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Literature review



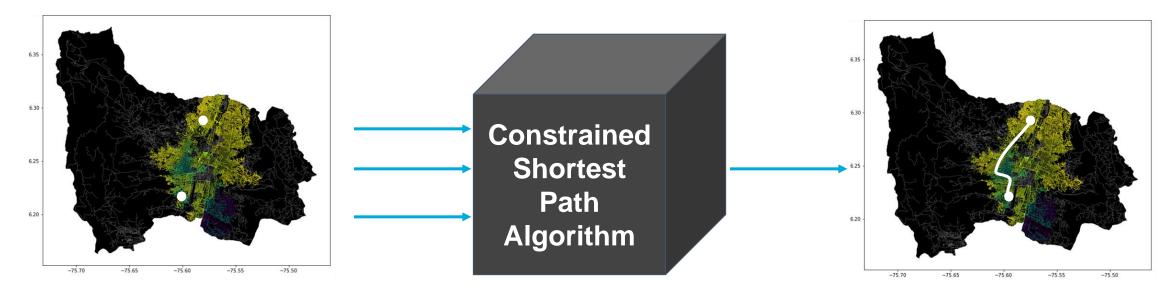
Mauricio
Toro
Data preparation





## **Problem Statement**





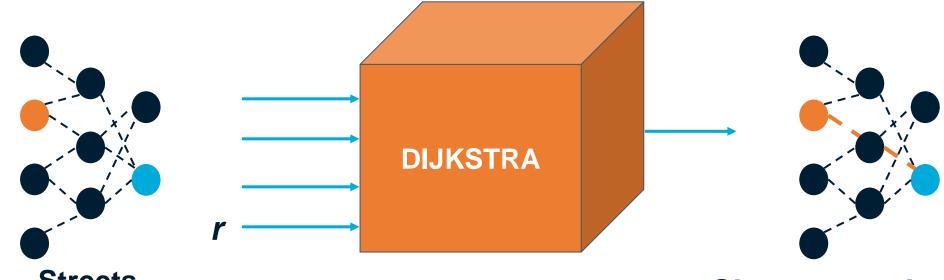
Streets of Medellín, Origin and Destination

Constrained
Shortest
Paths



## **First Algorithm**





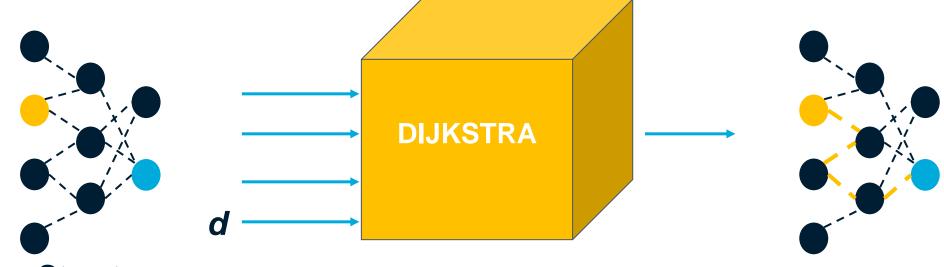
Streets of Medellín, Origin and Destination

Shortest path without exceeding a weighted-average risk of harassment *r* 



## **Second Algorithm**





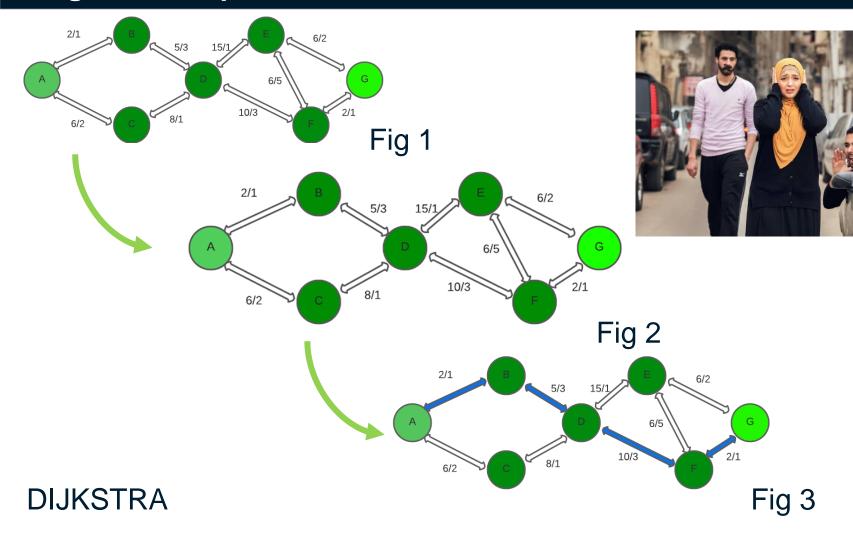
Streets of Medellín, Origin and Destination

Path with the lowest weighted-average risk of harassment without exceeding a distance *d* 



# **Algorithm Explanation**





```
FOR EACH node n in graph
    SET dist[n] TO infinity
   SET prev[n] TO unknown
   ADD n TO Q
ENDFOR
dist[source] = 0
WHILE Q is not empty:
    SET u TO node from Q with minimal dist[u]
   REMOVE u FROM O
   FOR EACH neighbor node n of u
        IF n in Q THEN
            temp = dist[u] + length between u and n
            IF temp < dist[v] THEN</pre>
                dist[v] = temp
                prev[v] = unknown
            ENDIF
       ENDIF
   ENDFOR
ENDWHILE
```

We can observe a graph in fig1 containing two values for each path from one node to another, that of the distance and that of the harassment. In fig2 we UNIVERSIDAD observe the operation of the Dijkstra algorithm, we can see how the algorithm search for each node the shortest path in its adjacent nodes and when it finds 📃 🔼 it, it adds them to the final path. This action is repeated until reaching the goal node. In fig 3 we can see the shortest path defined by the algorithm



## **Algorithm Complexity**



	Time Complexity	Memory Complexity
DIJKSTRA WITH MODIFIED WEIGHTS	O(( V + E ) log  V )	O(V <sup>2</sup> )



On the table are the time and memory complexity of the algorithms where V is the number of nodes, E is the number of edges And the memory complexity it is the one showed there because we used it for every single street (we don't use a binary monticule or a balanced binary tree).



#### **Shortest Path Results**



Origin	Destination	Shortest distance (meters)
Universidad EAFIT	Universidad de Medellín	6142m
Universidad de Antioquia	Universidad Nacional	815 m
Universidad Nacional	Universidad Luis Amigó	1469 m

Shortest distance obtained.



#### **Lowest Risk Results**



Origin	Destination	Weighted-average risk of harassment
Universidad EAFIT	Universidad de Medellín	0.642
Universidad de Antioquia	Universidad Nacional	0.605
Universidad Nacional	Universidad Luis Amigó	0.599

Lowest weighted-average risk of harassment obtained .



#### **Algorithm Execution Times**













13.5 seconds









11.5 seconds









12.49 seconds



#### **Future Work Directions**



#### **Databases**

Add other variables

Data managing

# **Project 1**

A Web
Application

Interface design

# **Software Eng.**

• • • • • • Multiplatform design

Cloud functionality

# **Project 2**

Web-app • development

User management and market launch



