CS 4Z03 - Functional Programming, in Application to Interactive Web Interfaces for Discrete Mathematics Education

Peter Santos

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This is a modified tautology checker that checks if two propositional statements are equal in their evaluation. It comes from the Tautology checker example in section 10.4 of *Programming in Haskell*, by Graham Hutton. This propositional checker no longer evaluates if a proposition a tautology, though two extra lines of code could enable this feature.

To begin, **Data.Set** is imported to allow for a more intuitive way of manipulating the lists that will be used in determining whether a proposition is equivalent to another, which will be discussed later.

```
module PropChecker where
import qualified Data.Set as Set
import Data.Set (Set)
```

A new data type is now defined which represents common propositional logic connectives such as Or, And, Implies, etc.

```
type Var = Char

data Prop = Const Bool

| Var Var

| Not Prop

| And Prop Prop

| Or Prop Prop

| Imply Prop Prop

| Equiv Prop Prop

| Equal Prop Prop

deriving (Show, Eq)
```

Subst will act kind of like a substitution since it doesn't really substitute variables, but rather is type that identifies what Bools should be used for what variables.

```
type Subst = Assoc Char Bool
```

Assoc acts lookup table for variables to bools, although this a general definition which only requires two different types, that may or may not be Char and Bool.

```
\label{eq:type} \begin{split} \mathbf{type} & \ \mathsf{Assoc} \ k \ v = [(k,v)] \\ & \ \mathsf{type} \ \mathsf{Rel} \ k \ v = \mathsf{Set} \ (k,v) \\ & \ \mathsf{--type} \ \mathsf{Fct} \ k \ v = \mathsf{Map.Map} \ k \ v \end{split}
```

find will take a variable(k) and table of variables to bools and find the first bool value for k and return it.

```
find :: Eq k \Rightarrow k \rightarrow Assoc k v \rightarrow v find k t = head [v | (k', v) \leftarrow t, k \equiv k']
```

Now that we have defined a way to associate bool values to variables, we can now use this in another function that recursively evaluates a proposition of type Prop. Using the find function it will find the appropriate bool value for a variable in a proposition and then evaluate the proposition once it has been constructed.

```
\begin{array}{lll} \text{eval} & :: \text{Subst} \rightarrow \text{Prop} \rightarrow \text{Bool} \\ \text{eval} \ \_ (\text{Const b}) & = b \\ \text{eval s } (\text{Var x}) & = \text{find x s} \\ \text{eval s } (\text{Not p}) & = \neg (\text{eval s p}) \\ \text{eval s } (\text{And p q}) & = \text{eval s p} \wedge \text{eval s q} \\ \text{eval s } (\text{Or p q}) & = \text{eval s p} \vee \text{eval s q} \\ \text{eval s } (\text{Imply p q}) & = \text{eval s p} \in \text{eval s q} \\ \text{eval s } (\text{Equiv p q}) & = \text{eval s (Imply p q}) \wedge \text{eval s (Imply q p)} \\ \text{eval s } (\text{Equal p q}) & = \text{eval s p} \equiv \text{eval s q} \\ \end{array}
```

We will use 'vars' to construct a list of variables, that are in a proposition. This function will produce duplicate variables, but that's okay for now, because we will remove them later.

```
\begin{array}{lll} \text{vars} & \text{:: Prop} \rightarrow [\text{Char}] \\ \text{vars} & (\text{Const}\_) & = [] \\ \text{vars} & (\text{Var } x) & = [x] \\ \text{vars} & (\text{Not } p) & = \text{vars } p \\ \text{vars} & (\text{And } p \ q) & = \text{vars } p + \text{vars } q \\ \text{vars} & (\text{Or } p \ q) & = \text{vars } p + \text{vars } q \\ \text{vars} & (\text{Equiv } p \ q) & = \text{vars } p + \text{vars } q \\ \text{vars} & (\text{Equal } p \ q) & = \text{vars } p + \text{vars } q \\ \text{vars} & (\text{Equal } p \ q) & = \text{vars } p + \text{vars } q \\ \text{vars} & (\text{Equal } p \ q) & = \text{vars } p + \text{vars } q \\ \text{vars} & (\text{Equal } p \ q) & = \text{vars } p + \text{vars } q \\ \text{vars} & (\text{Equal } p \ q) & = \text{vars } p + \text{vars } q \\ \text{vars} & (\text{Equal } p \ q) & = \text{vars } p + \text{vars } q \\ \text{vars} & (\text{Equal } p \ q) & = \text{vars } p + \text{vars } q \\ \text{vars} & (\text{Equal } p \ q) & = \text{vars } p + \text{vars } q \\ \text{vars} & (\text{Equal } p \ q) & = \text{vars } p + \text{vars } q \\ \text{vars} & (\text{Equal } p \ q) & = \text{vars } p + \text{vars } q \\ \text{vars} & (\text{Equal } p \ q) & = \text{vars } p + \text{vars } q \\ \text{vars} & (\text{Equal } p \ q) & = \text{vars } p + \text{vars } q \\ \text{vars} & (\text{Equal } p \ q) & = \text{vars } p + \text{vars } q \\ \text{vars} & (\text{Equal } p \ q) & = \text{vars } p + \text{vars } q \\ \text{vars} & (\text{Equal } p \ q) & = \text{vars } p + \text{vars } q \\ \text{vars} & (\text{Equal } p \ q) & = \text{vars } p + \text{vars } q \\ \text{vars} & (\text{Equal } p \ q) & = \text{vars } p + \text{vars } q \\ \text{vars} & (\text{Equal } p \ q) & = \text{vars } p + \text{vars } q \\ \text{vars} & (\text{Equal } p \ q) & = \text{vars } p + \text{vars } q \\ \text{vars} & (\text{Equal } p \ q) & = \text{vars } p + \text{vars } q \\ \text{vars} & (\text{Equal } p \ q) & = \text{vars } p + \text{vars } q \\ \text{vars} & (\text{Equal } p \ q) & = \text{vars } p + \text{vars } q \\ \text{vars} & (\text{Equal } p \ q) & = \text{vars } p + \text{vars } q \\ \text{vars} & (\text{Equal } p \ q) & = \text{vars } p + \text{vars } q \\ \text{vars} & (\text{Equal } p \ q) & = \text{vars } p + \text{vars } q \\ \text{vars} & (\text{Equal } p \ q) & = \text{vars } p + \text{vars } q \\ \text{vars} & (\text{Equal } p \ q) & = \text{vars } p + \text{vars } q \\ \text{vars} & (\text{Equal } p \ q) & = \text{vars } p + \text{vars } q \\ \text{vars} & (\text{Equal } p \ q) & = \text{vars } p + \text{vars } q \\ \text{vars} & (\text{Equal } p \ q) & = \text{vars } p + \text{va
```

The 'bools' function creates a complete truth table of all possible True and False combinations for a given number of variables.

```
\begin{array}{ll} \mathsf{bools} & :: \mathsf{Int} \to [[\mathsf{Bool}]] \\ \mathsf{bools} \ 0 & = [[]] \\ \mathsf{bools} \ (\mathsf{n}+1) = \mathsf{map} \ (\mathsf{False:}) \ \mathsf{bss} + \mathsf{map} \ (\mathsf{True:}) \ \mathsf{bss} \\ & \mathbf{where} \ \mathsf{bss} = \mathsf{bools} \ \mathsf{n} \\ \end{array}
```

This function filters out all the duplicates of a list, and is used to remove the duplicate variables from the list that is generated from 'vars'.

```
rmdups :: Eq a \Rightarrow [a] \rightarrow [a]

rmdups [] = []

rmdups (x : xs) = x : rmdups (filter (\not\equiv x) xs)
```

The function 'substs' will pair the variables produced from 'vars' with the list of bools generated from 'bools'.

```
substs :: Prop \rightarrow [Subst]
substs p = map (zip vs) (bools (length vs))
               where vs = rmdups (vars p)
type Rests
                                   = [Prop]
  -- This function will filter out all the props that don't satisfy the given
cleanSubst
                                   :: [\mathsf{Subst}] \to \mathsf{Prop} \to [\mathsf{Subst}]
cleanSubst [] p
cleanSubst (s : subs) p
                                   = if eval s p then
                                        s : cleanSubst subs p
                                      else
                                        cleanSubst subs p
  -- eval s p = flip eval p s = s 'eval' p = ('eval'p) s = (\lambda s \rightarrow \text{eval s p}) s
  -- flip cleanSubst p = filter (flip eval p)
  -- This will purge the substitution list of any substitutions that do not satisfy
  -- the given restrictions. ie, A \neq B
```

```
readySubst
                                   :: [\mathsf{Subst}] \to \mathsf{Rests} \to [\mathsf{Subst}]
readySubst subs []
                                    = subs
readySubst subs (p : ps)
                                    = readySubst (cleanSubst subs p) ps
  -- This will create the one part of the final substitution list
keepSubst :: [Subst] \rightarrow Set.Set (Char, Bool) \rightarrow [(Subst, Subst)]
keepSubst subs set1
                 = [(s, Set.toList set1) \mid s \leftarrow subs, set1 'Set.isSubsetOf' Set.fromList s]
  -- This will create the final substitution list. The first arguments should
  -- take a larger clean substitution list and a smaller substitution list
  -- if the last substitution list is of equal size, then is Equiv should be
  -- used instead
finalSubst
                                   :: [\mathsf{Subst}] \to [\mathsf{Subst}] \to [(\mathsf{Subst}, \mathsf{Subst})]
finalSubst subs1 subs2
                       = \mathsf{foldI} \; (\texttt{++}) \; [\; ] \; [\mathsf{keepSubst} \; \mathsf{subs1} \; (\mathsf{Set.fromList} \; \mathsf{s}) \; | \; \mathsf{s} \leftarrow \mathsf{subs2}]
  -- This function works with two propositions that have the same amount of
  -- variables. Boring, I know.
isEquiv
                                    :: \mathsf{Prop} \to \mathsf{Prop} \to \mathsf{Rests} \to \mathsf{Bool}
isEquiv p1 p2 r
                                    = if p1 \equiv p2 then
                                         True
                                      else
                                         and [eval s p1 \equiv eval s p2 | s \leftarrow subs]
                                            where subs = readySubst (substs p1) r
  -- This function works with two propositions that have an unequal amount of
  -- variables. The substitution list always has the propositions with more
  -- variables p1.
                                   :: \mathsf{Prop} \to \mathsf{Prop} \to [(\mathsf{Subst}, \mathsf{Subst})] \to \mathsf{Bool}
isEquiv'
isEquiv' p1 p2 subs
                                    = and [eval s1 p1 \equiv eval s2 p2 | (s1, s2) \leftarrow subs]
  -- Give propMachine two propositions and it's restrictions and it will tell you
  -- If they are equivalent
propMachine
                                   :: \mathsf{Prop} \to \mathsf{Prop} \to \mathsf{Rests} \to \mathsf{Bool}
propMachine p1 p2 r
                                    = if varLp1 > varLp2 then
                                         isEquiv' p1 p2 (finalSubst (cSub1) (substs p2))
                                      else if varLp2 > varLp1 then
                                         isEquiv' p2 p1 (finalSubst (cSub2) (substs p1))
                                      else
                                         isEquiv p1 p2 r
                                         where cSub1 = readySubst (substs p1) r
                                            cSub2 = readySubst (substs p2) r
                                            varLp1 = length (rmdups (vars p1))
                                            varLp2 = length (rmdups (vars p2))
  -- This will give the first instance of when two propositions disagree
  -- and format it into a string using disagreeM.
disagree
                       :: \mathsf{Prop} \to \mathsf{Prop} \to \mathsf{Rests} \to \mathsf{String}
disagree p1 p2 r
                       = if varLp1 > varLp2 then
                                   disagreeM (fstFalsePair p1 p2 (finalSubst (cSub1) (substs p2)))
                          else if varLp2 > varLp1 then
                                   disagreeM (fstFalsePair p2 p1 (finalSubst (cSub2) (substs p1)))
                          else
                                   disagreeM (fstFalsePair p1 p2 (finalSubst cSub1 cSub2))
                                   where cSub1 = readySubst (substs p1) r
                                      cSub2 = readySubst (substs p2) r
                                      varLp1 = length (rmdups (vars p1))
                                     varLp2 = length (rmdups (vars p2))
  -- Takes a substitution and translates to a string.
disagreeM
                       :: Maybe Subst \rightarrow String
```

```
disagreeM (Just s) = disagreeM' s
disagreeM Nothing = "your proposition is evaluated, it violates the given " ++
                                                     "restrictions. "
disagreeM'
                                                     :: Subst \rightarrow String
disagreeM'((v,b):[]) = v:" = " + (show b) + ", your proposition doesn't " + (show b) + ", your proposition doesn't " + (show b) + ", your proposition doesn't " + (show b) + ", your proposition doesn't " + (show b) + ", your proposition doesn't " + (show b) + ", your proposition doesn't " + (show b) + ", your proposition doesn't " + (show b) + ", your proposition doesn't " + (show b) + ", your proposition doesn't " + (show b) + ", your proposition doesn't " + (show b) + ", your proposition doesn't " + (show b) + ", your proposition doesn't " + (show b) + ", your proposition doesn't " + (show b) + 
                                                                  "correctly describe the situation. "
disagreeM'((v,b):subs) = v:" = " + (show b) + ", " + disagreeM' subs
fstFalse
                                                     :: Maybe (Subst, Subst) \rightarrow Maybe Subst
fstFalse subPair
                                                                 case subPair of
                                                                        Just (s1, s2) \rightarrow Just (rmdups (s1 + s2))
                                                                        Nothing → Nothing
fstFalsePair
                                               :: \mathsf{Prop} \to \mathsf{Prop} \to [(\mathsf{Subst}, \mathsf{Subst})] \to \mathsf{Maybe} \, \mathsf{Subst}
fstFalsePair p1 p2 s = case (safeHead (takeWhile (\lambda(s1,s2) \rightarrow eval\ s1\ p1 \not\equiv eval\ s2\ p2)\ s)) of
                                                           Just subPair → fstFalse (Just subPair)
                                                           Nothing \rightarrow fstFalse Nothing
safeHead (h: _{-}) = Just h
safeHead _ = Nothing
      -- Test propositions
p1 :: Prop
p1 = And (Var 'A') (Not (Var 'A'))
p2 :: Prop
p2 = Imply (And (Var 'A') (Var 'B')) (Var 'A')
p3 :: Prop
p3 = Imply (Var 'A') (And (Var 'A') (Var 'B'))
p4 :: Prop
p4 = Imply (And (Var 'A') (Imply (Var 'A') (Var 'B'))) (Var 'B')
p5 :: Prop
p5 = Or(Not(Var'A'))(Var'A')
p6 :: Prop
\mathsf{p6} = \; \; \mathsf{Equiv} \; (\mathsf{Imply} \; (\mathsf{Var} \; \mathsf{`A'}) \; (\mathsf{Var} \; \mathsf{`B'})) \; (\mathsf{Imply} \; (\mathsf{Not} \; (\mathsf{Var} \; \mathsf{`B'})) \; (\mathsf{Not} \; (\mathsf{Var} \; \mathsf{`A'})))
p7 :: Prop
p7 = Equiv (Not (And (Var 'A') (Var 'B'))) (Or (Not (Var 'A')) (Not (Var 'B')))
p8 :: Prop
p8 = Imply (Var 'A') (Or (Var 'B') (Var 'C'))
p9 :: Prop
p9 = \  \, \mathsf{Or}\,\,(\mathsf{Imply}\,\,(\mathsf{Var}\,\,{}^{\backprime}\!\mathtt{A}^{\backprime})\,\,(\mathsf{Var}\,\,{}^{\backprime}\!\mathtt{B}^{\backprime}))\,\,(\mathsf{Imply}\,\,(\mathsf{Var}\,\,{}^{\backprime}\!\mathtt{A}^{\backprime})\,\,(\mathsf{Var}\,\,{}^{\backprime}\!\mathtt{C}^{\backprime}))
      -- Ladies or Tigers example
p10 :: Prop
\mathsf{p10} = \mathsf{Equiv} \; (\mathsf{Var} \; \texttt{`a'}) \; (\mathsf{Var} \; \texttt{`d'})
p11 :: Prop
p11 = And (Or (Var 'a') (Var 'b')) (Or (Var 'c') (Var 'd'))
r1 :: Prop
r1 = Not (Equal (Var 'a') (Var 'c'))
r2 :: Prop
r2 = Not (Equal (Var 'b') (Var 'd'))
s1 :: Prop
s1 = And (Var 'a') (Var 'd')
s2 :: Prop
s2 = p11
s1s2 :: Prop
```

```
s1s2 = Not (Equiv (s1) (s2))
s3 :: Prop
s3 = And (Var 'b') (Var 'c')
```

1 PropParser

This is my propostional parser using Parsec (parsec2 package)

```
module PropParser where
import Text.ParserCombinators.Parsec -- :set -ignore-package parsec-3.1.0
import Text.ParserCombinators.Parsec.Expr
import qualified Text.ParserCombinators.Parsec.Token as T
import Text.ParserCombinators.Parsec.Language (haskellDef)
import Text.ParserCombinators.Parsec.Char
import PropChecker
import Char
  -- The Parser
mainParser = do whiteSpace
            \mathsf{e} \leftarrow \mathsf{expr}
            eof
            return e
       = buildExpressionParser table term
        <? > "expression"
term0 = parens expr
        <|>var
        <?> "simple proposition"
term = do
       t \leftarrow term0
       whiteSpace
       return t
table :: OperatorTable Char () Prop
       = [[prefix "~" Not]
           [binary "&" And AssocLeft, binary "v" Or AssocLeft]
           [binary "=>" Imply AssocLeft, binary "<=>" Equiv AssocNone]
binary name fun assoc = Infix (\mathbf{do} {reservedOp name; whiteSpace; return fun}) assoc
prefix name fun = Prefix (do \{reservedOp name; whiteSpace; return fun \})
postfix name fun = Postfix (do { reservedOp name; whiteSpace; return fun })
isVar
                    :: \mathsf{Char} \to \mathsf{Bool}
isVar c
                    = isAlpha c \wedge c \not\equiv 'v'
                    :: Parser Prop
var
                    = fmap Var $ satisfy isVar
var
evalProp
                  :: \mathsf{String} \to \mathsf{String} \to \mathsf{Rests} \to \mathsf{String}
evalProp x1 x2 rs
             = case (parse mainParser "" x1) of
                   Left err1 \rightarrow show err1
                   Right p1 \rightarrow case (parse mainParser "" x2) of
                            Left err2 \rightarrow show err2
                            Right p2 → show (propMachine p1 p2 rs)
  -- This will convert the first string to a Prop and allow disagree to work
evalDisagree
                    :: \mathsf{String} \to \mathsf{String} \to \mathsf{Rests} \to \mathsf{String}
```

```
evalDisagree x1 x2 rs = case (parse mainParser "" x1) of
                                        Left err1 \rightarrow show err1
                                        Right p1 \rightarrow case (parse mainParser "" x2) of
                                           Left\ err2 \rightarrow show\ err2
                                           Right\ p2 \to disagree\ p1\ p2\ rs
   -- Restriction parser
getRests
                           :: [\mathsf{String}] \to \mathsf{Rests} \to \mathsf{Either} \; \mathsf{String} \; \mathsf{Rests}
getRests [] ps
                           = Right ps
\mathsf{getRests}\;(\mathsf{r}\,\mathsf{:}\,\mathsf{rs})\;\mathsf{ps} = \mathbf{case}\;(\mathsf{parse}\;\mathsf{mainParser}\;\mathsf{"}\,\mathsf{"}\;\mathsf{r})\;\mathbf{of}
                                        Left err \rightarrow Left "hello"
                                        \mathsf{Right}\; \mathsf{p} \to \mathsf{Right}\; [\mathsf{p}]
removeEmpty xs = filter (\not\equiv "") xs
   -- The lexer
lexer
                  = {\sf T.makeTokenParser} \; {\sf haskelIDef} \;
                 = T.lexeme
lexeme
                 = T.parens lexer
parens
natural
                 = T.natural lexer
\mathsf{reservedOp} = \mathsf{T}.\mathsf{reservedOp} \ \mathsf{lexer}
white {\sf Space} = {\sf T.whiteSpace} \ {\sf lexer}
```