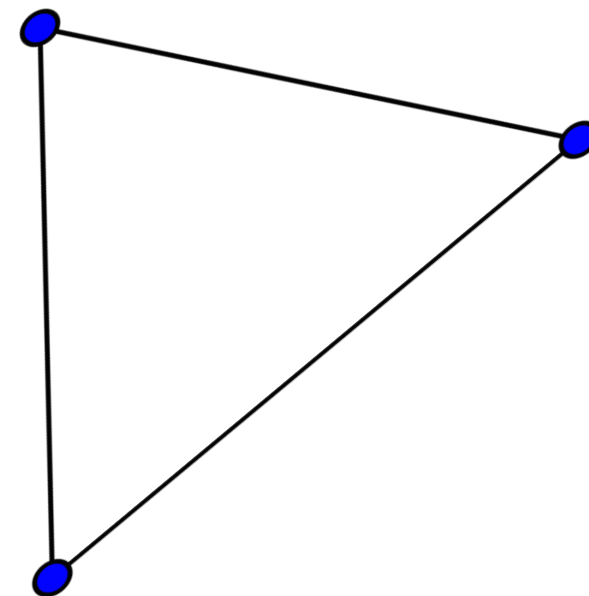


Graphs and Dijkstra's Shortest Path Algorithm

By
Arun Cyril Jose

Graphs

- Like Trees it is also a non-Linear Data Structure.
- Set of Vertices and Edges.
- Graphs Vs Trees

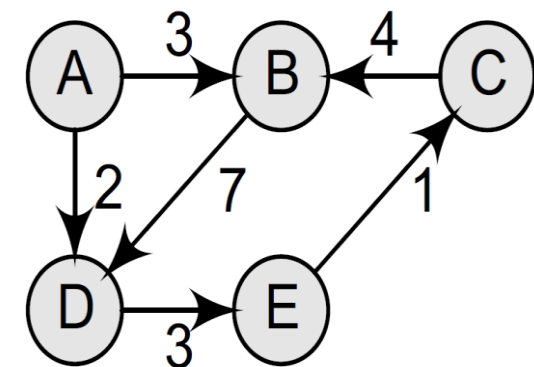
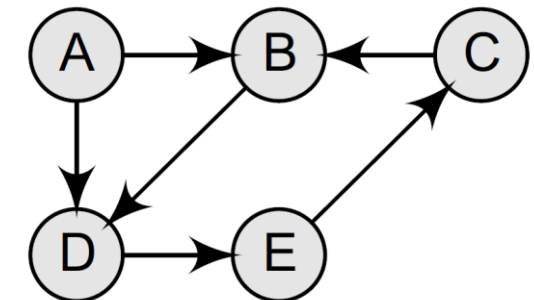
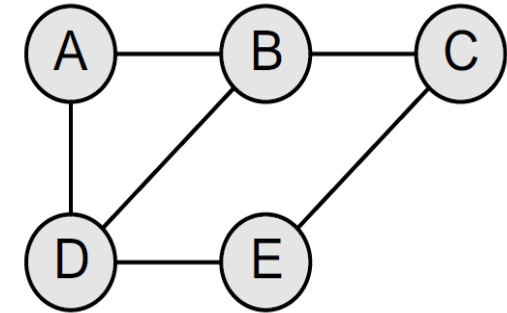


Some Applications of Graphs

- Airline Scheduling
- Computer Networks and Routing
- Directions in a map
- Plumbing/Hydraulic Systems
- Electrical Circuits.

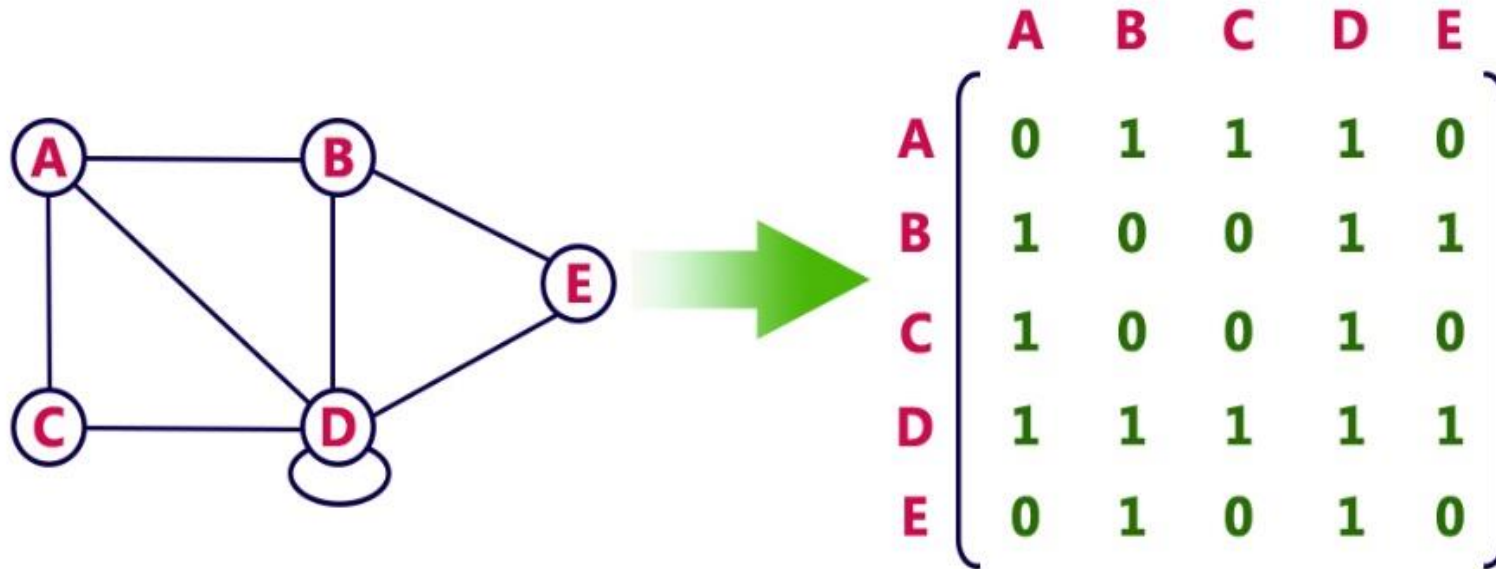
Graphs

- Used for Representing many-to-many relationships
- Two Forms:
 - **Directed** – Finite set of elements called Vertices (nodes) connected by a set of directed edges (arcs).
 - $G = \{V\} \{E\}$
 - $V = \{V_1, V_2, V_3 \dots V_n\}$ $E = \{ (V_{m1}, V_{n1}), (V_{m2}, V_{n2}) \dots (V_{mx}, V_{ny}) \}$
 - **Undirected** – Finite set of vertices connected by a set of undirected (bidirectional) edges.



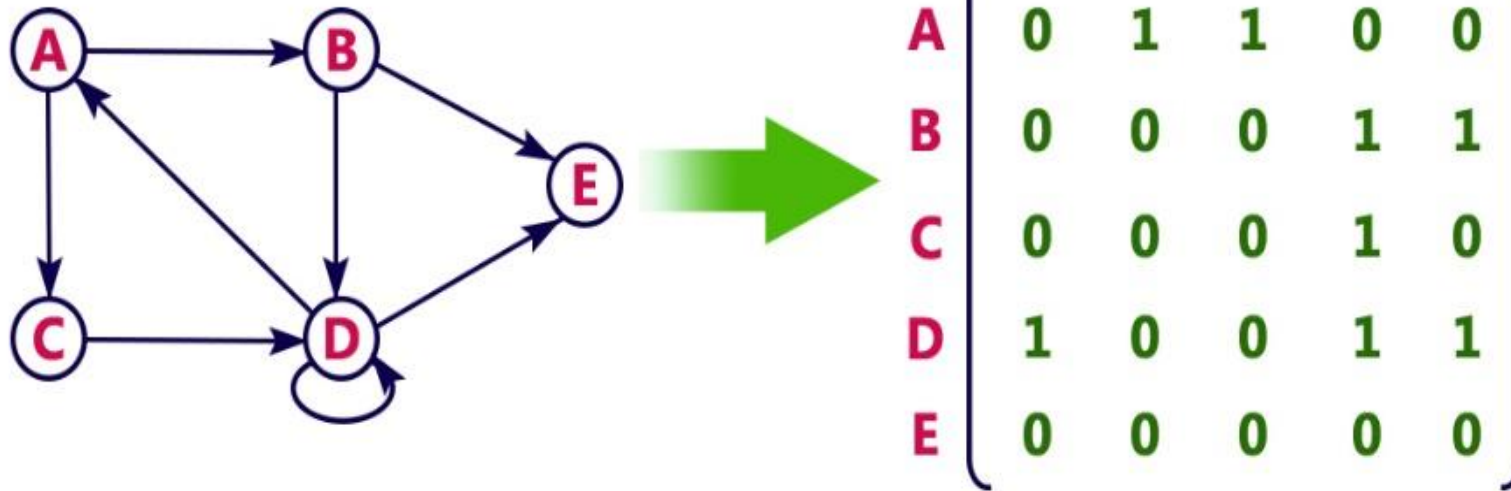
Graph Representations

- **Adjacency Matrix**



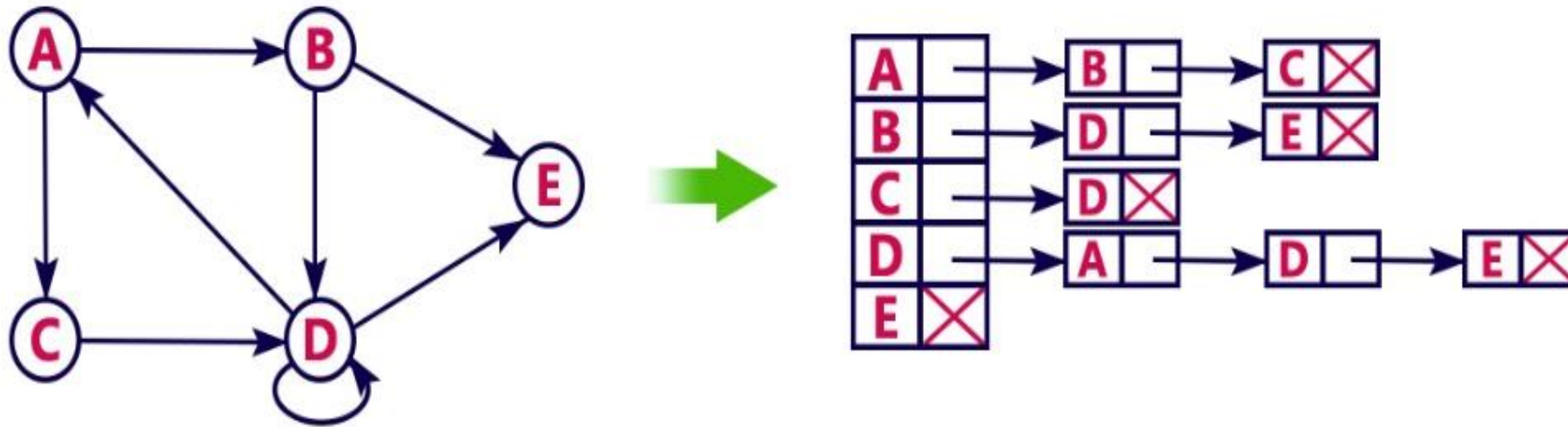
Graph Representations

- **Adjacency Matrix**



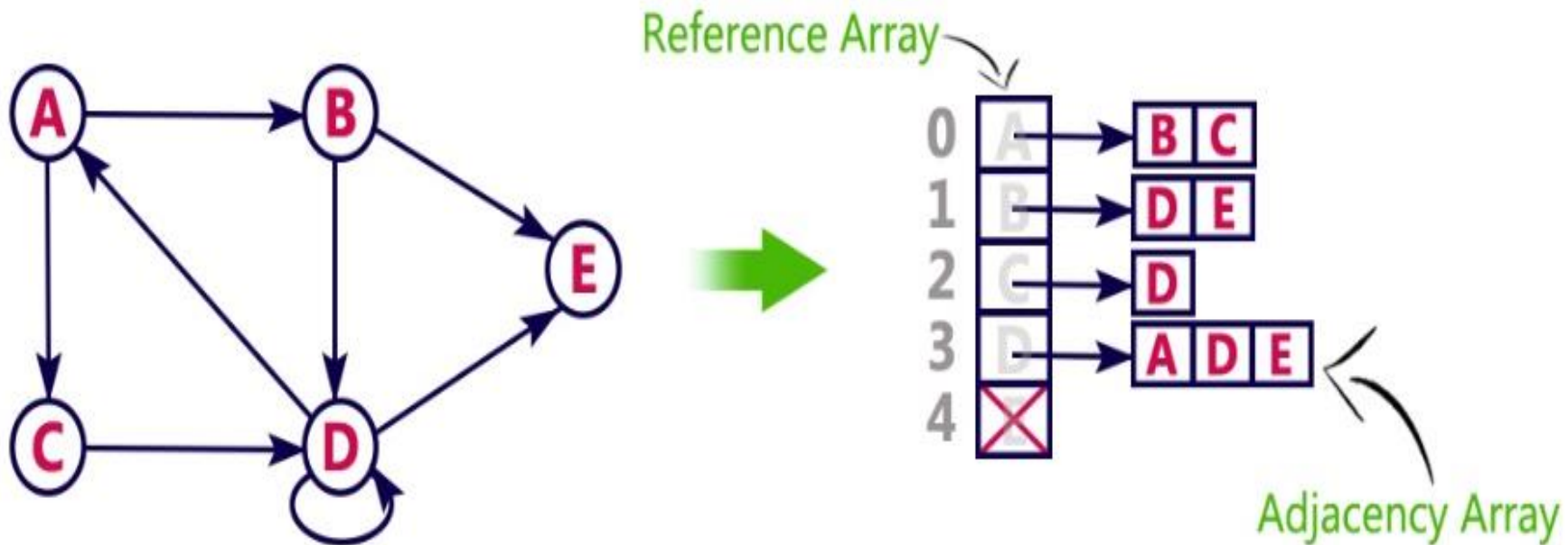
Graph Representations

- Adjacency List

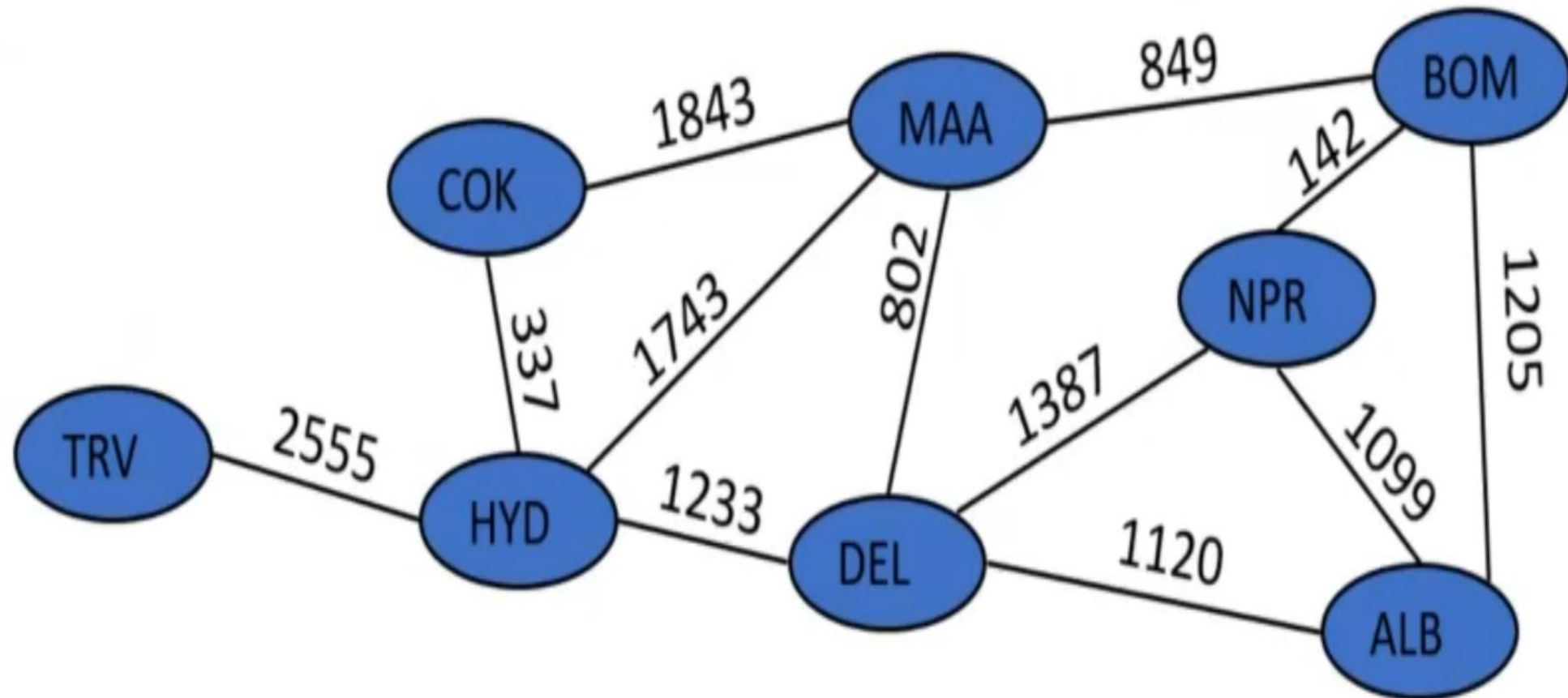


Graph Representations

- Adjacency List

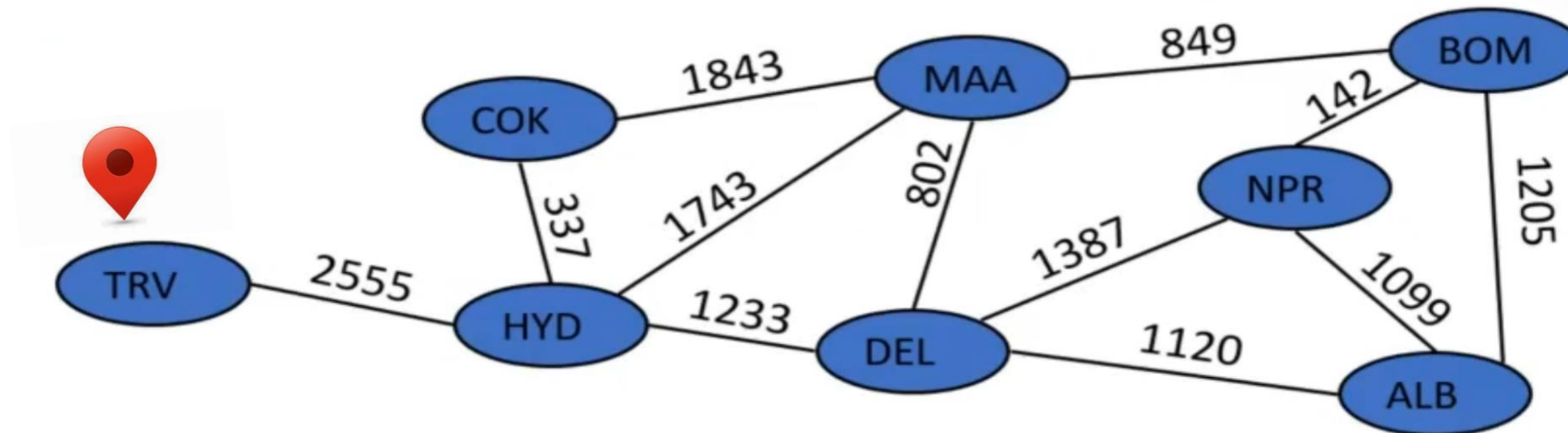


Weighted Graphs



Shortest Path Problem

- In the weighted graph and given two vertices **u** and **v**, we want to find a path of minimum total weight between **u** and **v**.
 - Length of the path is the sum of weights of its edges.
- Applications
 - Internet packet Routing, Flight reservations, Driving directions.

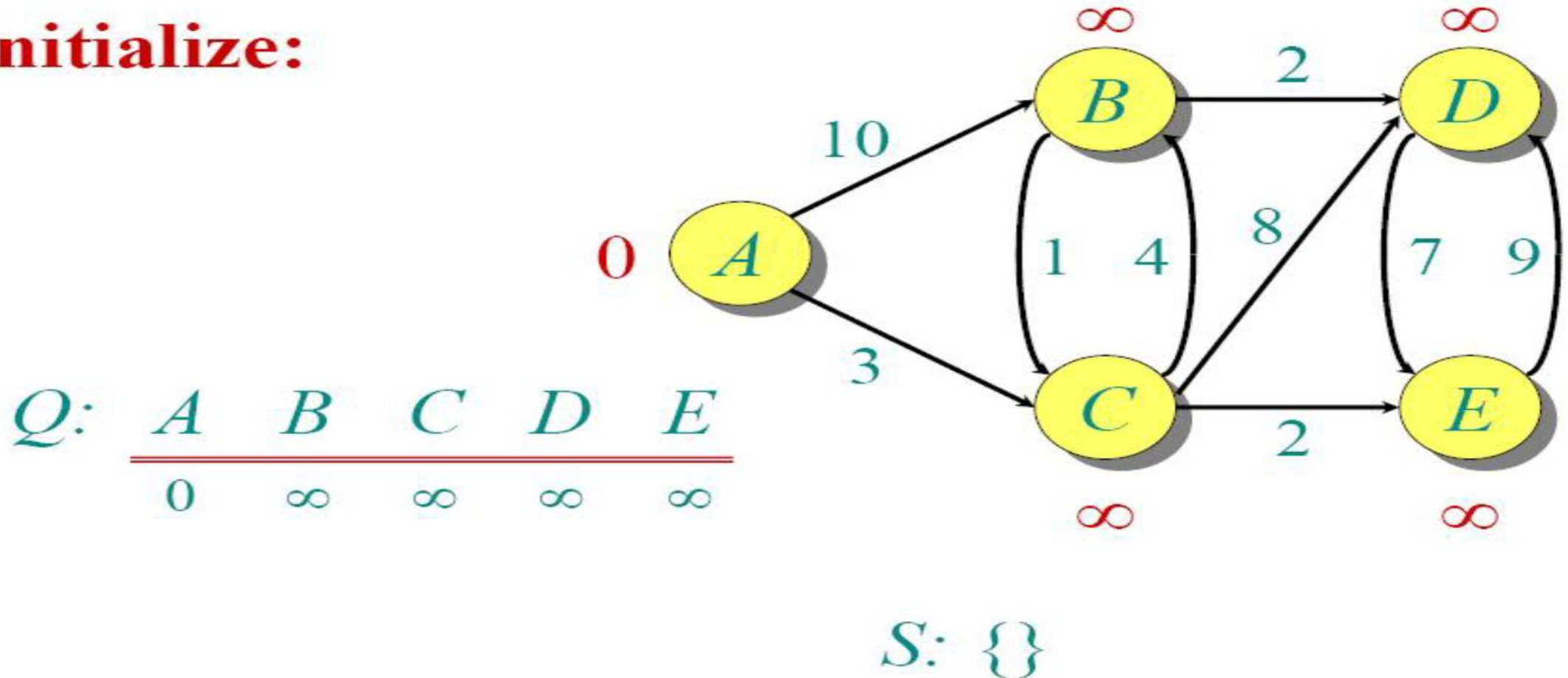


Why Shortest Path algorithms?

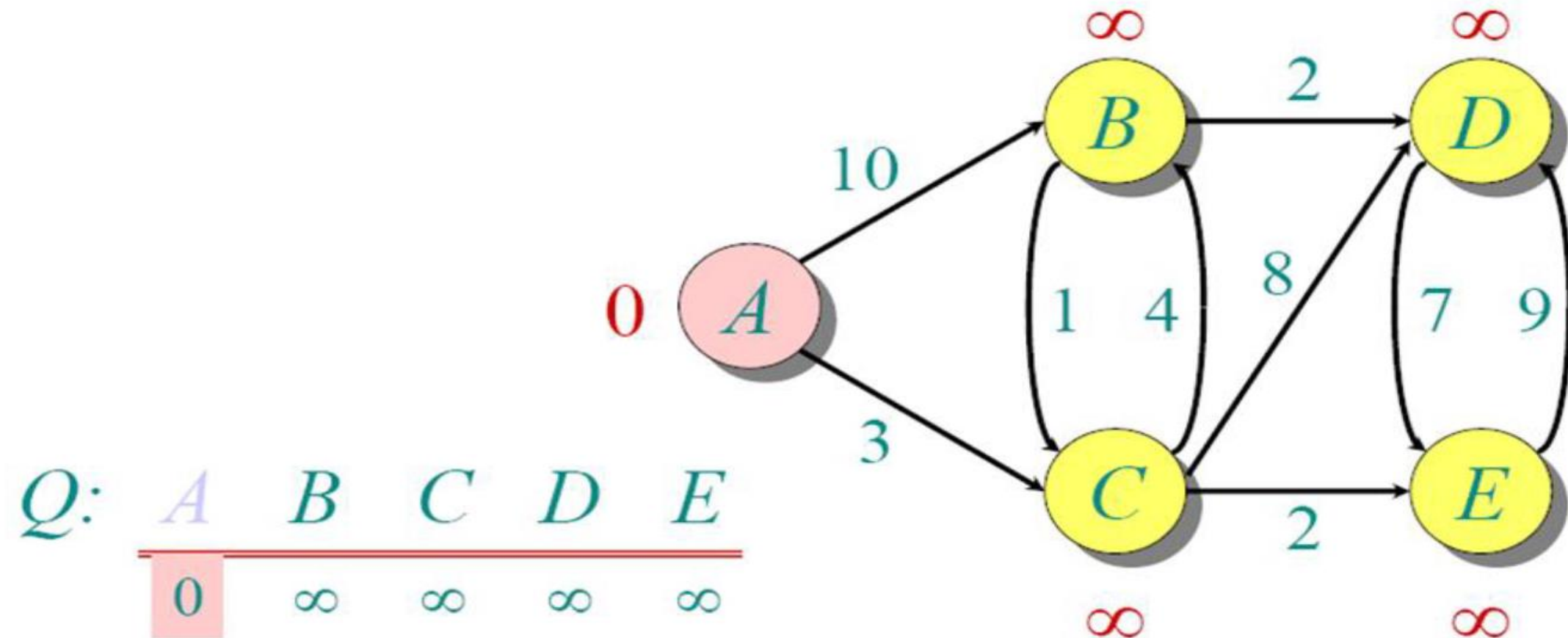
- Google Maps.
- Uber Ride Sharing.
- Routing computer Network.
- Almost all Logistical problems

Dijkstra's Shortest Path Algorithm

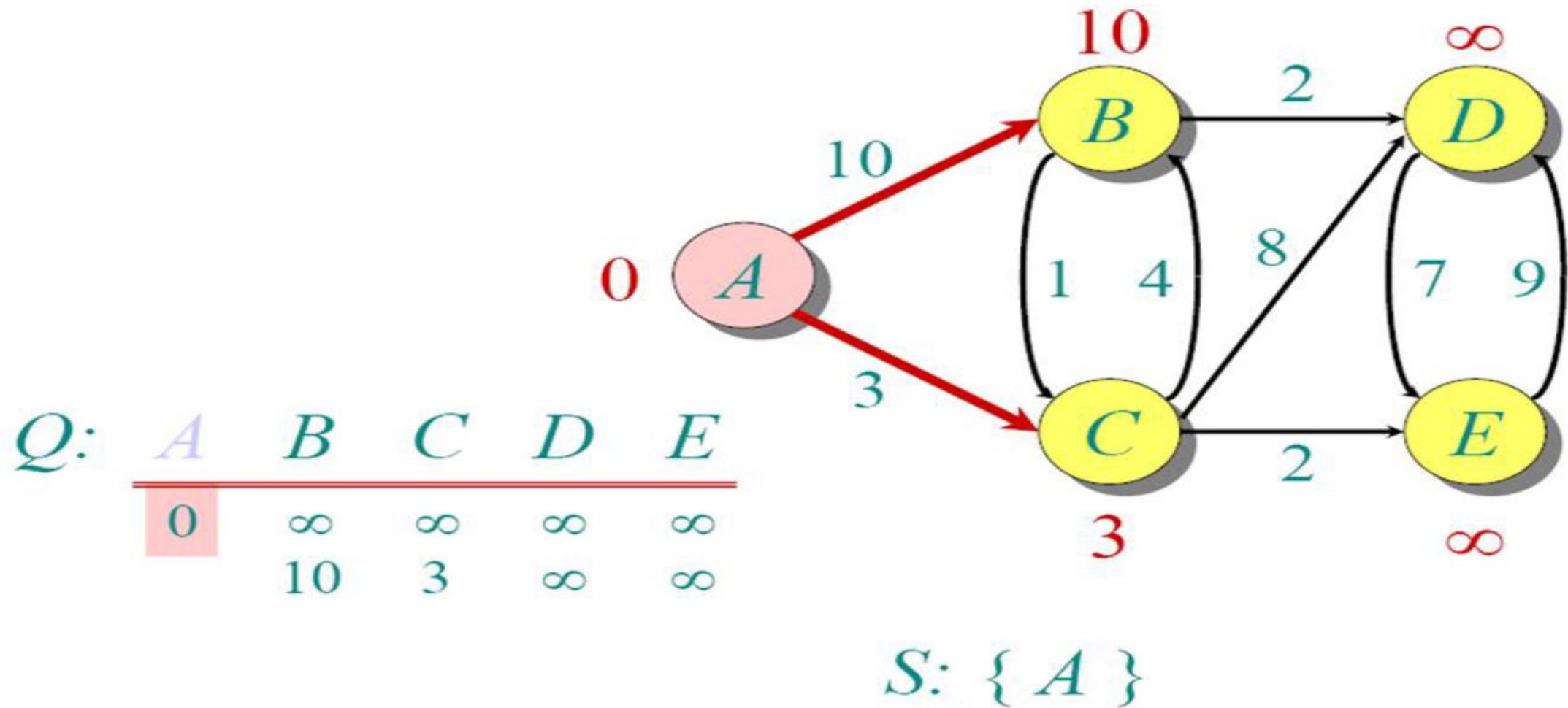
Initialize:



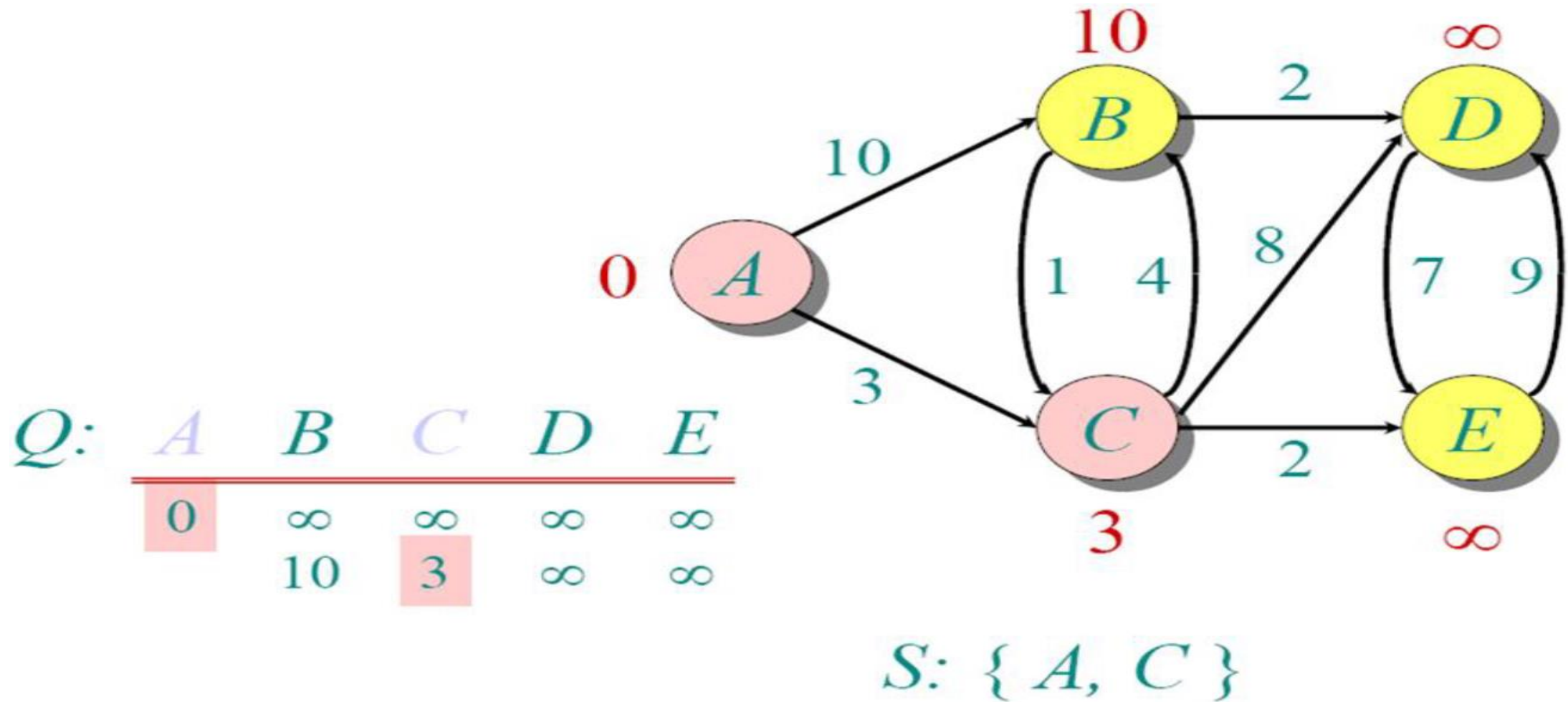
Dijkstra's Shortest Path Algorithm



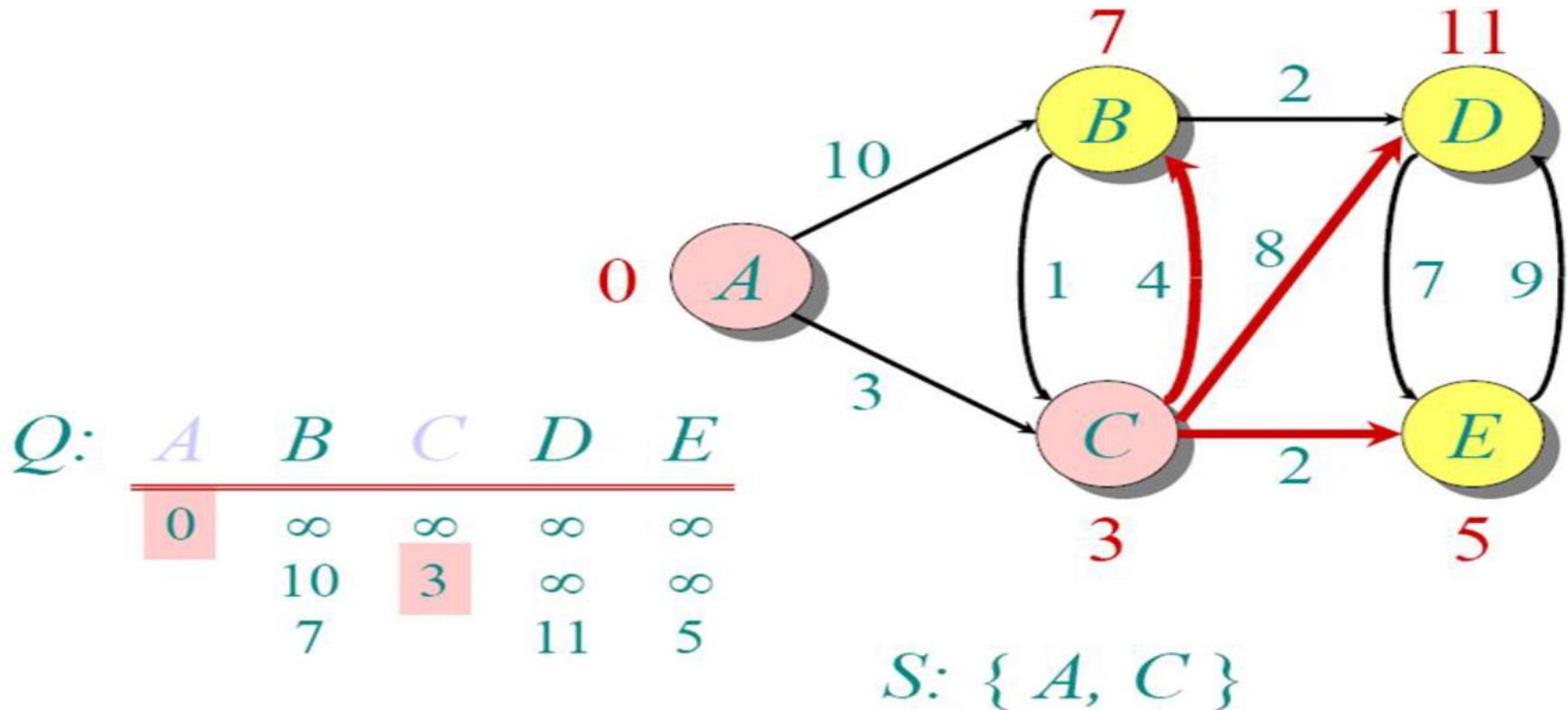
Dijkstra's Shortest Path Algorithm



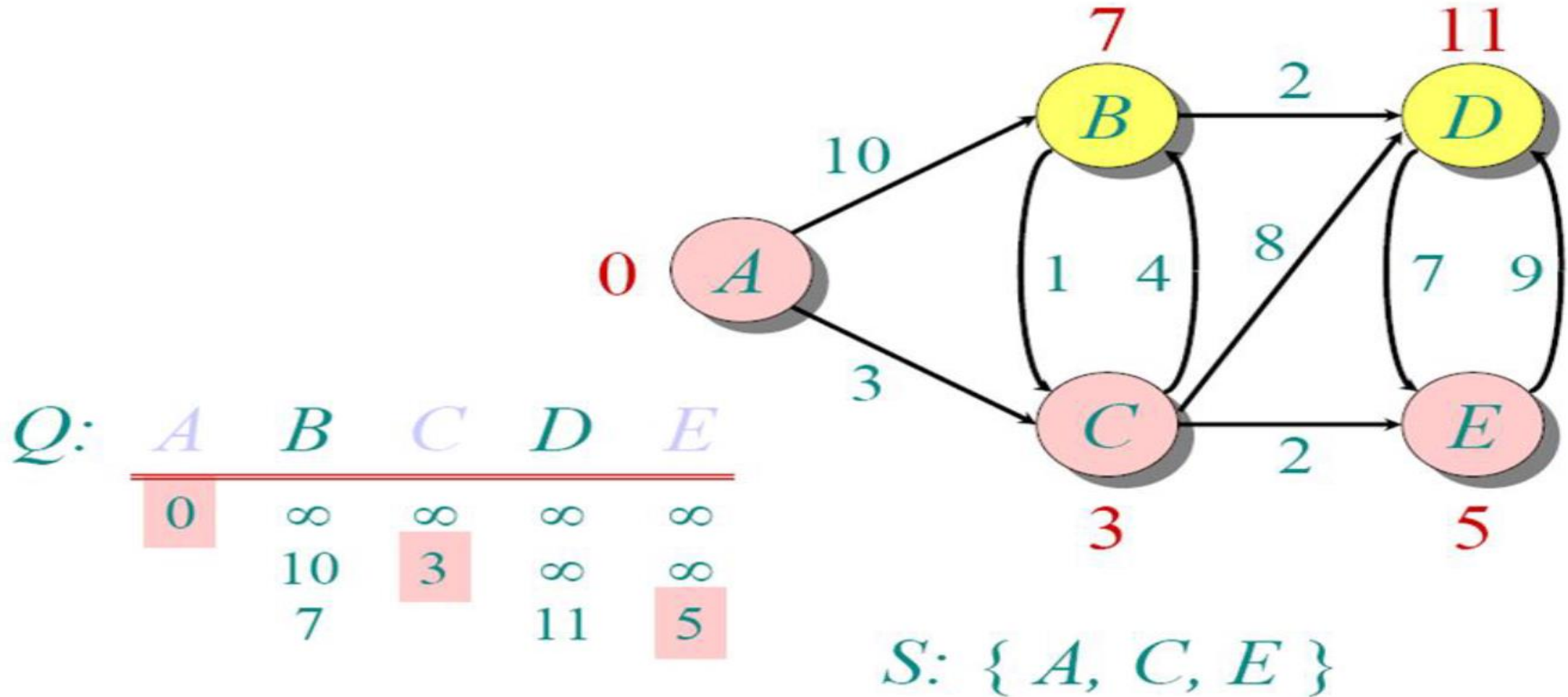
Dijkstra's Shortest Path Algorithm



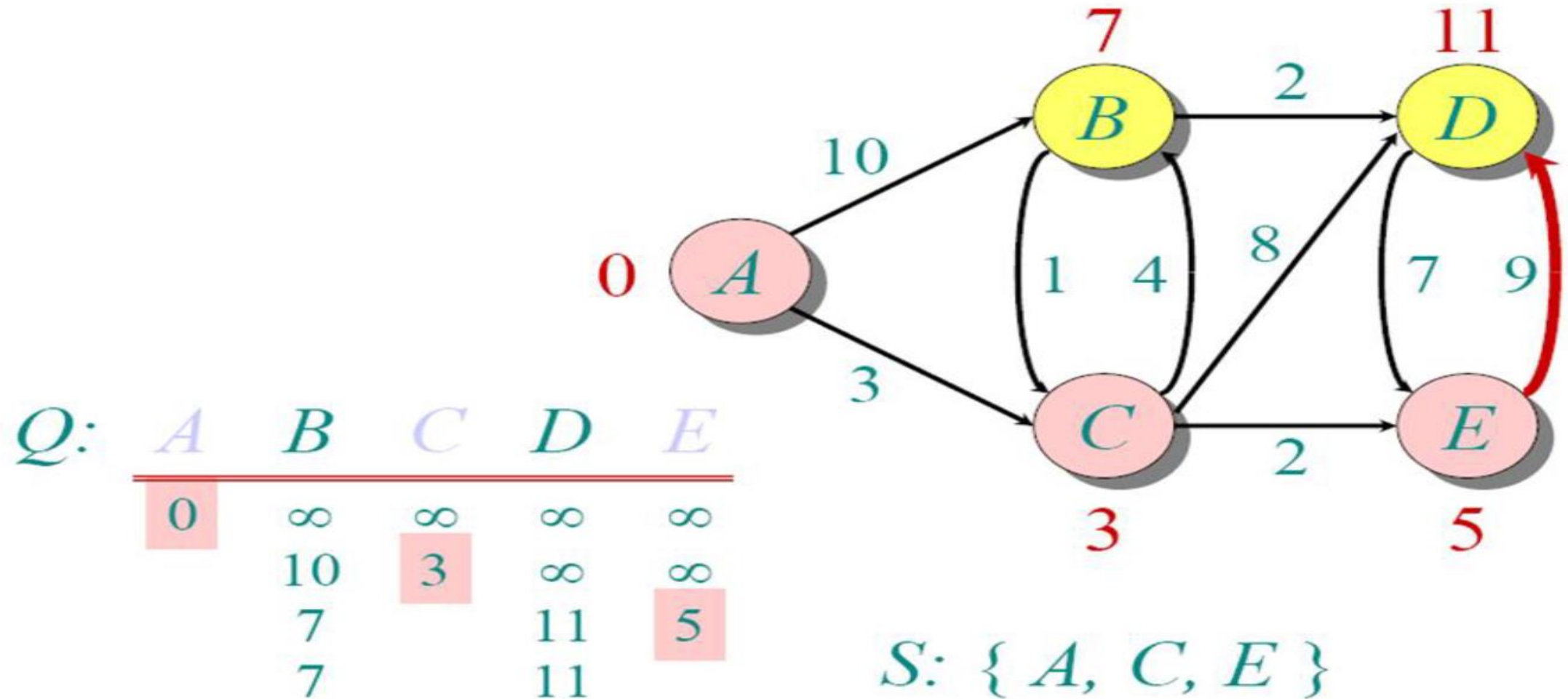
Dijkstra's Shortest Path Algorithm



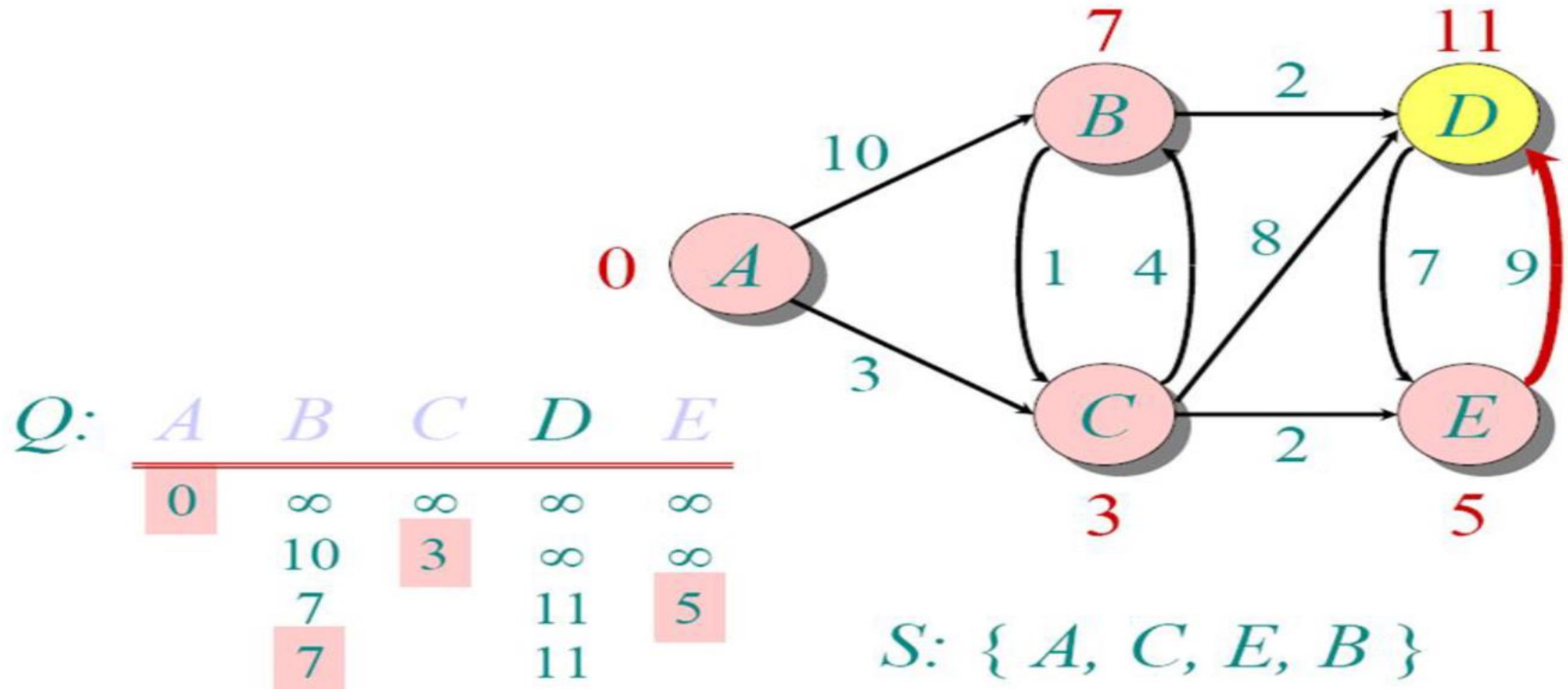
Dijkstra's Shortest Path Algorithm



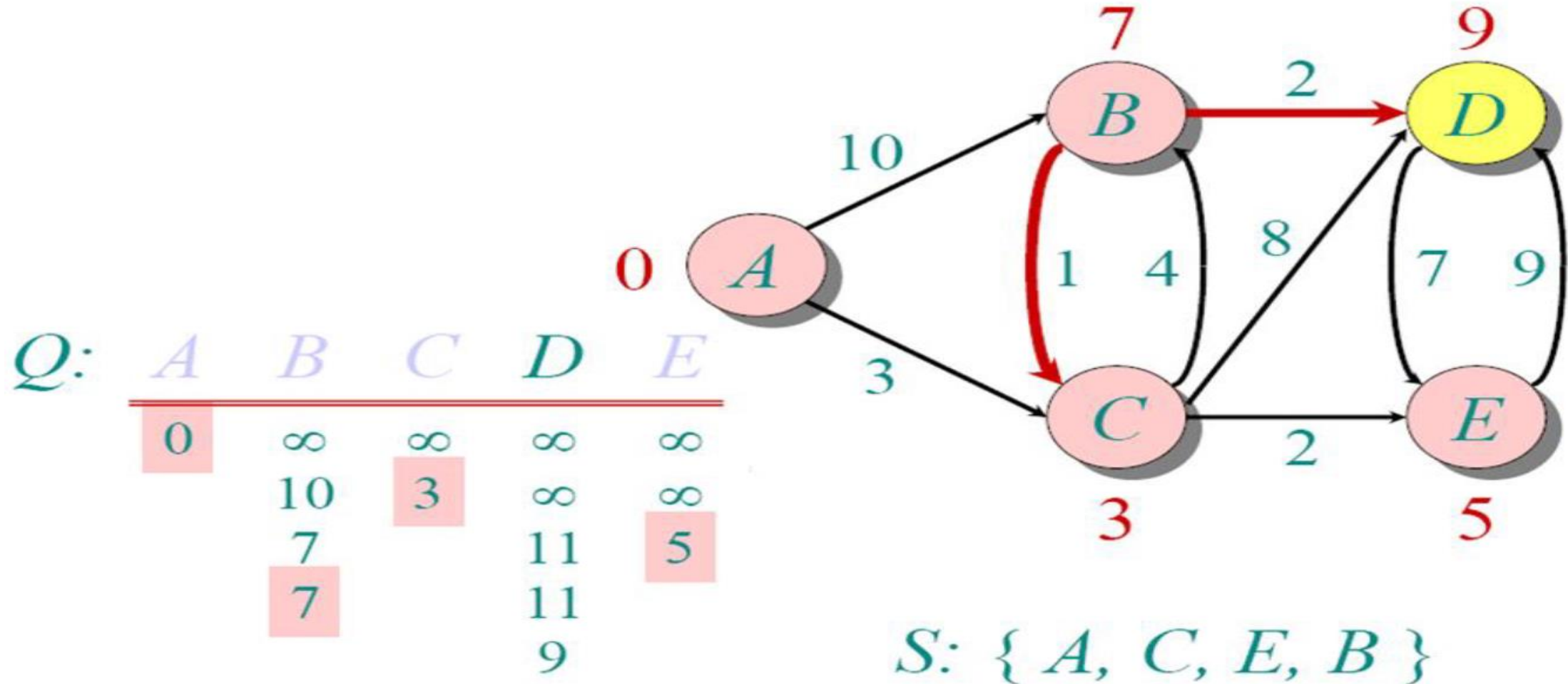
Dijkstra's Shortest Path Algorithm



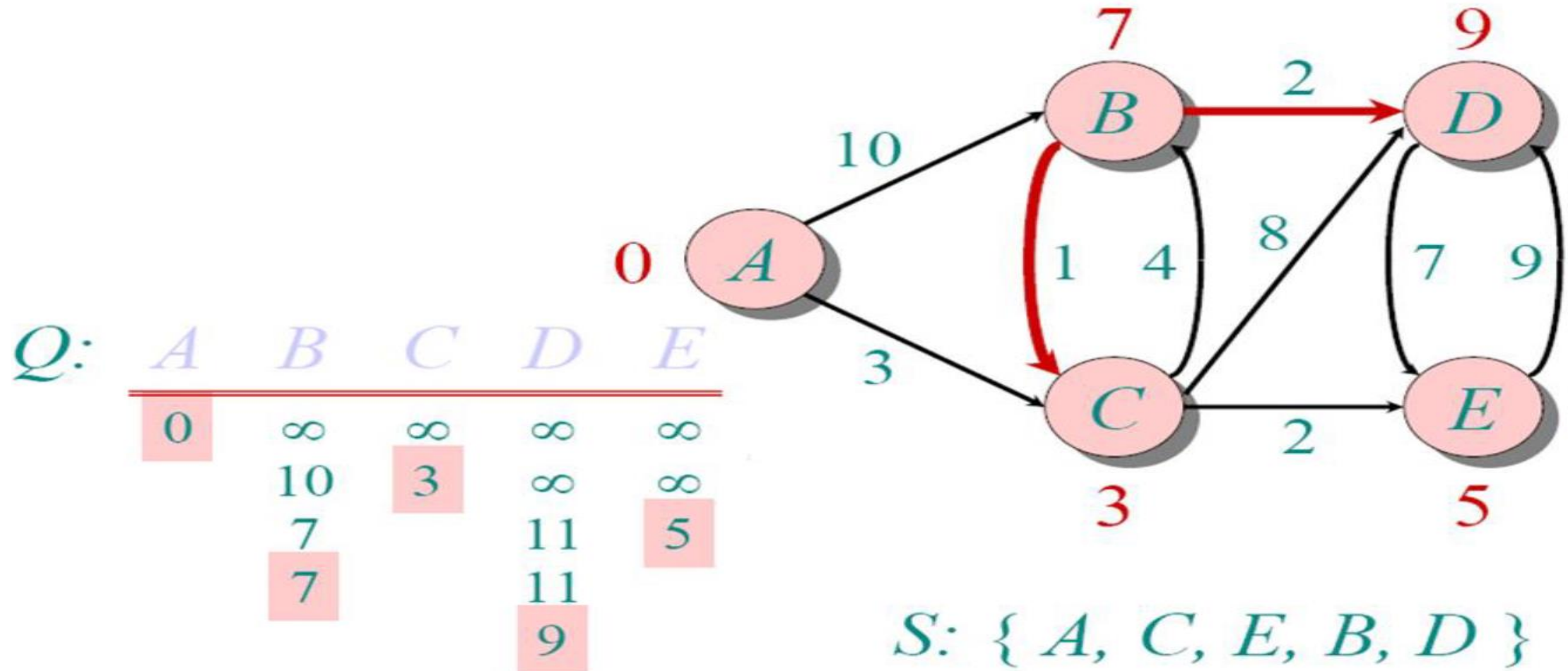
Dijkstra's Shortest Path Algorithm



Dijkstra's Shortest Path Algorithm



Dijkstra's Shortest Path Algorithm



Dijkstra's Shortest Path Algorithm

```
function dijkstra( $v_1, v_2$ ):  
  for each vertex  $v$ :                                // Initialize vertex info  
     $v$ 's cost := infinity.  
     $v$ 's previous := none.  
   $v_1$ 's cost := 0.  
   $pqueue$  := {all vertices, ordered by distance}.  
  
  while  $pqueue$  is not empty:  
     $v$  := remove vertex from  $pqueue$  with minimum cost.  
    mark  $v$  as visited.  
    for each unvisited neighbor  $n$  of  $v$ :  
       $cost$  :=  $v$ 's cost + weight of edge ( $v, n$ ).  
  
      if  $cost < n$ 's cost:  
         $n$ 's cost :=  $cost$ .  
         $n$ 's previous :=  $v$ .  
  
  reconstruct path from  $v_2$  back to  $v_1$ , following previous pointers.
```

Relax Operation

Dijkstra's Shortest Path Algorithm

function **dijkstra**(v_1, v_2):

 for each vertex v :

 // Initialize vertex info

v 's cost := infinity.

v 's previous := none.

v_1 's cost := 0.

$pqueue := \{\text{all vertices, ordered by distance}\}$.

while $pqueue$ is not empty:

$v :=$ remove vertex from $pqueue$ with minimum cost.

 mark v as visited.

 for each unvisited neighbor n of v :

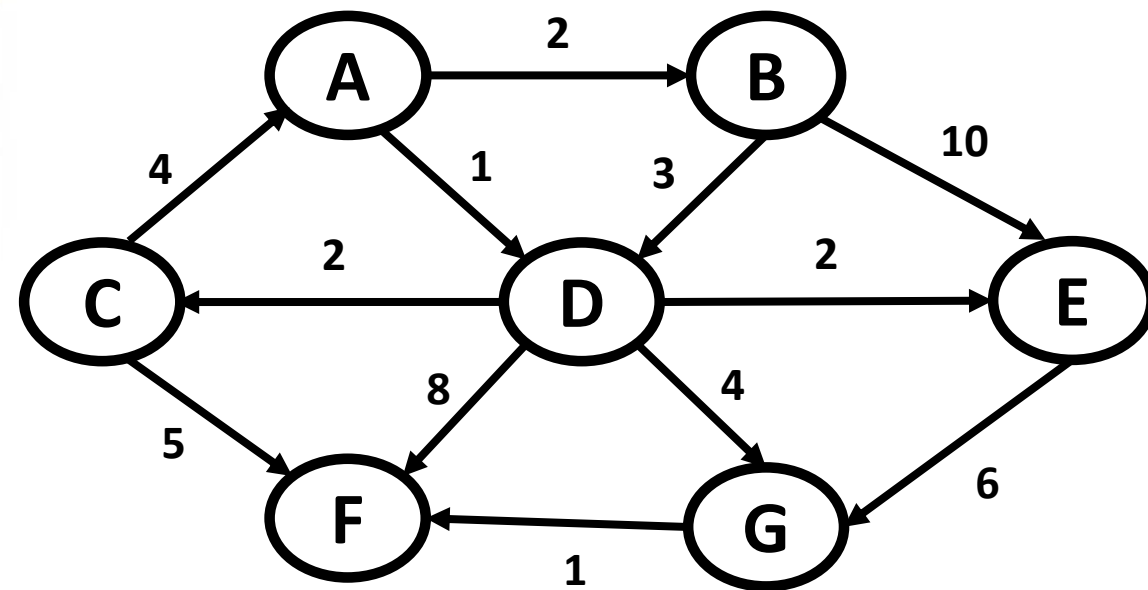
$cost := v$'s cost + weight of edge (v, n) .

 if $cost < n$'s cost:

n 's cost := $cost$.

n 's previous := v .

reconstruct path from v_2 back to v_1 , following previous pointers.



Dijkstra's Shortest Path Algorithm

function **dijkstra**(v_1, v_2):

 for each vertex v :

 // Initialize vertex info

v 's cost := infinity.

v 's previous := none.

v_1 's cost := 0.

$pqueue := \{\text{all vertices, ordered by distance}\}$.

while $pqueue$ is not empty:

$v := \text{remove vertex from } pqueue \text{ with minimum cost.}$ // A

 mark v as visited.

 for each unvisited neighbor n of v : // B, D

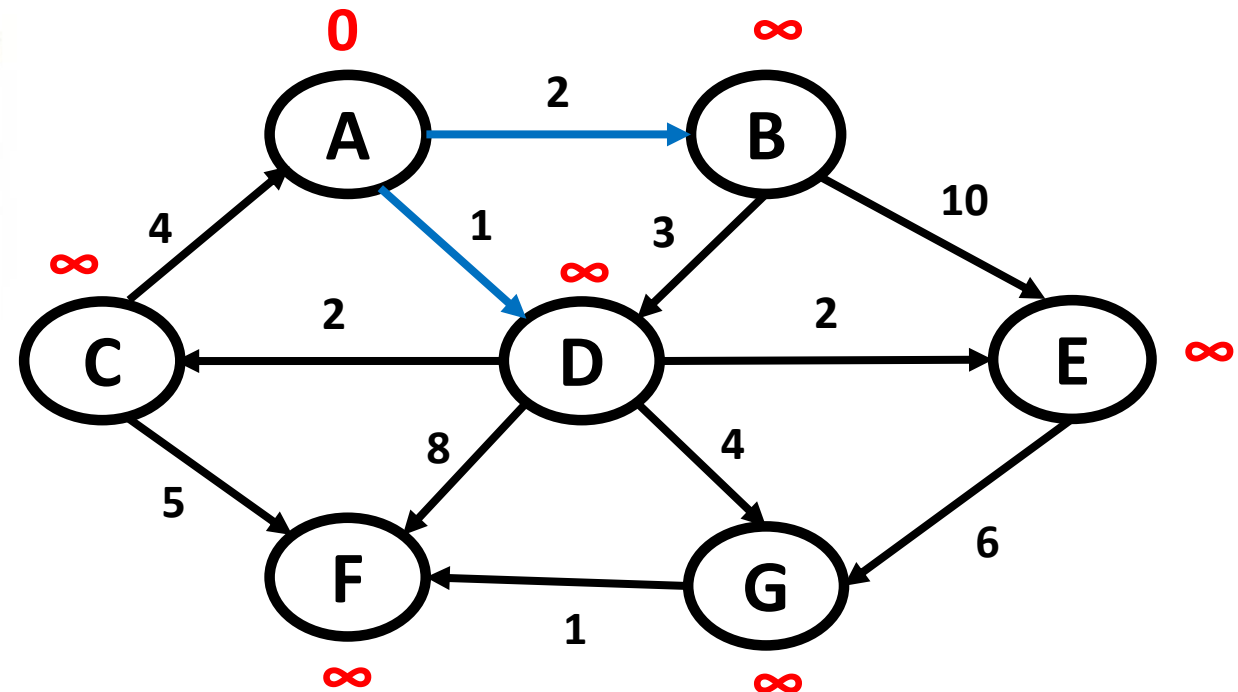
$cost := v$'s cost + weight of edge (v, n) .

 if $cost < n$'s cost:

n 's cost := $cost$.

n 's previous := v .

reconstruct path from v_2 back to v_1 , following previous pointers.



Priority Queue: [A: 0, B: ∞ , C: ∞ , D: ∞ , E: ∞ , F: ∞ , G: ∞]

Dijkstra's Shortest Path Algorithm

function **dijkstra**(v_1, v_2):

for each vertex v :

// Initialize vertex info

v 's cost := infinity.

v 's previous := none.

v_1 's cost := 0.

$pqueue := \{\text{all vertices, ordered by distance}\}$.

while $pqueue$ is not empty:

$v := \text{remove vertex from } pqueue \text{ with minimum cost.}$ // A

mark v as visited.

for each unvisited neighbor n of v : // B, D

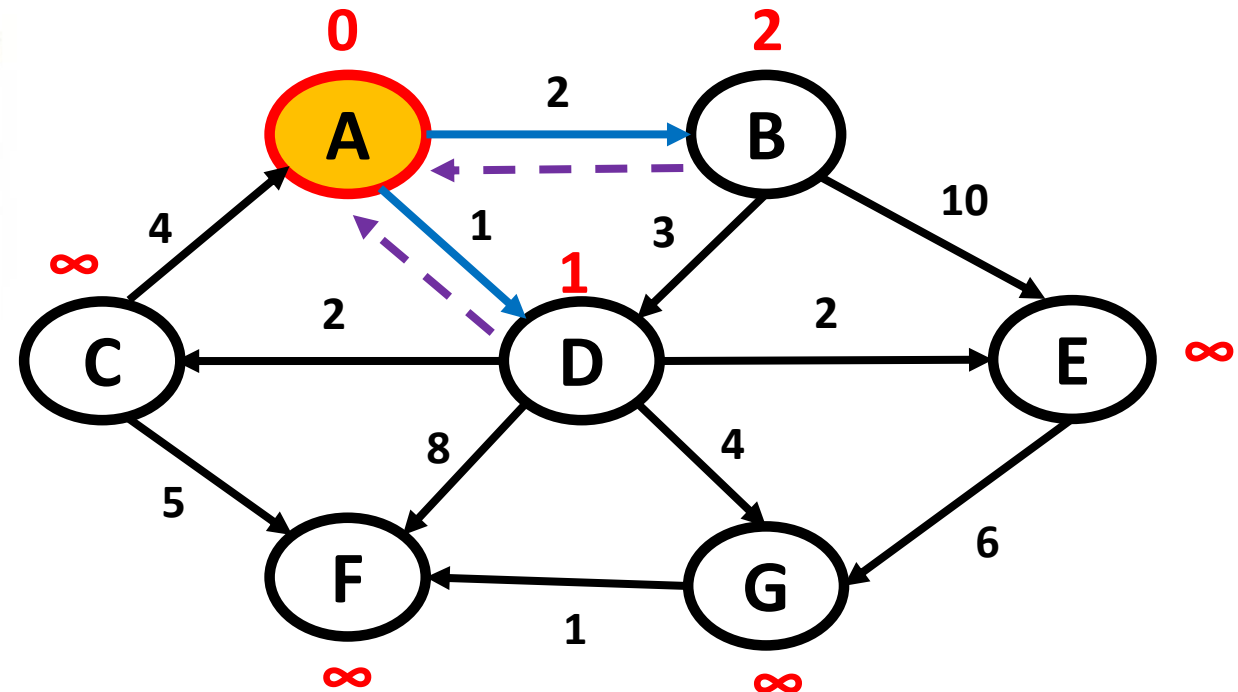
$cost := v$'s cost + weight of edge (v, n) .

if $cost < n$'s cost: // B's cost = $0 + 2$

n 's cost := $cost$. // D's cost = $0 + 1$

n 's previous := v .

reconstruct path from v_2 back to v_1 , following previous pointers.



Priority Queue: [B: 2, C: ∞ , D: 1, E: ∞ , F: ∞ , G: ∞]

Dijkstra's Shortest Path Algorithm

function **dijkstra**(v_1, v_2):

for each vertex v :

// Initialize vertex info

v 's cost := infinity.

v 's previous := none.

v_1 's cost := 0.

$pqueue := \{\text{all vertices, ordered by distance}\}$.

while $pqueue$ is not empty:

$v := \text{remove vertex from } pqueue \text{ with minimum cost.}$ // D

mark v as visited.

for each unvisited neighbor n of v : // C, E, F, G

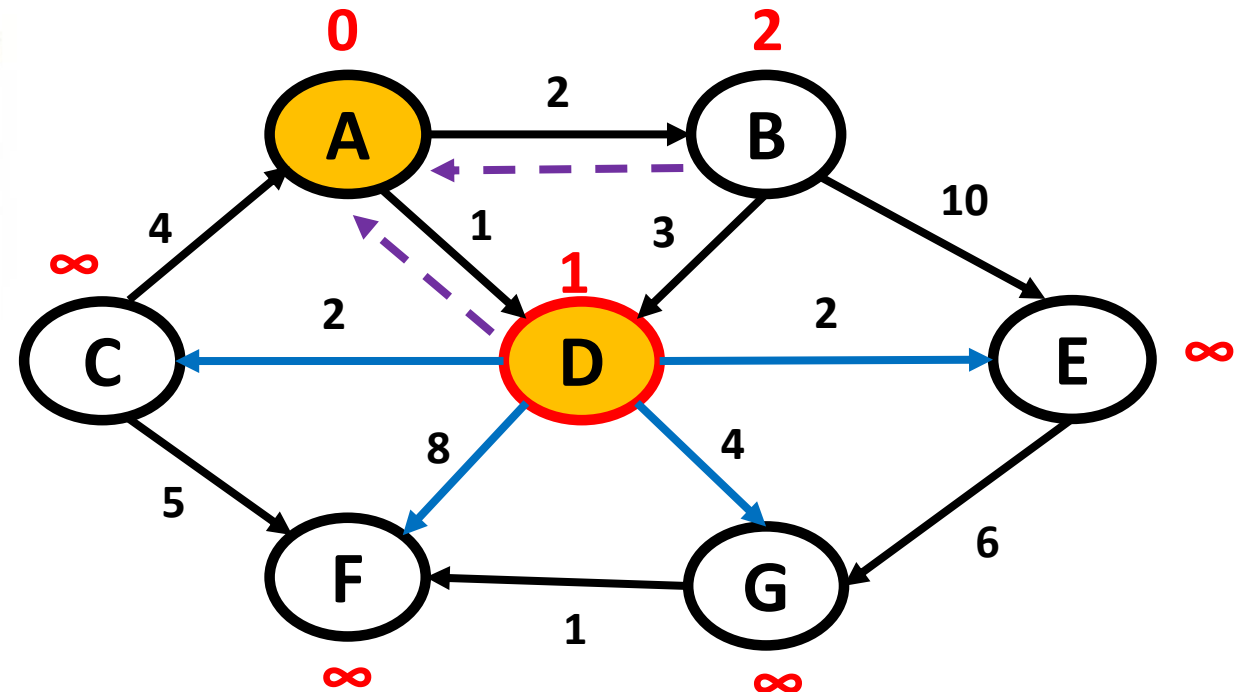
$cost := v$'s cost + weight of edge (v, n) .

if $cost < n$'s cost:

n 's cost := $cost$.

n 's previous := v .

reconstruct path from v_2 back to v_1 , following previous pointers.



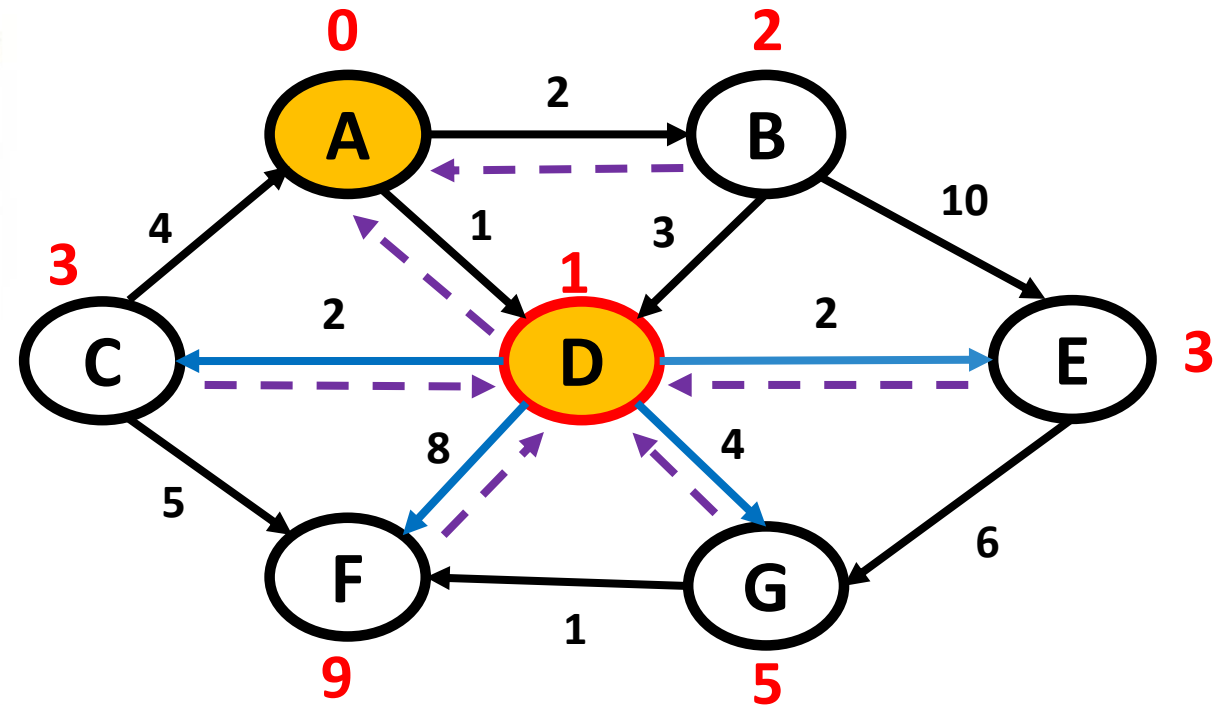
Priority Queue: [B: 2, C: ∞ , E: ∞ , F: ∞ , G: ∞]

Dijkstra's Shortest Path Algorithm

```
function dijkstra( $v_1, v_2$ ):
  for each vertex  $v$ :           // Initialize vertex info
     $v$ 's cost := infinity.
     $v$ 's previous := none.
   $v_1$ 's cost := 0.
   $pqueue := \{\text{all vertices, ordered by distance}\}$ .

  while  $pqueue$  is not empty:
     $v :=$  remove vertex from  $pqueue$  with minimum cost. // D
    mark  $v$  as visited.
    for each unvisited neighbor  $n$  of  $v$ : // C, E, F, G
       $cost := v$ 's cost + weight of edge ( $v, n$ ).
      // C's cost = 1 + 2
      if  $cost < n$ 's cost: // E's cost = 1 + 2
         $n$ 's cost :=  $cost$ . // F's cost = 1 + 8
         $n$ 's previous :=  $v$ . // G's cost = 1 + 4

  reconstruct path from  $v_2$  back to  $v_1$ , following previous pointers.
```



Priority Queue: [B: 2, C: 3, E: 3, F: 9, G: 5]

Dijkstra's Shortest Path Algorithm

function **dijkstra**(v_1, v_2):

for each vertex v :

// Initialize vertex info

v 's cost := infinity.

v 's previous := none.

v_1 's cost := 0.

$pqueue := \{\text{all vertices, ordered by distance}\}$.

while $pqueue$ is not empty:

$v := \text{remove vertex from } pqueue \text{ with minimum cost.}$ // B

mark v as visited.

for each unvisited neighbor n of v : // E

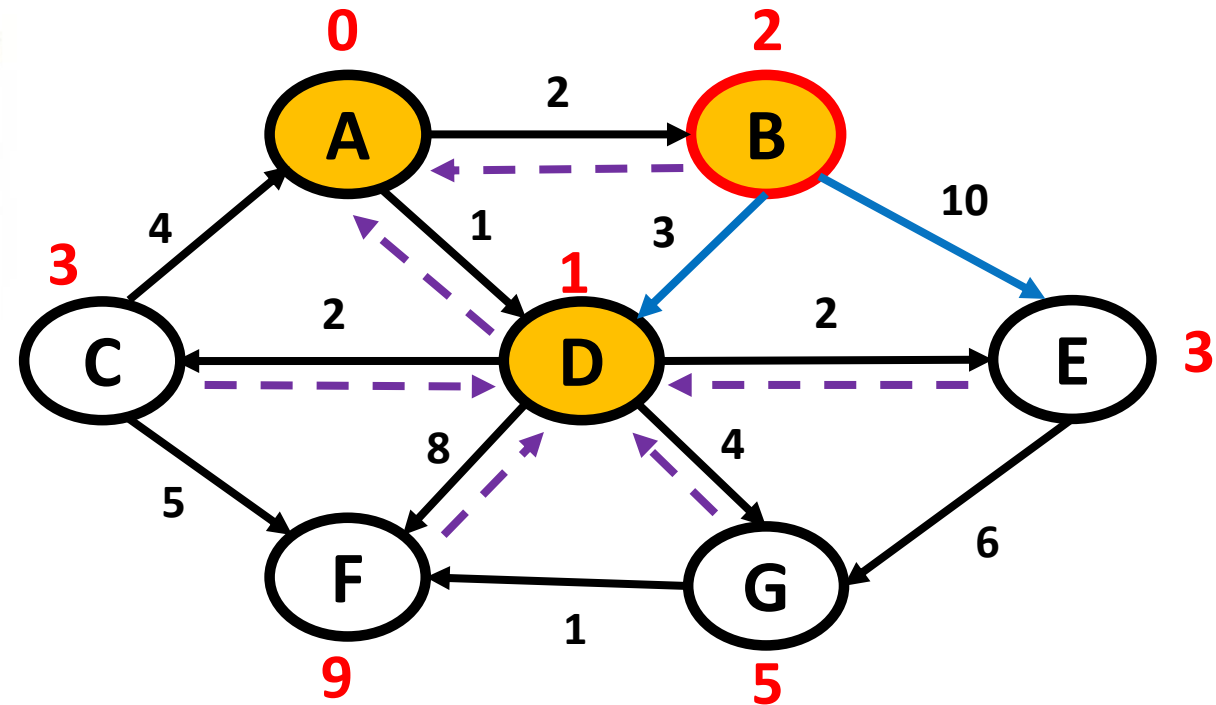
$cost := v$'s cost + weight of edge (v, n) .

if $cost < n$'s cost:

n 's cost := $cost$.

n 's previous := v .

reconstruct path from v_2 back to v_1 , following previous pointers.



Priority Queue: [C: 3, E: 3, F: 9, G: 5]

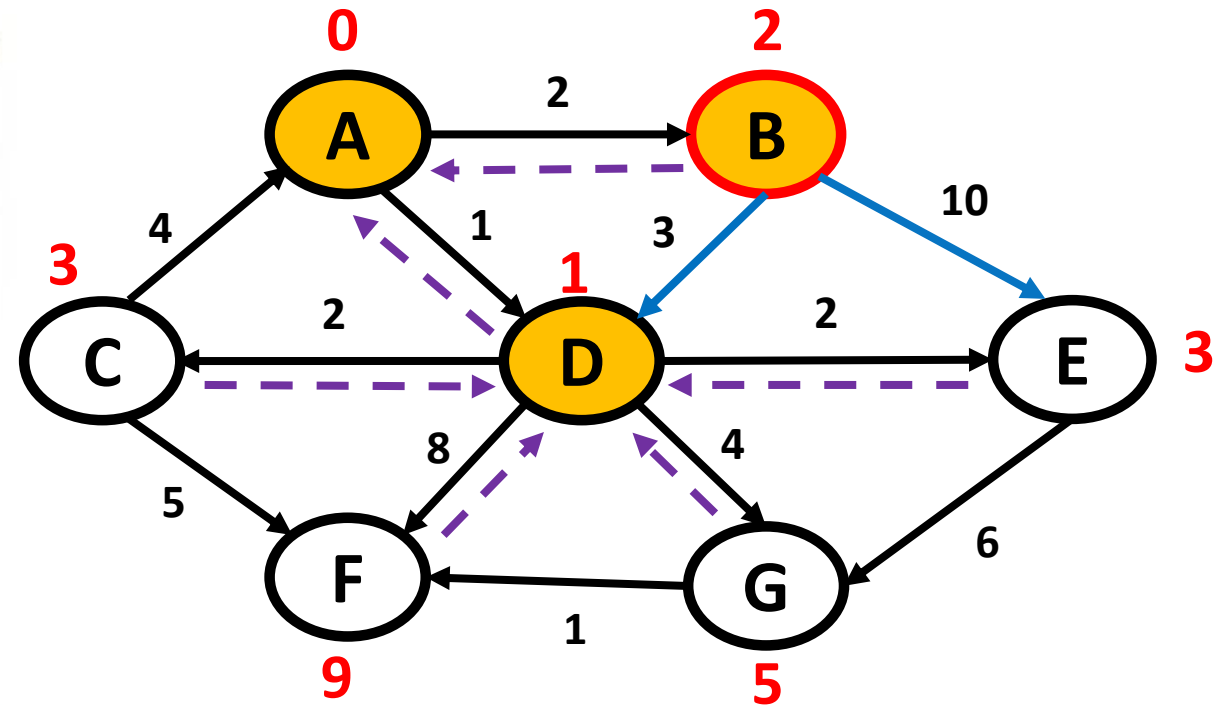
Dijkstra's Shortest Path Algorithm

```
function dijkstra( $v_1, v_2$ ):
  for each vertex  $v$ :           // Initialize vertex info
     $v$ 's cost := infinity.
     $v$ 's previous := none.
   $v_1$ 's cost := 0.
   $pqueue := \{\text{all vertices, ordered by distance}\}$ .

  while  $pqueue$  is not empty:
     $v := \text{remove vertex from } pqueue \text{ with minimum cost.}$  // B
    mark  $v$  as visited.
    for each unvisited neighbor  $n$  of  $v$ : // E
       $cost := v$ 's cost + weight of edge  $(v, n)$ . // 2 + 10

      if  $cost < n$ 's cost: // 12 > 3; False
         $n$ 's cost :=  $cost$ . // no cost change for E
         $n$ 's previous :=  $v$ .

  reconstruct path from  $v_2$  back to  $v_1$ , following previous pointers.
```



Priority Queue: [C: 3, E: 3, F: 9, G: 5]

Dijkstra's Shortest Path Algorithm

function **dijkstra**(v_1, v_2):

for each vertex v :

// Initialize vertex info

v 's cost := infinity.

v 's previous := none.

v_1 's cost := 0.

$pqueue := \{\text{all vertices, ordered by distance}\}$.

while $pqueue$ is not empty:

$v := \text{remove vertex from } pqueue \text{ with minimum cost.}$ // C

mark v as visited.

for each unvisited neighbor n of v : // F

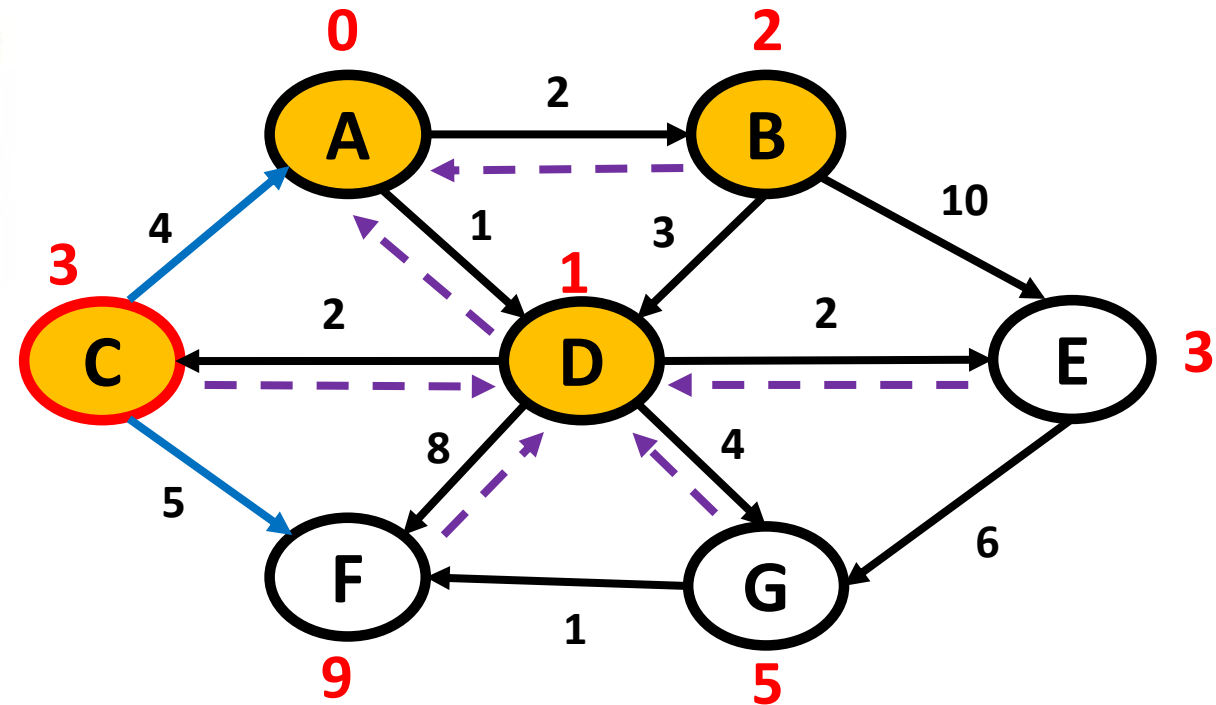
$cost := v$'s cost + weight of edge (v, n) .

if $cost < n$'s cost:

n 's cost := $cost$.

n 's previous := v .

reconstruct path from v_2 back to v_1 , following previous pointers.



Priority Queue: [E: 3, F: 9, G: 5]

Dijkstra's Shortest Path Algorithm

function **dijkstra**(v_1, v_2):

for each vertex v :

// Initialize vertex info

v 's cost := infinity.

v 's previous := none.

v_1 's cost := 0.

$pqueue := \{\text{all vertices, ordered by distance}\}$.

while $pqueue$ is not empty:

$v := \text{remove vertex from } pqueue \text{ with minimum cost.}$ // C

mark v as visited.

for each unvisited neighbor n of v : // F

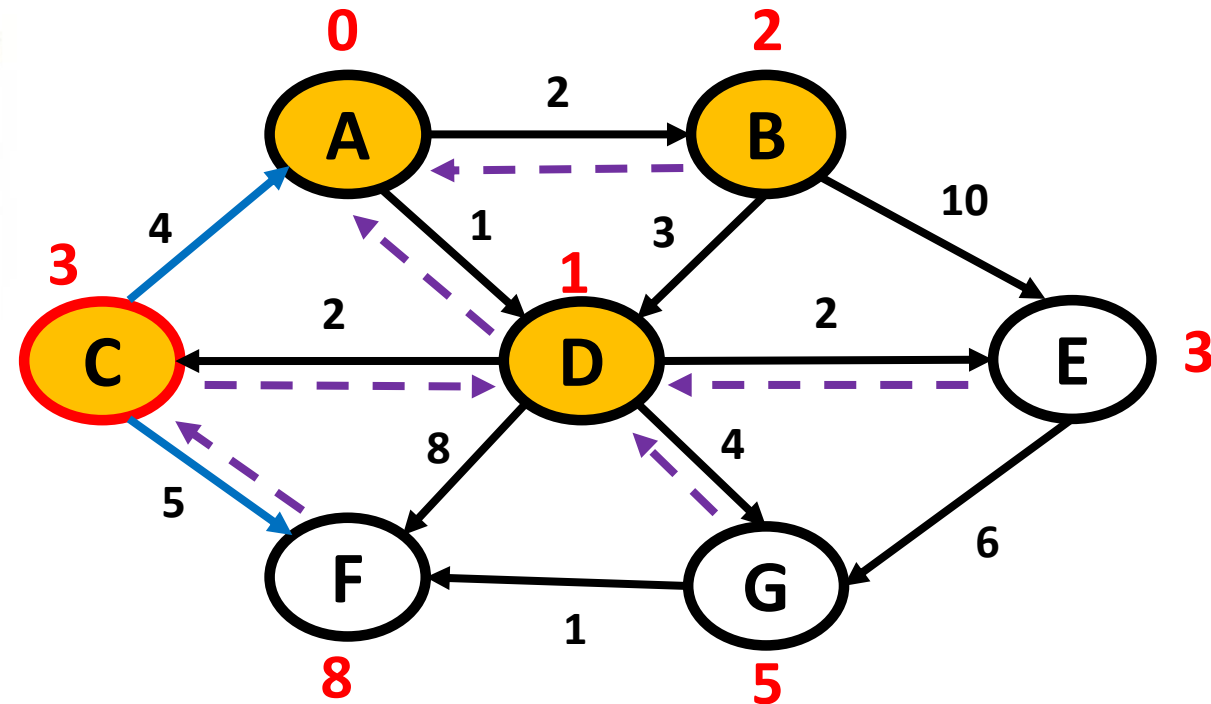
$cost := v$'s cost + weight of edge (v, n) . // $3 + 5$

if $cost < n$'s cost: // $8 < 9$; True

n 's cost := $cost$. // F's cost is updated to 8

n 's previous := v .

reconstruct path from v_2 back to v_1 , following previous pointers.



Priority Queue: [E: 3, F: 8, G: 5]

Dijkstra's Shortest Path Algorithm

function **dijkstra**(v_1, v_2):

for each vertex v :

// Initialize vertex info

v 's cost := infinity.

v 's previous := none.

v_1 's cost := 0.

$pqueue := \{\text{all vertices, ordered by distance}\}$.

while $pqueue$ is not empty:

$v := \text{remove vertex from } pqueue \text{ with minimum cost.}$ // E

mark v as visited.

for each unvisited neighbor n of v : // G

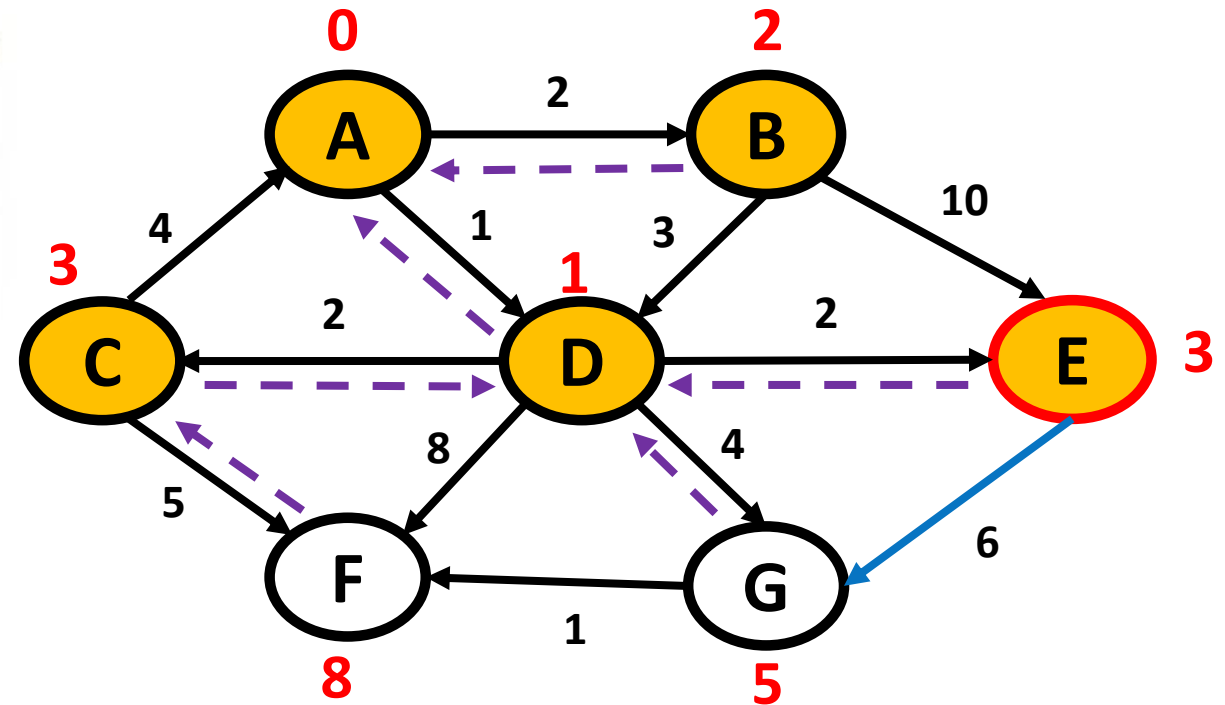
$cost := v$'s cost + weight of edge (v, n) .

if $cost < n$'s cost:

n 's cost := $cost$.

n 's previous := v .

reconstruct path from v_2 back to v_1 , following previous pointers.



Priority Queue: [F: 8, G: 5]

Dijkstra's Shortest Path Algorithm

function **dijkstra**(v_1, v_2):

for each vertex v :

// Initialize vertex info

v 's cost := infinity.

v 's previous := none.

v_1 's cost := 0.

$pqueue := \{\text{all vertices, ordered by distance}\}$.

while $pqueue$ is not empty:

$v := \text{remove vertex from } pqueue \text{ with minimum cost. // E}$

mark v as visited.

for each unvisited neighbor n of v : // G

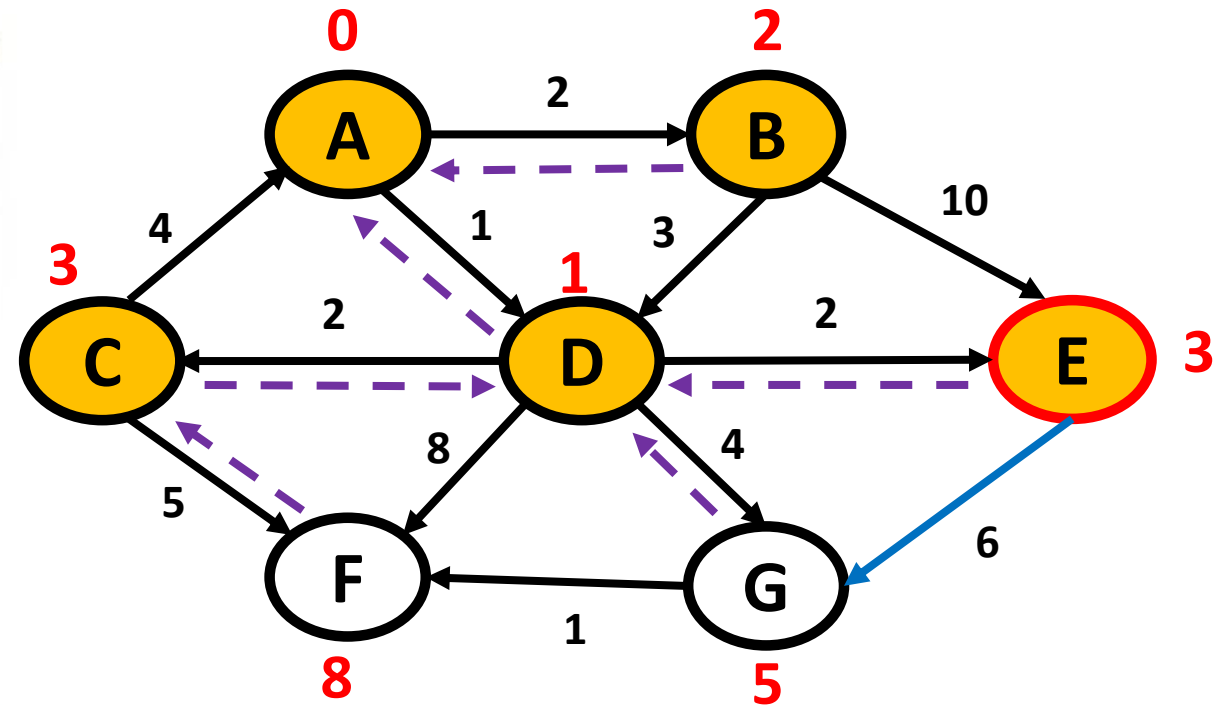
$cost := v$'s cost + weight of edge (v, n) . // $3 + 6$

if $cost < n$'s cost: // $9 > 5$; False

n 's cost := $cost$. // no cost change for G

n 's previous := v .

reconstruct path from v_2 back to v_1 , following previous pointers.



Priority Queue: [F: 8, G: 5]

Dijkstra's Shortest Path Algorithm

function **dijkstra**(v_1, v_2):

for each vertex v :

// Initialize vertex info

v 's cost := infinity.

v 's previous := none.

v_1 's cost := 0.

$pqueue := \{\text{all vertices, ordered by distance}\}$.

while $pqueue$ is not empty:

$v := \text{remove vertex from } pqueue \text{ with minimum cost.}$ // G

mark v as visited.

for each unvisited neighbor n of v : // F

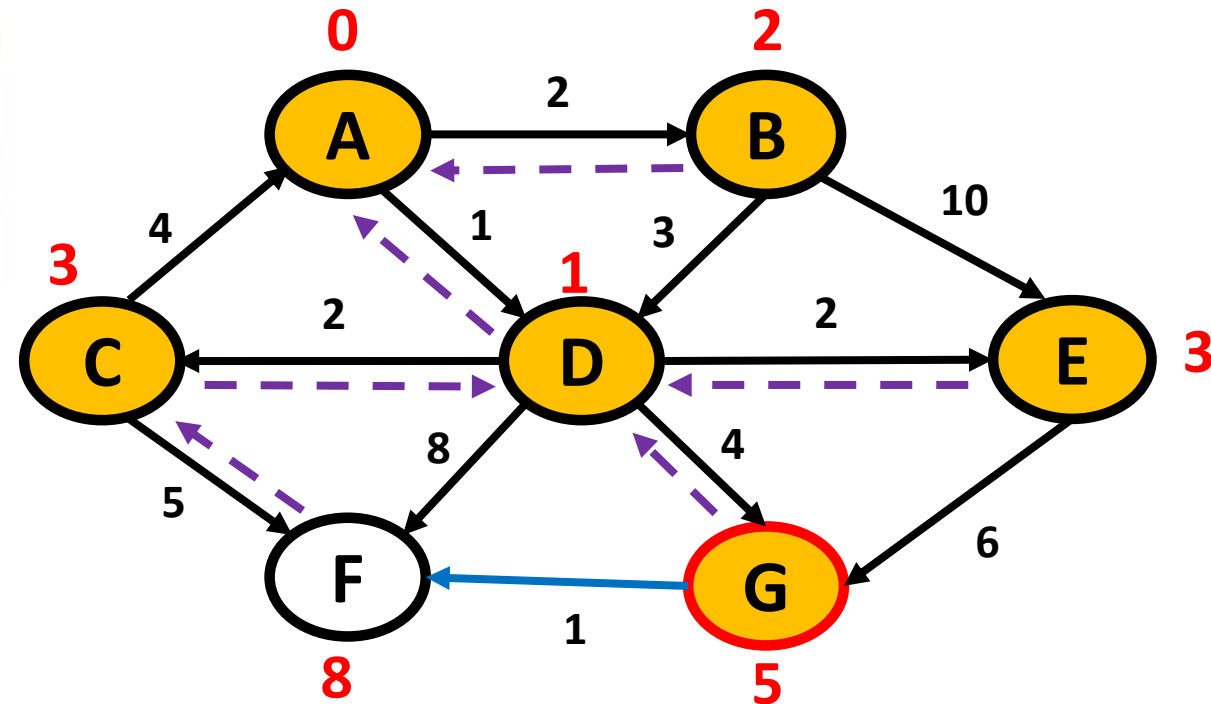
$cost := v$'s cost + weight of edge (v, n) .

if $cost < n$'s cost:

n 's cost := $cost$.

n 's previous := v .

reconstruct path from v_2 back to v_1 , following previous pointers.



Priority Queue: [F: 8]

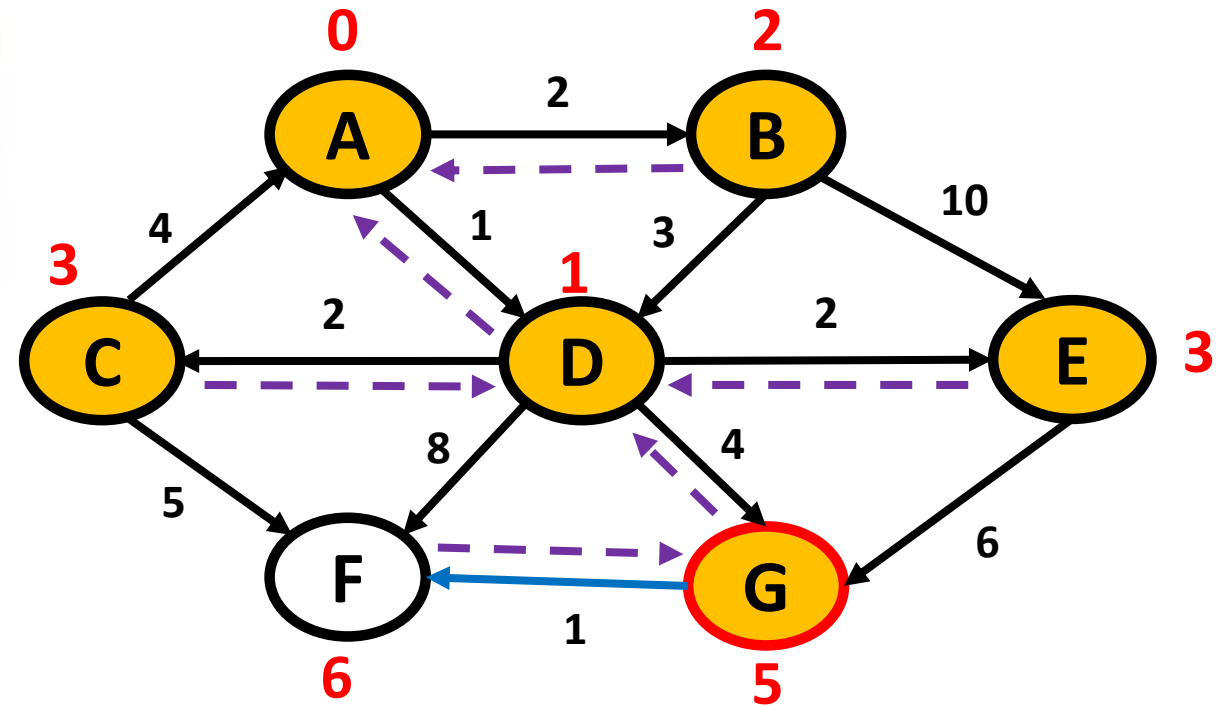
Dijkstra's Shortest Path Algorithm

```
function dijkstra( $v_1, v_2$ ):
  for each vertex  $v$ :           // Initialize vertex info
     $v$ 's cost := infinity.
     $v$ 's previous := none.
   $v_1$ 's cost := 0.
   $pqueue := \{\text{all vertices, ordered by distance}\}$ .

  while  $pqueue$  is not empty:
     $v := \text{remove vertex from } pqueue \text{ with minimum cost.}$  // G
    mark  $v$  as visited.
    for each unvisited neighbor  $n$  of  $v$ : // F
       $cost := v$ 's cost + weight of edge  $(v, n)$ . // 5 + 1

      if  $cost < n$ 's cost: // 6 < 8; True
         $n$ 's cost :=  $cost$ . // F's cost is updated to 6
         $n$ 's previous :=  $v$ .

  reconstruct path from  $v_2$  back to  $v_1$ , following previous pointers.
```



Priority Queue: [F: 6]

Dijkstra's Shortest Path Algorithm

function **dijkstra**(v_1, v_2):

for each vertex v :

// Initialize vertex info

v 's cost := infinity.

v 's previous := none.

v_1 's cost := 0.

$pqueue := \{\text{all vertices, ordered by distance}\}$.

while $pqueue$ is not empty:

$v := \text{remove vertex from } pqueue \text{ with minimum cost.}$ // F

mark v as visited.

for each unvisited neighbor n of v : // none

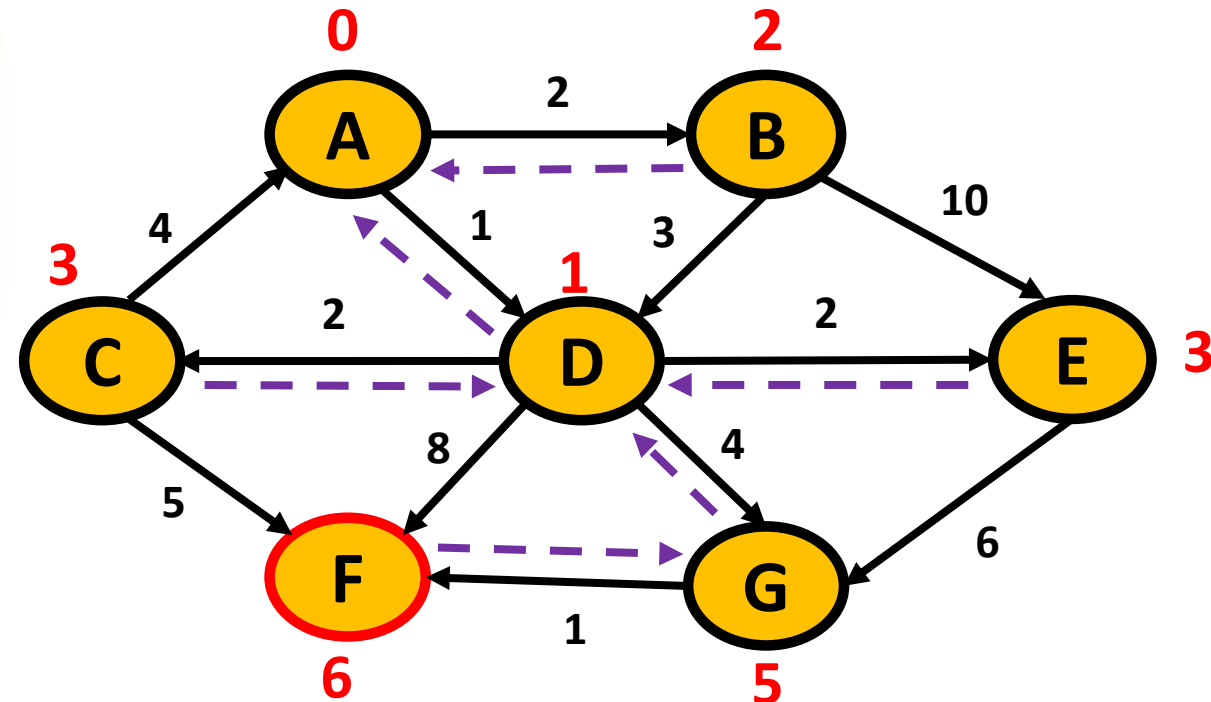
$cost := v$'s cost + weight of edge (v, n) .

if $cost < n$'s cost: // no costs change

n 's cost := $cost$.

n 's previous := v .

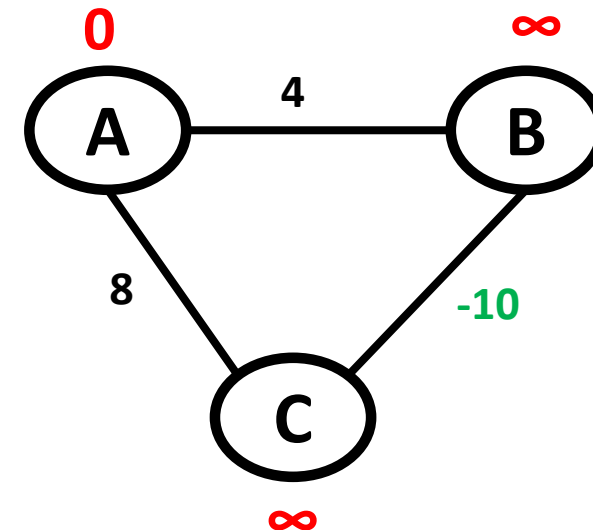
reconstruct path from v_2 back to v_1 , following previous pointers.



Priority Queue: []

Dijkstra's Shortest Path Algorithm

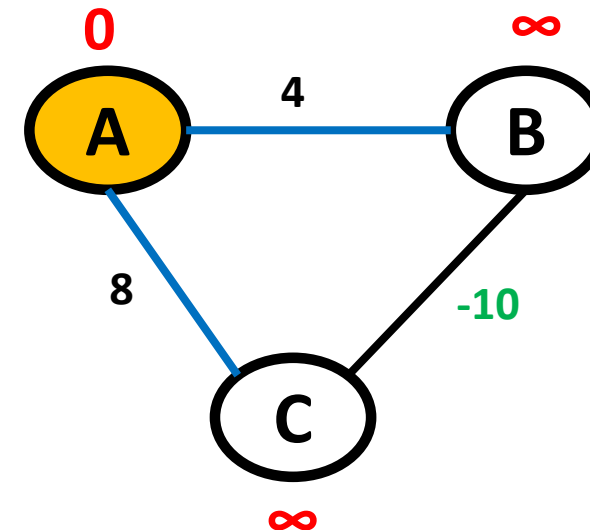
- Does Dijkstra's Shortest Path Algorithm works for **negative edges**?



Priority Queue: [A: 0, B: ∞ , C: ∞]

Dijkstra's Shortest Path Algorithm

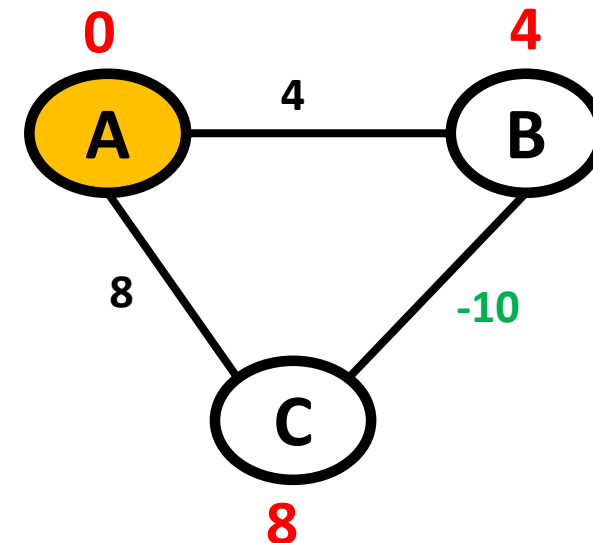
- Does Dijkstra's Shortest Path Algorithm works for **negative edges**?



Priority Queue: [B: ∞ , C: ∞]

Dijkstra's Shortest Path Algorithm

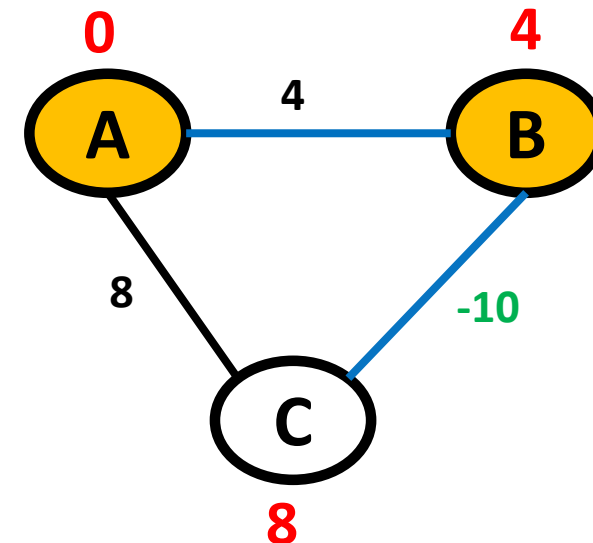
- Does Dijkstra's Shortest Path Algorithm works for **negative edges**?



Priority Queue: [B: 4, C: 8]

Dijkstra's Shortest Path Algorithm

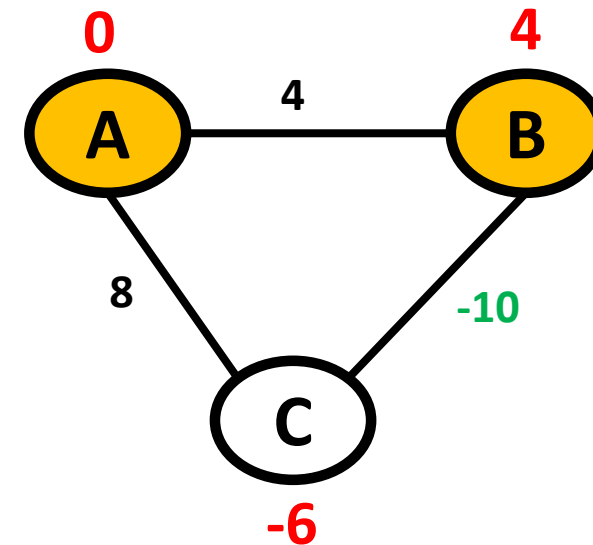
- Does Dijkstra's Shortest Path Algorithm work for **negative edges**?



Priority Queue: [C: 8]

Dijkstra's Shortest Path Algorithm

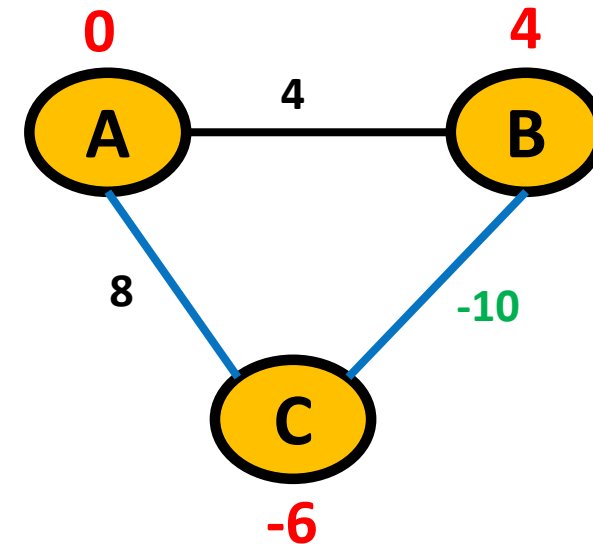
- Does Dijkstra's Shortest Path Algorithm works for **negative edges**?



Priority Queue: [C: -6]

Dijkstra's Shortest Path Algorithm

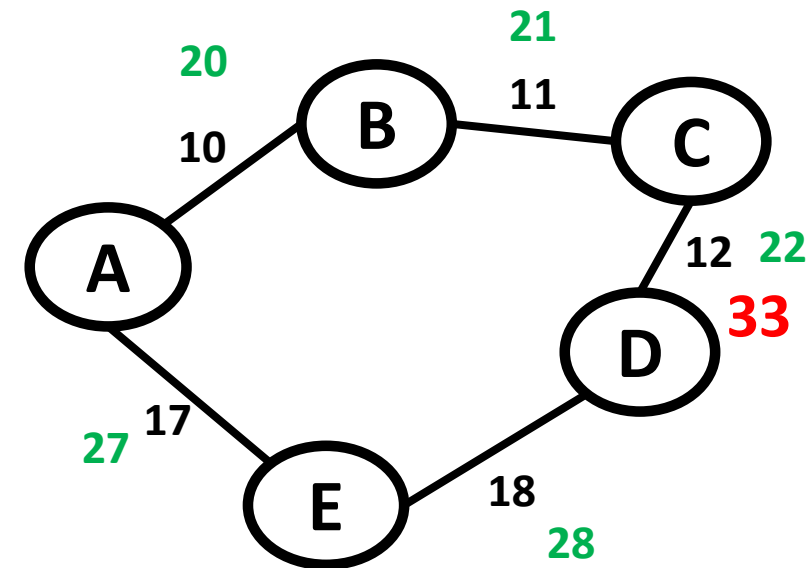
- Does Dijkstra's Shortest Path Algorithm works for **negative edges**?



Priority Queue: []

Dijkstra's Shortest Path Algorithm

- Does the shortest path changes if we **add a weight to all edges** of the original Graph?



Dijkstra's Shortest Path Algorithm

function **dijkstra**(v_1, v_2):

 for each vertex v :

v 's cost := infinity.

v 's previous := none.

v_1 's cost := 0.

$pqueue := \{\text{all vertices, ordered by distance}\}.$ } $V \log V$

Total Complexity = $E \log V + V \log V$

 while $pqueue$ is not empty:

$v := \text{remove vertex from } pqueue \text{ with minimum cost}.$ } $O(\log V)$ } V

 mark v as visited.

 for each unvisited neighbor n of v :

$cost := v$'s cost + weight of edge (v, n) .

 if $cost < n$'s cost:

n 's cost := $cost$.

n 's previous := v .

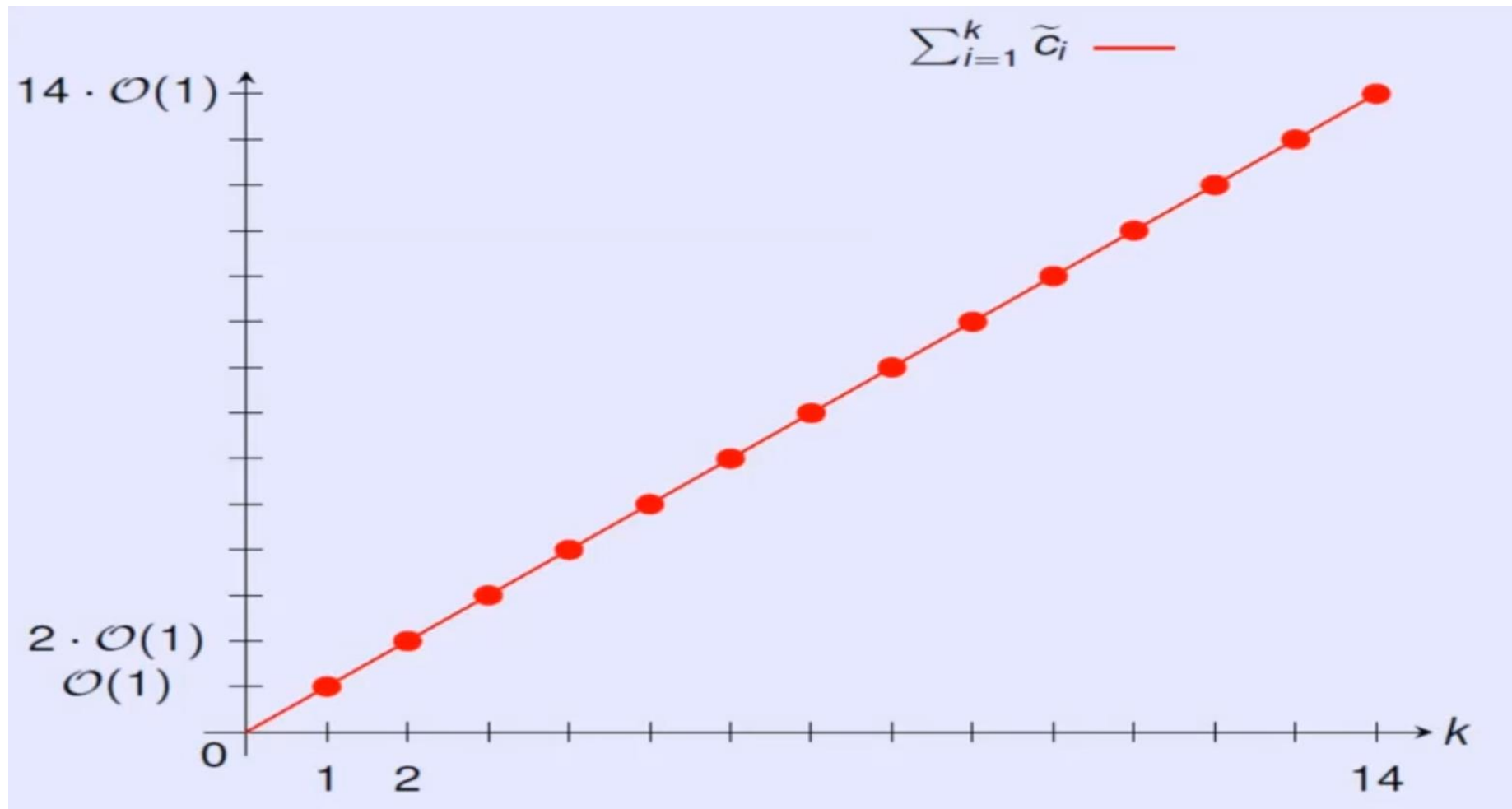
Relax Operation

$O(\log V)$

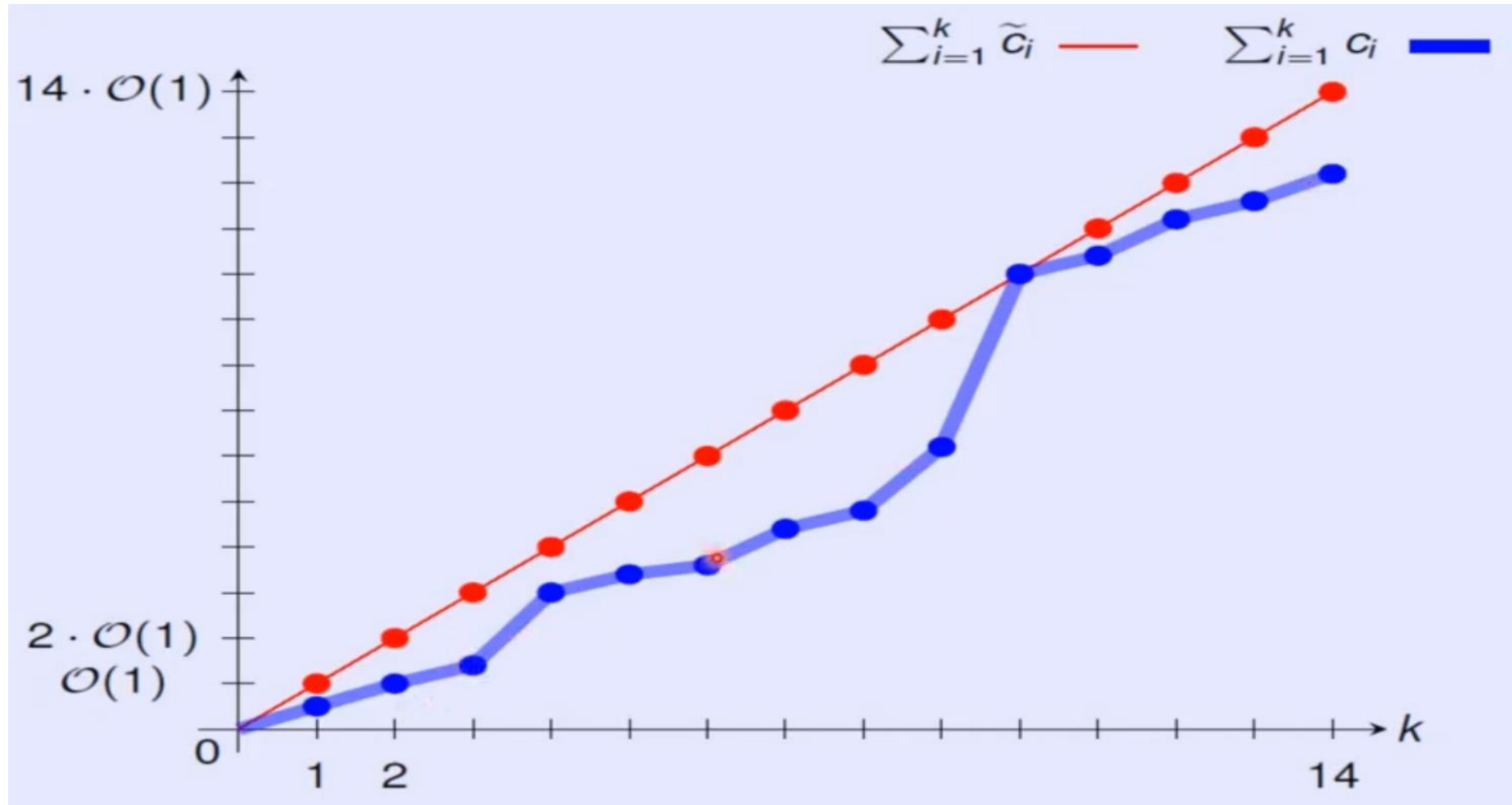
E

 reconstruct path from v_2 back to v_1 , following previous pointers.

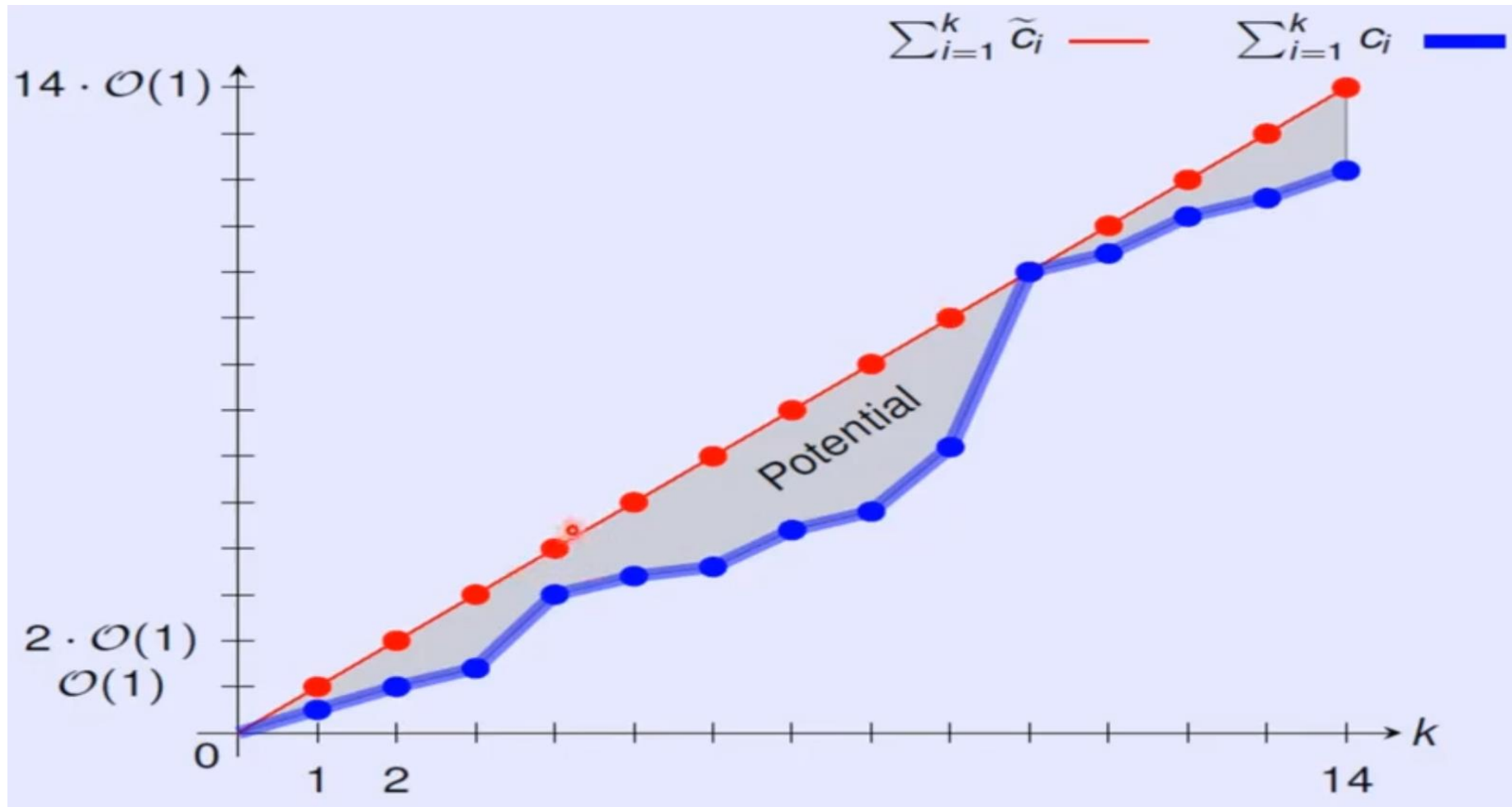
Actual Vs Amortised Cost



Actual Vs Amortised Cost



Actual Vs Amortised Cost



Actual Vs Amortised Cost

| PQ implementation | Insert | Extract - Min | Decrease Key | Total |
|-------------------|----------|---------------|--------------|----------------|
| Unordered array | 1 | V | 1 | V^2 |
| Binary Heap | $\log V$ | $\log V$ | $\log V$ | $E \log V$ |
| Fibonacci Heap | 1^* | $\log V^*$ | 1^* | $E + V \log V$ |