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### File Structure.

A file structure should be according to a required format that the OS can understand.

- A file has a certain defined structure according to its type.
- A Text file is a sequence of characters organised into lines.
- A source file is a sequence of procedures & functions.
- An Object file is a sequence of bytes organised into blocks that are understandable by the Machine.
- When OS define different file structure, it also contains the code to support this file structure. Unix, MSdos support minimum number of file structure.

### File Type.

File type refers to the ability of the OS to distinguish different types of file such as text files, source files & binary files etc.



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Many OS support many types of files. OS like MS dos & Unix have the following types of files

1. Ordinary files
  - These are the files that contain user information
  - These may have text, databases or executable programs.
  - The user can apply various operations on such files like add, modify, delete, or even remove the entire file
2. Directory files
  - These files contain list of file names & other information related to these files

3. Special files
  - These files are also known as device files
  - These files represent physical resource like disc, terminals, printers, networks, tape drive etc.

Special files are of 2 types

1. Character Special files

Data is handled character by character as in case of terminals or printers

2. Block Special files

Data is handled in blocks as in case of disc & tapes.



User and System programmer view of file system.

## 1. User view

The user view of the computer refers to the interface being used. Such systems are designed for one user to monopolize its resources, to maximize the work that the user is performing in that case, the OS is designed mostly for ease of use, with some attention paid to performance, and men paid to resource utilization.

## 2. System view

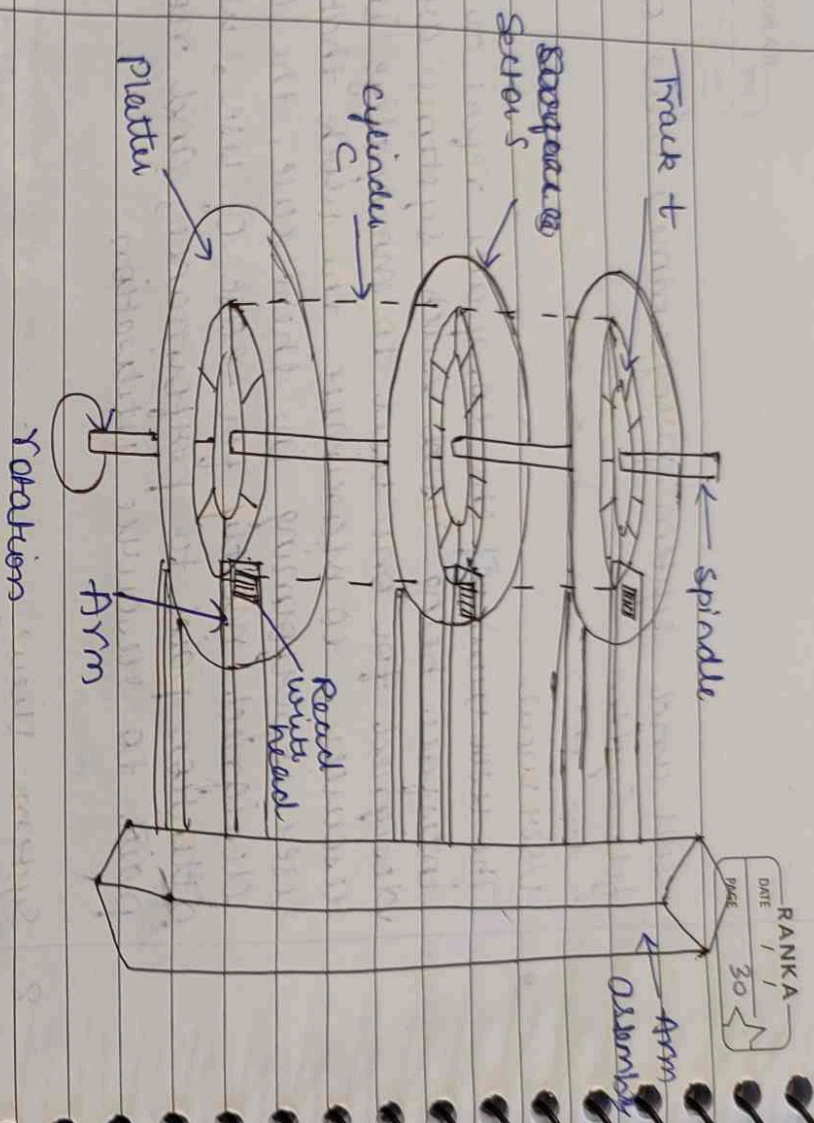
OS can be viewed as a resource allocator also. A computer system consist of many resources like hardware & software that must be managed efficiently. The OS act as the manager of the resources, receives & manages conflicting requests, control execution of programs etc.

## DISC ORGANISATION

A hard disc is a memory storage device which looks like the below picture.

The disc is divided into tracks each track is further divided into sectors. The point to note here is that outer tracks are bigger in size than the inner tracks but they contain the same number of sectors and have





### Hard disc

equal storage capacity this is because the storage density is high in sectors of the inner tracks whereas the bits are sparsely arranged in sectors of the outer tracks. Some space of every sector is used for formatting so the actual capacity of a sector is less than the given capacity.

→ Read write (R-W) head moves over the rotating hard disc. It is this read write head that performs all the read & write operations on the disc so the position of the read write head is a major concern. To perform a read or write operation on a memory location we need to place the



R-W head Over that position

Some important terms are

1. Seek time

The time taken by the R-W head to reach the desired track from its current position

2. Rotational latency

Time taken by the sector to come under the R-W head

3. Data transfer time

Time taken to transfer the required amount of data. It depends upon the rotational speed.

4. Controller time

The processing time taken by the controller

5. Average access time

Seek time + average rotation latency + Data transfer time + Controller time

## DIFFERENT MODULES OF A FILE SYSTEM

The basic file system level

works directly with the device drivers in terms of retrieving & storing raw blocks of data, without any consideration for what is in each block. Depending on the system blocks may be referred to with a single block number or with head sector cylinder combination



The file organisation module.

Knows about files and their logical blocks and how they map to physical blocks on the disc. In addition to translating from logical to physical blocks, the file organisation module also maintains the list of free blocks and allocates free blocks to file as needed.

The logical file System

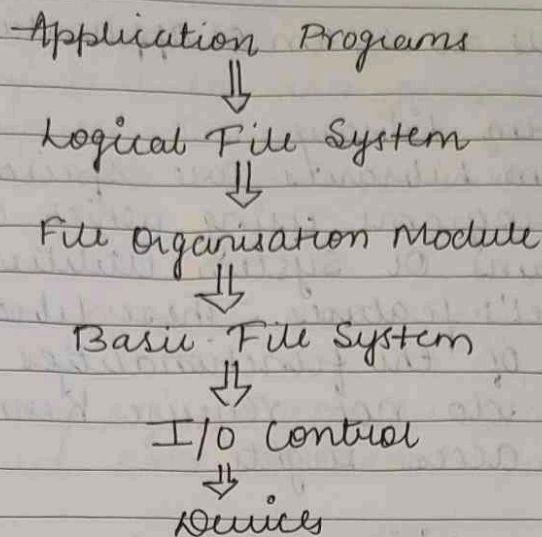
Deals with all of the meta data associated with the file i.e. everything about the file except the data itself. This level manages the directory structure and the mapping of file names to file control block (FCB), FCB's which contain all of the meta data as well as block number information for finding the data on the disc.

The layered approach to file systems

means that much of the code can be used uniformly for a wide variety of different file systems, and only certain layers need to be file system specific.

Common file systems in use include the Unix file system, Windows's file system etc.





### Layered File System.

## FILE SYSTEM IN LINUX AND WINDOWS

### • ~~File~~ LINUX

Linux is one of popular version of UNIX OS. It is open source as its source code is freely available. It is free to use. LINUX was designed considering UNIX compatibility its functionality list is quite similar to that of UNIX.

### Components of LINUX

Linux OS has primarily 3 components

#### 1. Kernel

Kernel is the core part of LINUX it is responsible for all major activities of this OS. It consists of various modules and it interacts directly with the underlying hardware. Kernel provides the required abstraction to high low level hardware.



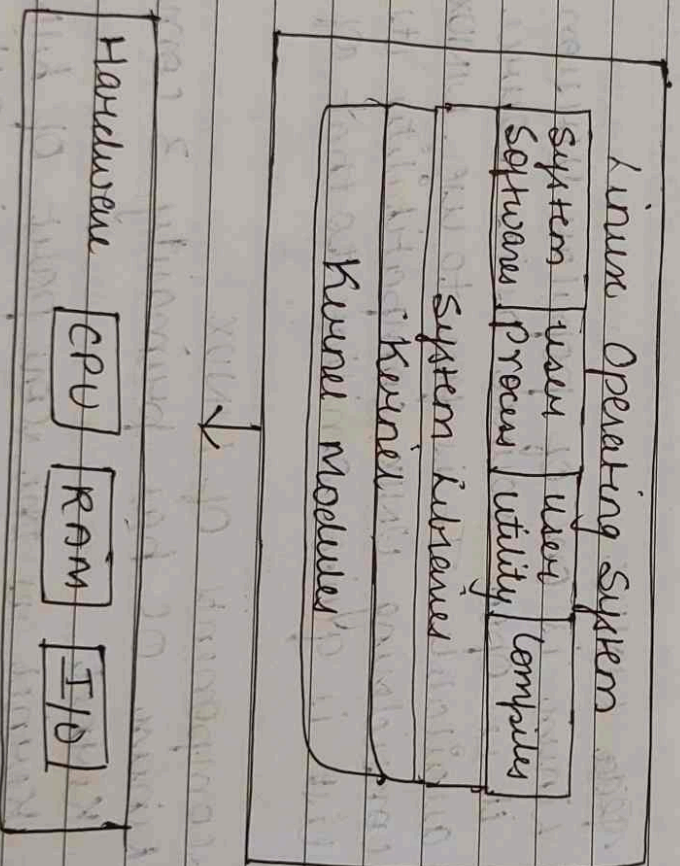
relates to system or application programs.

2.

System library  
System libraries are special functions or programs using which application programs or system utilities ~~access~~ access kernel's features. These libraries implement most of the functionalities of the OS and do not require kernel module's code access rights.

3

System Utility  
System utility programs are responsible to do specialised, individual level tasks.



Linux Operating System.



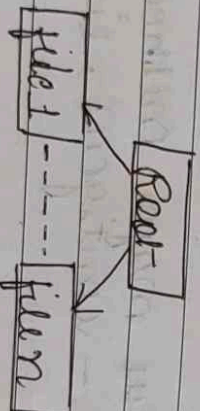
## DIRECTORY SYSTEM

A directory is a location for storing files on your computer. Directories are found in a hierarchical file system such as LINUX, MS-DOS, UNIX etc.

1. A collection of notes containing information about all files
2. A directory system can be classified into single level and hierarchical directory system

Single level directory system

In this type of directory system, there is a root directory which has all files it has directories. Advantage of SIDS is that it is easy to find a file in the directory. this type of directory system is used in cameras & phones.



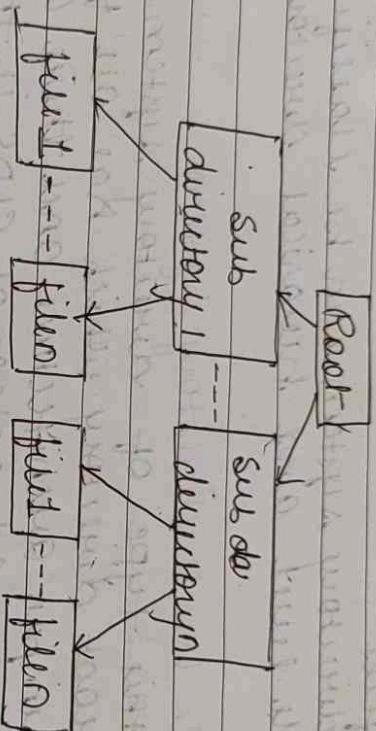
## Directory System

### Hierarchical Directory System

In a hierarchical directory system, files are grouped together to form a subdirectory at the top of the hierarchy is the root directory & then there are sub directories which has files



Advantage of HDS is that users can be provided access to a sub directory rather than the entire directory. It provides a better structure to file system also managing millions of files is easy with this architecture. Personal computers use hierarchical directory system for managing files.



### Hierarchical Directory system

#### DISC SPACE ALLOCATION METHOD

There are three major methods of storing files on disc - contiguous, linked and indexed.

#### Contiguous Allocation

- contiguous allocation requires that all blocks of a file be copped together contiguously.
- Performance is very fast because reading successive blocks of the same file generally requires no movement of the disc.



Fragmentation is loss of memory

RANK A  
DATE / /  
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heads or almost one small step to the next adjacent cylinder.

- Storage allocation involves the same issues for the allocation of contiguous blocks of memory (first fit, fragmentation problems etc.)
  - The distinction is that the high time penalty required for moving the disc head from spot to spot may now justify the benefits of keeping files contiguously when possible.
- Even file systems that do not by default store files contiguously can benefit from certain utilities that compact the disc and make all files contiguous in the process.
- Problems can arise when files grow or if the exact size of a file is unknown at creation time.
- Overestimation of the files final size increases external fragmentation and waste disc space.
- Underestimation of the files may require that a file be moved or a pointer aborted if the file grows beyond its originally allocated space.
- At a file grows slowly over a long time period and the total final space must



be allocated initially, then a lot of space becomes unusable before the file fills the space.

- A variation is to allocate file space in large contiguous runs, called extents when a file outgrows its original extent then an additional one is allocated for eg:- An extent may be the size of a cylinder-track or even cylinders, assigned on an appropriate track or cylinder boundary the high performance file system variably uses extents to optimize performance,

Count-				
0	1	2	3	
mail				
4	5	6	7	
list				
8	9	10	11	
block				
12	13	14	15	
16	17	18	19	
Image				
20	21	22	23	
24	25	26	27	
Page				
28	29	30	31	

file	start	length
count	0	2
mail	6	1
list	10	2
block	14	4
Image	22	5

Contiguous Allocation of Disc Space.



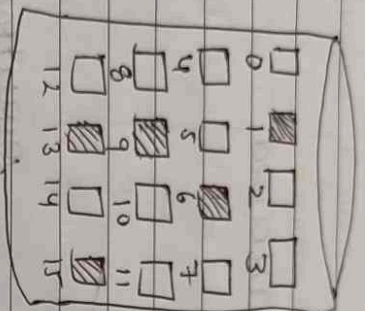
Advantage :-

- easy to implement
- excellent read performance and short access

Disadvantage :-

- since will become fragmented / fragmentation occurs
- difficult to grow file.

Linked Allocation



Directory	
file	start
A	1
B	6
C	9
D	15

- Since files can be stored as linked list, with the expense of the storage space consumed by each link

- Linked Allocation involves no external fragmentation does not require pre known file size & allows files to grow dynamically at any time

- Unfortunately linked allocation is only efficient for sequential access files as random access requires starting at the beginning of the list for each new location access