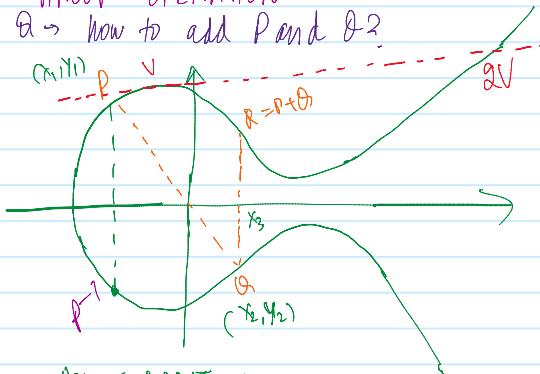
$$y = \pm \sqrt{x^3 + ax + b}$$

for a DLP we need a cyclic group. For a group we need

- (a) a set of elements
- (b) a group operation that fulfills the group laws

GROUP OPERATION

Q: how to add P and Q?



POINT ADDITION

$$P + Q = R$$

POINT DOUBLING?

$$\text{if } P = Q \Rightarrow P + P = 2P$$

ELLIPTIC CURVE LOGARITHMIC PROBLEM

We obtain immediately a discrete logarithmic problem:

Definition 9.2.1 Elliptic Curved Discrete Logarithm Problem (ECDLP)

Given is an elliptic curve E . We consider a primitive element P and another element T . The DL problem is finding the integer d , where $1 \leq d \leq \#E$, such that:

$$\underbrace{P + P + \dots + P}_{d \text{ times}} = dP = T.$$

↑ two of steps I
made in the worse
case from
starting point (9.2)
P to end at T.

e.g.: $P = (2, 5)$ generator
 $T = (16, 4) = dP$.
 $\underline{\text{ie}} \quad (16, 4) = d(2, 5) \quad d?$
 This is NOT EASY TO FIND.

Notes: (1) the ECDLP $d (= k_{p\text{prv}})$ is an integer but $T = (k_{p\text{pub}})$ is point on curve \in a group element.

Q: group cardinality of E ??

Theorem 9.2.2 Hasse's theorem

Given an elliptic curve E modulo p , the number of points on the curve is denoted by $\#E$ and is bounded by:

$$p+1-2\sqrt{p} \leq \#E \leq p+1+2\sqrt{p}.$$

finding exact $\#E$ is computationally difficult.

All EC protocols rely on the hardness of the ECDLP

If the EC is chosen carefully, the best known algorithm for computing the ECDLP requires $\approx \sqrt{p}$ steps.

EC DIFFIE-HELLMAN KEY EXCHANGE (ECDH)

→ straightforward $aP \rightarrow bP \rightarrow D \in \mathbb{Z}_p$

I Phase set up:

$$E: y^2 \equiv x^3 + ax + b \pmod{p}$$

primitive element $P = (x_p, y_p)$

II Phase protocol

Alice $\leftarrow E, p \rightarrow \text{Bob}$

$$a = k_{PA} \in \{2, 3, \dots, \#E-1\} \quad b = k_{PB} \in \{2, 3, \dots, \#E-1\}$$

$$P+Q = R$$

POINT DOUBLING?

$$\Leftrightarrow \text{what is } V+V = 2V$$

$\Theta \Rightarrow$ Analytical expression for the group operation?

Idea: Given $E: y^2 \equiv x^3 + ax + b$,
 $P = (x_1, y_1)$, $Q = (x_2, y_2)$.

1) Find the line PQ :

$$\text{lets say } y = Sx + M$$

$$\text{then } "S = E"$$

$$(Sx + M)^2 \equiv x^3 + ax + b$$

degree 3 $\Rightarrow 3$ solutions,

but we already have 2

solutions namely x_1, x_2 .

Elliptic Curve Point Addition and Point Doubling

$$x_3 = s^2 - x_1 - x_2 \bmod p$$

$$y_3 = s(x_1 - x_3) - y_1 \bmod p$$

where

$\rightarrow (y_2 - y_1)(x_2 - x_1)^{-1}$ from extended euclidean algorithm (d)

$$s = \begin{cases} \frac{y_2 - y_1}{x_2 - x_1} \bmod p & ; \text{if } P \neq Q \text{ (point addition)} \\ \frac{3x_1^2 + a}{2y_1} \bmod p & ; \text{if } P = Q \text{ (point doubling)} \end{cases}$$

$$\begin{array}{l} \text{Alice: } a = k_{\text{pub}, A} \in \{2, 3, \dots, \#E-1\} \quad b = k_{\text{pub}, B} \in \{2, 3, \dots, \#E-1\} \\ A = k_{\text{pub}, A} P \\ = a \cdot P \\ = (x_A, y_A) \end{array} \quad \begin{array}{l} \text{Bob: } B = k_{\text{pub}, B} \cdot bP \\ = (x_B, y_B) \end{array}$$

$$\begin{array}{c} A \xrightarrow{\quad} \\ aB = (x_{AB}, y_{AB}) \end{array} \quad \begin{array}{c} B \xrightarrow{\quad} \\ bA = (x_{AB}, y_{AB}) \end{array}$$

USE CASE:

MESSAGE M:

$$\text{AES}(M)$$

$$\times_{AB}$$

$$\xrightarrow{\quad C \quad} \text{AES}_{X_{AB}}^{-1}(C)$$

Proof of correctness:

Alice computes:

$$a \cdot b \leq a(bP) = abP$$

Bob computes:

$$bA = b(aP) = abP$$

Alice

$$\text{choose } a = k_{\text{pr}, A} = 3$$

$$A = k_{\text{pub}, A} = 3P = (10, 6)$$

Bob

$$\text{choose } b = k_{\text{pr}, B} = 10$$

$$B = k_{\text{pub}, B} = 10P = (7, 11)$$

$$T_{AB} = aB = 3(7, 11) = (13, 10)$$

$\xrightarrow{\quad A \quad}$

$$T_{AB} = bA = 10(10, 6) = (13, 10)$$

$\Theta \Rightarrow$ how to compute aP \leq $P+P+\dots+P$

$\xrightarrow{\quad a \text{ times} \quad}$

$$\text{eg: } 26P = \overbrace{(11010)_2}^{\text{LEFT TO RIGHT}} P$$

Step:

$$\emptyset \quad P = 1_2 P$$

$$1_A \quad P + P = 2_2 P = 10_2 P \quad D$$

$$1_B \quad 2P + P = 3P = 11_2 P \quad A$$

$$2_A \quad 3P + 3P = 6P = 110_2 P \quad D$$

$$3_A \quad 6P + 6P = 12P = 1100_2 P \quad D$$

$$3_B \quad 12P + P = 13P = 1101_2 P \quad A$$

$$4_A \quad 13P + 13P = 11010_2 P \quad D$$

\approx

$\Theta \Rightarrow$ as per the group laws, what is the NEUTRAL ELEMENT.

$$P + Z = P \quad \forall P \in E$$

We define a "point at infinity" \mathcal{O}

property 3 of DEF 4.37

$$P + \mathcal{O} = P \quad \forall P \in E$$

$$\text{property 4} \rightarrow P + (-P) = \mathcal{O}$$

$$\Rightarrow -P \text{ of } P = (x, y) \text{ is}$$

by definition

$$-P \equiv (x, -y)$$

\checkmark
the INVERSE
of group
operation

the MINUS
we know
(just mod p)

example:

$$E: y^2 \equiv x^3 + 2x + 2 \bmod 17$$

we want to double $P = (5, 1)$

$$2P = P + P = (5, 1) + (5, 1) = (b_1, b_2)$$

$$S = \frac{3x_1^2 + a}{2y_1} = (2 \cdot 5)^{-1} (5, 5 + 2) \pmod{17}$$

$$= 2^{-1} \cdot 9 \pmod{17}$$

$$= 9 \cdot 9 \pmod{17}$$

$$= 13$$

$$x_2 = s^2 - x_1 - x_2 = 13^2 - 5 - 5 \pmod{17}$$

$$= 169 - 10 \pmod{17}$$

$$= 13$$

$$\begin{aligned}x_3 &= s^2 - x_1 - x_2 = 13^2 - 5 - 5 \bmod 17 \\&= 159 \bmod 17 \\&= 6 \bmod 17\end{aligned}$$

$$\begin{aligned}y_3 &= s(x_1 - x_3) - y_1 = 13(5-6) - 1 \\&= -14 \bmod 17 \\&\equiv 3 \bmod 17\end{aligned}$$

note: $(6, 3)$ lies on the curve

The points on an elliptic curve including \mathcal{O} , have CYCLIC SUBGROUPS.
under certain conditions all points
on EC form a CYCLIC GROUP.

Example 9.5. We want to find all points on the curve:

$$E : y^2 \equiv x^3 + 2 \cdot x + 2 \pmod{17}$$

It happens that all points on the curve form a cyclic group and that the order is $\#E = 19$. For this specific curve the group order is a prime and, according to Theorem 8.2.4, every element is primitive.

As in the previous example we start with the primitive element $P = (5, 1)$. We compute now all "powers" of P . More precisely, since the group operation is addition, we compute $P, 2P, \dots, \#E P$. Here is a list of the elements that we obtain:

$$\begin{array}{ll}2P = (5, 1) + (5, 1) = (6, 3) & 11P = (13, 10) \\3P = 2P + P = (10, 6) & 12P = (0, 11) \\4P = (3, 1) & 13P = (16, 4) \\5P = (9, 16) & 14P = (9, 1) \\6P = (16, 13) & 15P = (3, 16) \\7P = (0, 6) & 16P = (10, 11) \\8P = (13, 7) & 17P = (6, 14) \\9P = (7, 6) & 18P = (5, 16) \\10P = (7, 11) & 19P = \mathcal{O}\end{array}$$

From now on, the cyclic structure becomes visible since:

$$\begin{aligned}20P &= 19P + P = \mathcal{O} + P = P \\21P &= 2P \\&\vdots\end{aligned}$$

INFINITY POINT

It is also instructive to look at the last computation above, which yielded:

$$18P + P = \mathcal{O}.$$