

Binary Tree and Traversal

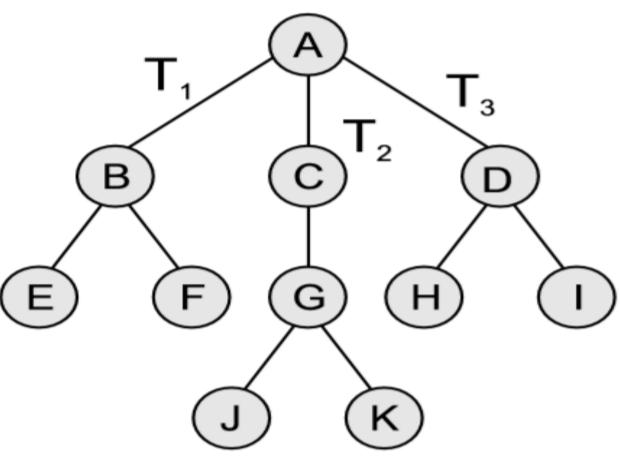
BY

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Trees

Root node





Trees

- Connected acyclic undirected graph.
- A Graph in which any two vertices are connected by exactly one path.



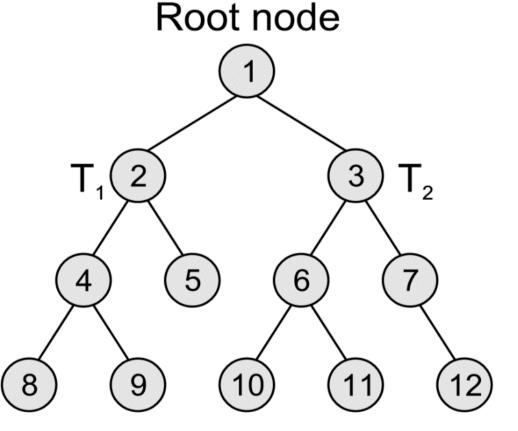
Trees

- Root: Root node R is the topmost node in the tree.
- Leaf node: A node that has no children is called the leaf node or the terminal node.
- **Sub-trees:** If the root node R is not NULL, then the trees T1, T2, and T3 are called the sub-trees of R.
- Path: A sequence of consecutive edges is called a path.
- Ancestor node: An ancestor of a node is any predecessor node on the path from root to that node.



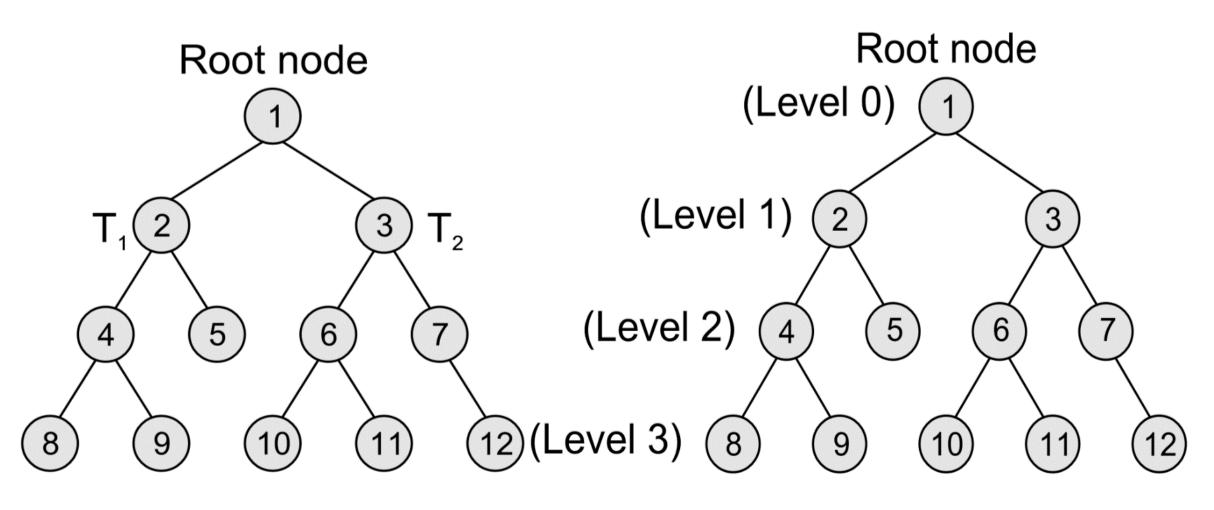
Binary Tree

- DS defined as a collection of elements called nodes.
- Topmost element is called the **root node**, and each node has 0, 1, or at the most 2 children.





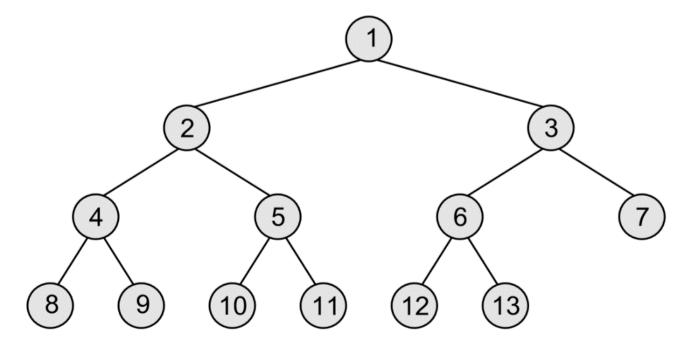
Binary Trees





Complete Binary Tree

- Binary Tree that satisfies these two properties:
- Every level, except possibly the last, is completely filled.
- All nodes appear as far left as possible.
- In a complete binary tree T_n, there are exactly n nodes and level r of T can have at most 2^r nodes.



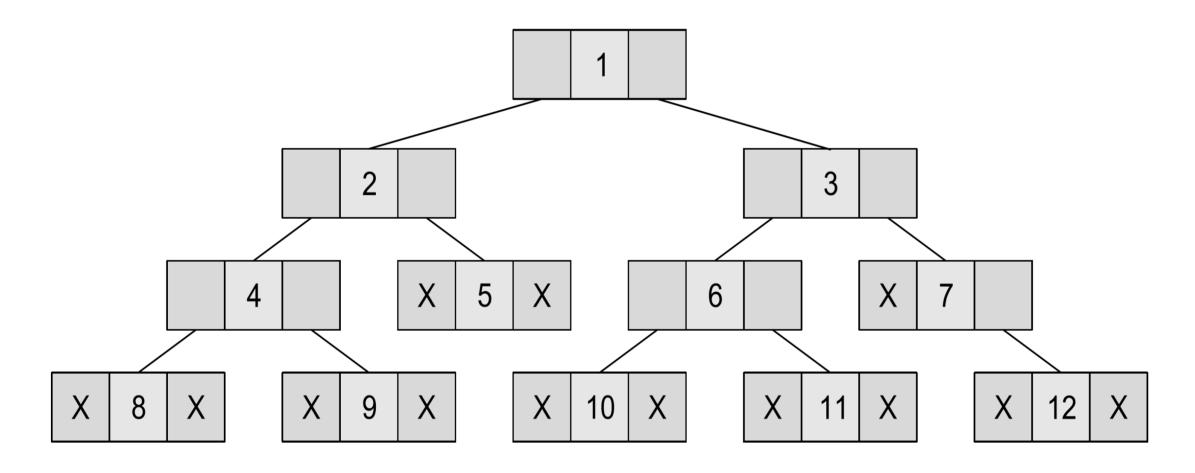


Binary Tree Representation

```
struct node {
    struct node *left;
    int data;
    struct node *right;
};
```

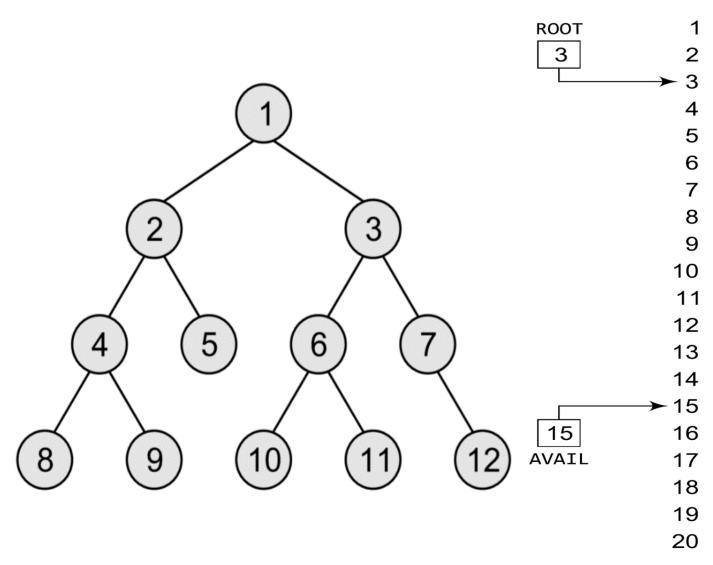


Binary Tree Representation





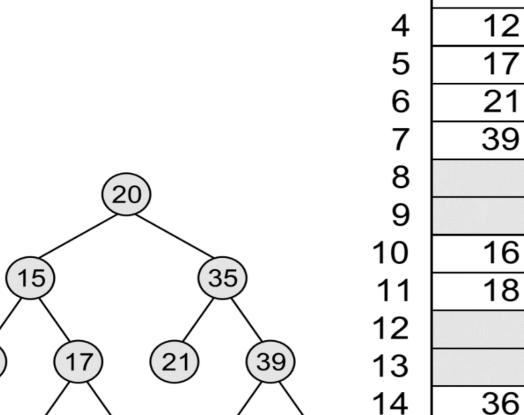




LEFT	DATA	RIGHT
–1	8	–1
–1	10	–1
5	1	8
9	2	14
20	3 4	11
1	4	12
–1	7	18
–1	9	–1
–1	5	–1
–1	11	–1
–1	12	–1
2	6	16

Binary Tree Representation

- 1D array.
- Root is stored in the first location
 - Here it is TREE[1]
- The children of a node stored in location K will be stored in locations (2 \times K) and (2 \times K+1)
- Maximum size of the array TREE is given as (2^h-1), where h is the height of the tree.



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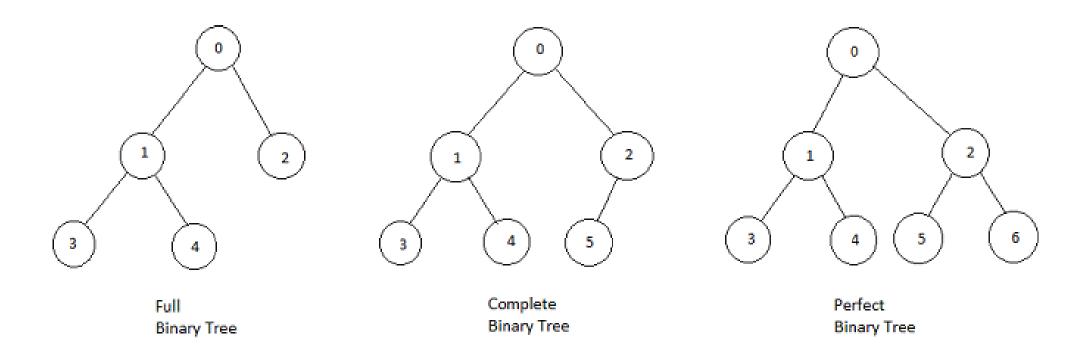
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Information Technology

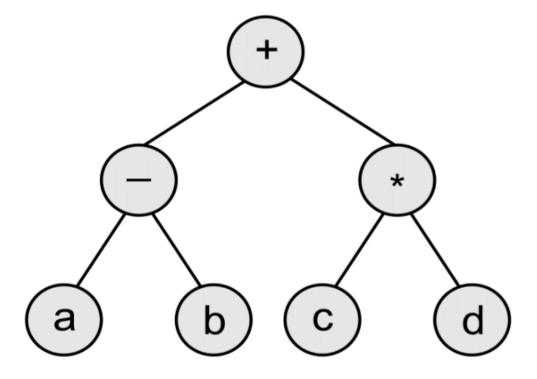


Binary Tree



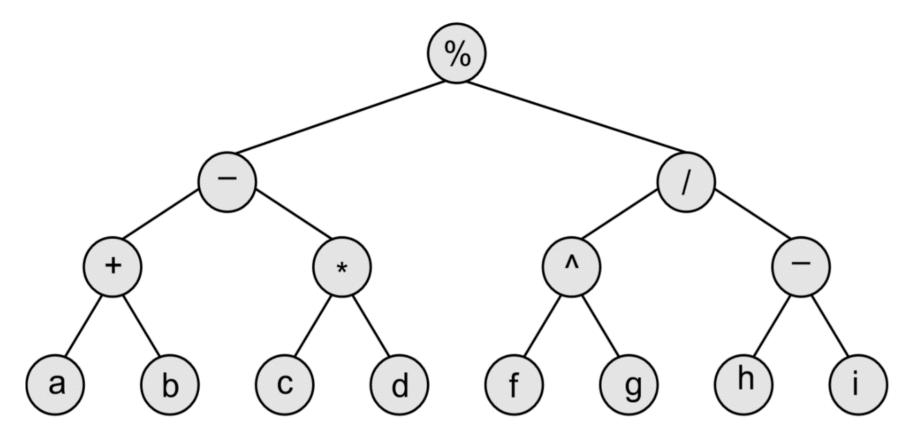


- Binary Tree can be used to store algebraic expression.
- (a b) + (c * d)



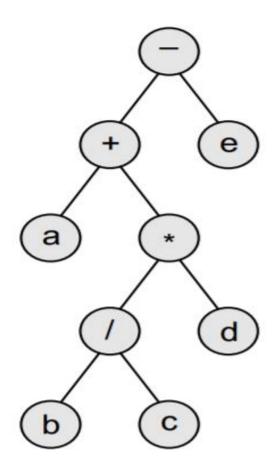


• $((a + b) - (c * d)) % ((f ^ g) / (h - i))$

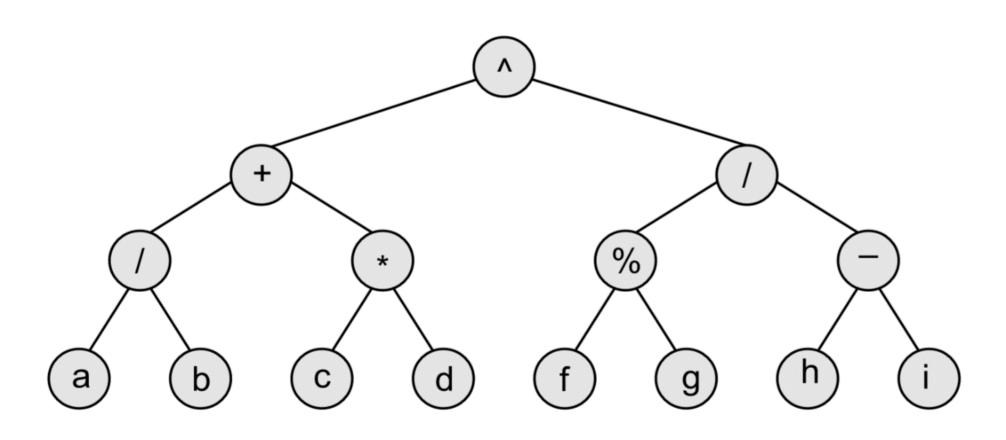




• a + b / c * d - e



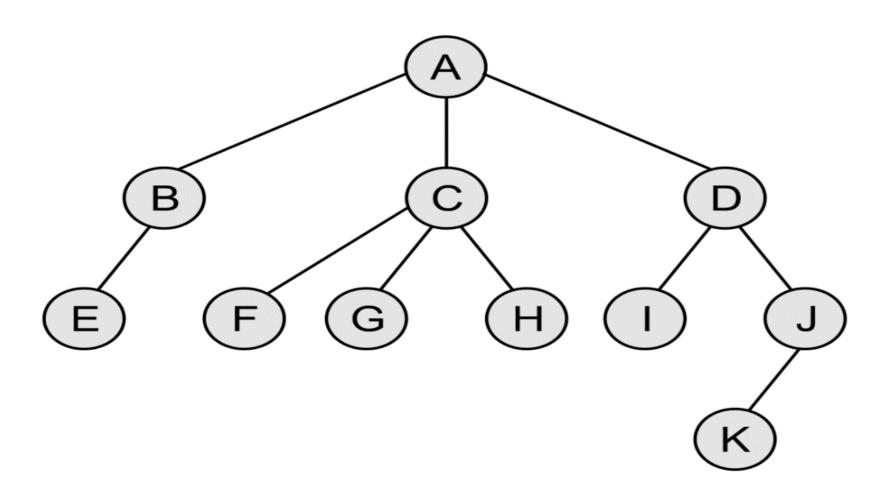




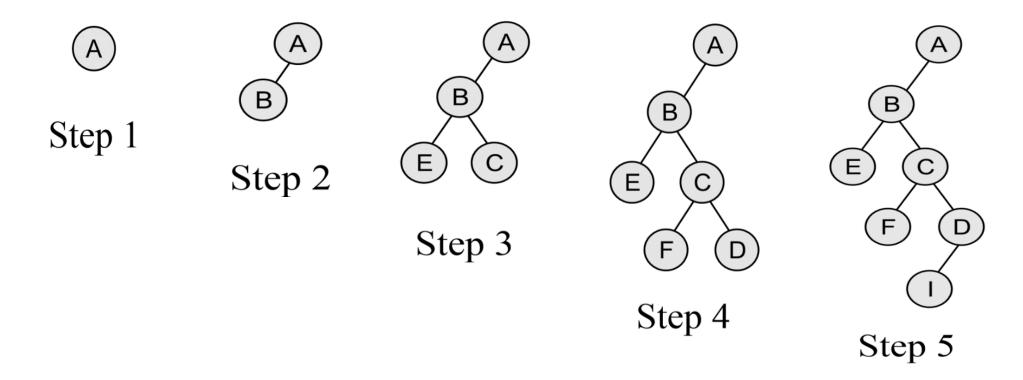


- Rule 1: Root of the binary tree = Root of the general tree
- Rule 2: Left child of a node = Leftmost child of the node in the binary tree in the general tree
- Rule 3: Right child of a node in the binary tree = Right sibling of the node in the general tree

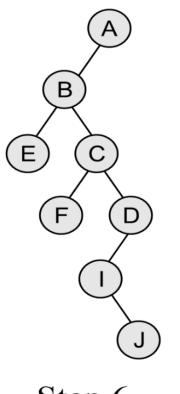




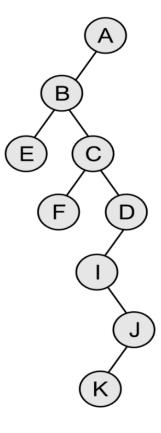




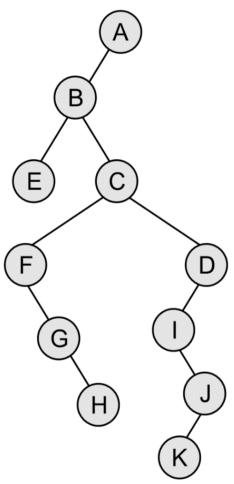




Step 6



Step 7



Step 8



Binary Tree Traversal

 Process of visiting each node in the tree exactly once in a systematic way.

Nonlinear data structure can be traversed in many different ways.

Algorithms differ in the order in which the nodes are visited.

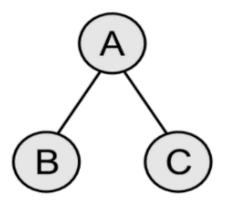


- Following operations are performed recursively at each node:
 - 1. Visiting the root node.
 - 2. Traversing the left sub-tree.
 - 3. Traversing the right sub-tree.



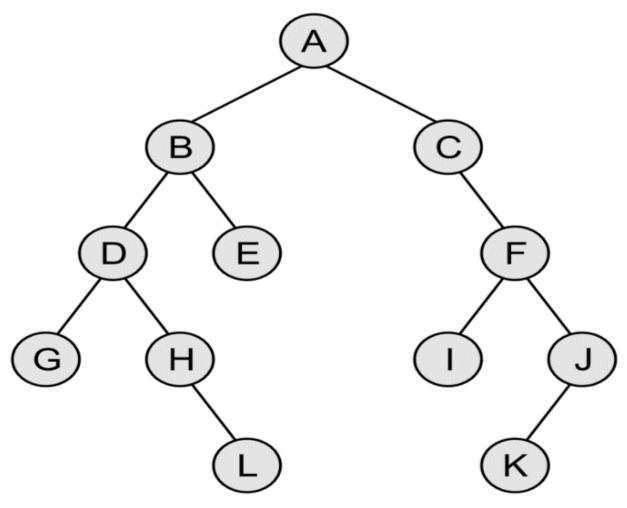
• Pre-order Traversal: A, B, C

• Depth-first traversal.

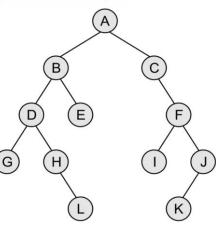


- Root node is accessed prior to any other nodes in the left and right sub-trees.
- NLR traversal algorithm (Node-Left-Right).





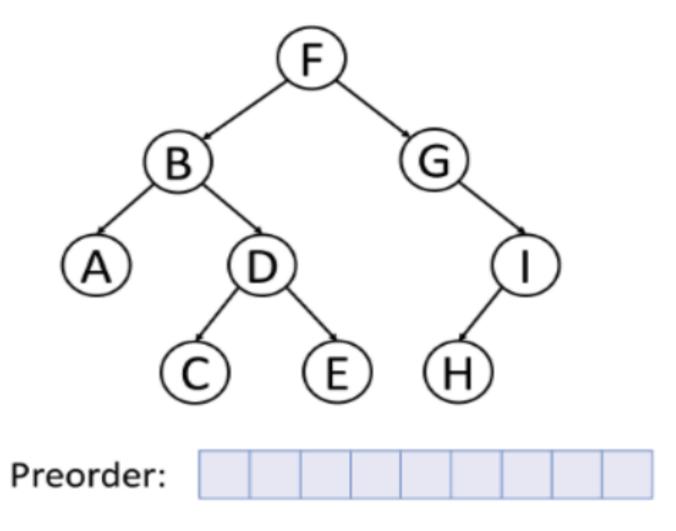




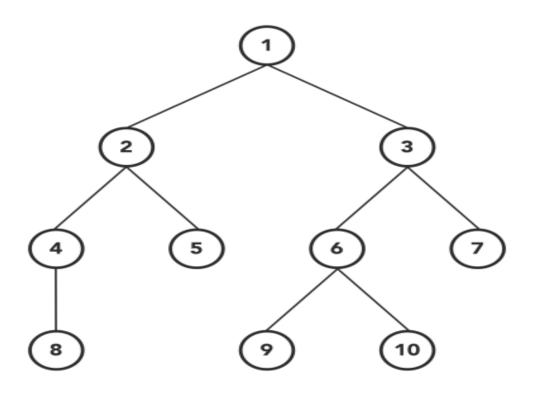




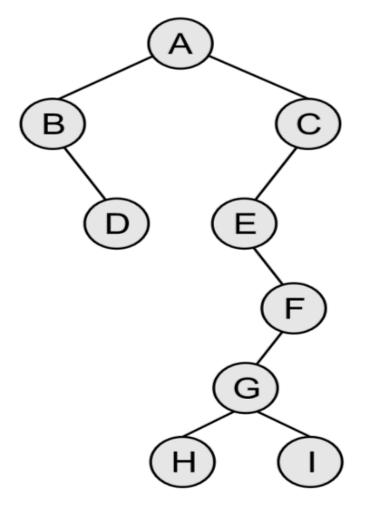






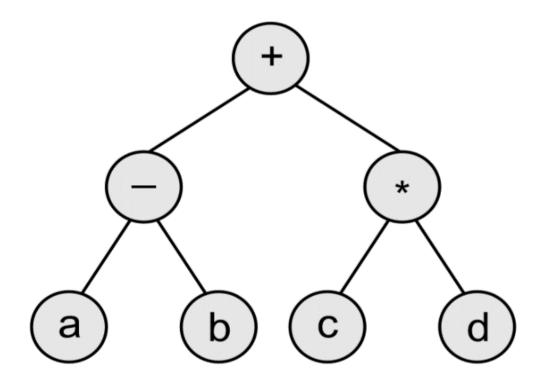






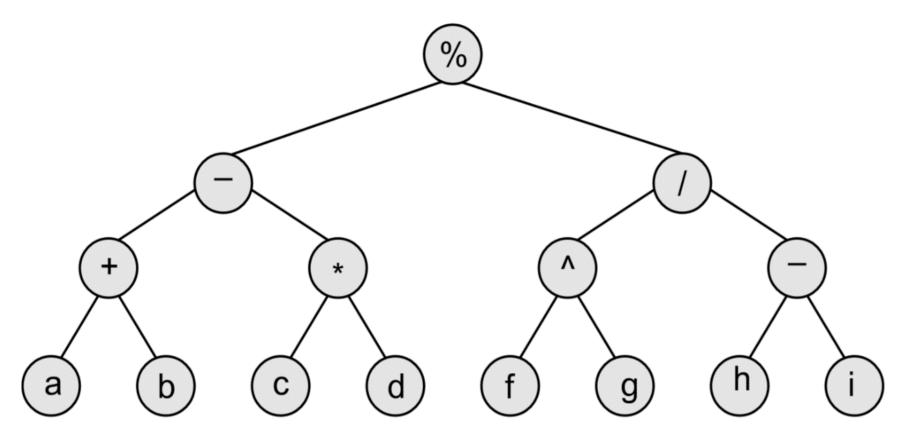


• (a - b) + (c * d)



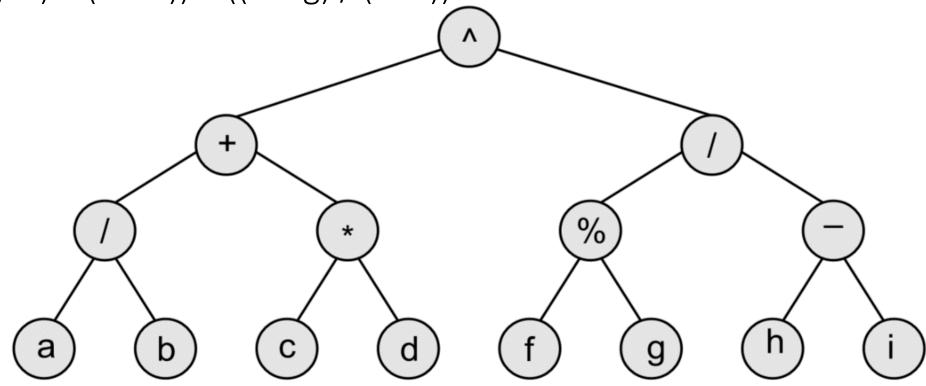


• $((a + b) - (c * d)) \% ((f ^ g) / (h - i))$





• ((a / b) + (c * d)) ^ ((f % g) / (h - i))



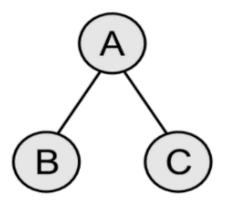


- Following operations are performed recursively at each node:
 - 1. Traversing the left sub-tree.
 - 2. Visiting the root node.
 - 3. Traversing the right sub-tree.



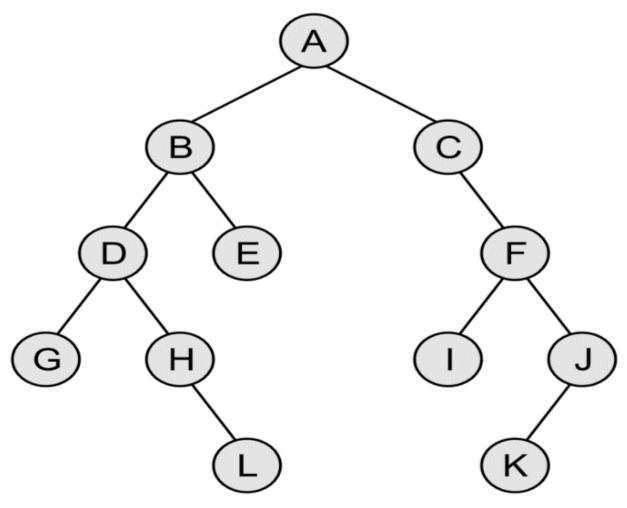
• In-order Traversal: B, A, C

Symmetric traversal.

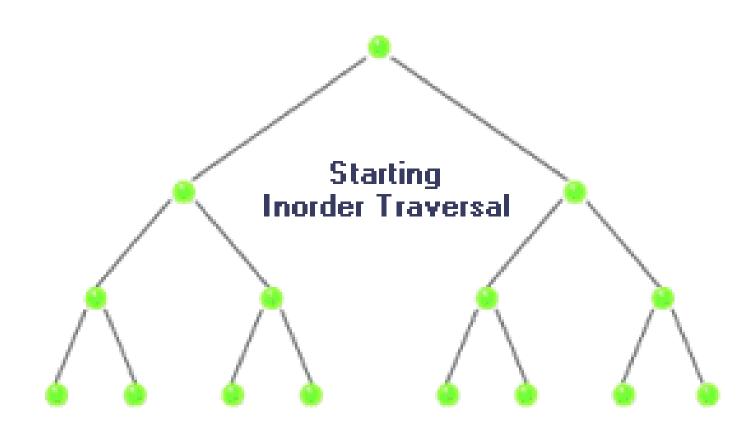


- Left sub-tree is always traversed before the root node and the right sub-tree.
- LNR traversal algorithm.

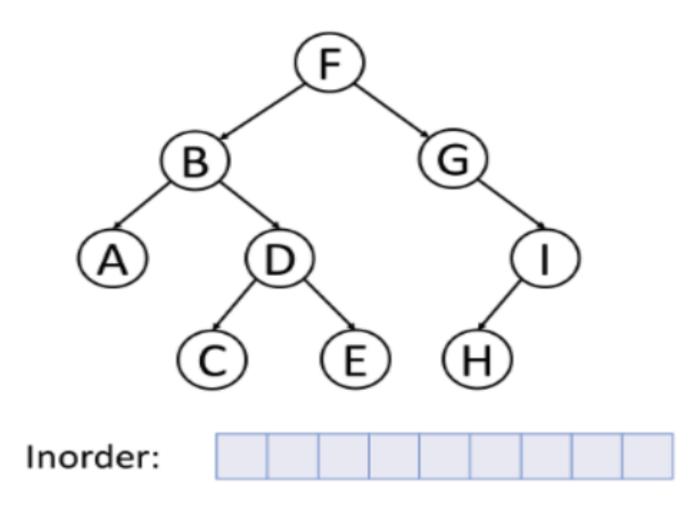




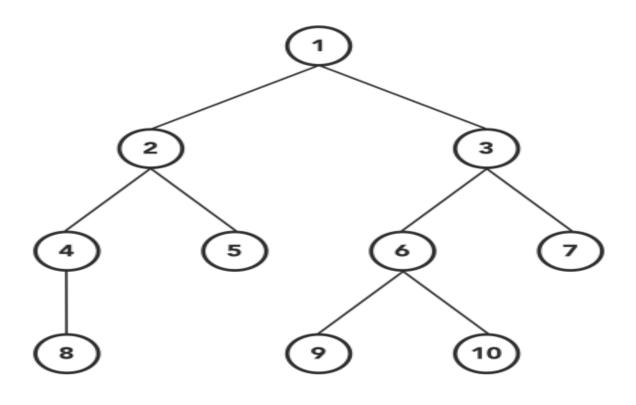




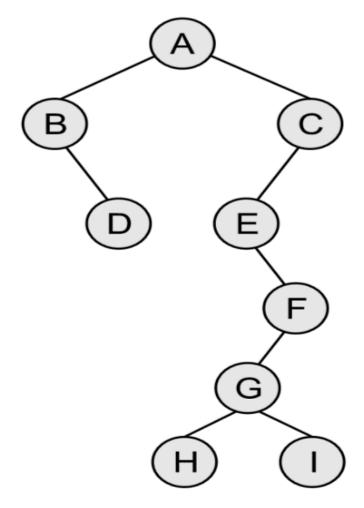








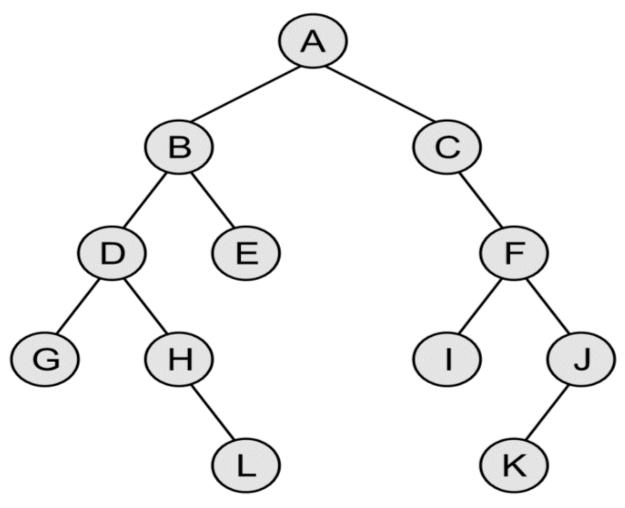




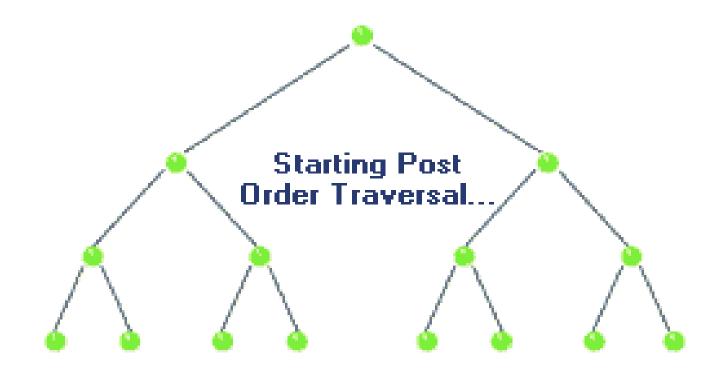


- Following operations are performed recursively at each node:
 - 1. Traversing the left sub-tree.
 - 2. Traversing the right sub-tree.
 - 3. Visiting the root node.

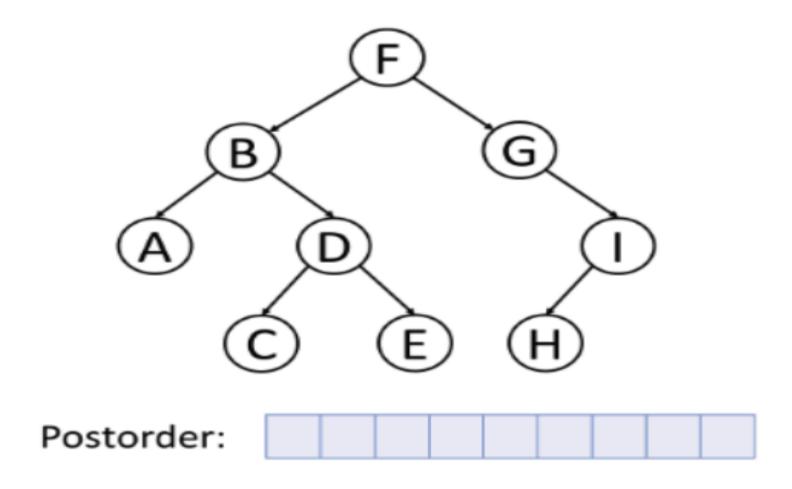




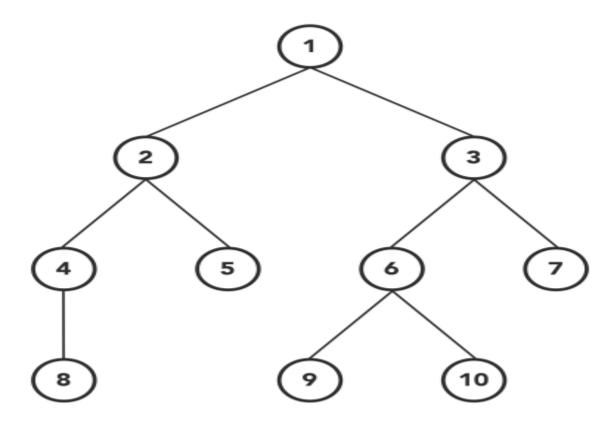




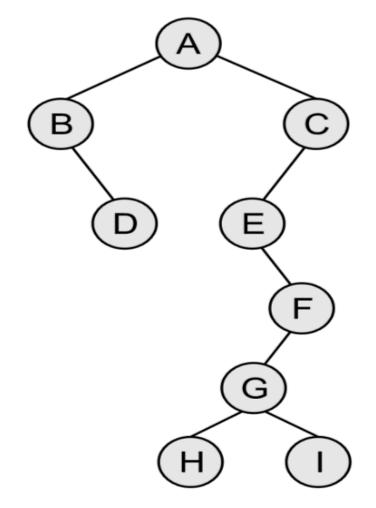








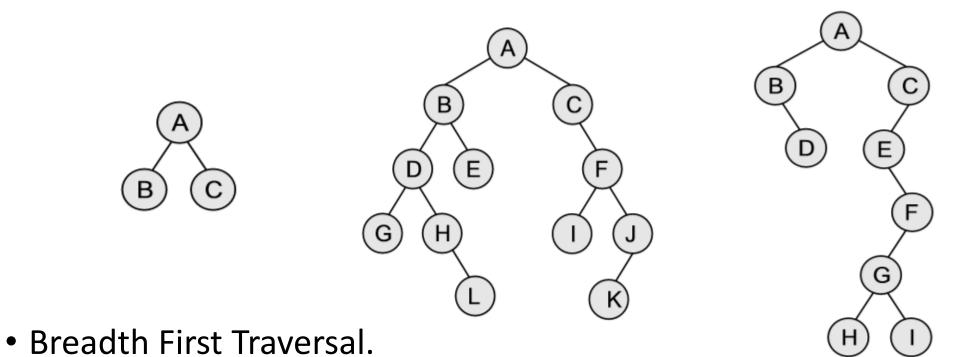






Binary Tree Traversal: Level-order Traversal

All the nodes at a level are accessed before going to the next level.





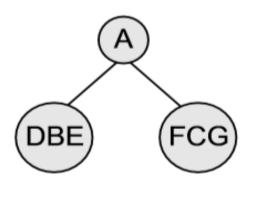
- At least two traversal results.
- One must be the in-order traversal and the second can be either preorder or post-order traversal.

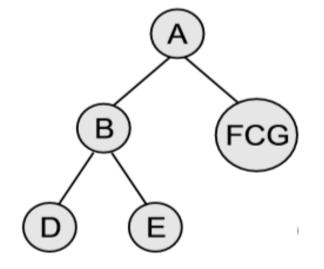


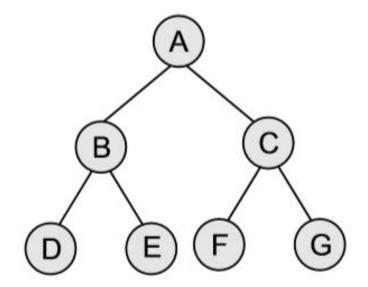
- Step 1: Use the pre-order/post-order sequence to determine the root node of the tree.
- Step 2: Elements on the left side of the root node in the in-order traversal sequence form the left sub-tree of the root node. Elements on the right side of the root node in the in-order traversal sequence form the right sub-tree of the root node.
- Step 3: Recursively select each element from pre-order/post-order traversal sequence and create its left and right sub-trees from the in-order traversal sequence.



• In-order Traversal: D B E A F C G Pre-order Traversal: A B D E C F G









- In-order Traversal: D B H E I A F J C G
- Post order Traversal: D H I E B J F G C A

