



Tecnológico de Monterrey

Tecnológico de Monterrey - Campus Monterrey
School of Engineering and Sciences
Engineering in Computational Technologies
Analysis and Design of Advanced Algorithms

Class Activity 3: Coin Change through Memoization or Tabulation

Group: 607
Team #3
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Gold Mine - Tabulation

https://colab.research.google.com/drive/1kRMVjIz5v4-58qJouW5aF5BnfX8_xgFW?usp=sharing

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GoldMine_Homework3.py X

DynamicProgramming > GoldMine_Homework3.py > max_coins_tabulation

```
1 # Analysis and Design of Advanced Algorithms
2 # Group #607
3 # Team 3
4 # Dr. Katie Brodhead
5
6 # Santiago Quintana Moreno A01571222
7 # Miguel Ángel Álvarez Hermida A01722925
8
9 # ----- COIN COLLECTING - TABULATION BOTTOM UP -----
10
11
12 # Import List, Tuple, and Optional for type annotations.
13 # List and Tuple are used to specify the types of the grid and path.
14 # Optional is used for the previous cell, which can be None.
15 from typing import List, Tuple, Optional
16
17 def max_coins_tabulation(C: List[List[int]]) -> Tuple[int, List[Tuple[int, int]]]:
18     if not C or not C[0]:
19         return 0, []
20
21     # Get the dimensions of the grid
22     n, m = len(C), len(C[0])
23     # dp[i][j] will store the maximum coins collected to reach cell (i, j)
24     dp = [[0] * m for _ in range(n)]
25     # prev[i][j] will store the previous cell (i, j) in the optimal path to (i, j)
26     prev: List[List[Optional[Tuple[int, int]]]] = [[None] * m for _ in range(n)]
27
28     for i in range(n):
29         for j in range(m):
30             # Initialize the best value and previous cell for dp[i][j]
31             best_val = 0
32             best_prev = None
33
34             # If we can come from the cell above, consider its value
35             if i > 0:
36                 best_val = dp[i - 1][j]
37                 best_prev = (i - 1, j)
38             # If coming from the left cell yields a better value, update
39             if j > 0 and dp[i][j - 1] > best_val:
40                 best_val = dp[i][j - 1]
41                 best_prev = (i, j - 1)
42
```

powershell X

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Max coins: 4

One optimal path: [(0, 1), (1, 1), (1, 2), (1, 3), (2, 3), (2, 4), (3, 4), (4, 4)]

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GoldMine_Homework3.py ×

DynamicProgramming > GoldMine_Homework3.py > max_coins_tabulation

```
17 def max_coins_tabulation(C: List[List[int]]) -> Tuple[int, List[Tuple[int, int]]]:
43     # Store the maximum coins collected to reach (i, j)
44     dp[i][j] = best_val + C[i][j]
45     # Store the previous cell in the optimal path
46     prev[i][j] = best_prev
47
48 # Reconstruct the path from (n-1, m-1) to (0, 0) using the prev array
49 path: List[Tuple[int, int]] = []
50 i, j = n - 1, m - 1
51 while True:
52     path.append((i, j)) # Add current cell to the path
53     if prev[i][j] is None: # If there's no previous cell, we've reached the start
54         break
55     i, j = prev[i][j] # Move to the previous cell in the optimal path
56 path.reverse() # Reverse to get path from start to end
57
58 # Return the maximum coins collected and the reconstructed path
59 return dp[n - 1][m - 1], path
60
61
62 # --- Example run ---
63 if __name__ == "__main__":
64     # Example approximating the slide: 1 means there's a coin in that cell.
65     C = [
66         [0, 0, 0, 0, 1],
67         [0, 1, 0, 1, 0],
68         [0, 0, 0, 1, 0],
69         [0, 0, 1, 0, 0],
70         [1, 0, 0, 0, 1],
71     ]
72     total, path = max_coins_tabulation(C)
73     print("Max coins:", total)
74     print("One optimal path:", path)
75
```

powershell ×

PS D:\1.SQM\1.UNIVERSIDAD\5. QUINTO SEMESTRE\2.Advanced Algorithms> & C:\Users\sant...
ty\AppData\Local\Microsoft\WindowsApps\python3.12.exe "d:/1.SQM/1.UNIVERSIDAD/5. Q
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Why tabulation (Bottom Up) solution?

In Python, this problem is better bottom-up through tabulation for this grid because it avoids recursion overhead and recursion-depth limits on large inputs. Tabulation has the same $O(nm)$ time complexity but uses predictable loops that are more cache-friendly and typically faster in practice. It also makes extras like path reconstruction easy with a predecessor table, and if you only need the value, you can shrink space to $O(m)$ with a sweep. Since memoization offers no asymptotic gain and adds call overhead plus stack risk, tabulation is the safer, simpler choice.

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